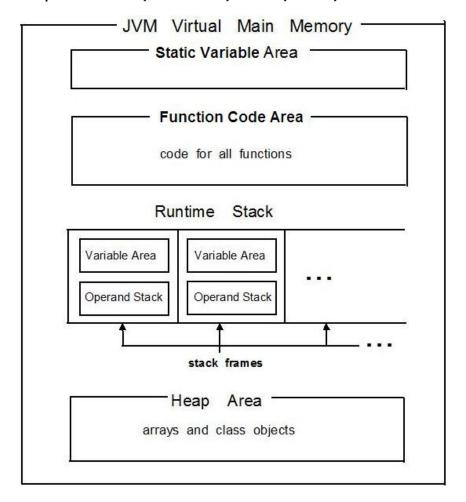
9/5/2010 Course Notes #3

**Java Virtual Machine (JVM)** is a member of the family of stack-based machines that use operand stacks for expression evaluations. They have no separate, conventional registers – or we could say operand stacks serve as high-level, virtual registers. The JVM has four main components:

- The static variable area that contains memory cells for static variables.
- The function code area that contains the instructions implementing functions.
- The runtime stack that controls function calls. Every function call pushes a new stack frame
  onto the runtime stack; it is popped when the function returns. Each stack frame contains a
  variable area and an operand stack. The variable area contains memory cells for the
  function's formal parameters and local variables. The operand stack is used to evaluate
  expressions that appear in the function body; operand stacks are sometimes called
  evaluation stacks.
- The heap area where arrays and class objects are dynamically allocated.



The Common Intermediate Language, used for example in Microsoft .NET Framework, is also based on a similar stack-based virtual machine.

Consider this example function where S represents a statement schema.

```
void example()
{
    int i = 0;
    while ( i < 100 )
    {
        S;
        i++;
    }
}</pre>
```

The following is an example JVM code that a compiler could generate. Every source function code is compiled to start from relative address 0. In the action description, "stack" refers to the operand stack. First, let's assume S is empty.

```
0 iconst_0 push int constant 0 onto stack
```

For a non-empty S, we simply insert JVM code for S after the if\_icmpge instruction and make necessary increments to instruction addresses.

```
push int constant 0 onto stack
0
     iconst 0
1
     istore 0
                        pop stack and store into address 0 in variable area (i=0)
2
     iload 0
                        push value at address 0 onto stack
      bipush 100 push int constant 100 onto stack if_icmpge X+14 pop top two values from stack; if stack[top-1] \geq stack[top] then go to X+14 (test if i < 100)
3
5
      . . .
      code for S
      iinc 0 1
X+8
                          increment value at address 0 by 1 (i++)
X+11 goto 2
X+14 return
```

Here X is the # of bytes occupied by the code for S.

As you might have guessed, any type of "load" instruction pushes a data item onto an operand stack, while any type of "store" instruction pops a data item from the stack and stores it at a location in the variable area.

Every JVM instruction occupies one byte, hence JVM "bytecode". An operand occupies one or more bytes depending on its instruction. Although the actual ---.class files contain JVM code in binary format, we can see their contents in text format by the JDK disassembler command javap. The binary contents of example.class compiled from the following example.java

can be disassembled to text format by the command "javap -c example" producing

```
Compiled from "example.java"
class example extends java.lang.Object{
static int j;
example();
 Code:
   0:
       aload 0
   1:
       invokespecial #1; //Method java/lang/Object."":() V
   4: return
static void example();
  Code:
   0: iconst_0
   1: istore 0
   2: iload_0
   3: bipush 100
   5:
       if_icmpge
                       2.2
   8: getstatic
                       #2; //Field j:I
   11: iload 0
   12: iadd
   13: putstatic
16: iinc  0, 1
                       #2; //Field j:I
   19: goto
               2
   22: return
```

## 9/5/2010 Course Notes #3

```
static {};
  Code:
    0:    iconst_0
    1:    putstatic  #2; //Field j:I
    4:    return
}
```