

# 10 reasons why YOU should care about programming competitions

Daniel Epstein, CSE Grad Student

What are Programming  
Competitions?



Rock, Paper, Scissors is a two player game, where each player simultaneously chooses one of the three items after counting to three. The game typically lasts a pre-determined number of rounds. The player who wins the most rounds wins the game. Given the number of rounds the players will compete, it is your job to determine which player wins after those rounds have been played.

The rules for what item wins are as follows:

- Rock always beats Scissors (Rock crushes Scissors)
- Scissors always beat Paper (Scissors cut Paper)
- Paper always beats Rock (Paper covers Rock)

## Input

The first value in the input file will be an integer  $t$  ( $0 < t < 1000$ ) representing the number of test cases in the input file. Following this, on a case by case basis, will be an integer  $n$  ( $0 < n < 100$ ) specifying the number of rounds of Rock, Paper, Scissors played. Next will be  $n$  lines, each with either a capital R, P, or S, followed by a space, followed by a capital R, P, or S, followed by a newline. The first letter is Player 1's choice; the second letter is Player 2's choice.

## Output

For each test case, report the name of the player ('Player 1' or 'Player 2') that wins the game, followed by a newline. If the game ends up in a tie, print 'TIE'.

## Sample Input

```
3
2
R P
S R
3
P P
R S
S R
1
P R
```

## Sample Output

```
Player 2
TIE
Player 1
```

Rock, Paper, Scissors is a two player game, where each player simultaneously chooses one of the three items after counting to three. The game typically lasts a pre-determined number of rounds. The player who wins the most rounds wins the game. Given the number of rounds the players will compete, it is your job to determine which player wins after those rounds have been played.

The rules for what item wins are as follows:

- Rock always beats Scissors (Rock crushes Scissors)
- Scissors always beat Paper (Scissors cut Paper)
- Paper always beats Rock (Paper covers Rock)

## Input

The first value in the input file will be an integer  $t$  ( $0 < t < 1000$ ) representing the number of test cases in the input file. Following this, on a case by case basis, will be an integer  $n$  ( $0 < n < 100$ ) specifying the number of rounds of Rock, Paper, Scissors played. Next will be  $n$  lines, each with either a capital R, P, or S, followed by a space, followed by a capital R, P, or S, followed by a newline. The first letter is Player 1's choice; the second letter is Player 2's choice.

## Output

For each test case, report the name of the player ('Player 1' or 'Player 2') that wins the game, followed by a newline. If the game ends up in a tie, print 'TIE'.

## Sample Input

```
3
2
R P
S R
3
P P
R S
S R
1
P R
```

## Sample Output

```
Player 2
TIE
Player 1
```

Rock, Paper, Scissors is a two player game, where each player simultaneously chooses one of the three items after counting to three. The game typically lasts a pre-determined number of rounds. The player who wins the most rounds wins the game. Given the number of rounds the players will compete, it is your job to determine which player wins after those rounds have been played.

The rules for what item wins are as follows:

- Rock always beats Scissors (Rock crushes Scissors)
- Scissors always beat Paper (Scissors cut Paper)
- Paper always beats Rock (Paper covers Rock)

## Input

The first value in the input file will be an integer  $t$  ( $0 < t < 1000$ ) representing the number of test cases in the input file. Following this, on a case by case basis, will be an integer  $n$  ( $0 < n < 100$ ) specifying the number of rounds of Rock, Paper, Scissors played. Next will be  $n$  lines, each with either a capital R, P, or S, followed by a space, followed by a capital R, P, or S, followed by a newline. The first letter is Player 1's choice; the second letter is Player 2's choice.

## Output

For each test case, report the name of the player ('Player 1' or 'Player 2') that wins the game, followed by a newline. If the game ends up in a tie, print 'TIE'.

## Sample Input

```
3
2
R P
S R
3
P P
R S
S R
1
P R
```

## Sample Output

```
Player 2
TIE
Player 1
```

Rock, Paper, Scissors is a two player game, where each player simultaneously chooses one of the three items after counting to three. The game typically lasts a pre-determined number of rounds. The player who wins the most rounds wins the game. Given the number of rounds the players will compete, it is your job to determine which player wins after those rounds have been played.

The rules for what item wins are as follows:

- Rock always beats Scissors (Rock crushes Scissors)
- Scissors always beat Paper (Scissors cut Paper)
- Paper always beats Rock (Paper covers Rock)

## Input

The first value in the input file will be an integer  $t$  ( $0 < t < 1000$ ) representing the number of test cases in the input file. Following this, on a case by case basis, will be an integer  $n$  ( $0 < n < 100$ ) specifying the number of rounds of Rock, Paper, Scissors played. Next will be  $n$  lines, each with either a capital R, P, or S, followed by a space, followed by a capital R, P, or S, followed by a newline. The first letter is Player 1's choice; the second letter is Player 2's choice.

## Output

For each test case, report the name of the player ('Player 1' or 'Player 2') that wins the game, followed by a newline. If the game ends up in a tie, print 'TIE'.

## Sample Input

```
3
2
R P
S R
3
P P
R S
S R
1
P R
```

## Sample Output

```
Player 2
TIE
Player 1
```

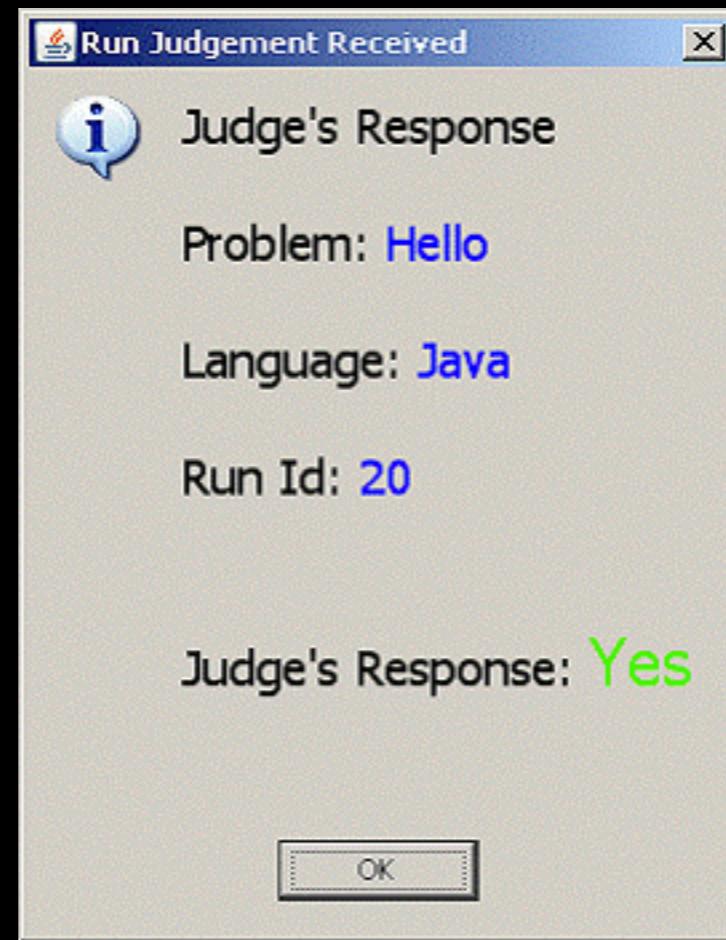
```

import java.io.*;
import java.util.*;

public class Main {
    public static void main(String[] args) {
        Scanner in = new Scanner(System.in);
        int t = in.nextInt();
        for(int a = 0; a < t; a++) {
            int n = in.nextInt();
            int p1score = 0, p2score = 0;
            for(int i = 0; i < n; i++) {
                char p1 = in.next().charAt(0);
                char p2 = in.next().charAt(0);
                if(p1 == 'R') {
                    if(p2 == 'P') {
                        p2score++;
                    } else if(p2 == 'S') {
                        p1score++;
                    }
                } else if(p1 == 'P') {
                    if(p2 == 'R') {
                        p1score++;
                    } else if(p2 == 'S') {
                        p2score++;
                    }
                } else { // 'S'
                    if(p2 == 'R') {
                        p2score++;
                    } else if(p2 == 'P') {
                        p1score++;
                    }
                }
            }
            System.out.println(p1score==p2score?"TIE":(p1score>p2score?"Player 1":"Player 2"));
        }
    }
}

```

1 try, <5 minutes



No - Compilation Error

No - Runtime Exception

No - Time Limit Exceeded

No - Wrong Answer

No - See Contest Staff

No - Compilation Error

No - Runtime Exception

No - Time Limit Exceeded  
~10,000,000 operations

No - Wrong Answer

No - See Contest Staff

# ACM-ICPC World Finals

June 22 - June 26

# 2014

Ekaterinburg

host Ural Federal University

## world finals ↓

- Schedule
- Activities
- Local Information
- Teams
- World Finals Rules
- On-Site Registration
- Video/Photo Coverage
- World Finals Results
- Past Problems
- Fact Sheet
- Prog. Environment

## regionals ↓

- Regional Finder
- Upcoming Regionals
- Regional Results
- Regional Rules
- Getting Involved
- Starting a Regional
- Free ACM Membership

## compete ↓

- Preparation
- Policies & Procedures
- FAQs
- The Problems

## community ↓

- IBM
- Upsilon Pi Epsilon
- ACM
- Fact Sheet
- History
- Contacts

## Participate in Regionals Now!

Find a regional contest

I am a  
[Coach](#)   [Contestant](#)   [Volunteer](#)

Think. Create. Solve.



ICPCNews

Connect.  
(and click for more!)



Share your story.  
[#ICPC2014](#)



IBM

event  
sponsor



**world finals** ↗  
Schedule  
Activities  
Local Information  
Teams  
World Finals Rules  
On-Site Registration  
Video/Photo Coverage  
World Finals Results  
Past Problems  
Fact Sheet  
Prog. Environment

**regionals** ↗  
Regional Finder  
Upcoming Regionals  
Regional Results  
Regional Rules  
Getting Involved  
Starting a Regional  
Free ACM Membership

**compete** ↗  
Preparation  
Policies & Procedures  
FAQs  
The Problems

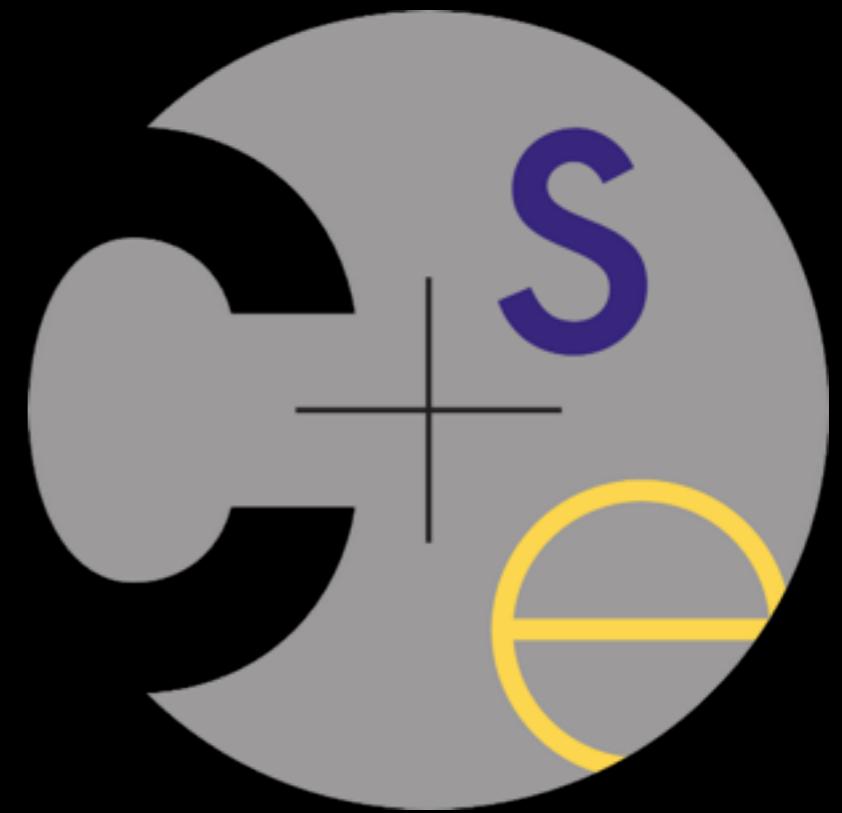
**community** ↗  
IBM  
Upsilon Pi Epsilon  
ACM  
Fact Sheet  
History  
Contacts



Registration opens 12.20.10



Who am I?





Michigan Region









**acm**

2011 World Finals  
International Collegiate  
Programming Contest

**IBM.**

event  
sponsor



**A - Light Purple**

**G - Yellow**

**B – Dark Purple**

**H - Silver**

**C - Blue**

**I – Black**

**D - RED**

**J - Pink**

**E - White**

**K- Orange**

**F - Green**

| Team  | A | B | C | D | E | F  | G | H | I | J | K | S |
|---|---|---|---|---|---|----|---|---|---|---|---|---|
| Zhejiang University                               | 4 | 1 | 1 |   | 2 |    | 4 | 1 |   | 1 | 1 |   |
| University of Michigan at Ann Arbor               | 2 | 1 | 2 |   | 3 |    | 1 | 3 |   | 2 | 1 |   |
| tsinghua University                               | 1 |   | 2 |   | 1 | 2  |   | 2 |   | 2 | 1 |   |
| . Petersburg State University                     | 1 | 3 | 1 |   | 1 |    | 1 |   | 1 | 1 | 1 |   |
| zhny Novgorod State University                    | 2 | 2 | 1 |   | 1 |    | 2 | 2 |   | 1 | 1 |   |
| aratov State University                           | 2 | 2 | 1 |   | 4 | 10 |   | 1 |   | 1 | 1 |   |
| rich-Alexander-University Erlangen-Nuremberg      | 1 |   | 1 |   | 1 |    | 3 | 1 | 6 | 1 |   |   |
| netsk National University                         | 1 |   | 1 | 1 | 1 |    | 2 | 2 |   | 4 | 1 |   |
| ellenian University in Krakow                     | 5 | 3 | 1 |   | 2 | 5  |   | 2 |   | 1 | 1 |   |
| cow State University                              | 3 |   | 2 |   | 1 | 3  | 3 | 2 |   | 2 | 3 |   |
| State University                                  | 2 |   | 1 | 2 | 3 |    | 6 | 4 |   | 1 | 1 |   |
| iversity of Waterloo                              | 2 |   | 1 |   | 2 |    | 5 | 5 | 4 | 3 |   |   |
| Ia V.I. Vernadsky National University             | 1 |   | 1 |   | 1 |    | 9 | 1 |   | 1 | 1 |   |
| nal Taiwan University                             |   | 6 | 1 |   | 3 |    | 2 | 1 |   | 1 | 1 |   |
| rsity of Warsaw                                   | 2 | 4 | 2 |   | 8 | 2  |   | 3 |   | 3 | 3 |   |
| burg State University of IT, Mechanics and Optics | 2 |   | 2 |   | 4 |    | 5 | 2 |   | 1 |   |   |
| ng Technological University                       | 1 |   | 1 |   | 5 |    | 7 | 3 |   | 1 |   |   |
| sidad de Buenos Aires - FCEN                      | 1 |   | 2 |   | 3 |    |   | 1 |   | 1 |   |   |
| anced Institute of Science and Technology         | 1 |   | 1 |   | 5 | 4  |   | 1 |   | 2 |   |   |
| University  | 3 |   | 1 |   | 1 | 2  |   | 3 |   | 5 |   |   |

| Team  | A | B | C | D | E  | F | G | H | I | J | K | S |
|---|---|---|---|---|----|---|---|---|---|---|---|---|
| Zhejiang University                               | 4 | 1 | 1 | 2 |    | 4 | 1 |   | 1 | 1 |   |   |
| University of Michigan at Ann Arbor               |   |   | 2 | 3 | 1  | 3 | 1 | 3 | 2 | 2 | 1 |   |
| tsinghua University                               | 1 |   | 2 | 1 | 2  |   | 2 |   | 2 | 2 | 1 |   |
| . Petersburg State University                     | 1 | 3 | 1 | 1 |    |   | 1 |   | 1 | 1 | 1 |   |
| zhny Novgorod State University                    | 2 | 2 | 1 | 1 |    | 2 | 2 |   | 1 | 1 | 1 |   |
| aratov State University                           | 2 | 2 | 1 | 4 | 10 |   | 1 |   | 1 | 1 | 1 |   |
| rich-Alexander-University Erlangen-Nuremberg      | 1 |   | 1 | 1 |    | 3 | 1 | 6 | 1 | 4 | 1 |   |
| netsk National University                         | 1 |   | 1 | 1 | 1  | 2 | 2 |   | 4 | 1 |   |   |
| ellenian University in Krakow                     | 5 | 3 | 1 | 2 | 5  |   | 2 |   | 1 | 1 | 1 |   |
| cow State University                              | 3 |   | 2 | 1 | 3  | 3 | 2 |   | 2 | 2 | 3 |   |
| State University                                  | 2 |   | 1 | 2 | 3  |   | 6 | 4 |   | 1 | 1 |   |
| iversity of Waterloo                              | 2 |   | 1 | 2 |    |   | 5 | 5 | 4 | 3 | 3 |   |
| Ia V.I. Vernadsky National University             | 1 |   | 1 | 1 |    | 9 | 1 |   | 1 | 1 | 1 |   |
| nal Taiwan University                             |   | 6 | 1 | 3 |    | 2 | 1 |   | 1 | 1 | 1 |   |
| rsity of Warsaw                                   | 2 | 4 | 2 | 8 | 2  |   | 3 |   | 3 | 3 | 3 |   |
| burg State University of IT, Mechanics and Optics | 2 |   | 2 | 4 |    | 5 | 2 |   | 1 | 1 | 1 |   |
| ng Technological University                       | 1 |   | 1 | 5 |    | 7 | 3 |   | 1 | 1 | 1 |   |
| sidad de Buenos Aires - FCEN                      | 1 |   | 2 | 3 |    |   | 1 |   | 1 | 1 | 1 |   |
| anced Institute of Science and Technology         | 1 |   | 1 | 5 | 4  |   | 1 |   | 2 | 2 | 2 |   |
| University  | 3 |   | 1 | 1 | 2  |   | 3 |   | 3 | 5 | 5 |   |



Cassandra  
Daniel Epstein  
University of Oregon  
Microsoft

2011  
Mid-Atlantic Regional  
Robotics Competition

lenovo

Microsoft





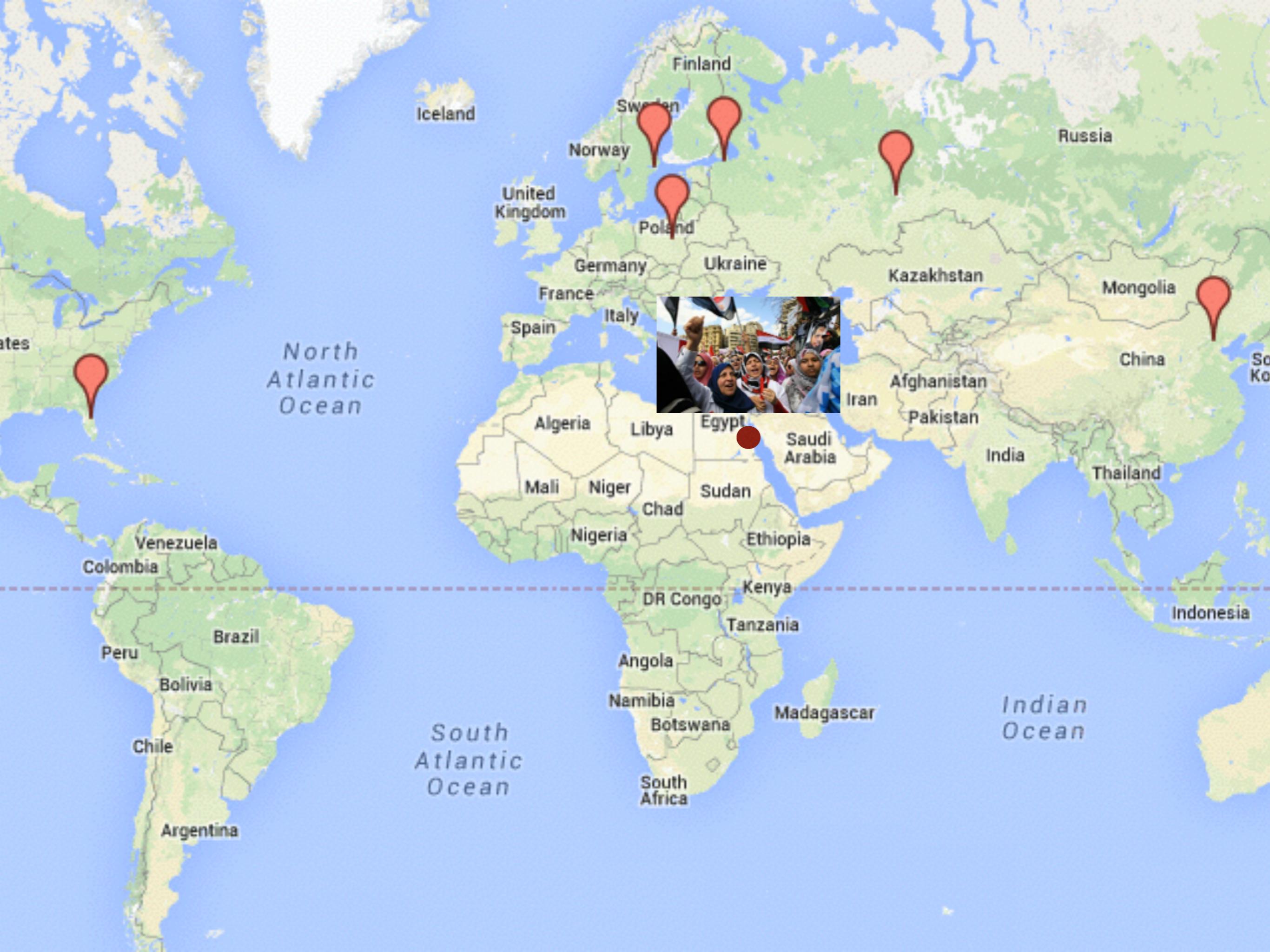
So... why compete?

#1: Travel to exotic  
locations









North  
Atlantic  
Ocean

South  
Atlantic  
Ocean

Indian  
Ocean





# #2: Free stuff!

U Virginia

2011 World Finals  
Orlando



acm International Collegiate  
Programming Contest

IBM

event  
sponsor

ACM-ICPC  
2009

World Finals

Stockholm,  
Sweden



acm  
International  
Computer  
Programming  
Contest

IBM

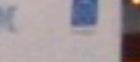
Official Sponsor



IBM  
Software  
Group



IBM



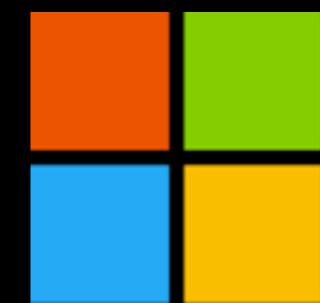


#3: Looks great on a  
resume





Google



Microsoft

#4: Learn to play nice  
with others



  Allowed on Timeline ▾

I have nothing more to say... — with Daniel Epstein.

[Tag Photo](#) [Add Location](#)

[Like](#) · [Comment](#) · [Stop Notifications](#) · [Share](#)

 [REDACTED] and [REDACTED] like this.

 **Daniel Epstein** "fight" isn't the right word here... and it's 3v3 instead of 6v6. [REDACTED] and [REDACTED] I choose you!

November 4, 2011 at 3:34pm · Like

 [REDACTED] Wow, guys. Wow...

November 4, 2011 at 3:35pm · Like

 [REDACTED] Oh, and: IT'S ON!!!!

November 4, 2011 at 3:35pm · Like

 [REDACTED] So yeah, I could probably have put more thought into the picture, but some of us are busy trying to do other things like getting ready for victory.

November 4, 2011 at 3:38pm · Like

 **Daniel Epstein** My team is already level 100. Nothing more to be done.

November 4, 2011 at 4:12pm · Like · 1

 Write a comment... 

# Former Roommate



Current Roommates

I have nothing more to say... — with Daniel Epstein.

Like · Comment · Stop Notifications · Share

**Daniel Epstein** "fight" isn't the right word here... and it's 3v3 instead of 6v6. [REDACTED] and [REDACTED] I choose you!

November 4, 2011 at 3:34pm · Like

[REDACTED] Wow, guys. Wow... November 4, 2011 at 3:35pm · Like

[REDACTED] Oh, and: IT'S ON!!!! November 4, 2011 at 3:35pm · Like

[REDACTED] So yeah, I could probably have put more thought into the picture, but some of us are busy trying to do other things like getting ready for victory.

November 4, 2011 at 3:38pm · Like

**Daniel Epstein** My team is already level 100. Nothing more to be done.

November 4, 2011 at 4:12pm · Like · 1

Write a comment...

A screenshot of a Facebook post by Daniel Epstein. The post says "I have nothing more to say... — with Daniel Epstein." and includes a photo of two people. Below the post are several comments from other users. A blue arrow points from the text "Epstein wants to fight!" in the Mario screenshot up to the first comment on the Facebook post. Another blue arrow points from the "Former Roommate" header down to the Facebook post area.

# #5: Learn to debug

List1

1

2

3

List1

- 1
- 2
- 3

List2

- 6
- 1
- 3
- 5
- 2
- 8
- 9

List1

1 2 3

List2

6 5 8 9

# Find the bug!

```
import java.io.*;
import java.util.*;

public class iterate {

    public static void main(String[] args) {
        ArrayList<Integer> list1 = new ArrayList<Integer>(Arrays.asList(new Integer[]{1, 2, 3}));
        ArrayList<Integer> list2 = new ArrayList<Integer>(Arrays.asList(new Integer[]{6, 1, 3, 5, 2, 8, 9}));
        // Remove all occurrences of elements in list1 from list2
        for(int i=0;i<list2.size();i++) {
            for(int j=0;j<list1.size(); j++) {
                if(list2.get(i) == list1.get(j)) {
                    list2.remove(i);
                    break;
                }
            }
        }
        System.out.println(list2);
    }
}
```

# Find the bug!

```
import java.io.*;
import java.util.*;

public class iterate {

    public static void main(String[] args) {
        ArrayList<Integer> list1 = new ArrayList<Integer>(Arrays.asList(new Integer[]{1, 2, 3}));
        ArrayList<Integer> list2 = new ArrayList<Integer>(Arrays.asList(new Integer[]{6, 1, 3, 5, 2, 8, 9}));
        // Remove all occurrences of elements in list1 from list2
        for(int i=0;i<list2.size();i++) {
            for(int j=0;j<list1.size(); j++) {
                if(list2.get(i) == list1.get(j)) {
                    list2.remove(i);
                    break;
                }
            }
        }
        System.out.println(list2);
    }
}
```



#6: Ace technical  
interviews

Bessie has gone on a trip, and she's riding a roller coaster! Bessie really likes riding the roller coaster, but unfortunately she often gets dizzy.

The roller coaster has a number of distinct sections that Bessie rides in order. At the beginning of the ride, Bessie's dizziness and fun levels are both at 0. For each section of the roller coaster, Bessie can either keep her eyes open or keep them closed (and must keep them that way for the whole section). If she keeps her eyes open for a section, her total fun increases by a Fun factor for that section, and her dizziness increases by a Dizziness factor for that section. However, if she keeps her eyes closed for the section, her total fun will not change, but her dizziness will decrease by a value that's constant for the entire roller coaster. (Note that her dizziness can never go below 0.)

If, at any point, Bessie's dizziness is above a certain limit, Bessie will get sick. Write a program to find the maximum amount of fun Bessie can have without getting sick.

## Input

There will be several test cases in the input. Each test case will begin with a line with three integers:

*N K L*

Where  $N (1 \leq N \leq 1,000)$  is the number of sections in this particular roller coaster,  $K (1 \leq K \leq 500)$  is the amount that Bessie's dizziness level will go down if she keeps her eyes closed on any section of the ride, and  $L (1 \leq L \leq 300,000)$  is the limit of dizziness that Bessie can tolerate - if her dizziness ever becomes larger than  $L$ , Bessie will get sick, and that's not fun!

Each of the next  $N$  lines will describe a section of the roller coaster, and will have two integers:

*F D*

Where  $F (1 \leq F \leq 20)$  is the increase to Bessie's total fun that she'll get if she keeps her eyes open on that section, and  $D (1 \leq D \leq 500)$  is the increase to her dizziness level if she keeps her eyes open on that section. The sections will be listed in order. The input will end with a line with three 0s.

## Output

For each test case, output a single integer, representing the maximum amount of fun Bessie can have on that roller coaster without exceeding her dizziness limit. Print each integer on its own line with no spaces. Do not print any blank lines between answers.

Bessie has gone on a trip, and she's riding a roller coaster! Bessie really likes riding the roller coaster, but unfortunately she often gets dizzy.

The roller coaster has a number of distinct sections that Bessie rides in order. At the beginning of the ride, Bessie's dizziness and fun levels are both at 0. For each section of the roller coaster, Bessie can either keep her eyes open or keep them closed (and must keep them that way for the whole section). If she keeps her eyes open for a section, her total fun increases by a Fun factor for that section, and her dizziness increases by a Dizziness factor for that section. However, if she keeps her eyes closed for the section, her total fun will not change, but her dizziness will decrease by a value that's constant for the entire roller coaster. (Note that her dizziness can never go below 0.)

If, at any point, Bessie's dizziness is above a certain limit, Bessie will get sick. Write a program to find the maximum amount of fun Bessie can have without getting sick.

## Input

There will be several test cases in the input. Each test case will begin with a line with three integers:

$N \ K \ L$

Where  $N (1 \leq N \leq 1,000)$  is the number of sections in this particular roller coaster,  $K (1 \leq K \leq 500)$  is the amount that Bessie's dizziness level will go down if she keeps her eyes closed on any section of the ride, and  $L (1 \leq L \leq 300,000)$  is the limit of dizziness that Bessie can tolerate - if her dizziness ever becomes larger than  $L$ , Bessie will get sick, and that's not fun!

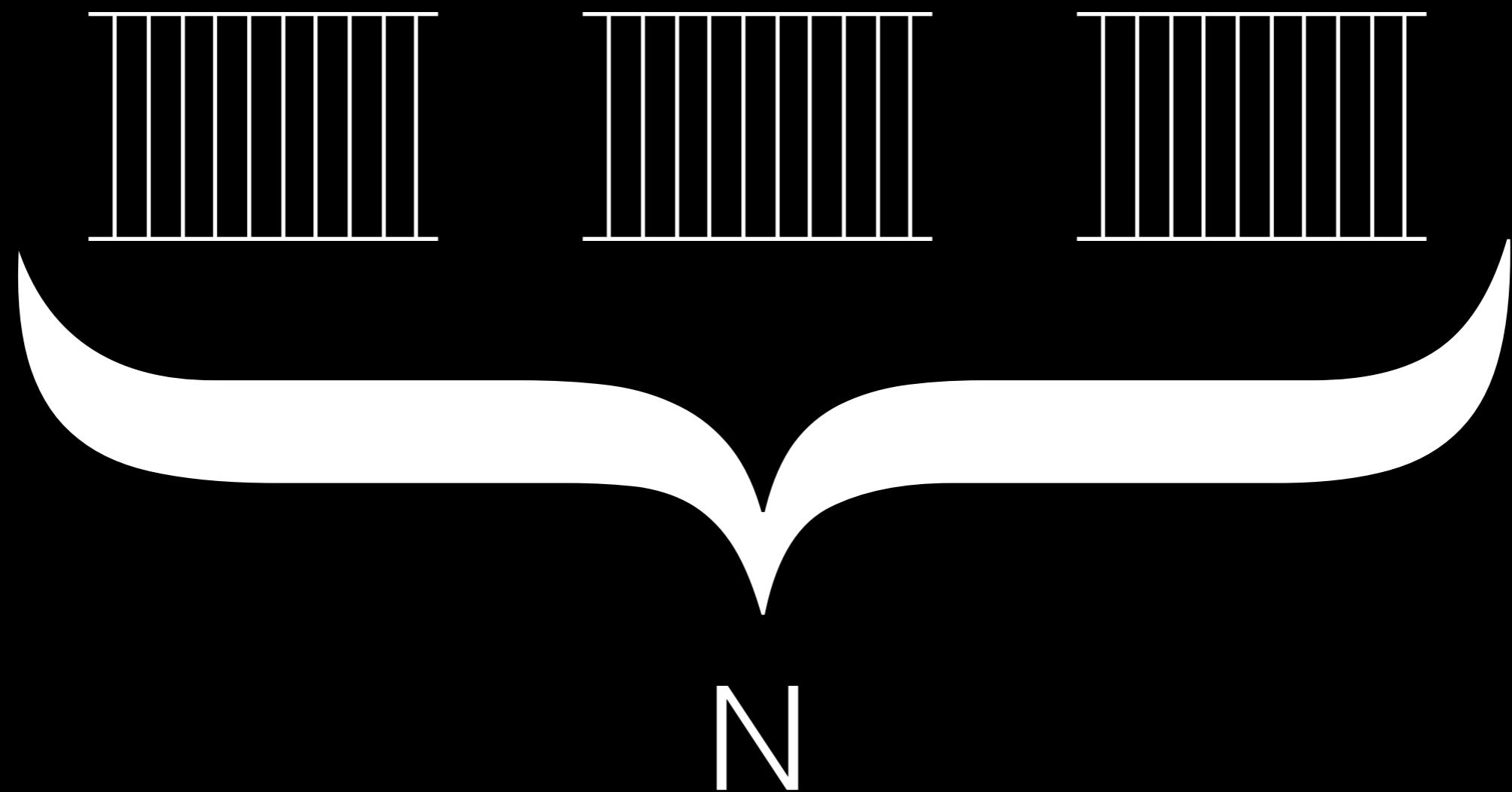
Each of the next  $N$  lines will describe a section of the roller coaster, and will have two integers:

$F \ D$

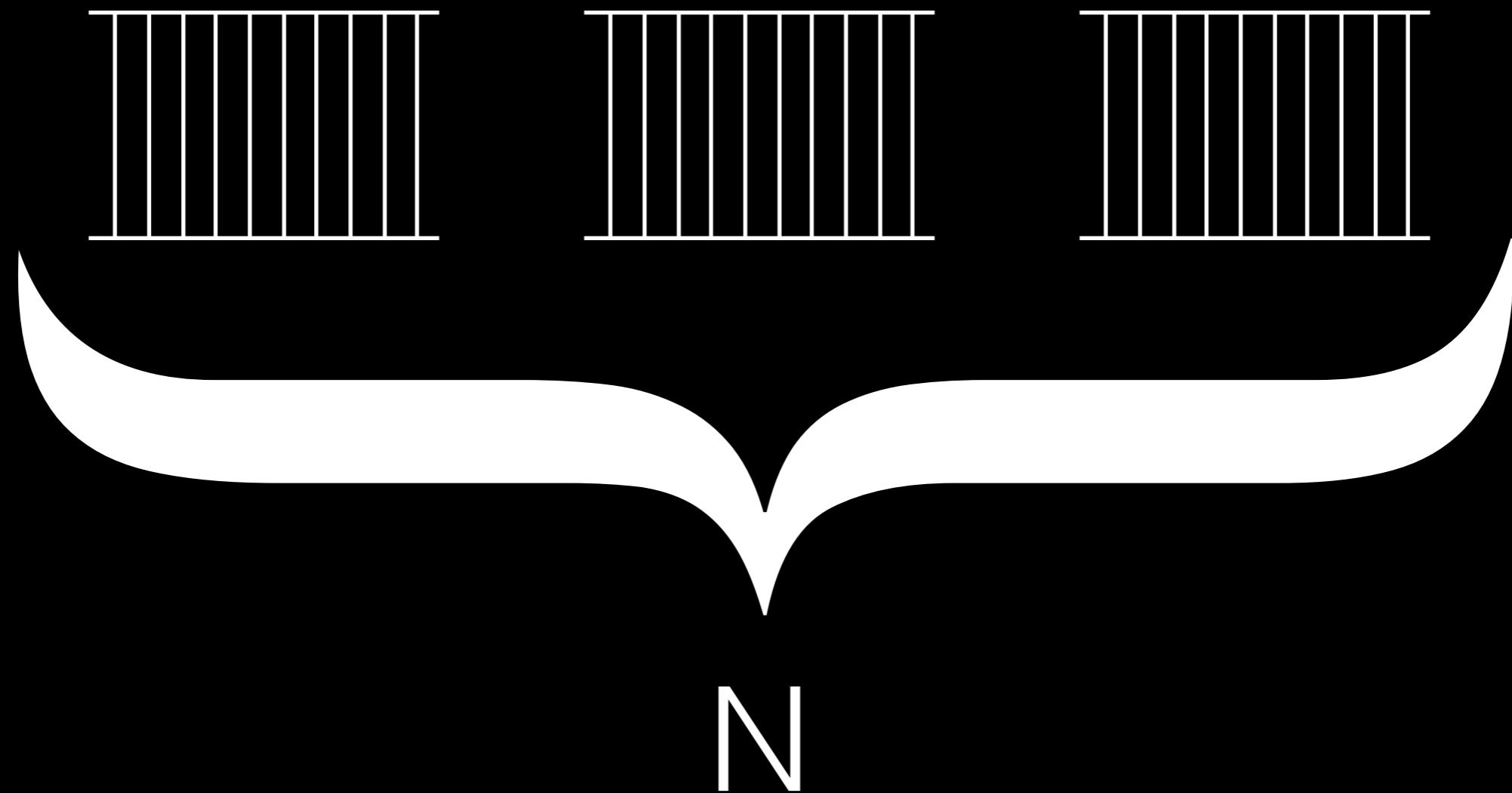
Where  $F (1 \leq F \leq 20)$  is the increase to Bessie's total fun that she'll get if she keeps her eyes open on that section, and  $D (1 \leq D \leq 500)$  is the increase to her dizziness level if she keeps her eyes open on that section. The sections will be listed in order. The input will end with a line with three 0s.

## Output

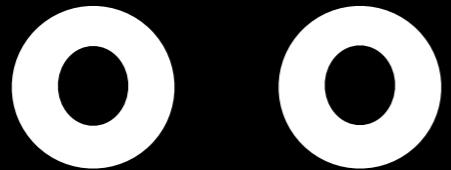
For each test case, output a single integer, representing the maximum amount of fun Bessie can have on that roller coaster without exceeding her dizziness limit. Print each integer on its own line with no spaces. Do not print any blank lines between answers.



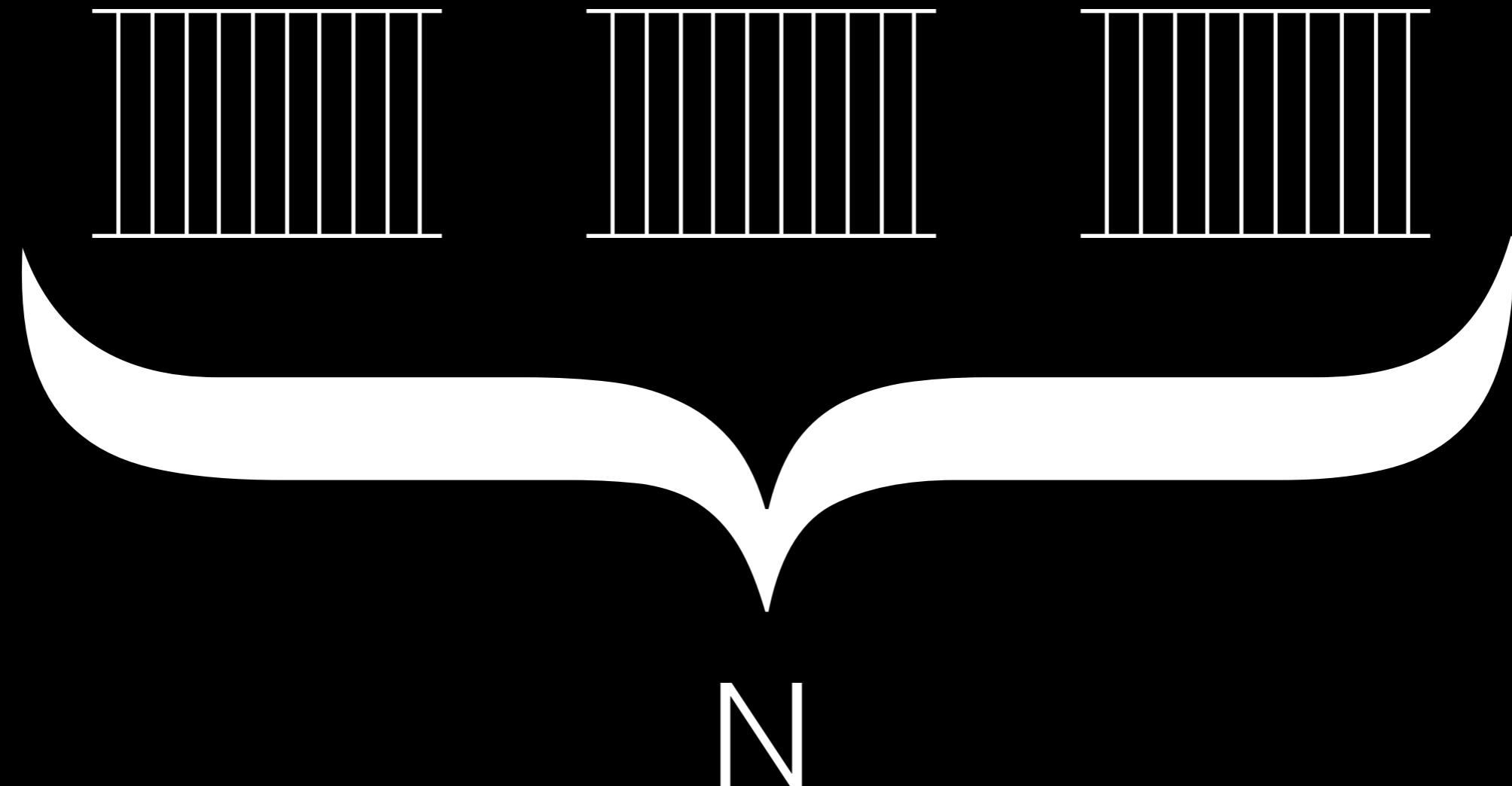
Fun=X, Dizzy=Y



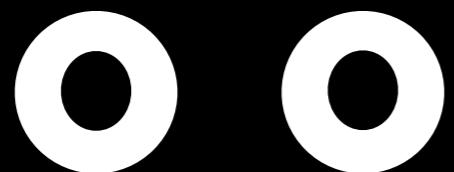
Fun=X, Dizzy=Y



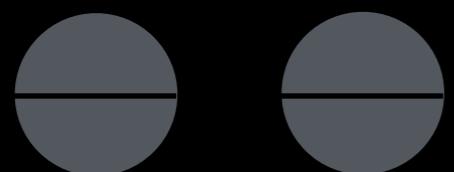
Fun+=F[i], Dizzy+=D[i]



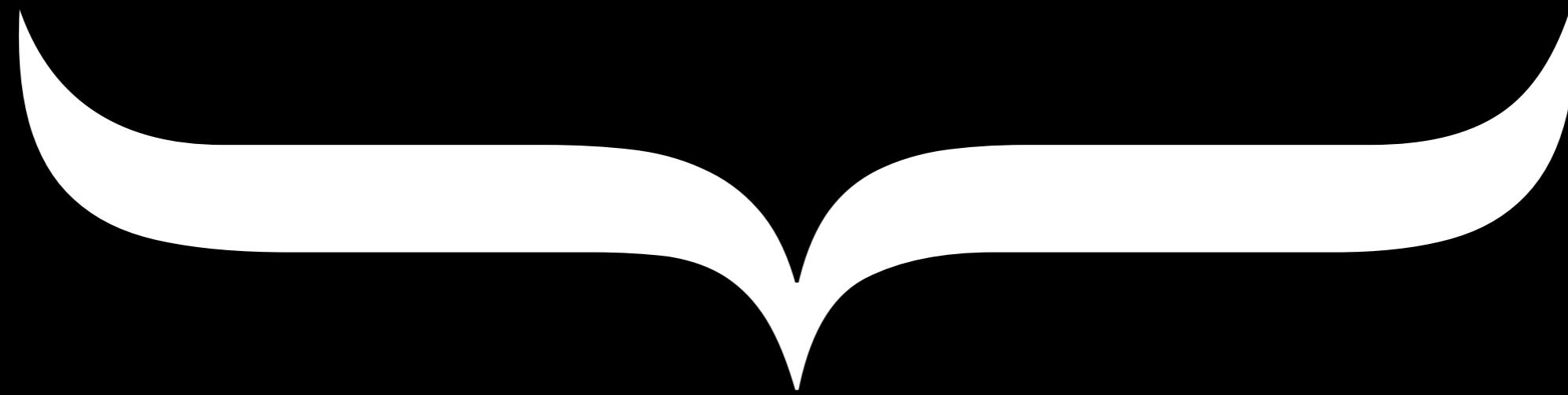
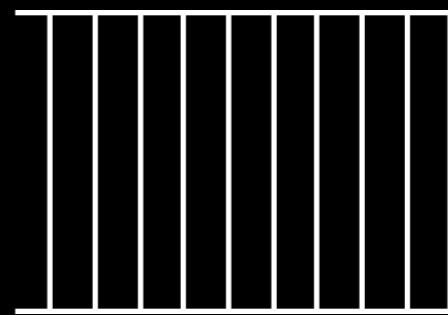
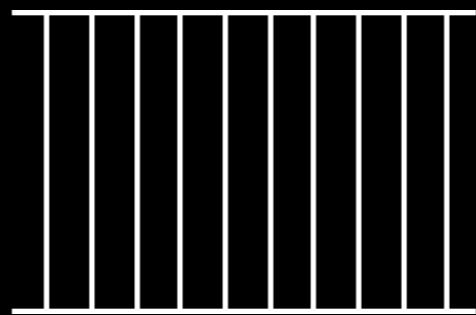
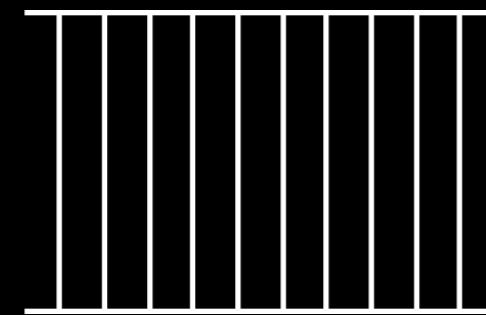
Fun=X, Dizzy=Y



Fun+=F[i], Dizzy+=D[i]

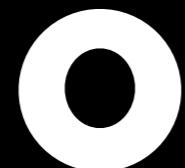
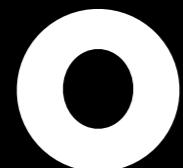


Dizzy-=K

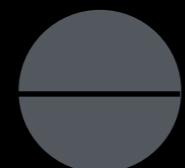
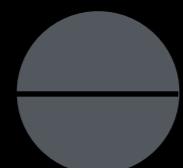


N

$\text{Fun} = X$ ,  $\text{Dizzy} = Y$

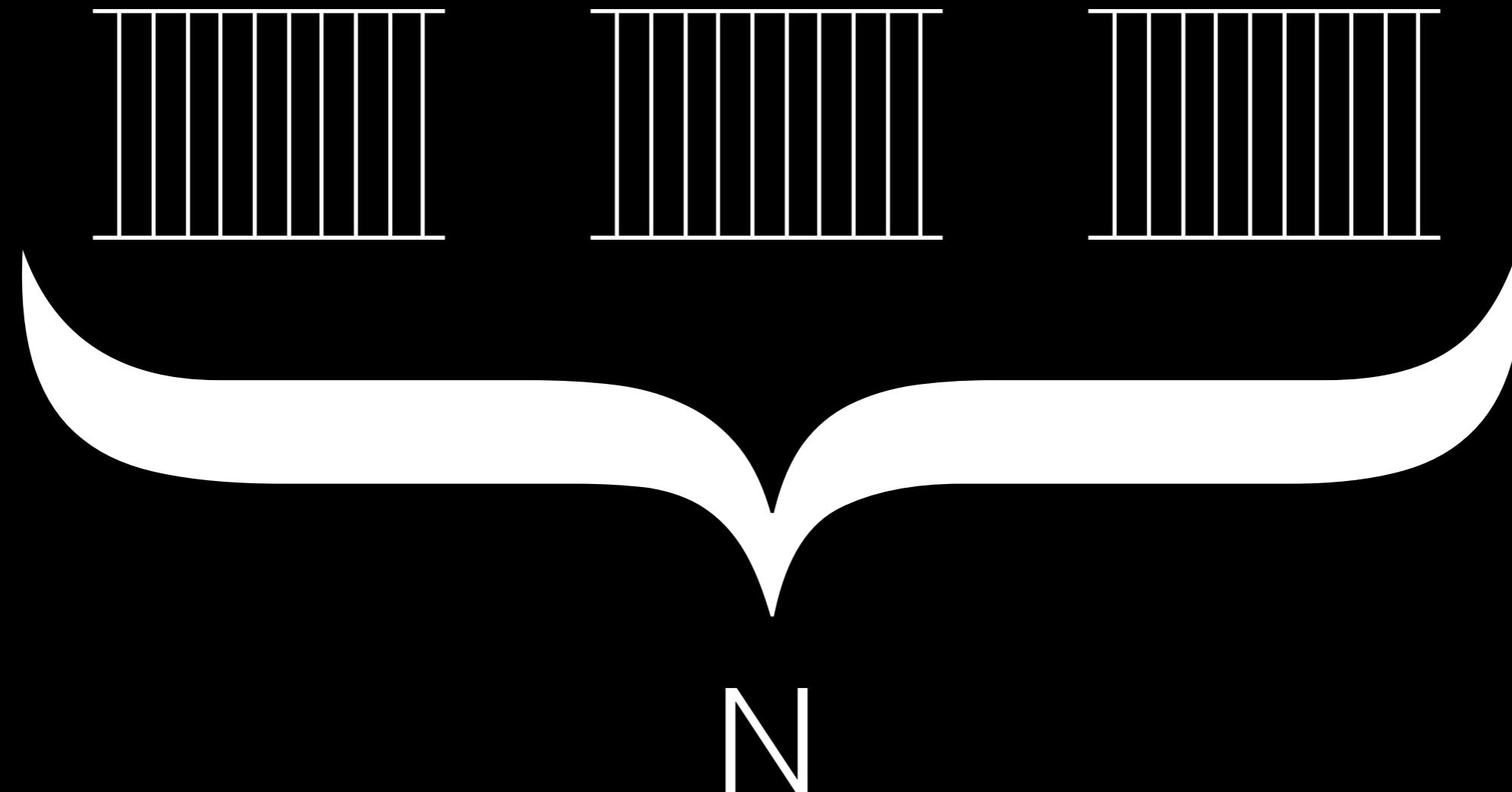


$\text{Fun} += F[i]$ ,  $\text{Dizzy} += D[i]$



$\text{Dizzy} -= K$

$\text{Dizzy} \leq L$   
at all times



```
import java.io.*;
import java.util.*;

public class Main {
    public static void main(String[ ] args) {
        Scanner in = new Scanner(System.in);
        while(true) {
            int N = in.nextInt(), K = in.nextInt(), L = in.nextInt();
            if(N==0 && K==0 && L==0) {
                break;
            }
            int[ ] fun = new int[N], dizzy = new int[N];
            for(int i=0; i<N; i++) {
                fun[i] = in.nextInt();
                dizzy[i] = in.nextInt();
            }
            // Now what?
        }
    }
}
```

```

import java.io.*;
import java.util.*;

public class Main {
    public static void main(String[ ] args) {
        Scanner in = new Scanner(System.in);
        while(true) {
            int N = in.nextInt(), K = in.nextInt(), L = in.nextInt();
            if(N==0 && K==0 && L==0) {
                break;
            }
            int[ ] fun = new int[N], dizzy = new int[N];
            for(int i=0; i<N; i++) {
                fun[i] = in.nextInt();
                dizzy[i] = in.nextInt();
            }
            // Now what?
        }
    }
}

```

$dp[n][l]$  = maximum possible fun as of section  $n$  with dizziness less than  $l$   
 $dp[n][l] = \max(dp[n-1][l - dizzy[n]] + fun[n], dp[n-1][l+k])$

**O(N\*L)**

```

import java.io.*;
import java.util.*;

public class Main {
    public static void main(String[ ] args) {
        Scanner in = new Scanner(System.in);
        while(true) {
            int N = in.nextInt(), K = in.nextInt(), L = in.nextInt();
            if(N==0 && K==0 && L==0) {
                break;
            }
            int[ ] fun = new int[N], dizzy = new int[N];
            for(int i=0; i<N; i++) {
                fun[i] = in.nextInt();
                dizzy[i] = in.nextInt();
            }
            // Now what?
        }
    }
}

```

$dp[n][l]$  = maximum possible fun as of section  $n$  with dizziness less than  $l$   
 $dp[n][l] = \max(dp[n-1][l - dizzy[n]] + fun[n], dp[n-1][l+k])$

**O( $N*L$ )**

Where  $N (1 \leq N \leq 1,000)$  is the number of sections in this particular roller coaster,  $K (1 \leq K \leq 500)$  is the amount that Bessie's dizziness level will go down if she keeps her eyes closed on any section of the ride, and  $L (1 \leq L \leq 300,000)$  is the limit of dizziness that Bessie can tolerate - if her dizziness ever becomes larger than  $L$ , Bessie will get sick, and that's not fun!

$1,000 * 300,000 = 300,000,000 \rightarrow$  WAY TOO HIGH!!

## Input

There will be several test cases in the input. Each test case will begin with a line with three integers:

*N K L*

Where  $N (1 \leq N \leq 1,000)$  is the number of sections in this particular roller coaster,  $K (1 \leq K \leq 500)$  is the amount that Bessie's dizziness level will go down if she keeps her eyes closed on any section of the ride, and  $L (1 \leq L \leq 300,000)$  is the limit of dizziness that Bessie can tolerate - if her dizziness ever becomes larger than  $L$ , Bessie will get sick, and that's not fun!

Each of the next  $N$  lines will describe a section of the roller coaster, and will have two integers:

*F D*

Where  $F (1 \leq F \leq 20)$  is the increase to Bessie's total fun that she'll get if she keeps her eyes open on that section, and  $D (1 \leq D \leq 500)$  is the increase to her dizziness level if she keeps her eyes open on that section.

## Input

There will be several test cases in the input. Each test case will begin with a line with three integers:

$N K L$

Where  $N (1 \leq N \leq 1,000)$  is the number of sections in this particular roller coaster,  $K (1 \leq K \leq 500)$  is the amount that Bessie's dizziness level will go down if she keeps her eyes closed on any section of the ride, and  $L (1 \leq L \leq 300,000)$  is the limit of dizziness that Bessie can tolerate - if her dizziness ever becomes larger than  $L$ , Bessie will get sick, and that's not fun!

Each of the next  $N$  lines will describe a section of the roller coaster, and will have two integers:

$F D$

Where  $F (1 \leq F \leq 20)$  is the increase to Bessie's total fun that she'll get if she keeps her eyes open on that section, and  $D (1 \leq D \leq 500)$  is the increase to her dizziness level if she keeps her eyes open on that section.

## Input

There will be several test cases in the input. Each test case will begin with a line with three integers:

*N K L*

Where  $N (1 \leq N \leq 1,000)$  is the number of sections in this particular roller coaster,  $K (1 \leq K \leq 500)$  is the amount that Bessie's dizziness level will go down if she keeps her eyes closed on any section of the ride, and  $L (1 \leq L \leq 300,000)$  is the limit of dizziness that Bessie can tolerate - if her dizziness ever becomes larger than  $L$ , Bessie will get sick, and that's not fun!

Each of the next  $N$  lines will describe a section of the roller coaster, and will have two integers:

*F D*

Where  $F (1 \leq F \leq 20)$  is the increase to Bessie's total fun that she'll get if she keeps her eyes open on that section, and  $D (1 \leq D \leq 500)$  is the increase to her dizziness level if she keeps her eyes open on that section.

$dp[n][f] = \text{minimum dizziness with fun level } f$

$dp[n][f] = \min(dp[n-1][f - \text{fun}[n]] + \text{dizzy}[n], dp[n-1][f] - k)$

**O(N\*F)**

$1,000 * 20 = 20,000 \rightarrow \text{Totally reasonable!}$

```

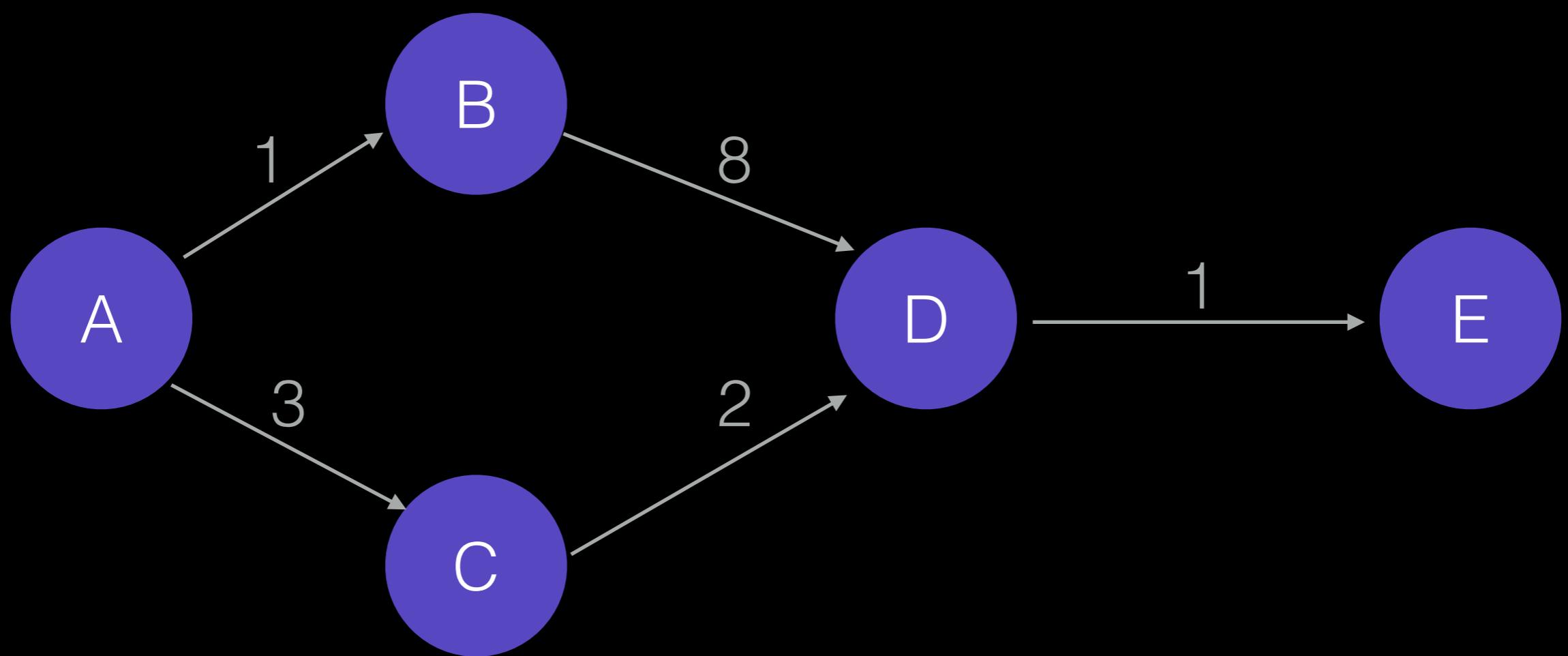
import java.io.*;
import java.util.*;

public class Main {
    public static void main(String[] args) {
        Scanner in = new Scanner(System.in);
        while(true) {
            int N = in.nextInt(), K = in.nextInt(), L = in.nextInt();
            if(N==0 && K==0 && L==0) {
                break;
            }
            int[] fun = new int[N], dizzy = new int[N];
            int maxF = 0;
            for(int i=0; i<N; i++) {
                fun[i] = in.nextInt();
                dizzy[i] = in.nextInt();
                maxF += fun[i];
            }
            // Now what?
            int[][] dp = new int[N][maxF+1];
            for(int i=0; i<N; i++) {
                for(int j=0; j<=maxF; j++) {
                    dp[i][j] = Integer.MAX_VALUE;
                }
            }
            dp[0][0] = 0; //Populate with the first values
            if(dizzy[0] <= L) {
                dp[0][fun[0]] = dizzy[0];
            }
            for(int n=1; n<N; n++) {
                for(int f=0; f<=maxF; f++) {
                    if(dp[n-1][f] != Integer.MAX_VALUE) {
                        dp[n][f] = Math.max(dp[n-1][f] - K, 0); // Dizziness can't go below 0
                    }
                    if(f - fun[n] >= 0 && dp[n-1][f-fun[n]] != Integer.MAX_VALUE && dp[n-1][f-fun[n]] + dizzy[n] <= L) {
                        dp[n][f] = Math.min(dp[n-1][f - fun[n]] + dizzy[n], dp[n][f]);
                    }
                }
            }
            for(int i=maxF; i>=0; i--) {
                if(dp[N-1][i] != Integer.MAX_VALUE) {
                    System.out.println(i);
                    break;
                }
            }
        }
    }
}

```

2 tries, ~40 minutes (off-by-one error)

#7: Deepen  
understanding of  
language semantics



# Find the bug!

```
import java.io.*;
import java.util.*;

public class dijkstra {

    public static void main(String[] args) {
        Node a = new Node();
        Node b = new Node();
        Node c = new Node();
        Node d = new Node();
        Node e = new Node();
        a.edges.put(b, 1);
        a.edges.put(c, 3);
        b.edges.put(d, 8);
        c.edges.put(d, 2);
        d.edges.put(e, 1);
        ArrayList<Node> allNodes = new ArrayList<Node>(Arrays.asList(new Node[]{e, d, c, b, a}));
        dijkstra(a, allNodes);
        System.out.printf("Distance to Node e is: %d\n", e.distance);
    }

    public static void dijkstra(Node root, ArrayList<Node> allNodes) {
        PriorityQueue<Node> q = new PriorityQueue<Node>();
        root.distance = 0;
        //Add all nodes to the priority queue
        for(Node n : allNodes) {
            q.add(n);
        }

        while(q.size() > 0) {
            Node u = q.poll();
            for(Node n : u.edges.keySet()) {
                n.distance = Math.min(n.distance, u.distance + u.edges.get(n));
            }
        }
    }
}

class Node implements Comparable<Node> {
    public HashMap<Node, Integer> edges = new HashMap<Node, Integer>();
    public int distance = Integer.MAX_VALUE;

    public int compareTo(Node o) {
        return (distance < o.distance) ? -1 : ((distance == o.distance) ? 0 : 1);
    }
}
```

# Find the bug!

```
import java.io.*;
import java.util.*;

public class dijkstra {

    public static void main(String[] args) {
        Node a = new Node();
        Node b = new Node();
        Node c = new Node();
        Node d = new Node();
        Node e = new Node();
        a.edges.put(b, 1);
        a.edges.put(c, 3);
        b.edges.put(d, 8);
        c.edges.put(d, 2);
        d.edges.put(e, 1);
        ArrayList<Node> allNodes = new ArrayList<Node>(Arrays.asList(new Node[]{e, d, c, b, a}));
        dijkstra(a, allNodes);
        System.out.printf("Distance to Node e is: %d\n", e.distance);
    }

    public static void dijkstra(Node root, ArrayList<Node> allNodes) {
        PriorityQueue<Node> q = new PriorityQueue<Node>();
        root.distance = 0;
        //Add all nodes to the priority queue
        for(Node n : allNodes) {
            q.add(n);
        }

        while(q.size() > 0) {
            Node u = q.poll();
            for(Node n : u.edges.keySet()) {
                n.distance = Math.min(n.distance, u.distance + u.edges.get(n));
            }
        }
    }

    class Node implements Comparable<Node> {
        public HashMap<Node, Integer> edges = new HashMap<Node, Integer>();
        public int distance = Integer.MAX_VALUE;

        public int compareTo(Node o) {
            return (distance < o.distance) ? -1 : ((distance == o.distance) ? 0 : 1);
        }
    }
}
```

Doesn't re-order  
the priority queue

#8: Learn clever  
optimizations

Competition has been fierce among the cellphone providers in the town of Eastern WestField. Several companies have been running ads in which they each claim to provide more coverage to the town than do any of their competitors.

AetherTech Telecommunications suspects that they, in fact, really do cover a larger portion of the town than do their competitors. They would like to compute their actual coverage so thay they can advertise the true amount.

The metropolitan area of EasternWestField will be modeled as a circle around the town center. The area covered by each cell tower can also be modeled as a circle, centered on the location of the tower, indicating the area within which a the tower provides sufficient signal strength for a phone to connect (e.g., one bar on the average phone's signal strength meter).

All cell towers to be considered will have their centers within the metropolitan area, though their area of coverage may extend beyond the border of the metropolitan area. The areas covered by different towers can overlap, though for economic reasons no two towers have been constructed so close to one another that the center of one would lie within the coverage area of the other.

Compute the fraction of the metropolitan area within which a cell phone would be covered by at least one tower.

## Input

Input consists of multiple data sets. Each data set begins with a line containing a single integer  $N$ ,  $1 \leq N \leq 25$ . This line is followed by  $N$  lines, each containing three floating point numbers  $x$ ,  $y$ ,  $r$ . These give the  $(x,y)$  coordinates of the circle's center and the radius of that circle, respectively. The first circle in the list denotes the metropolitan area of Eastern Westfield. Each of the remaining  $N-1$  circles describes a cell tower.

The final data set is followed by a line containing a zero.

## Output

For each data set, compute the fraction of the metropolitan area covered by the cell towers (as a number in the range  $0.00\dots1.00$  and print a line containing that fraction as a real number presented to 2 decimal points precision.

## Sample Input

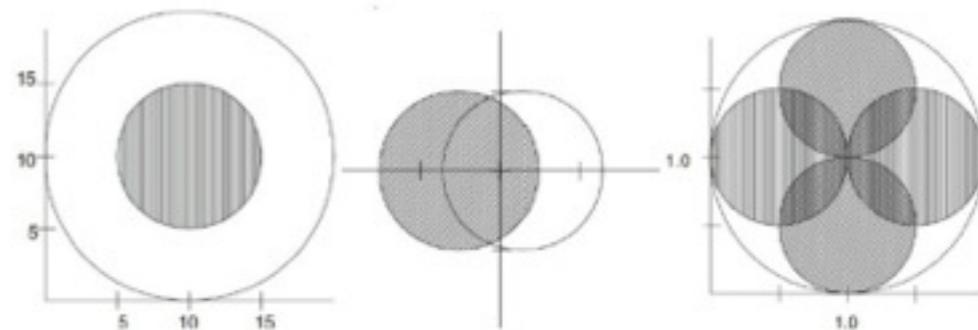
```
2
10 10 10
10 10 5
2
0.404 0.0 1.0
-0.404 0.0 1.0
5
0 0 2
1 0 1
0 1 1
```

```
-1 0 1
0 -1 1
0
```

## Sample Output

```
0.25
0.50
0.82
```

Note: The images corresponding to the sample input are:



Competition has been fierce among the cellphone providers in the town of Eastern WestField. Several companies have been running ads in which they each claim to provide more coverage to the town than do any of their competitors.

AetherTech Telecommunications suspects that they, in fact, really do cover a larger portion of the town than do their competitors. They would like to compute their actual coverage so thay they can advertise the true amount.

The metropolitan area of EasternWestField will be modeled as a circle around the town center. The area covered by each cell tower can also be modeled as a circle, centered on the location of the tower, indicating the area within which a the tower provides sufficient signal strength for a phone to connect (e.g., one bar on the average phone's signal strength meter).

All cell towers to be considered will have their centers within the metropolitan area, though their area of coverage may extend beyond the border of the metropolitan area. The areas covered by different towers can overlap, though for economic reasons no two towers have been constructed so close to one another that the center of one would lie within the coverage area of the other.

Compute the fraction of the metropolitan area within which a cell phone would be covered by at least one tower.

## Input

Input consists of multiple data sets. Each data set begins with a line containing a single integer  $N$ ,  $1 \leq N \leq 25$ . This line is followed by  $N$  lines, each containing three floating point numbers  $x$ ,  $y$ ,  $r$ . These give the  $(x,y)$  coordinates of the circle's center and the radius of that circle, respectively. The first circle in the list denotes the metropolitan area of Eastern Westfield. Each of the remaining  $N-1$  circles describes a cell tower.

The final data set is followed by a line containing a zero.

## Output

For each data set, compute the fraction of the metropolitan area covered by the cell towers (as a number in the range  $0.00\dots1.00$  and print a line containing that fraction as a real number presented to 2 decimal points precision.

## Sample Input

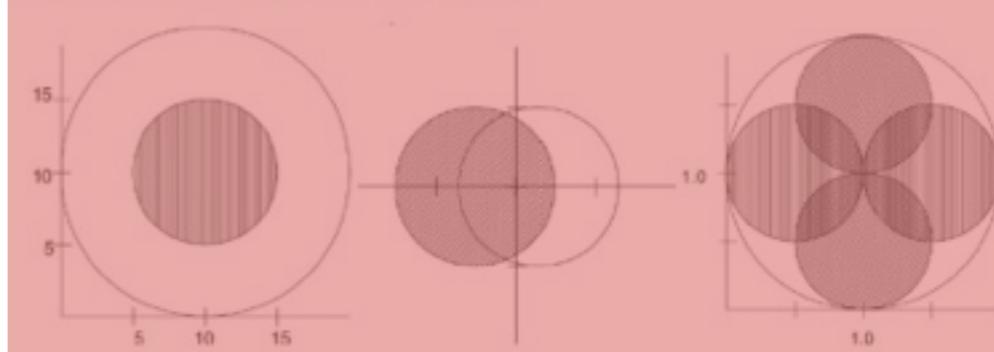
```
2
10 10 10
10 10 5
2
0.404 0.0 1.0
-0.404 0.0 1.0
5
0 0 2
1 0 1
0 1 1
```

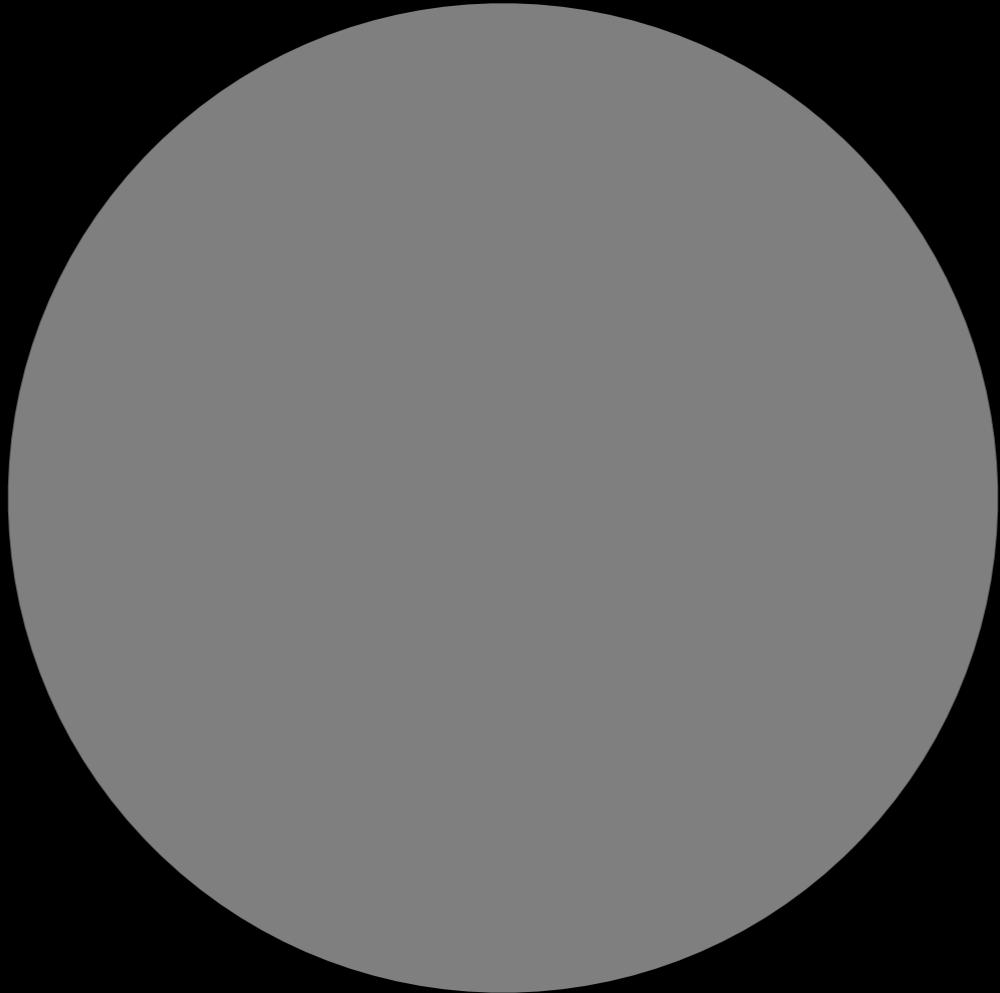
```
-1 0 1
0 -1 1
0
```

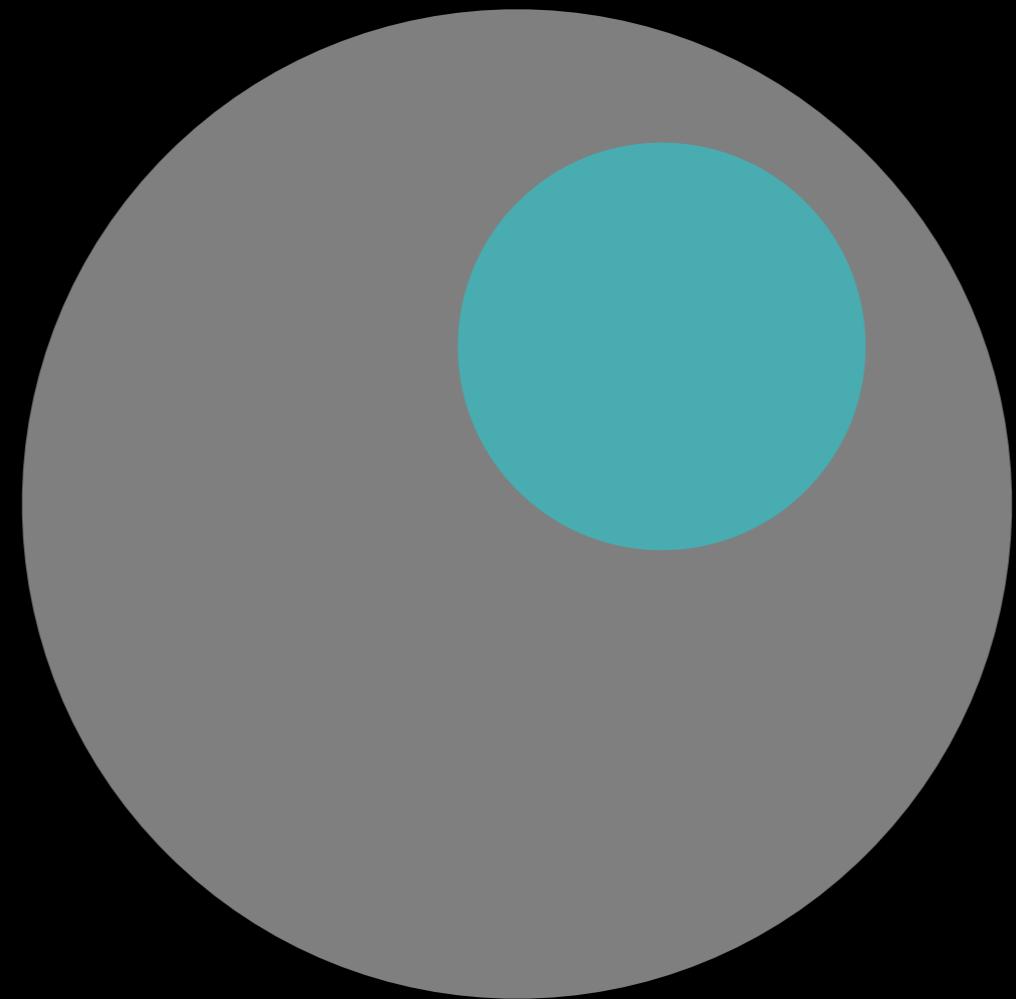
## Sample Output

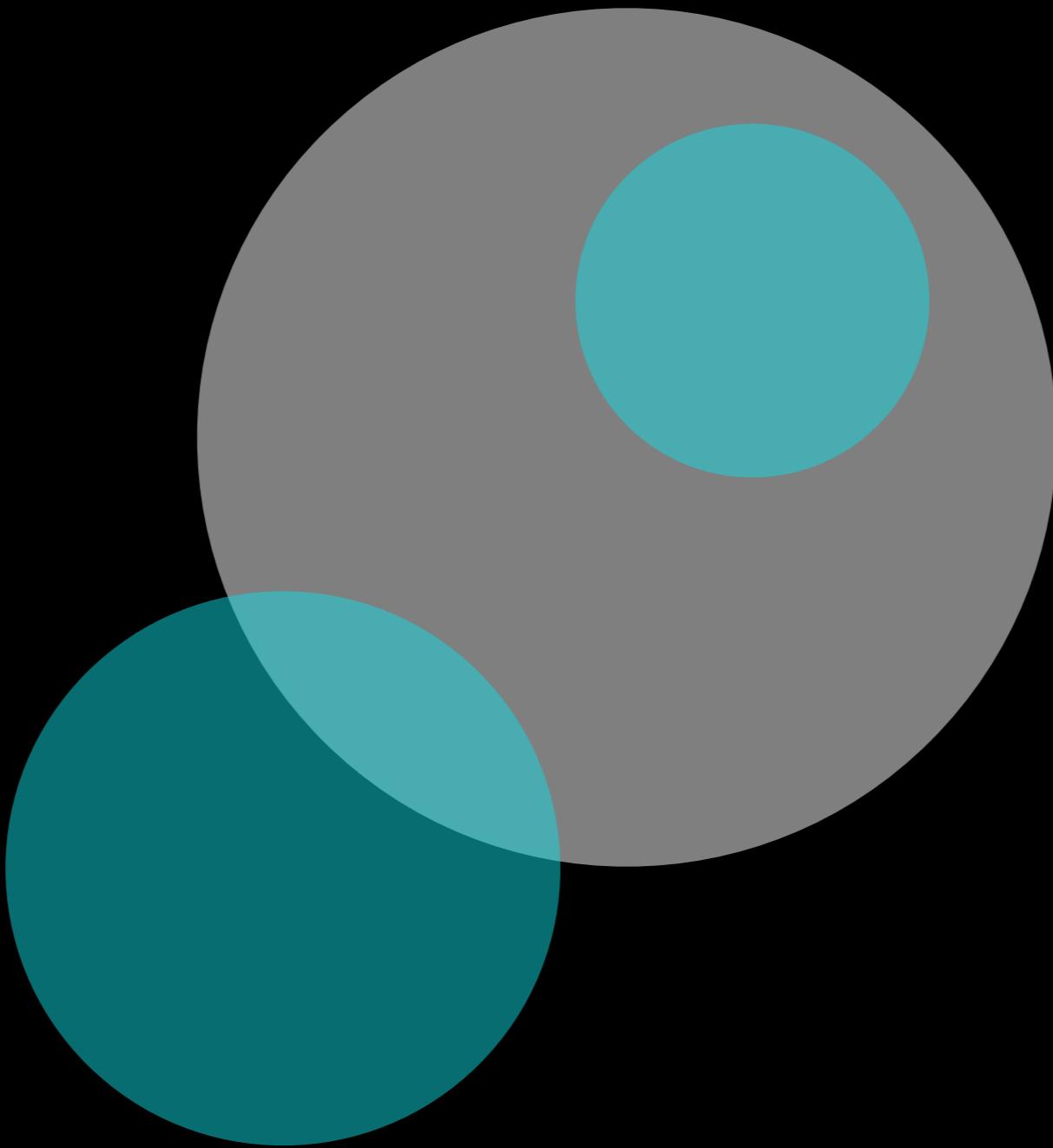
```
0.25
0.50
0.82
```

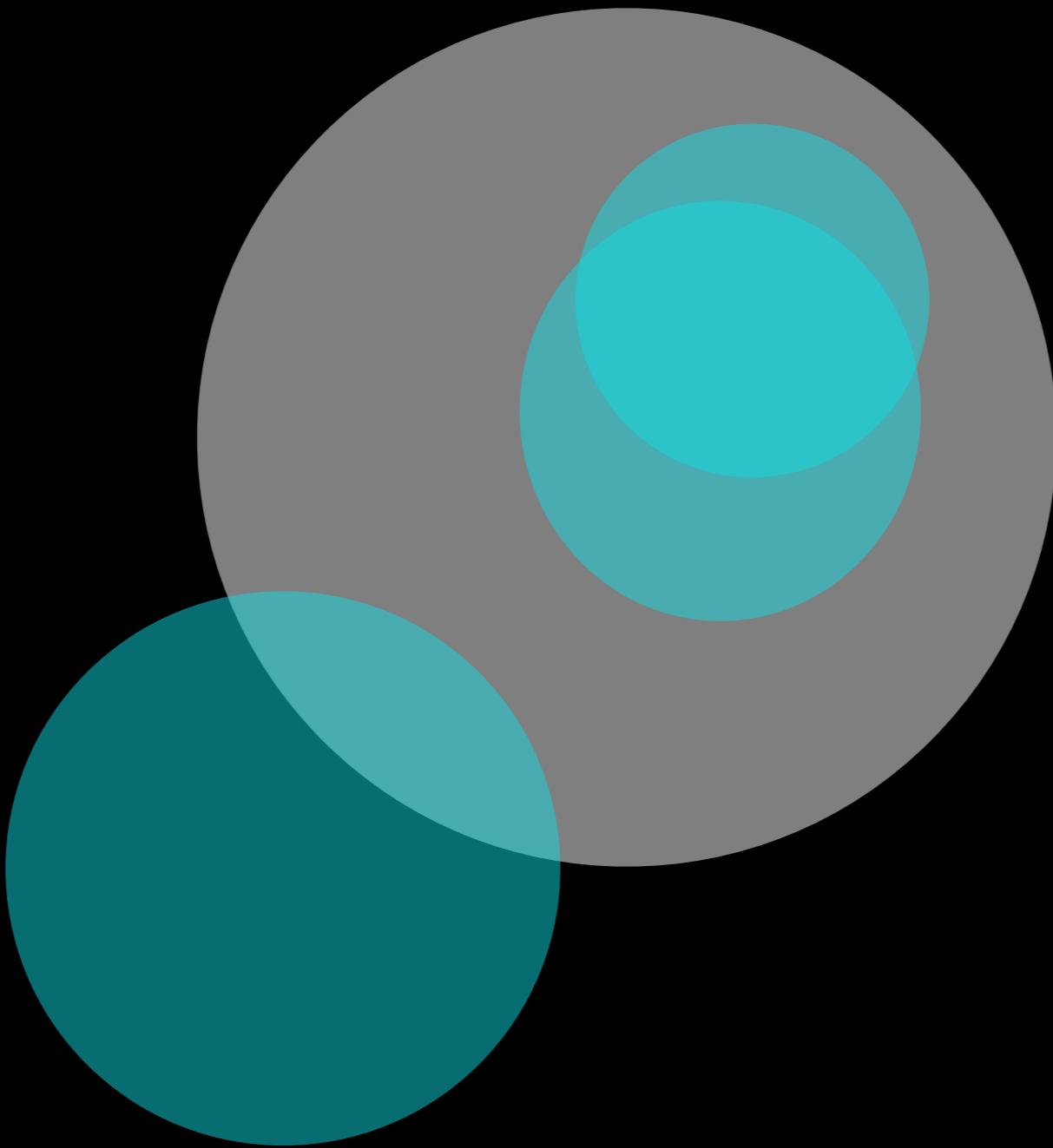
Note: The images corresponding to the sample input are:

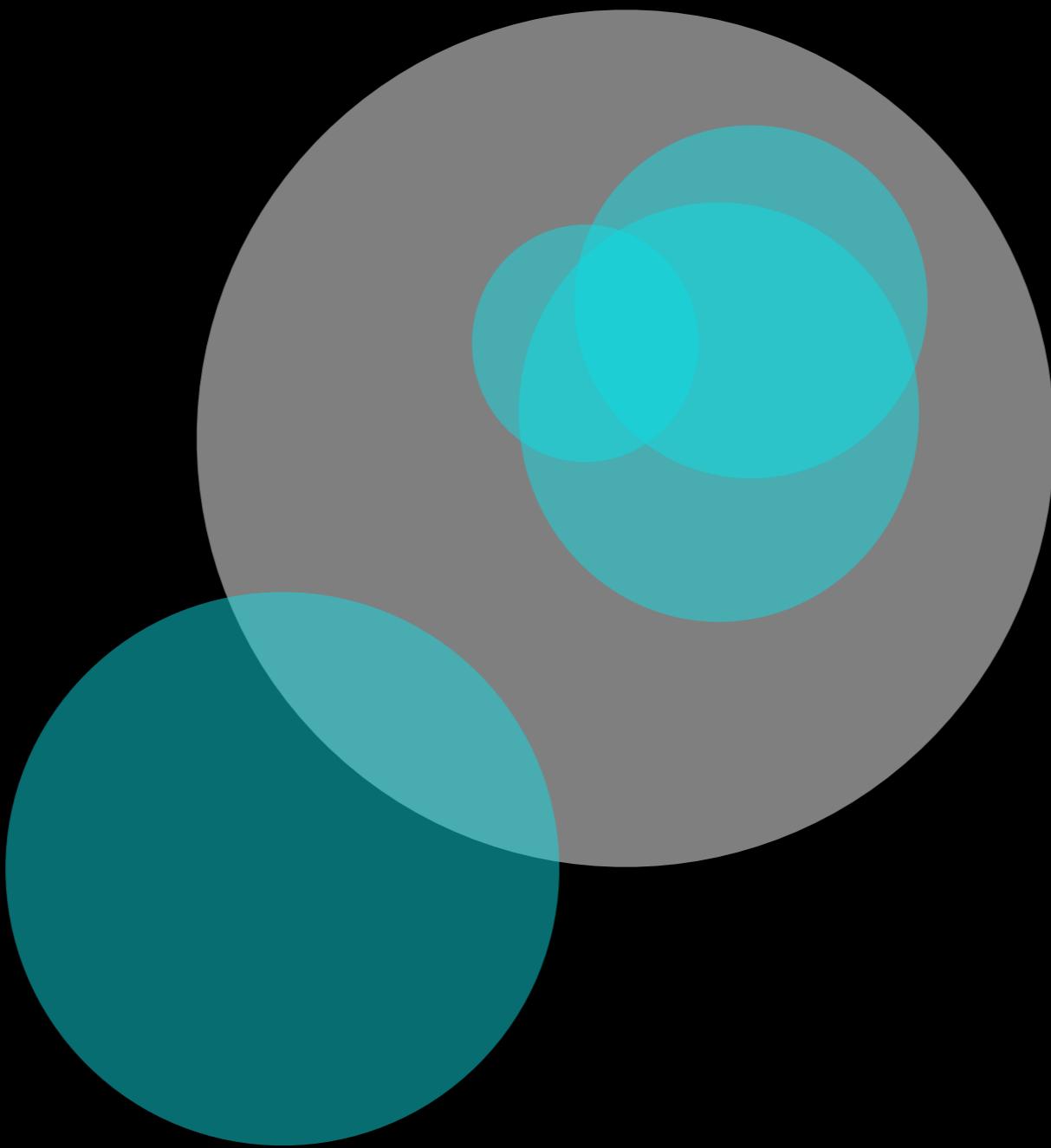












Competition has been fierce among the cellphone providers in the town of Eastern WestField. Several companies have been running ads in which they each claim to provide more coverage to the town than do any of their competitors.

AetherTech Telecommunications suspects that they, in fact, really do cover a larger portion of the town than do their competitors. They would like to compute their actual coverage so thay they can advertise the true amount.

The metropolitan area of EasternWestField will be modeled as a circle around the town center. The area covered by each cell tower can also be modeled as a circle, centered on the location of the tower, indicating the area within which a the tower provides sufficient signal strength for a phone to connect (e.g., one bar on the average phone's signal strength meter).

All cell towers to be considered will have their centers within the metropolitan area, though their area of coverage may extend beyond the border of the metropolitan area. The areas covered by different towers can overlap, though for economic reasons no two towers have been constructed so close to one another that the center of one would lie within the coverage area of the other.

Compute the fraction of the metropolitan area within which a cell phone would be covered by at least one tower.

## Input

Input consists of multiple data sets. Each data set begins with a line containing a single integer  $N$ ,  $1 \leq N \leq 25$ . This line is followed by  $N$  lines, each containing three floating point numbers  $x$ ,  $y$ ,  $r$ . These give the  $(x,y)$  coordinates of the circle's center and the radius of that circle, respectively. The first circle in the list denotes the metropolitan area of Eastern Westfield. Each of the remaining  $N-1$  circles describes a cell tower.

The final data set is followed by a line containing a zero.

## Output

For each data set, compute the fraction of the metropolitan area covered by the cell towers (as a number in the range  $0.00\dots1.00$  and print a line containing that fraction as a real number presented to 2 decimal points precision.

## Sample Input

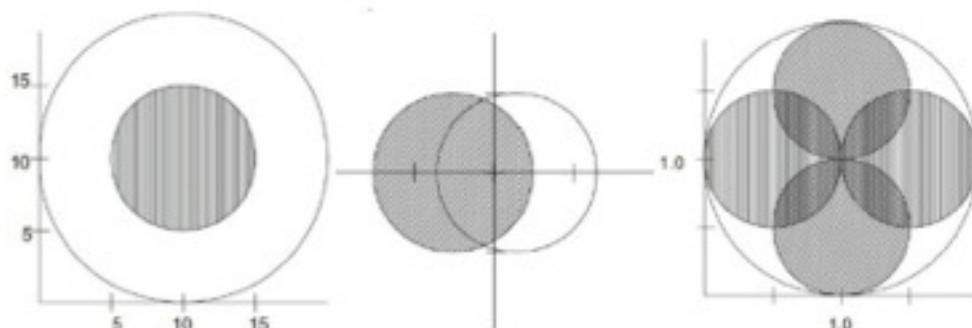
```
2
10 10 10
10 10 5
2
0.404 0.0 1.0
-0.404 0.0 1.0
5
0 0 2
1 0 1
0 1 1
```

```
-1 0 1
0 -1 1
0
```

## Sample Output

```
0.25
0.50
0.82
```

Note: The images corresponding to the sample input are:



Competition has been fierce among the cellphone providers in the town of Eastern WestField. Several companies have been running ads in which they each claim to provide more coverage to the town than do any of their competitors.

AetherTech Telecommunications suspects that they, in fact, really do cover a larger portion of the town than do their competitors. They would like to compute their actual coverage so thay they can advertise the true amount.

The metropolitan area of EasternWestField will be modeled as a circle around the town center. The area covered by each cell tower can also be modeled as a circle, centered on the location of the tower, indicating the area within which a the tower provides sufficient signal strength for a phone to connect (e.g., one bar on the average phone's signal strength meter).

All cell towers to be considered will have their centers within the metropolitan area, though their area of coverage may extend beyond the border of the metropolitan area. The areas covered by different towers can overlap, though for economic reasons no two towers have been constructed so close to one another that the center of one would lie within the coverage area of the other.

Compute the fraction of the metropolitan area within which a cell phone would be covered by at least one tower.

## Input

Input consists of multiple data sets. This line is followed by the coordinates of the centers of the metropolita

For each data set, compute the fraction of the metropolitan area covered by the cell towers (as a number in the range 0.00...1.00 and print a line containing that fraction as a real number presented to 2 decimal points precision.

The final data set is followed by a line containing a zero.

## Output

For each data set, compute the fraction of the metropolitan area covered by the cell towers (as a number in the range 0.00...1.00 and print a line containing that fraction as a real number presented to 2 decimal points precision.

## Sample Input

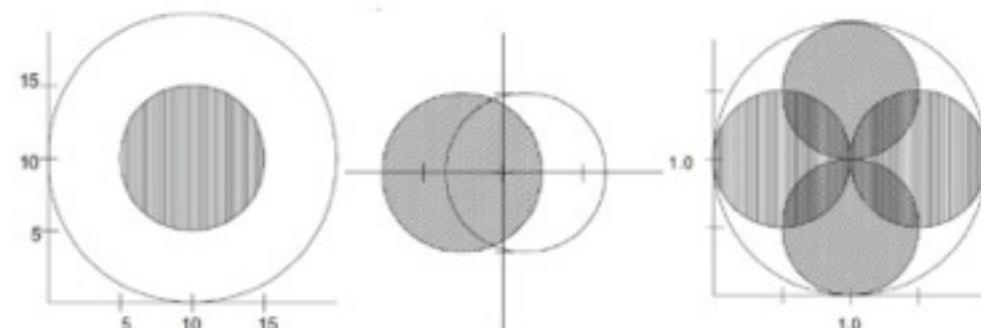
```
2
10 10 10
10 10 5
2
0.404 0.0 1.0
-0.404 0.0 1.0
5
0 0 2
1 0 1
0 1 1
```

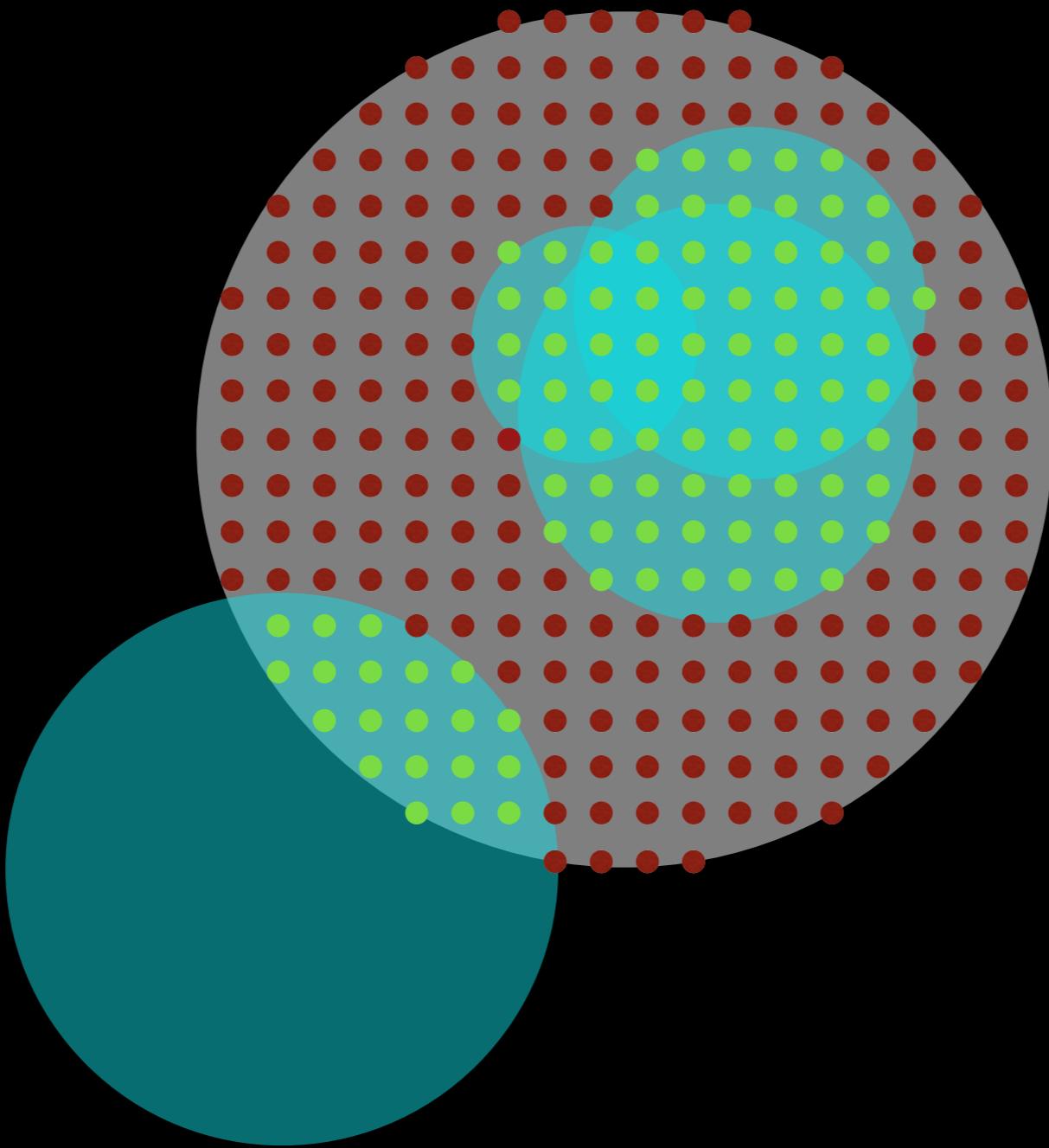
```
-1 0 1
0 -1 1
0
```

## Sample Output

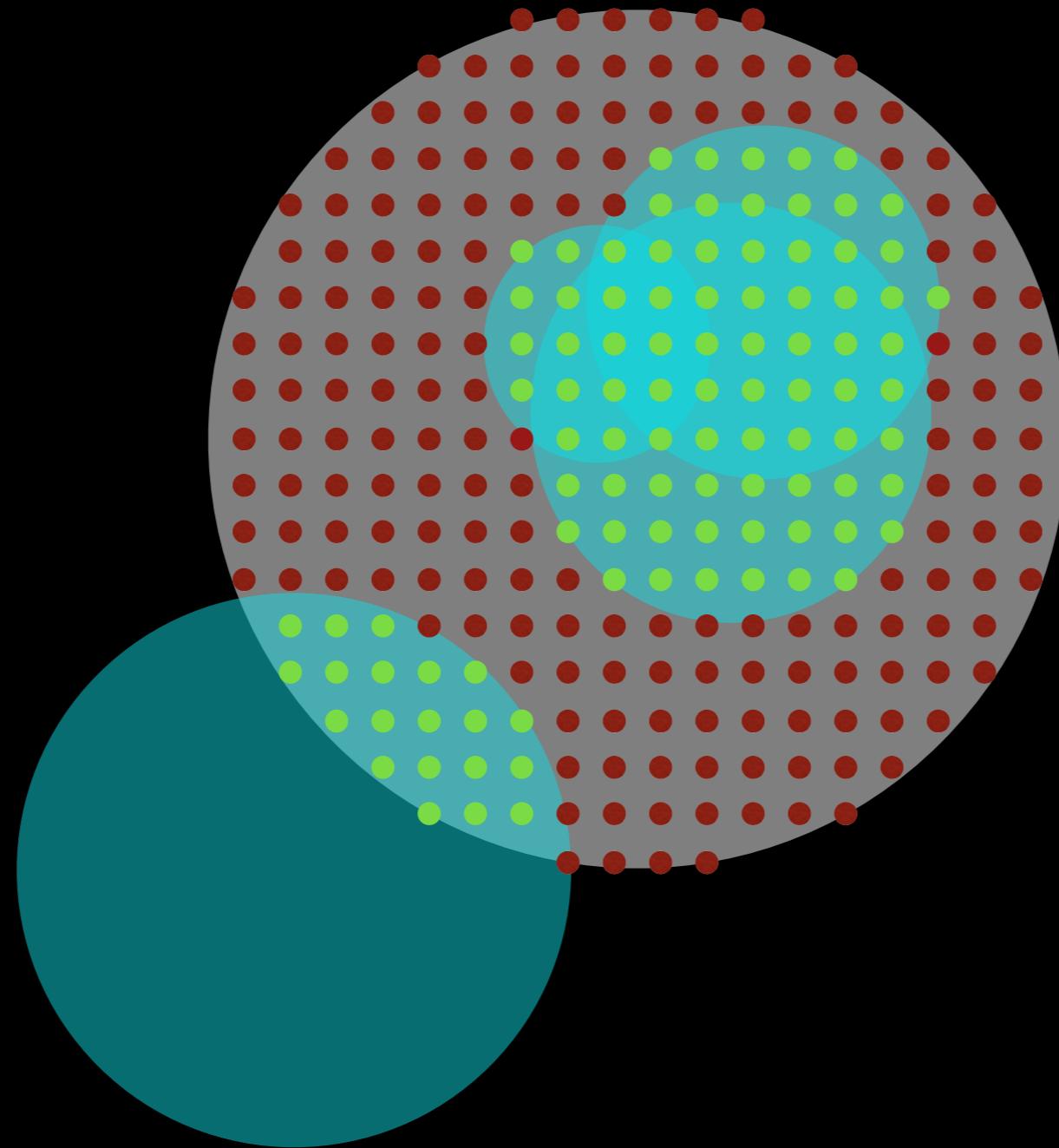
```
0.25
0.50
0.82
```

Note: The images corresponding to the sample input are:





$500 * 500 = 250,000$  random samples  
worst case check every circle,  $N \leq 25$   
 $250,000 * 25 = 6,250,000 \rightarrow$  Manageable!



```

import java.util.*;
import java.io.*;

public class cells {
    public static final double SAMPLE = 1000;

    public static void main(String[] args) {
        Scanner in = new Scanner(System.in);
        while(true) {
            int N = in.nextInt();
            if(N==0)
                break;
            double X = in.nextDouble(), Y = in.nextDouble(), R = in.nextDouble();
            double[] x = new double[N-1], y = new double[N-1], r = new double[N-1];
            for(int i=0;i<N-1;i++) {
                x[i] = in.nextDouble();
                y[i] = in.nextDouble();
                r[i] = in.nextDouble();
            }
            int inRange = 0;
            int covered = 0;
            for(double a = X-R; a<=X+R; a+= R/SAMPLE)
                for(double b = Y-R; b<=Y+R; b+= R/SAMPLE) {
                    if(distSQ(X, Y, a, b) <= R*R) {
                        inRange++;
                        for(int i=0;i<N-1;i++) {
                            if(distSQ(x[i],y[i], a, b) <= r[i]*r[i]) {
                                covered++;
                                break;
                            }
                        }
                    }
                }
            System.out.printf("%.02f\n", ((double)covered)/inRange);
        }
    }

    public static double distSQ(double x1, double y1, double x2, double y2) {
        double dx = x1-x2, dy=y1-y2;
        return dx*dx+dy*dy;
    }
}

```

1 try, <10 minutes

#9: Makes you  
appreciate algorithms,  
data structures

## Problem J

### Oil Skimming

Thanks to a certain "green" resources company, there is a new profitable industry of oil skimming. There are large slicks of crude oil floating in the Gulf of Mexico just waiting to be scooped up by enterprising oil barons. One such oil baron has a special plane that can skim the surface of the water collecting oil on the water's surface. However, each scoop covers a 10m by 20m rectangle (going either east/west or north/south). It also requires that the rectangle be completely covered in oil, otherwise the product is contaminated by pure ocean water and thus unprofitable!

Given a map of an oil slick, the oil baron would like you to compute the maximum number of scoops that may be extracted. The map is an NxN grid where each cell represents a 10m square of water, and each cell is marked as either being covered in oil or pure water.

#### **Input**

The input starts with an integer  $K$  ( $1 \leq K \leq 100$ ) indicating the number of cases. Each case starts with an integer  $N$  ( $1 \leq N \leq 600$ ) indicating the size of the square grid. Each of the following  $N$  lines contains  $N$  characters that represent the cells of a row in the grid. A character of '#' represents an oily cell, and a character of '.' represents a pure water cell.

#### **Output**

For each case, one line should be produced, formatted exactly as follows: "Case X: M" where X is the case number (starting from 1) and M is the maximum number of scoops of oil that may be extracted.

| <b>Sample Input</b>  | <b>Output for the Sample Input</b> |
|--|------------------------------------|
| 1<br>6<br>.....<br>.##...<br>.##...<br>....#.<br>....##<br>..... | Case 1: 3                          |

## Problem J

### Oil Skimming

Thanks to a certain "green" resources company, there is a new profitable industry of oil skimming. There are large slicks of crude oil floating in the Gulf of Mexico just waiting to be scooped up by enterprising oil barons. One such oil baron has a special plane that can skim the surface of the water collecting oil on the water's surface. However, each scoop covers a 10m by 20m rectangle (going either east/west or north/south). It also requires that the rectangle be completely covered in oil, otherwise the product is contaminated by pure ocean water and thus unprofitable!

Given a map of an oil slick, the oil baron would like you to compute the maximum number of scoops that may be extracted. The map is an NxN grid where each cell represents a 10m square of water, and each cell is marked as either being covered in oil or pure water.

#### Input

The input starts with an integer  $K$  ( $1 \leq K \leq 100$ ) indicating the number of cases. Each case starts with an integer  $N$  ( $1 \leq N \leq 600$ ) indicating the size of the square grid. Each of the following  $N$  lines contains  $N$  characters that represent the cells of a row in the grid. A character of '#' represents an oily cell, and a character of '.' represents a pure water cell.

#### Output

For each case, one line should be produced, formatted exactly as follows: "Case X: M" where X is the case number (starting from 1) and M is the maximum number of scoops of oil that may be extracted.

| Sample Input   | Output for the Sample Input |
|--|-----------------------------|
| 1<br>6<br>.....<br>.##...<br>.##...<br>....#.<br>....##<br>..... | Case 1: 3                   |

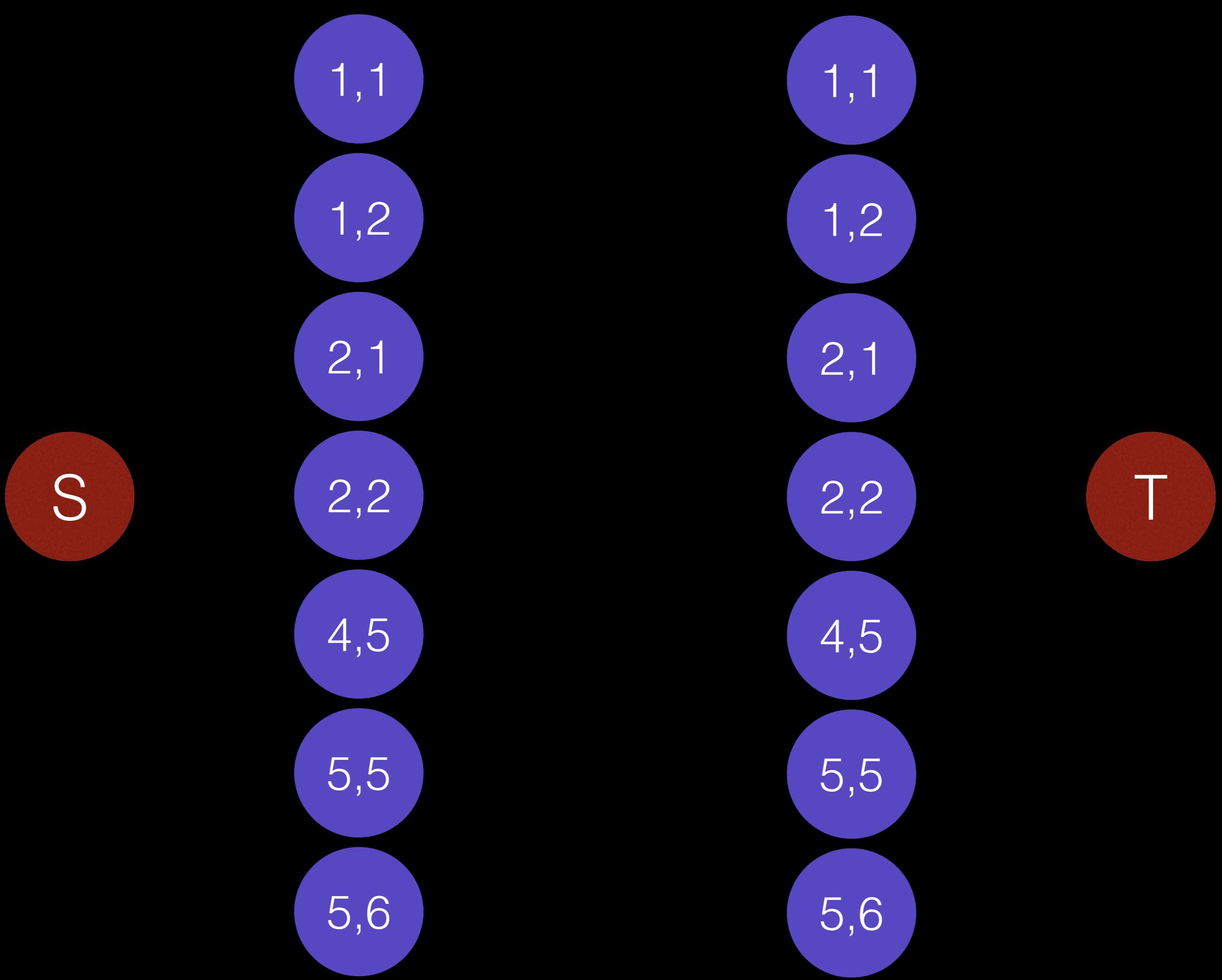
|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| . | . | . | . | . | . | . |
| . | # | # | . | . | . | . |
| . | # | # | . | . | . | . |
| . | . | . | . | # | . | . |
| . | . | . | . | # | # | . |
| . | . | . | . | . | . | . |

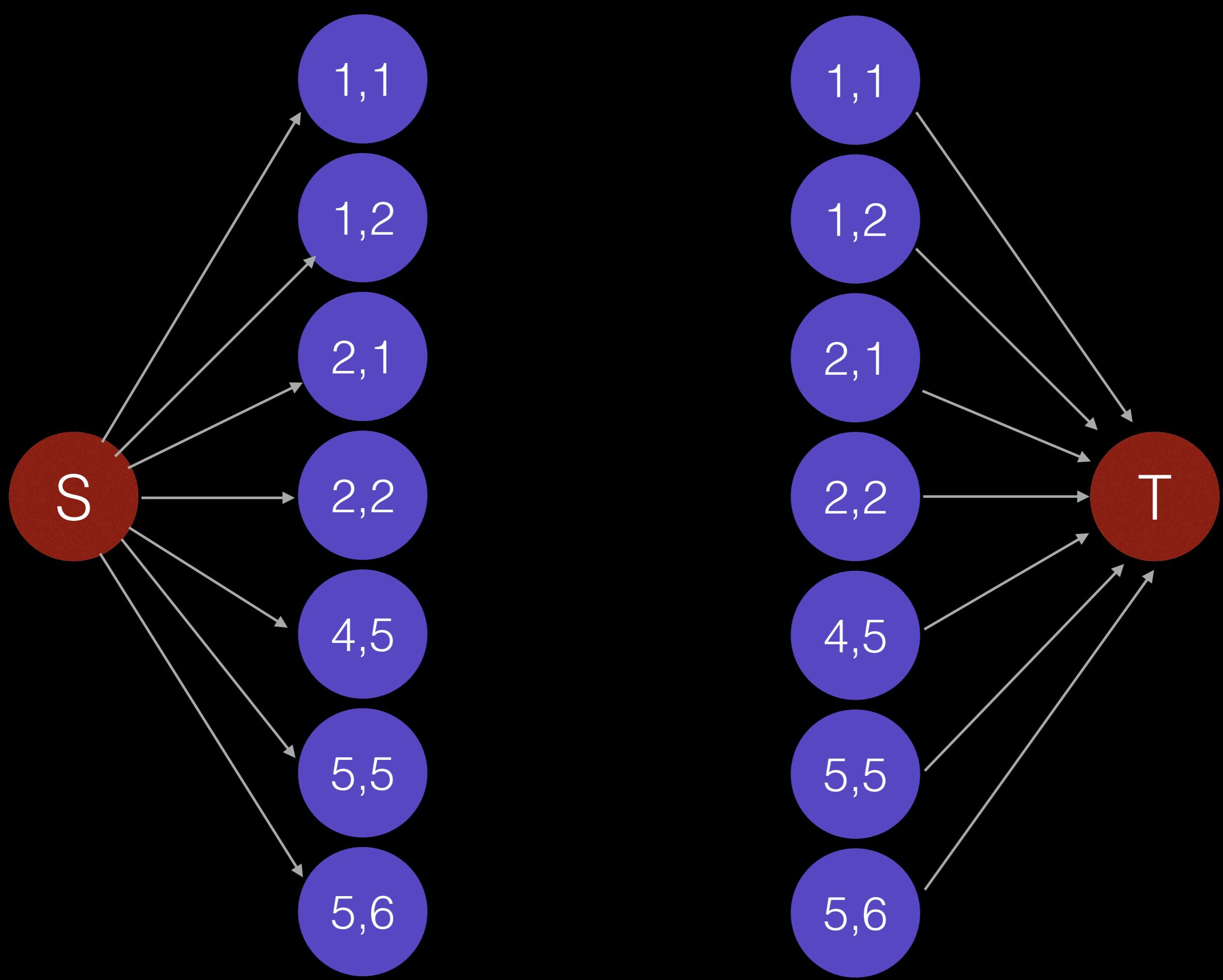
• • • • • •  
• # # • • •  
• # # • • •  
• • • • # •  
• • • • # #  
• • • • • •

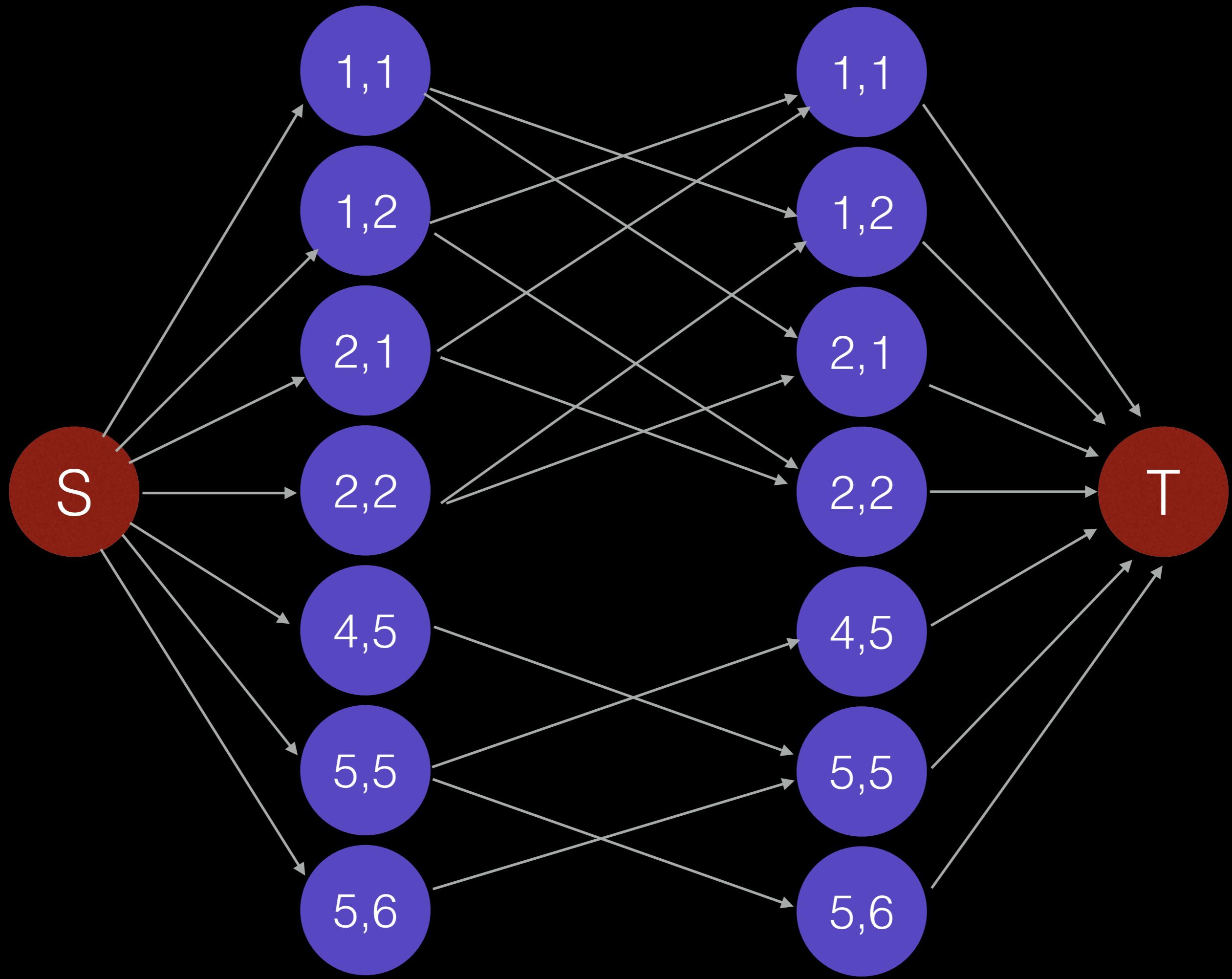
This is actually a max  
flow problem!

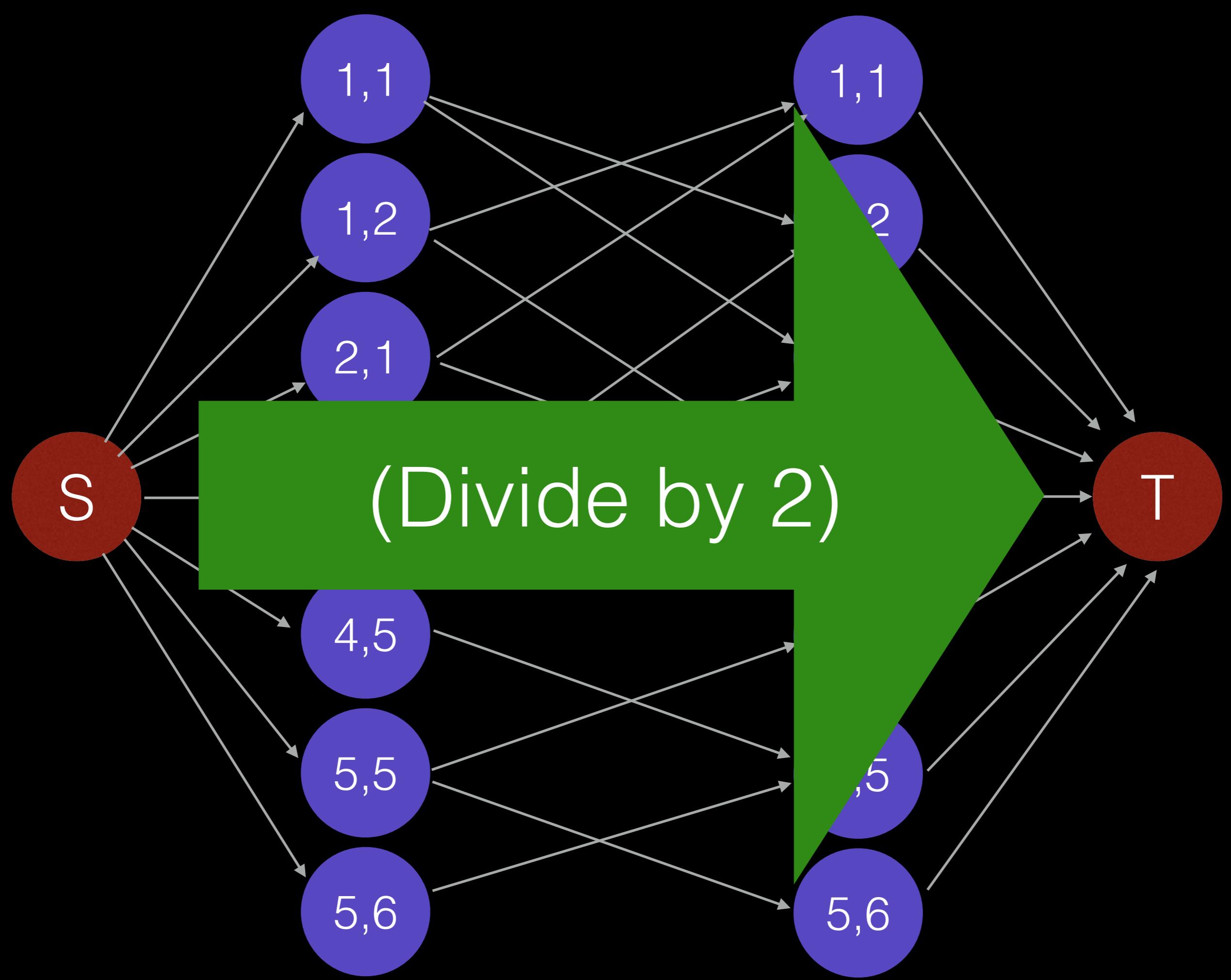
S

T









```

import java.util.*;
import java.io.*;

public class oil {

    static HashMap<String, Node> nodes = new HashMap<String, Node>();

    public static void main(String[] args) {
        Scanner in = new Scanner(System.in);

        int numCases = in.nextInt();
        for (int c = 1; c <= numCases; c++) {
            nodes.clear();

            int N = in.nextInt();
            boolean[][] grid = new boolean[N][N];
            for (int i = 0; i < N; i++) {
                String s = in.next();
                for (int j = 0; j < N; j++)
                    grid[i][j] = s.charAt(j)=='#';
            }

            Node source = new Node();
            Node sink = new Node();
            source.isSourceOrSink = true;
            sink.isSourceOrSink = true;
            int numNodes = 2;

            Node[][] aNodes = new Node[N][N];
            Node[][] bNodes = new Node[N][N];
            for (int i = 0; i < N; i++) {
                for (int j = 0; j < N; j++) {
                    if (grid[i][j]) {
                        Node n = new Node();
                        new Edge(n, source, 1, false);
                        aNodes[i][j] = n;
                        n = new Node();
                        new Edge(sink, n, 1, false);
                        bNodes[i][j] = n;
                        numNodes += 2;
                    }
                }
            }
            source.height = numNodes;
            source.excess = Integer.MAX_VALUE;
            for (int i = 0; i < N; i++) {
                for (int j = 0; j < N; j++) {
                    if (grid[i][j]) {
                        if (i - 1 >= 0 && grid[i - 1][j])
                            new Edge(bNodes[i-1][j], aNodes[i][j], 1, false);
                        if (j - 1 >= 0 && grid[i][j - 1])
                            new Edge(bNodes[i][j-1], aNodes[i][j], 1, false);
                        if (i + 1 < N && grid[i + 1][j])
                            new Edge(bNodes[i+1][j], aNodes[i][j], 1, false);
                        if (j + 1 < N && grid[i][j + 1])
                            new Edge(bNodes[i][j+1], aNodes[i][j], 1, false);
                    }
                }
            }
            PreflowPush(source, sink);
            System.out.printf("Case %d: %d\n", c, sink.excess/2);
        }
    }

    public static void PreflowPush(Node source, Node sink) {
        // source.height should be the number of nodes, sink.height should be 1
        LinkedList<Node> q = new LinkedList<Node>();
        q.add(source);
        while (!q.isEmpty()) {
            Node node = q.remove();
            int minHeight = Integer.MAX_VALUE;
            for (int i = 0; i < node.neighbors.size(); i++) {
                if (node.neighbors.get(i).dest.height < minHeight
                    && node.neighbors.get(i).remaining() > 0)
                    minHeight = node.neighbors.get(i).dest.height;
            }
            if (minHeight != Integer.MAX_VALUE && minHeight >= node.height)
                node.height = minHeight + 1;
            for (int i = 0; i < node.neighbors.size(); i++) {
                if (node.neighbors.get(i).dest.height < node.height) {
                    int pushedFlow = node.neighbors.get(i).remaining();
                    if (pushedFlow > node.excess)
                        pushedFlow = node.excess;
                    if (pushedFlow > 0) {
                        node.neighbors.get(i).flow += pushedFlow;
                        node.neighbors.get(i).back.flow -= pushedFlow;
                        node.neighbors.get(i).dest.excess += pushedFlow;
                        node.excess -= pushedFlow;
                        if (!node.neighbors.get(i).dest.isSourceOrSink)
                            q.add(node.neighbors.get(i).dest);
                        if (node.excess <= 0)
                            break;
                    }
                }
            }
            if (node.excess > 0 && !node.isSourceOrSink)
                q.add(node);
        }
    }

    class Node {
        public ArrayList<Edge> neighbors = new ArrayList<Edge>();
        public int height = 1;
        public int excess = 0;
        public boolean isSourceOrSink = false;
    }

    class Edge {
        public Node dest;
        public Edge back;
        public int capacity;
        public int flow = 0;

        public Edge(Node dest, Node source, int capacity, boolean isBack) {
            this.dest = dest;
            this.capacity = capacity;
            source.neighbors.add(this);
            if (!isBack) {
                Edge backEdge = new Edge(source, this.dest, 0, true);
                this.back = backEdge;
                backEdge.back = this;
            }
        }

        public int remaining() {
            return capacity - flow;
        }
    }
}

```

1 try, ~1 hour (copying/understanding max flow code)

#10: Hang out with  
friends and write bad  
code



```

public void calcPoints() {
    points[0] = dc[1]*1;
    points[1] = dc[2]*2;
    points[2] = dc[3]*3;
    points[3] = dc[4]*4;
    points[4] = dc[5]*5;
    points[5] = dc[6]*6;
    points[6] = sum;
    points[7] = isXOfAKind(3) ? sum : 0;
    points[8] = isXOfAKind(4) ? sum : 0;
    points[9] = isXOfAKind(5) ? 50 : 0;
    points[10] = straightOfLength(4) ? 25 : 0;
    points[11] = straightOfLength(5) ? 35 : 0;
    points[12] = isFullHouse() ? 40 : 0;
}
private boolean isXOfAKind(int x) {
    for(int n=0;n<7;n++)
        if(dc[n]>=x)
            return true;
    return false;
}
private boolean straightOfLength(int len) {
    if(len==4) {
        return (dc[1]>0 && dc[2]>0 && dc[3]>0 && dc[4]>0) ||
               (dc[2]>0 && dc[3]>0 && dc[4]>0 && dc[5]>0) ||
               (dc[3]>0 && dc[4]>0 && dc[5]>0 && dc[6]>0);
    }
    if(len==5) {
        return (dc[1]>0 && dc[2]>0 && dc[3]>0 && dc[4]>0 && dc[5]>0) ||
               (dc[2]>0 && dc[3]>0 && dc[4]>0 && dc[5]>0 && dc[6]>0);
    }
    return false;
}
private boolean isFullHouse() {
    return (dice[0] == dice[1] && dice[2] == dice[4]) ||
           (dice[0] == dice[2] && dice[3] == dice[4]);
}
}

```

How can I help you?

# How can I help you?

Daniel Epstein

[depstein@cs.washington.edu](mailto:depstein@cs.washington.edu)

<https://github.com/depstein/programming-competitions>