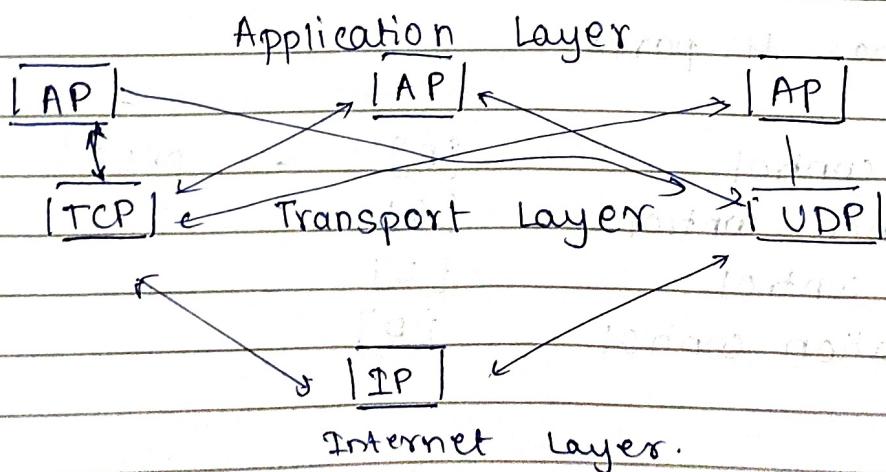


## Module 3 - The Transport Layer.

- The transport layer is the 4th layer within a computer network.
- It provides a variety of services but the main role of the transport layer is to provide communication services to application processes running on different hosts.
- The transport layer provides a logical connection between application processes from different hosts since none of them are physically connected to each other.
- The transport layer has two protocols, - TCP and UDP.
- Each of the applications in the application layer has the ability to send a message using either TCP or UDP. So the applications can read and write to the transport layer. Hence we say this communication is a two-way process.



## \* Services of Transport Layer.

Services provided by transport layers are divided into 5 categories:-

- 1] End - to - end delivery [E]
- 2] Addressing [A] [ADRFAE]
- 3] Reliable delivery [R]
- 4] Flow control [F]
- 5] Multiplexing [D]

### \* End - to - end delivery.

The transport layer transmits the entire message from source to destination. Providing end-to-end delivery of the message.

### \* Reliable delivery.

Retransmission of lost and damaged packets is providing reliability of message delivery.

It has 4 parts:

- 1] Error control [E] → 'DBLS'
- 2] Sequence control [S]
- 3] Loss control [L]
- 4] Duplication control [D]

## \* Error control.

- No transmission is error-free. So the transport layer protocols are designed to provide error-free transmission.
- The error-free experience is the main role of the data link layer but it only ensures node-to-node error correction. It doesn't ensure end-to-end error-free reliance.
- If there is an error in the router, the data-link layer cannot detect it; it can only detect errors in the beginning and at the end of the link. But the transport layer makes sure the message arrives correctly.

## \* Sequence control.

- This aspect makes sure that the packets received from the upper layers are accessible to the lower layers.
- Also the segments can be re-arranged according to the priority provided by the user.

### \* Loss control.

- Makes sure all the messages sent from the source are received at the destination, not some of them. Every single message is given a sequence number, allowing to identify any missing segment.

### \* Duplication control.

- Guarantees that data being received at the destination is not duplicated preventing wastage of storage space. Sequence numbering can also identify lost packets.

### \* Flow control.

- Used to prevent the sender from overwhelming the receiver.
- If the receiver is overloaded with packets, then it discards the packets which need retransmission.
- It uses the sliding window protocol that makes data transmission more efficient as well as controls the flow of data.

## \* Multiplexing

→ Can be of two types :-

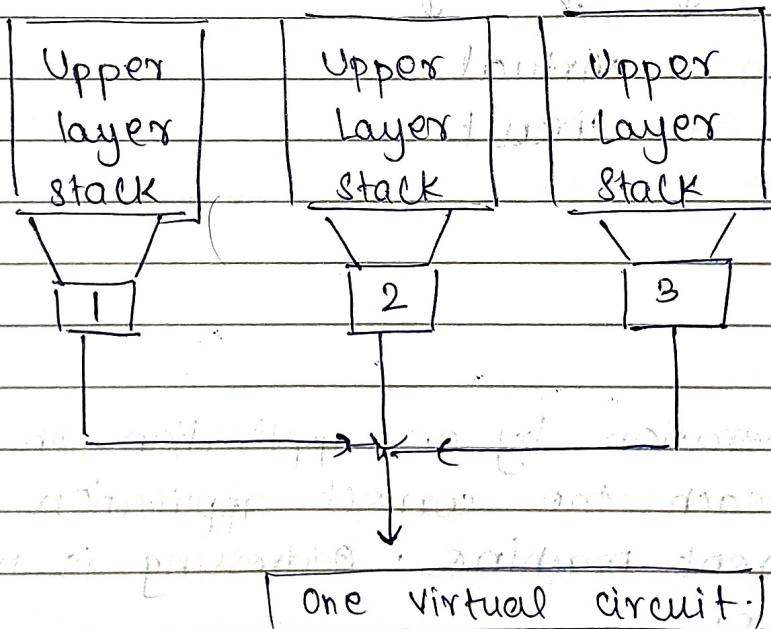
1] Upward M.P

2] Downward M.P.

→ Upward Multiplexing

This means multiple transport layer connections use the same network connection.

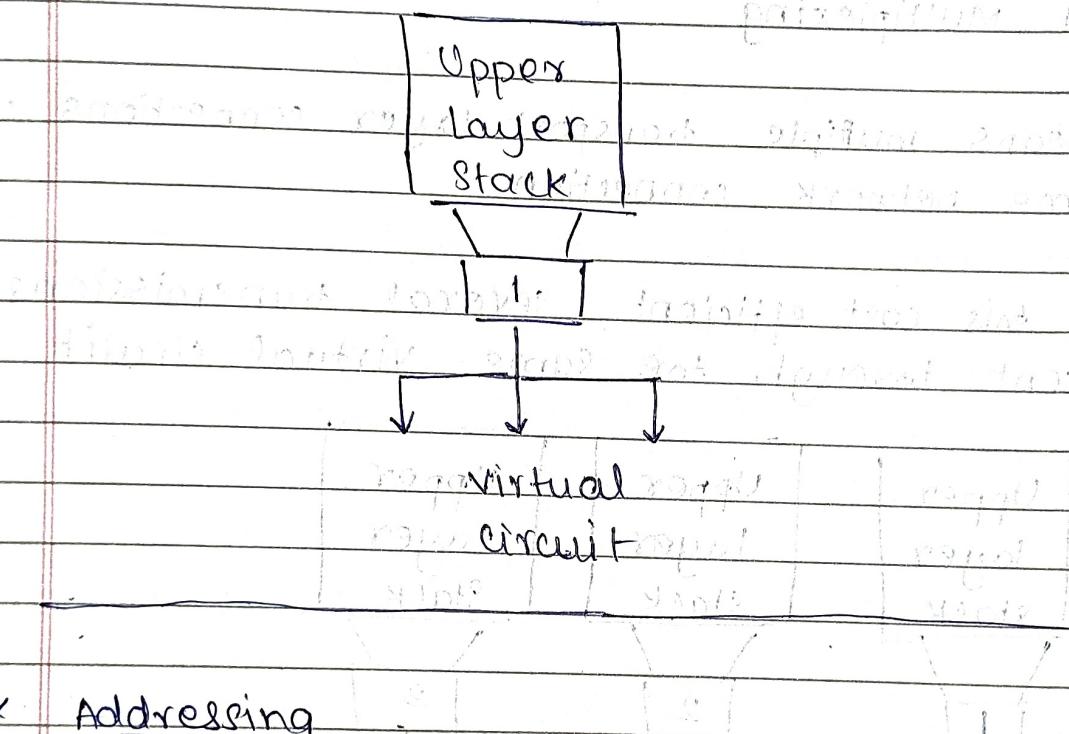
Making this cost efficient, several transmissions are sent through the same virtual circuit



## → Downward Multiplexing:

When one transport layer connection uses multiple network connections.

This allows to split connection between several paths to improve accuracy and throughput.



## \* Addressing

- Data generation by an application on one machine should reach the correct application on same or different machine. Addressing is helpful in these cases.
- The TL provides the user address which is specified as the station or port.

This port variable represents a TS user or a station called Transport service access point (TSAP).

- \* Transport layer protocols.
- \* UDP

→ UDP (User Data gram Protocol) is a simple protocol providing non-sequential transmission functionality.

→ UDP is a connectionless protocol.

→ UDP focuses on speed and size efficiency rather than reliability and security.

→ This type of protocol adds transport-level address, checksum error control and length info to the data from the upper layer.

- \* format of UDP packets.

Source port address (16-bits)	Dest port address (16-bits)
Total length. (16-bits)	Checksum (16-bits)
Data.	

In this diagram we see:

- 1) Source port address - Defines the address of application process that has delivered the message.
- 2) Destination port address - It defines the address of the application process which will receive the message.
- 3) Total length - defines the total length of the user datagram (16-bit)
- 4) checksum - Used in error detection.

#### \* Applications of UDP:

- 1) Lossless Data transmission.
- 2) Gaming, voice calls and video calls.
- 3) Multicasting and routing updates.
- 4) Fast applications.

### \* Disadvantages of UDP:

1] Doesn't provide sequencing or re-ordering of user datagrams.

2] Doesn't give specifications about which data packet is lost during loss control. There are no sequence numbers here.

### \* TCP (Transmission Control Protocol)

→ Tcp stands for Transmission Control Protocol.

→ It is a connection oriented protocol which means a connection is established between both ends of the transmission.

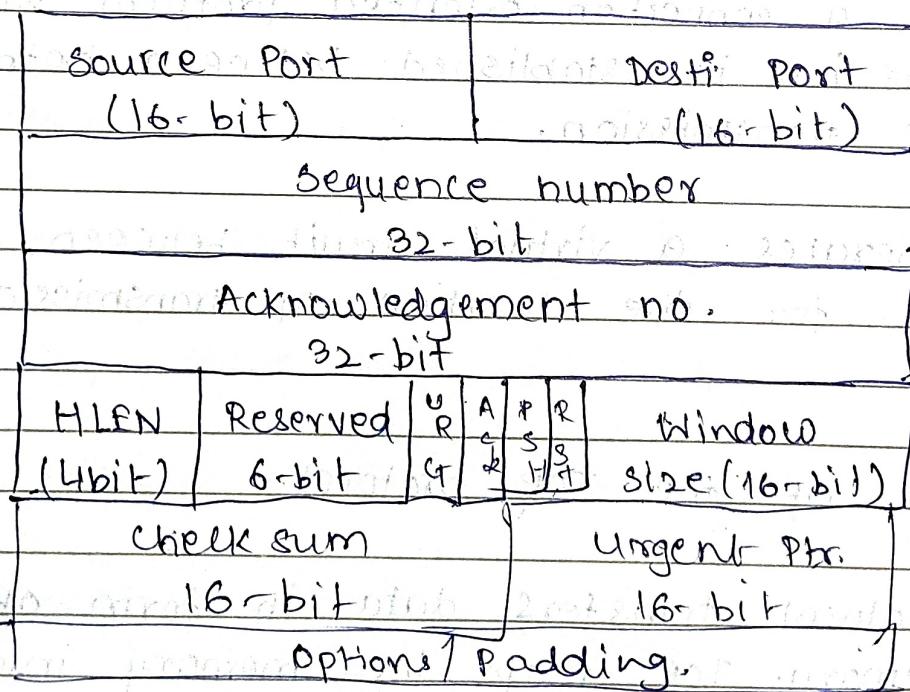
→ TCP generates a virtual circuit between sender and receiver, for the duration of transmission.

### \* Features of the TCP protocol.

1] TCP protocol transfers data in form of a stream of memory. TCP groups the memory into segments and passes it to the IP layer for transmission to the destination.

- 2] TCP assigns sequence numbers to each of its segments and upon reaching its destination, expects a positive acknowledgement. If this acknowledgement is not received, then retransmit the data.
- 3] Prevents the sender from overflowing the receiver.
- 4] Allows multiplexing for various situations.

#### \* TCP Segment Format:



Here,

- 1] Source Port Address:- Address of the application program of source computer.
- 2] Destination Port Address:- Address of the application program of destination computer.
- 3] Sequence Number:- Represents the position of a particular segment of data in a data stream.
- 4] Acknowledgement Number - Acknowledges data from other communicating devices. If ACK field is 1 then it specifies the sequence number that the receiver is expecting to receive.
- 5] Reserved - Reserved for future use.
- 6] Header Length (HLEN) - Specifies the size of the TCP header in 32-bit words, the minimum size of the header is 5 words, and max is 15 words.

Control Field:- Defines the use of the segment.

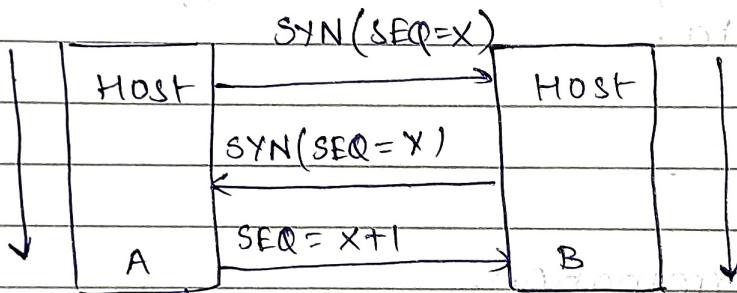
- a] URG - Indicates if msg is urgent
- b] ACK - Validate acknowledgement number.
- c] PSH - Informe that a higher throughput is needed.
- d] RST - resets the TCP connection if there is confusion in sequence number.
- e] SYN - Synchronizes the msg to 3 parts, request, confirmation, and acknowledgement.
- f] FIN - When sender has finished sending data.

→ Window size:- defines size of the window

## QUESTION

### \* TCP Connection Establishment

- Connection Establishment is performed by using a 3-way handshake mechanism. This synchronizes both ends of a network and helps agree on original sequence numbers.
- This mechanism also provides both the sides to transmit data and learn if the other side is able to communicate.



- 1) The requesting end (A) sends a SYN segment with initial seq. no. 'x'.
- 2) The server (HOST B) acknowledges his own SYN segment with seq. no. 'Y' and accepts the sender with seq. no.  $x+1$ .
- 3) An SYN consumes 1 seq. no., then the client acknowledges this SYN from server by accepting seq.  $+1$  ( $seq = x+1$ ,  $ACK = Y+1$ ).

### \* Data Transfer.

- After connection is established, bidirectional data transfer is possible.
- Message is divided to 4 segments, first 3 send both data and acknowledgement while the fourth segment sends only an acknowledgement as there is no more data to be sent.
- When receiver gets the package it increases the acknowledgement number by the length of received data.

### \* Connection Termination.

- The client TCP, after receiving the close command from the client process, sends a FIN segment which consumes only one sequence number.
- The server TCP, after receiving the FIN segment, informs the process of the situation and sends the second segment a FIN+ACK segment to confirm the receipt of the FIN segment.



The client segment then sends a lost msg.  
an ACK segment to confirm ACK segment from  
server TCP