SECRET KEY-TASK 3

```
File Actions Edit View Help

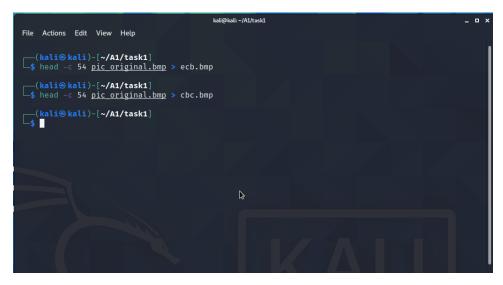
(kali@ kali)-[~/A1/task1]
$ openssl enc -aes-128-ecb -e -in pic original.bmp -out ecb.bin -K 00112233445566778889aabbccdde eff

(kali@ kali)-[~/A1/task1]
$ openssl enc -aes-128-cbc -e -in pic original.bmp -out cbc.bin -K 00112233445566778889aabbccdde eff -iv 0102030405060708 hex string is too short, padding with zero bytes to length

(kali@ kali)-[~/A1/task1]

(kali@ kali)-[~/A1/task1]
```

- First, we encrypt the file using 2 encrypted method as seen above
- Notice that, while encrypt using cbc, the message is too short compared to the key, that why we need some padding in order to match with the length of the key



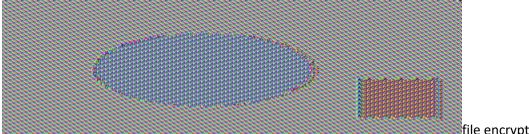
Then we preserve 54 bits header of the file, since they are needed to properly displayed

• Finally, we append the encrypted bits from offset 55 to the end of file



Outcome

- What we observe is the encrypted file using ecb looks similar to the original file.
 - This is because ecb using 1 to 1 mapping encryption methods with single key, The key can then be easily identify



file encrypted using ecb

- Encrypting file using cbc make it impossible to identify the original message
 - This is because cbc using 1 to 1 mapping but the key change with each block of the plaintext



file encrypted using cbc

SECRET KEY —TASK 7

```
devel@localhost:~/PrivateKeyTask7
File Edit View Search Terminal Help
         i++;
    }
if((tolower(*string1))!=0){
    result=1;
    return result;
           }
         } return result;
void word_processing(char *word, int len)
     int numWordAdd = 16-len;
while (numWordAdd--)
        word[len++] = '#';
     word[len++] = '#';
}
word[len] = '\0';
int i;
for (i = 0; word[i]; i++)
{
         word[i] = tolower(word[i]);
}
int main()
{
   /* Error */
strcat("%s","error");
return 0;
                  }
/* Buffer passed to EVP_EncryptFinal() must be after data just
* encrypted to avoid overwriting it.
*/
if (!EVP_EncryptFinal_ex(&ctx, outbuf + outlen, &tmplen))
,
                        /* Error */
strcat("%s","error");
return 0;
                   }
outlen += tmplen;
EVP_CIPHER_CTX_cleanup(&ctx);
                                                                                                                               78,9
                                                                                                                                               72%
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RSA-TASK 1

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File Actions Edit View Help

- (Mail@ Mail) - [/mut/hgfs/shared folder/LAB1/RSA1]
- (Mail@ Mail) - [/mut/hgfs/shared folder/LAB1/RSA1]
- (Mail@ Mail) - [/mut/hgfs/shared folder/LAB1/RSA1]
- S. /A.out
- Private Key - 3587A24598E5F2A21D8007D89D18CC50ABA5075BA19A33890FE7C28A9B49GAEB

- (Mail@ Mail) - [/mut/hgfs/shared folder/LAB1/RSA1]

- BY OFFENSIVE SECURITY

BY OFFENSIVE SECURITY
```

- Given p,q we can calculate phi of (q*p), by calculating (phi of q) = q-1 and (phi of p) = p-1
- then we know that m^phi =1 mod n

```
=> m^(phi*random +1) = m mod n
=> phi*random+1 = e*d
```

=> e*d = 1 mod phi

=> d = inverse modular of e mod phi

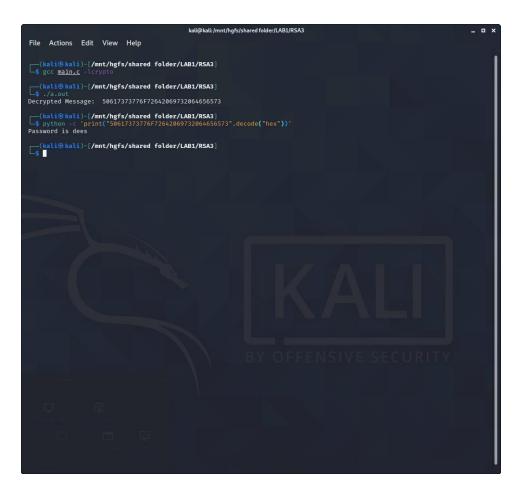
RSA-TASK 2



- Convert Ascii message to hex using command python -c 'print("A top secret!".encode("hex"))'
 We get message = 4120746f702073656372657421
- Then we encrypt it using message ^ e mod n
- Then we double check by decrypt the ciphertext and compare with original message
- The decryption using encrypted_message ^d mod n

RSA-TASK 3

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```



- The decryption using encrypted_message ^d mod n
- Then we get a result in hex value
- Then we convert it back to ascii using command

python -c 'print(message.encode("hex"))'