B-Link trees

Concurrency

- is hard
- Use latches (short-term "locks"), not strict-2PL locks
- Readers go fast: no latch acquisition, no waiting
- Implication: writers must leave tree in valid state at any point - readers don't respect latches!

Concurrency

- Processes go down and to the right
- Insertions expand the tree "rightward" from point of insertion
 - Might also expand parents "rightward" if recursive calls required

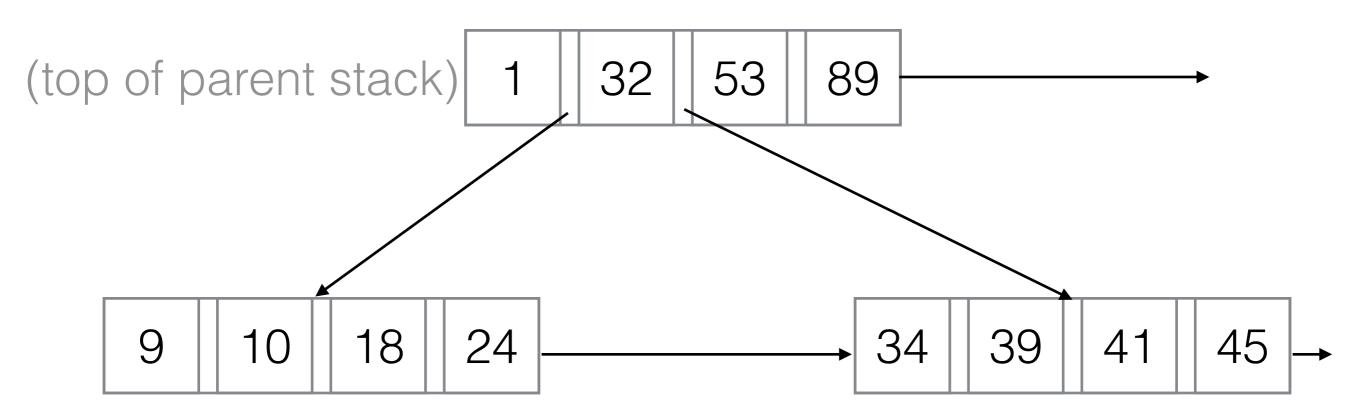
Reads

- Standard B-tree recursive read, but also check high-key
- while not leaf: check node
 - search item > high key? go right
 - otherwise, pick pointer as normal and go down

Writes

- No split same as normal insert (latch and insert to leaf)
- Splits can be recursive might need to set many latches!
 - Want to minimize # latches held at any point (will achieve maximum of 3)
 - Want to avoid holding latches during I/O, if possible

0. Search to get to the correct leaf node, keeping stack of parents



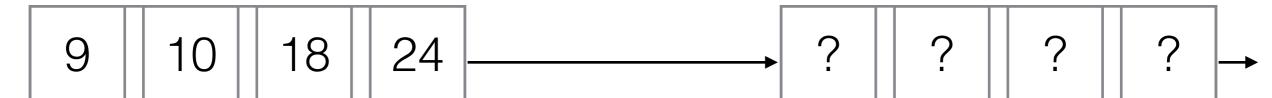
note: high keys not shown

Insert 15

1. Acquire latch on node to be split.

(top of parent stack) ? ? ?

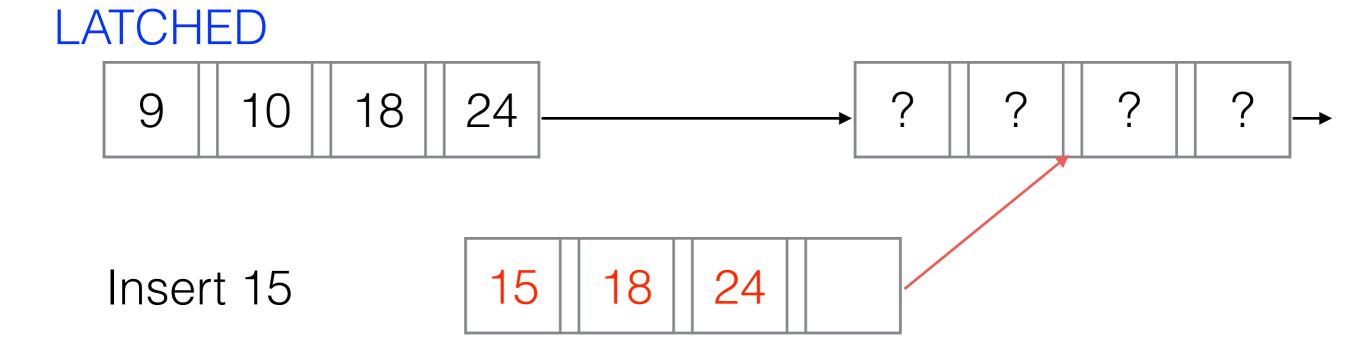
LATCHED



Insert 15

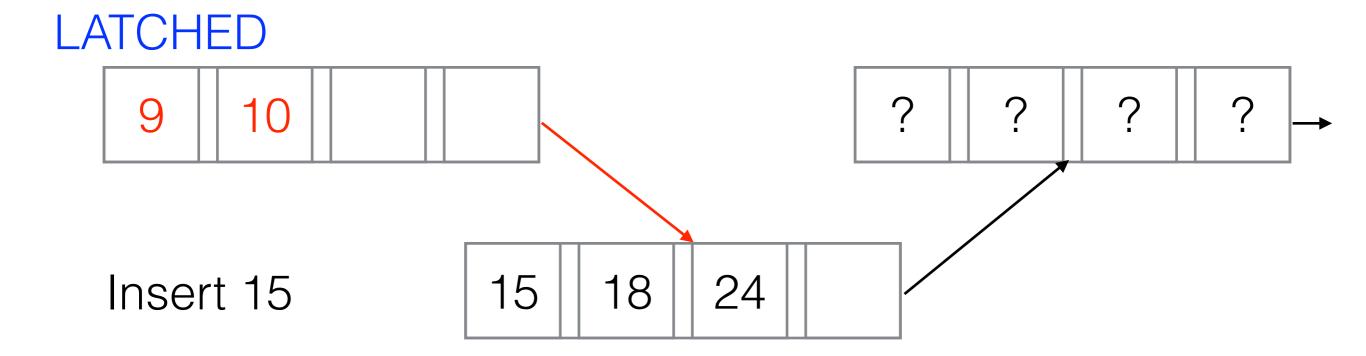
2. Create and write new twin node. Note that it is not yet visible! Nothing points to it.

(top of parent stack) ? ? ?



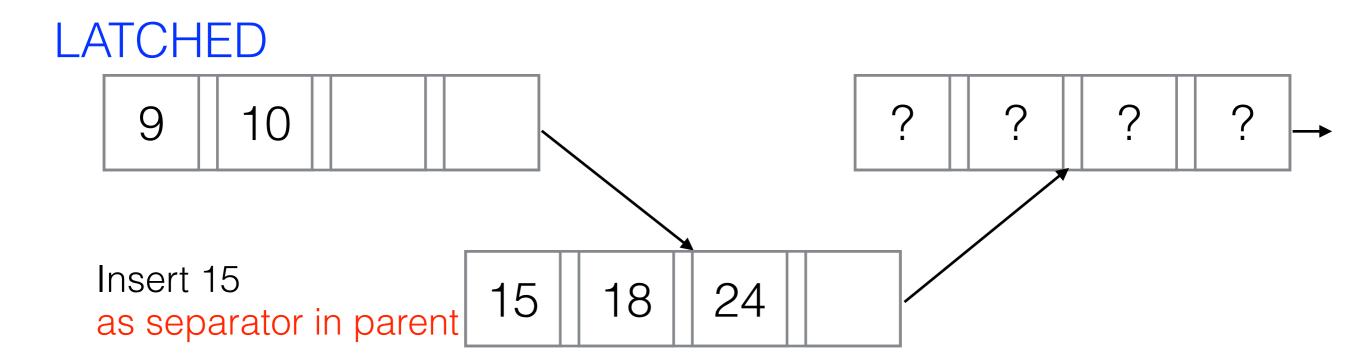
3. Write updated original twin node.

(top of parent stack) ? ? ?



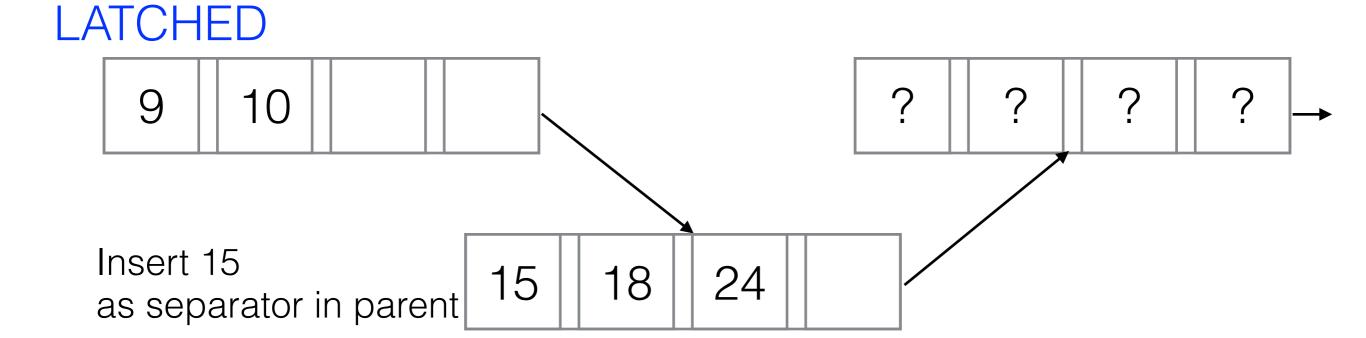
At this point, tree is "consistent", but if we stop now it we eventually degenerate into a linked list - need to add pointer from parent to new node!

(top of parent stack) ? ? ?

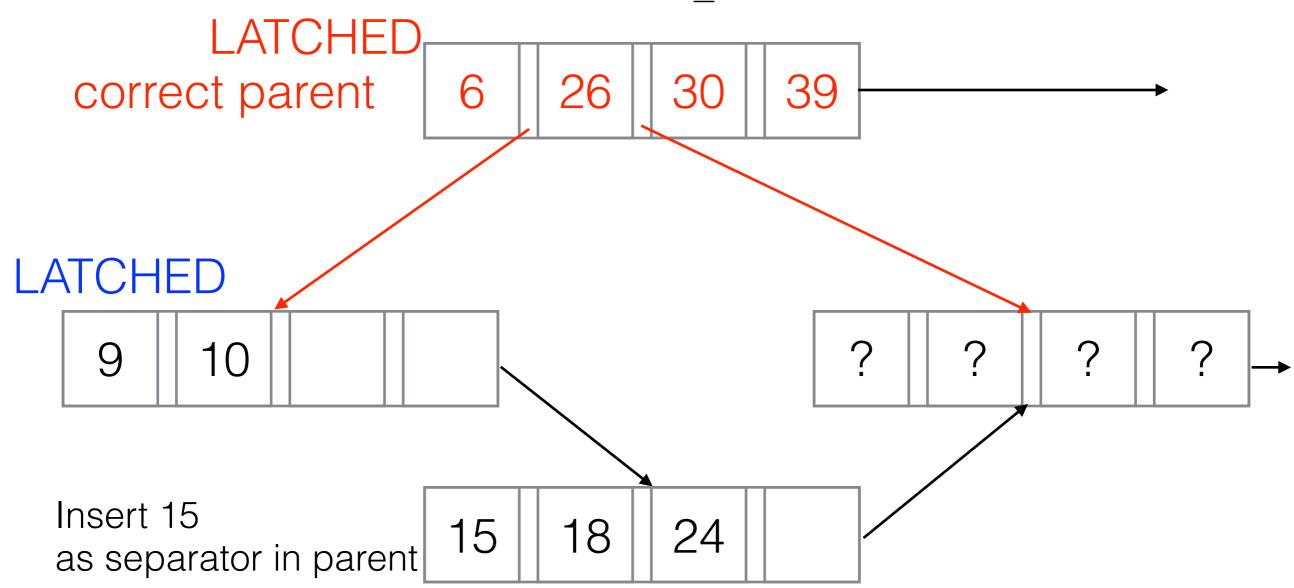


4. Re-read parent (found by popping off from our stack). Parent might have changed! Want *correct* parent pointing to our node.

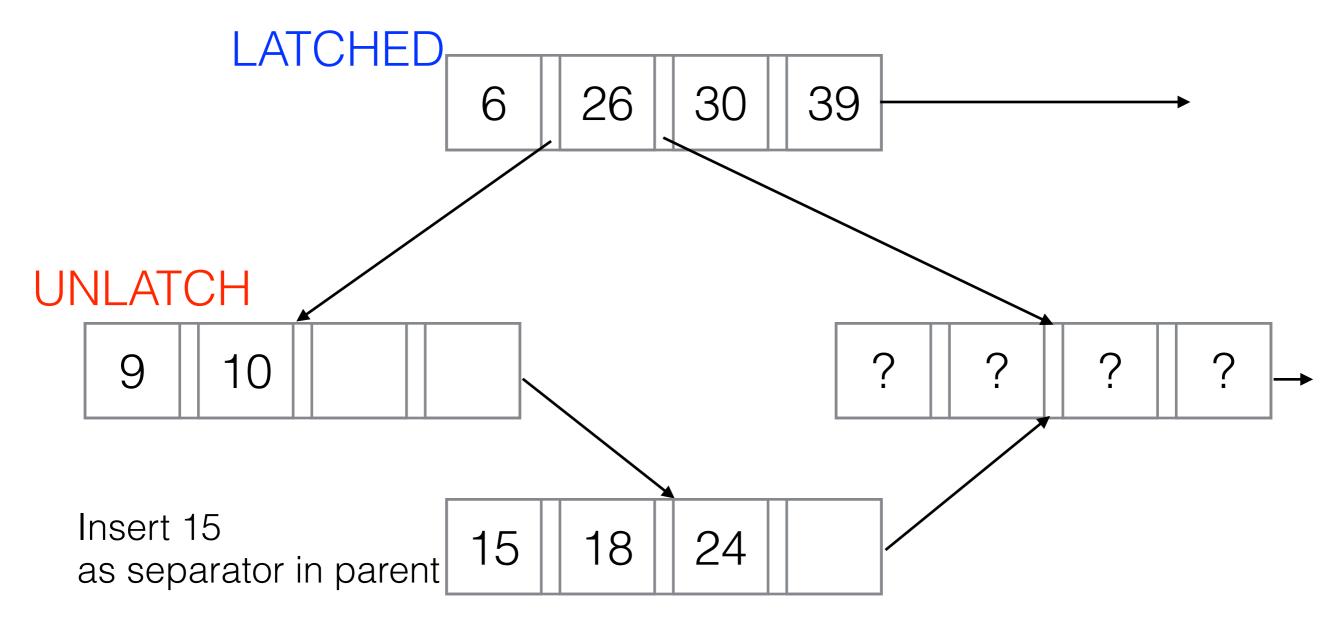




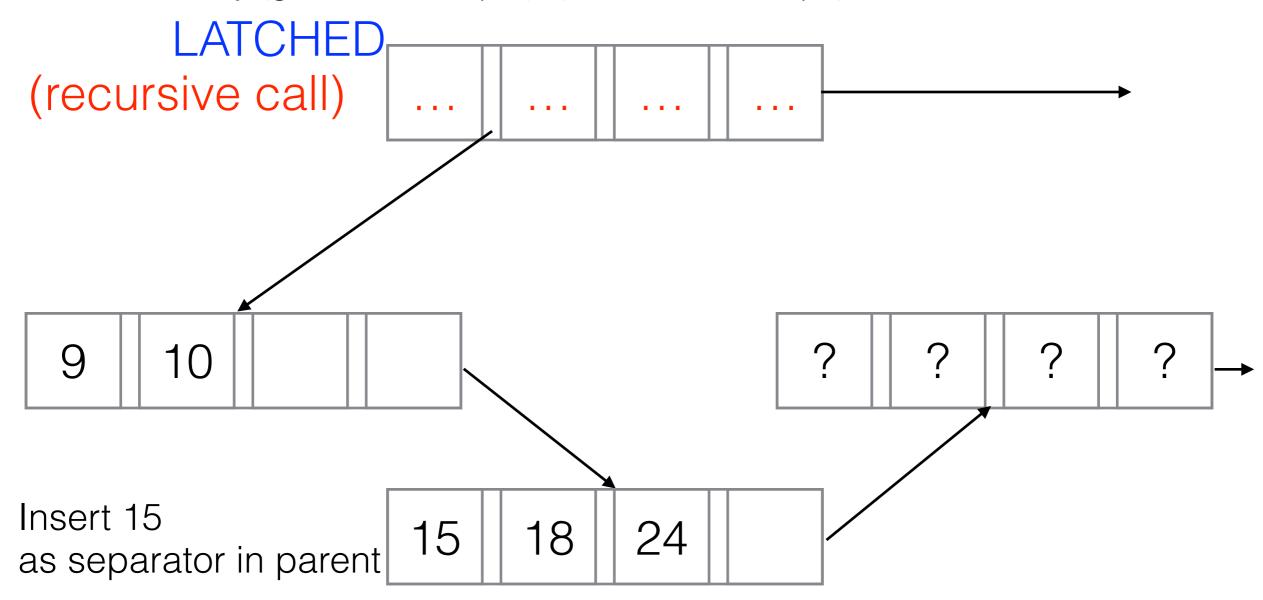
4a. While we haven't found parent pointing to us: move right. Acquire and release latches node by node. (cf. move_right() below for details.)



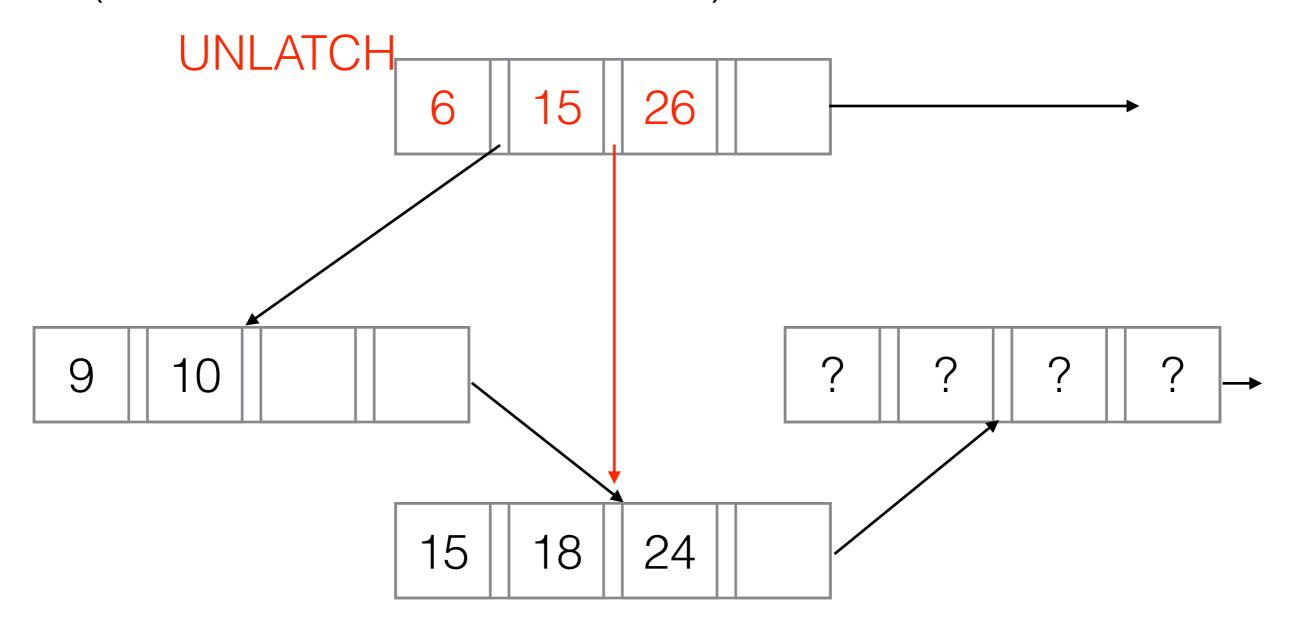
5. We can finally release our latch (why?).



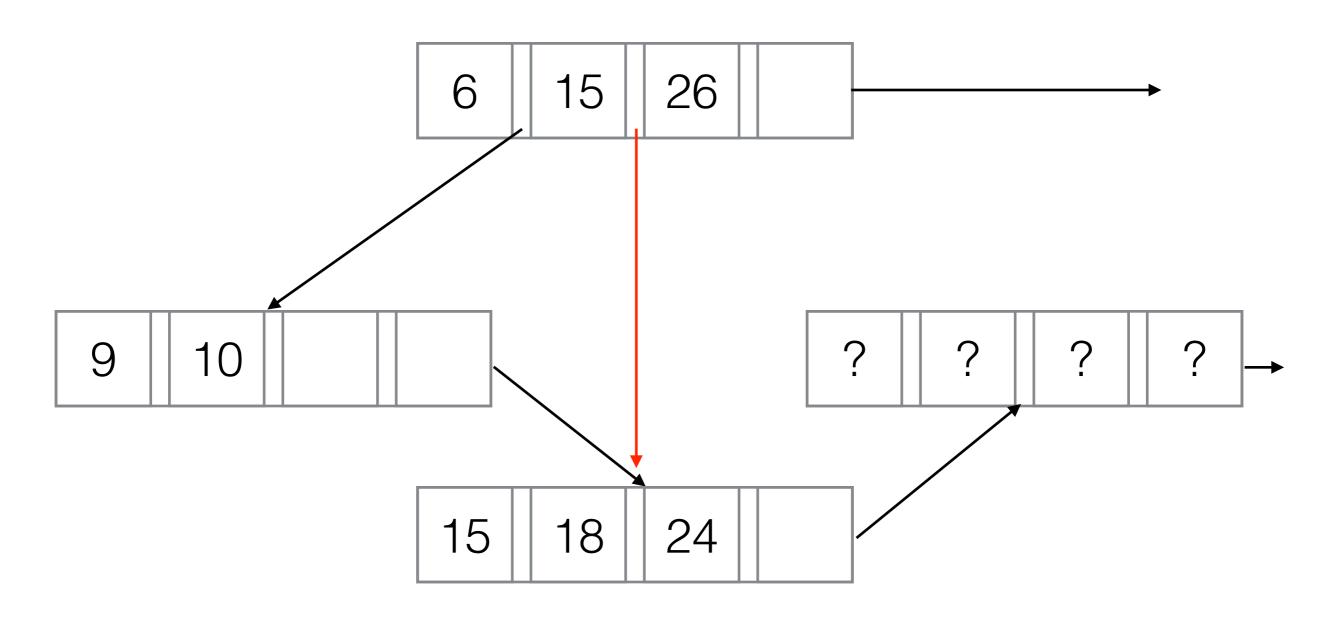
6. Insert separator (along with pointer) into parent. **May need to split that, too!** Call recursively (go back to step 1). (When do we stop?)



6. (Returned from recursive call)

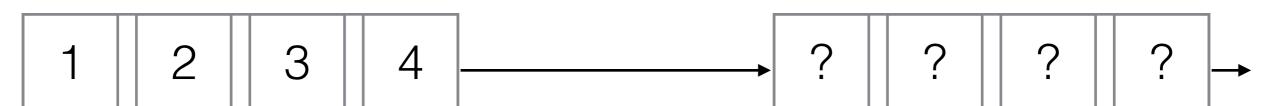


We're done (finally).



0. Acquire latch on node.

LATCHED



1. Latch right-pointer's node.



2. Unlatch old node.



3. Done? If not, go back to step 1.

