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Ambiguity \Rightarrow More than one parse tree

• An unambiguous BNF for $\langle E \rangle$

- Introduce precedence = * and / have higher precedence than + and -
- Introduce associativity = at the same precedence level, the operators associate to the right

$\langle E \rangle \Rightarrow \langle \text{term} \rangle \mid \langle \text{term} \rangle + \langle E \rangle \mid \langle \text{term} \rangle - \langle E \rangle$

$\langle E \rangle \Rightarrow \langle \text{term} \rangle + \langle E \rangle \Rightarrow \langle \text{term} \rangle + \langle \text{term} \rangle - \langle E \rangle \Rightarrow \langle \text{term} \rangle + \langle \text{term} \rangle - \langle \text{term} \rangle + \langle E \rangle$
 $\Rightarrow \dots$ terminate by $\langle E \rangle \Rightarrow \langle \text{term} \rangle$

Any expression is a sequence of terms separated by + / -

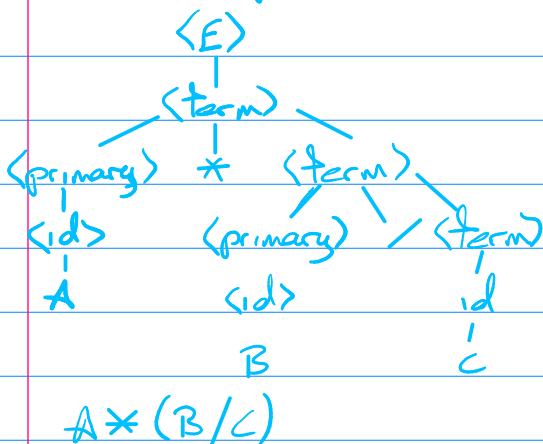
$\langle \text{term} \rangle \Rightarrow \langle \text{primary} \rangle \mid \langle \text{primary} \rangle * \langle \text{term} \rangle \mid \langle \text{primary} \rangle / \langle \text{term} \rangle$

$\langle \text{term} \rangle \Rightarrow \langle \text{prim} \rangle * \langle \text{term} \rangle \Rightarrow \langle \text{prim} \rangle * \langle \text{prim} \rangle / \langle \text{term} \rangle \Rightarrow \langle \text{prim} \rangle * \langle \text{prim} \rangle / \langle \text{prim} \rangle$
 $* \langle \text{term} \rangle \Rightarrow \dots$ terminated by $\langle \text{term} \rangle \Rightarrow \langle \text{prim} \rangle$

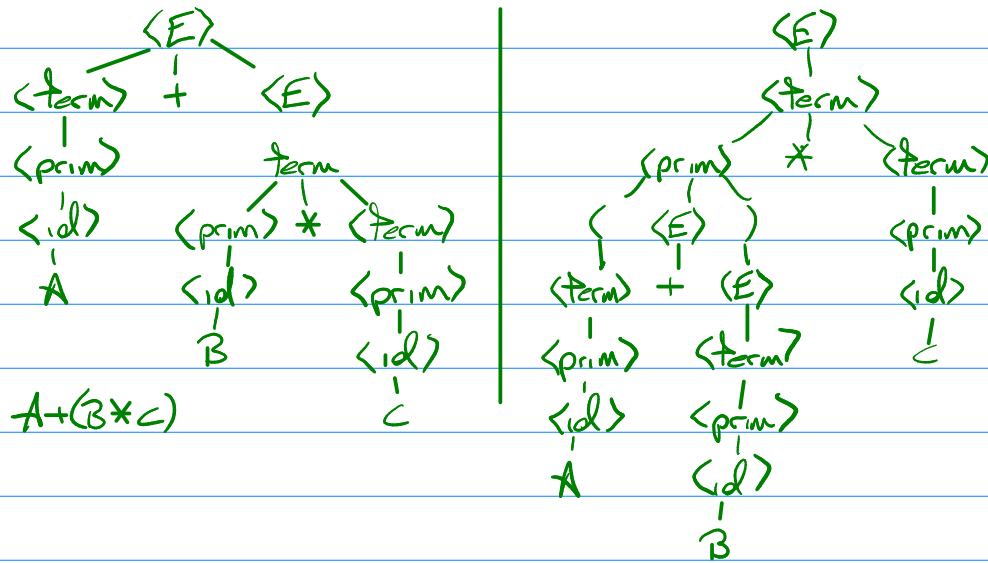
Any term is a sequence of primaries separated by * /

$\langle \text{primary} \rangle \Rightarrow \langle \text{id} \rangle \mid \langle \text{int} \rangle \mid \langle \text{float} \rangle \mid "(" \langle E \rangle ")"$

Construct the parse tree for $A * B / C$ Construct the parse tree for $A + B + C$



Construct the parse tree for $A+B \times C$



To get left associating at the same procedure level, use left recursion:

$\langle E \rangle \rightarrow \langle \text{term} \rangle \mid \langle E \rangle + \langle \text{term} \rangle \mid \langle E \rangle - \langle \text{term} \rangle$

$\langle \text{term} \rangle \rightarrow \langle \text{primary} \rangle \mid \langle \text{term} \rangle * \langle \text{primary} \rangle \mid \langle \text{term} \rangle / \langle \text{primary} \rangle$

To give higher precedence to $+/-$ over $* /$ (although not used in math)

$\langle E \rangle \rightarrow \langle \text{term} \rangle \mid \langle \text{term} \rangle * \langle \text{term} \rangle / \langle E \rangle$

$\langle \text{term} \rangle \rightarrow \langle \text{prim} \rangle \mid \langle \text{prim} \rangle + \langle \text{term} \rangle \mid \langle \text{prim} \rangle - \langle \text{term} \rangle$

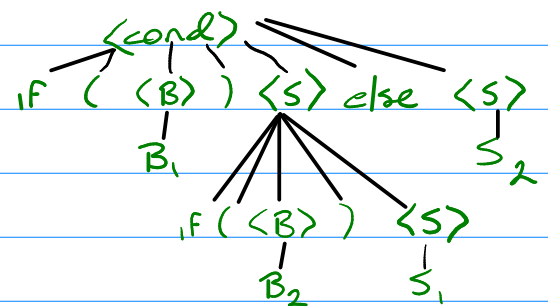
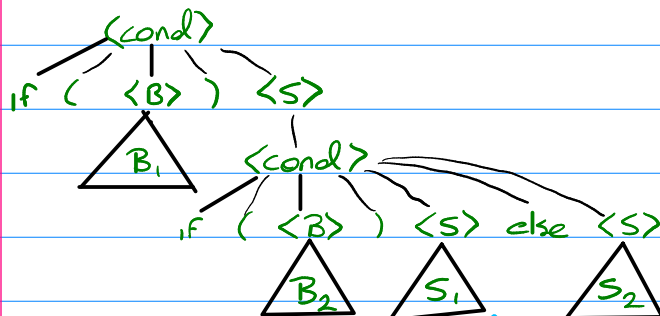
Conditionals

$\langle S \rangle \rightarrow \dots \mid \langle \text{cond} \rangle \mid \dots$

$\langle \text{cond} \rangle \rightarrow \text{if } "(" \langle B \rangle ")" \langle S \rangle \mid \text{if } "(" \langle B \rangle ")" \langle S \rangle \text{ else } \langle S \rangle$

Construct 2 parse trees for $\text{if } \langle B \rangle \text{ if } \langle B_2 \rangle S_1 \text{ else } S_2$

$\uparrow \quad \quad \uparrow \quad \quad \uparrow$
 $\quad ? \quad \quad ? \quad \quad$



Intermediate code would compute

$\text{if } (B_1) \{ \text{if } (B_2) S_1 \text{ else } S_2 \}$

This is the default grouping in C++/Java, and almost all languages.

Intermediate code would compute

$\text{if } (B_1) \{ \text{if } (B_2) S_1 \} \text{ else } S_2$

The following is an unambiguous BNF incorporating the rule:

each else matches with the closest preceding unmatched if

Introduce $\langle \text{matched } S \rangle$ for the statements containing no 1-branch conditional directly nested with other conditionals.

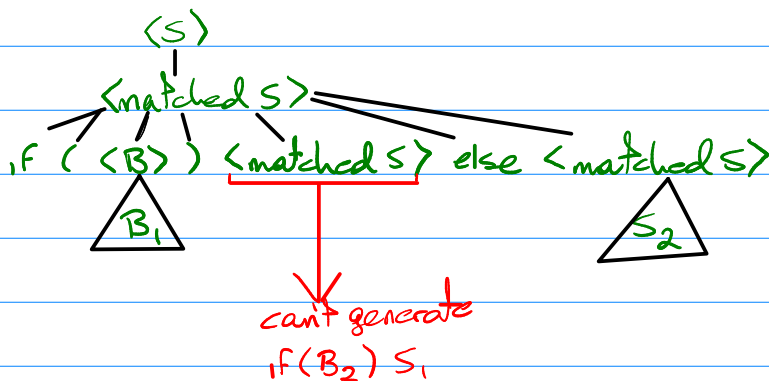
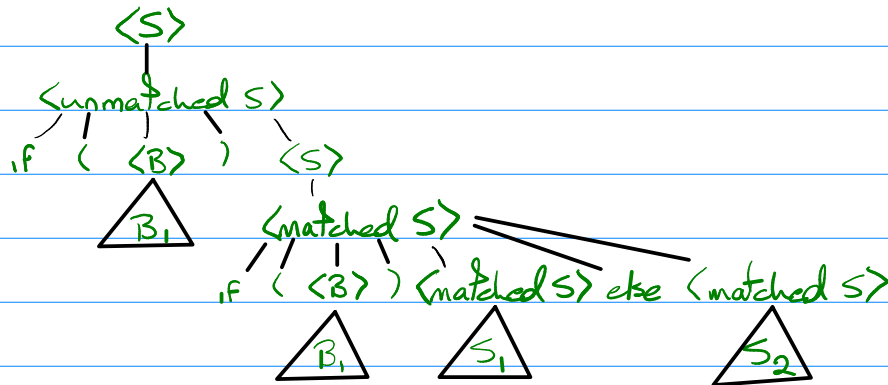
$\langle S \rangle \rightarrow \langle \text{matched } S \rangle \mid \langle \text{unmatched } S \rangle$

$\langle \text{matched } S \rangle \rightarrow \text{if } (\langle B \rangle) \langle \text{matched } S \rangle \text{ else } \langle \text{matched } S \rangle \mid \dots \text{ other kinds of}$

(assume S_1 and S_2 are simple statements, like assignments)

non-conditional
statements

$\langle \text{unmatched } S \rangle \rightarrow \text{if } (\langle B \rangle) \langle S \rangle \mid \text{if } (\langle B \rangle) \langle \text{matched } S \rangle \text{ else } \langle \text{unmatched } S \rangle$



This BNF is NOT used in practice.

- Not so easy to understand
- The parser cannot make the right choice without heavy look-ahead analysis.