

- ①
 - ② ~~F~~
 - ③ F, they are the same
 - ④ F, they should be the same (F 1D, T if they're 2D)
 - ⑤ F, Compaction defrags the heap
 - ⑥ T
 - ⑦ F, Syntax describes grammatical structures
 - ⑧ T
 - ⑨ T, because it is a pointer
 - ⑩ T
 - ⑪ T
 - ⑫ F, RTS controls function calls. ARs are simply function states
 - ⑬ F, Points to the end of the AR that called it.
 - ⑭ T
 - ⑮ F, Garbage Collection is for the heap, not the AR
 - ⑯ T
 - ⑰ T
 - ⑱ T, Static variables are fixed in one location
 - ⑲ F, Heap and RTS are separate
 - ⑳ F, Everything is static
 - ㉑ T
 - ㉒ F Binding = Type Binding and can't undo in statically-typed
 - ㉓ T, Function Table is used dynamic inheritance
 - ㉔ T, Everything is a function call

② ① 10-13-09 Notes - 2 branch conditional

- ③
- ④
- ⑤

- ③
- ④

④

- ⑤
 - ① Type checking @ compile time (statically)
 - ② Faster to code, more flexible
 - ③ Java / C++
 - ④ PHP, Python, JavaScript

⑥ ② I, II, III (C++)

③ I II III (Java)

- ④ I (local vars in functions)
- ⑤ II (formal params in functions)

⑥ III

⑦ I

- ⑧ II
- ⑨ II

⑩ III - Because of explicit memory calls

- ⑪ I
- ⑫ III

± Static = Variables that don't change locations

II Stack-Dynamic = Depends on runtime

III Heap-Dynamic = Anything in the heap.
Garbage collection

- ⑦ a) x
 b) Yes, $\text{Node } a = \text{new Node}()$;
 $\text{Node } m = \text{new Node}()$;
 $\text{Node } n$;

⑧ Exercise Set #4, question 5 (and notes from 10-29-09)

⑧ Base Address = $\text{address}([5]) = 10$
 Element Size = 2

Give the formula for $\text{address}([i, j])$

$[5, \dots, 15]$ $S = a_1, 15 = B_1$

$\text{address}([i, j]) = \text{BA} + \text{rank}([i, j]) \times \text{ES}$

$\text{address}([i, j]) = 10 + (i - 5) \times 2$

From: notes-10
 $\text{address}([i_1, \dots, i_n]) = \text{BA} + \text{rank}([i_1, \dots, i_n]) \times \text{ES}$

For 1D array
 $\text{rank}([i, j]) = i - a_1$

⑨ 2D array $[0 \dots 10, 5 \dots 20]$

Base Address = $\text{address}([0, 5]) = 10$

Element Size = 1

Compute $\text{address}([5, 15])$

Row Major = $\text{BA} + \text{ES} \times \text{rowSize} \times (i_1 - B_1) + \text{ES} \times (i_2 - B_2)$

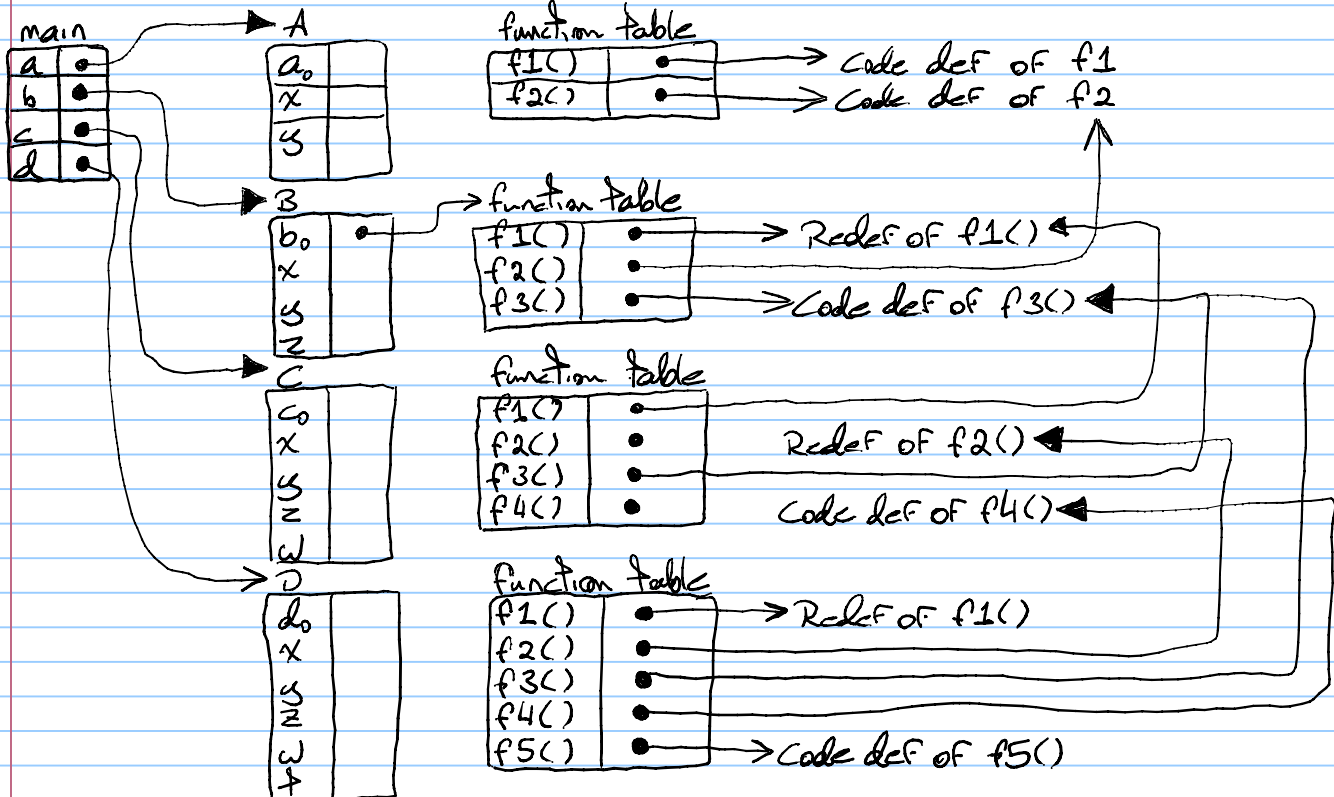
Column Major = $\text{BA} + \text{ES} \times (i_2 - B_2) \times \text{colSize} + (i_1 - B_1)$

⑩ $\text{base}([1]) = \text{BA}$ = the start address of the array $[a_1, \dots, b_1]$

$\text{base}([i_1, \dots, i_{k-1}])$ = the value of the cell @ $\text{address}([i_1, \dots, i_{k-1}])$

$\text{address}([i_1, \dots, i_k]) = \text{base}([i_1, \dots, i_{k-1}]) + (i_k - a_k) \times S$, where $S = \text{ES}$ if $k = n$, $S = AS$ if $k < n$

⑪ Notes from 11-12-09 (Person object, Name object, etc)



12) From Notes 12

- a) Root nodes - static variables, memory cells, ARs within memory cells, registers, which contain references to heap cells
- b) The nodes are root nodes and non-root nodes.
The links are the references to the nodes
- c) Reachable = Nodes which are allocated and can be reached by a chain of 1 or more references from one of the root nodes.
- Garbage = A node that isn't reachable. (i.e. it loses a reference to the root)

- 13) a) (b) • cons = constructs a list
- b) (2()) or (2) • car = returns the first parameter of the input
- c) (a) • cdr = returns the second parameter of the input
- d) (b) Check Exercise Set #8 (or every parameter excluding the first)
- e) (1 2 3)

14) a) (define (append L1 L2) From 12-08-09
 (if (null? L1) L2
 (cons(car L1) (append(cdr L1) L2))
))

b) (define (map f list) From 12-08-09
 (if (null? list) ()
 (cons(f(car list)) (map f(cdr list)))
))

15) length([], 0).
 length(H/T, Y) :- length(T, Y1), Y is Y1 + 1.

16)