

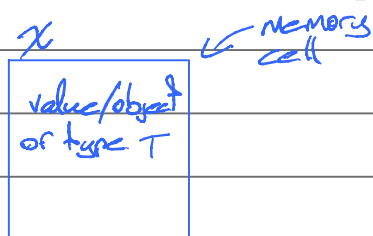
## Memory Mapping of Class Objects in Statically Typed OOPs

Let  $x$  be a variable declared to be

of type  $T$ .  $\{ \text{size}(T) = \text{size of the memory cell}$

$C++, C\#, Java, Eiffel, etc.$

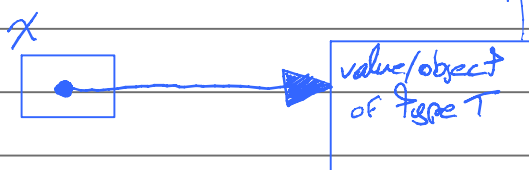
①  $T$  is of non-reference type  $\{ \text{allocated to one element of } T \}$



②  $T$  is of reference type  $\{ \text{usually inside heap}$

$\{ \text{eg. dynamically created by constructor functions} \}$

$\text{size}(x) = \text{size}(T)$



$\text{size}(T) = \text{the size of the memory cell allocated to a reference itself}$

In Java/Eiffel, all class types (including arrays types) belong to category 2

In  $C++$ , class types can be either 1 or 2.

$Cx$ ; // non-reference type

$C^*x$ ; // reference type

Java /  $C++$  with  $*$

class Person

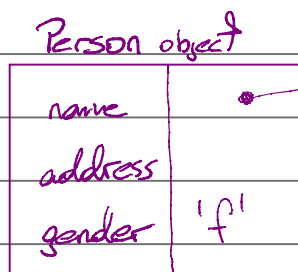
```
{
    Name name;
    Address address;
    char gender;
}
```

class Name

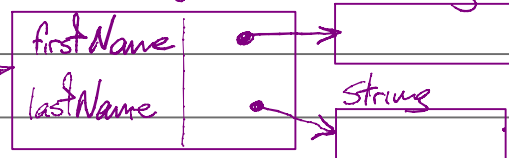
```
{
    String firstName;
    String lastName;
}
```

class Address

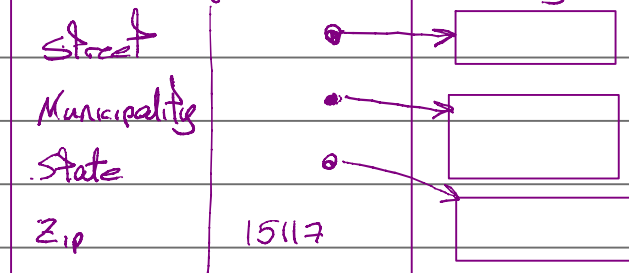
```
{
    String street;
    String municipality;
    String state;
    int zip;
}
```



Name object



Address object



## C++ without \*

Person object

name.firstName	
name.lastName	
⋮	
⋮	
⋮	

Memory Cell format for class objects without Dynamic Binding of function code  
(No inheritance polymorphism)

Base Address →

$x_1$

$x_2$

⋮

$x_m$

class X

{  $T_1 x_1;$

⋮

$T_m x_m;$

$f_1(\dots);$

⋮

$f_n(\dots);$

m instance level fields.

The compiler will call the code for  $f_j$  as it is defined in class X

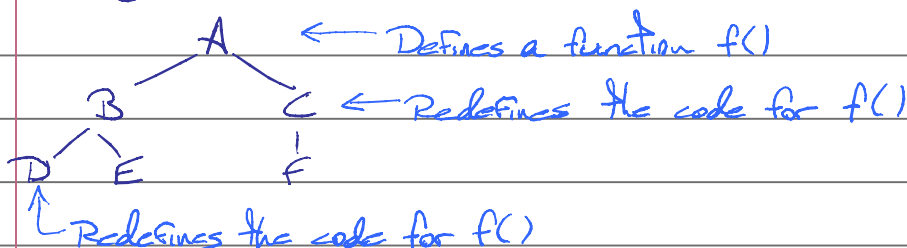
n instance level method functions

The functions  $f_j$  are statically bound at compilation time

$$\text{address}(x_i) = \text{Base Address} + \sum_{1 \leq k \leq i-1} \text{size}(x_k), \quad 1 \leq i \leq m$$

With Dynamic Binding of Function Code

Single-Inheritance Tree



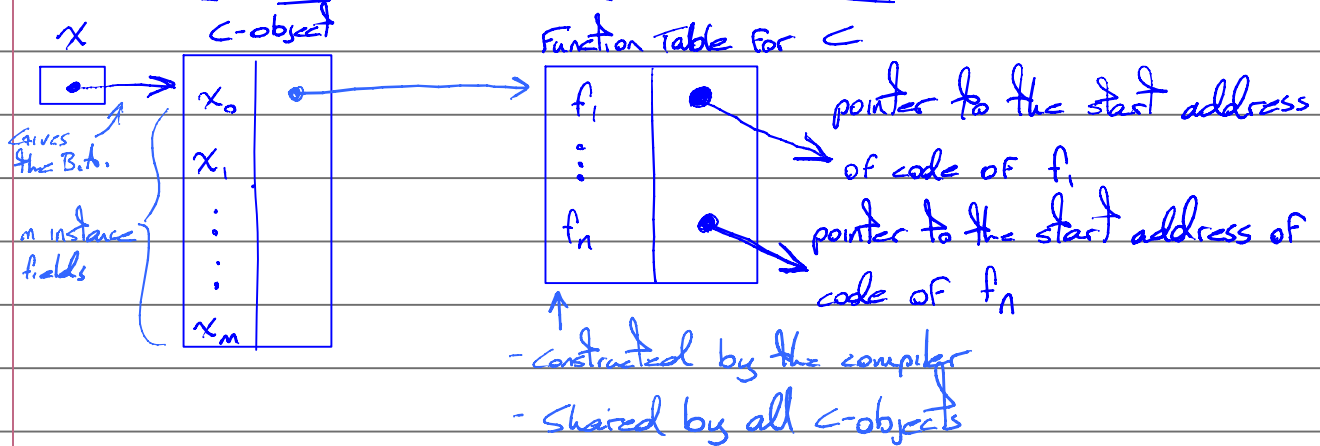
A a; // Java, Eiffel, etc. - dynamic binding is default

A\* a; // C++ - dynamic binding is enabled when reference type is used and  $f()$  is declared to be a virtual function

a.f(); // a may refer to A-object or object of any subclass of A  
a refers to A-object  $\Rightarrow f()$  defined in A will be called

$a \text{ refers to } B\text{-object} \Rightarrow f() \text{ defined in } A \text{ will be called}$   
 $a \text{ refers to } C\text{-object} \Rightarrow f() \text{ defined in } C$   
 $a \text{ refers to } D\text{-object} \Rightarrow f() \text{ defined in } D$   
 $a \text{ refers to } E\text{-object} \Rightarrow f() \text{ defined in } A$   
 $a \text{ refers to } F\text{-object} \Rightarrow f() \text{ defined in } C$

Use of function tables (method tables) is most common



$$\begin{aligned}
 \text{address}(x_i) &= \text{Base address (reference value in } x) + \text{offset}(i) \\
 &= \text{Base address (reference value in } x) + \sum_{0 \leq k \leq i-1} \text{size}(x_k); \quad 0 \leq i \leq m \\
 \text{address}(f_j) &= \text{Reference in } x_0 \text{ field} + \text{offset}(j) \\
 &= \text{Reference in } x_0 \text{ field} + \sum_{1 \leq k \leq j-1} \text{size}(f_k), \quad 1 \leq j \leq n
 \end{aligned}$$

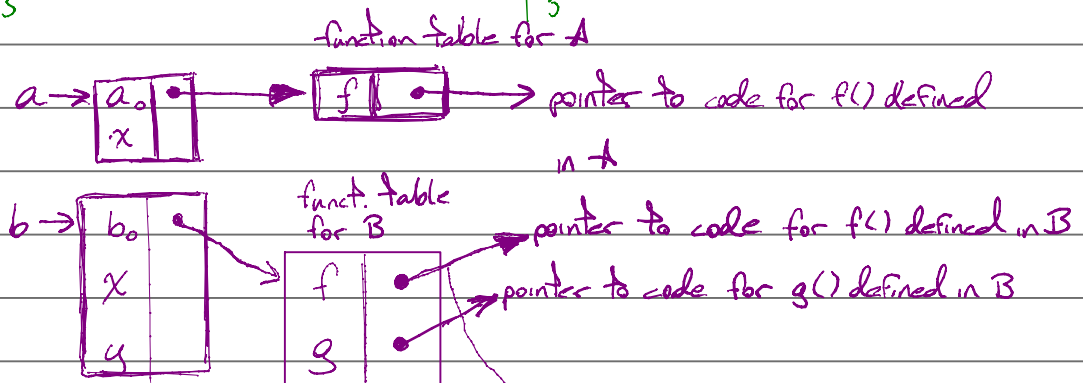
Use Java-like Notation

<pre> class A {   int x;   void f() {...} }         </pre>	<pre> A   B   C         </pre>	<pre> class B extends A {   int y;   void f() { /* f redefined */ }   void g() {...} }         </pre>	<pre> class C extends B {   int x;   void g() { /* g redefined */ }   void h() {...} }         </pre>
--	--------------------------------	---	---

void main()

```

{
  A a = new A();
  B b = new B();
  C c = new C();
}
  
```



$C \rightarrow$

$c_0$	
$x$	
$y$	
$z$	

function table for  $C$

$f$	
$g$	
$h$	