

TreeSearch User Guide

Version 0.9

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Abstract

The TreeSearch library abstracts the structure of a search tree in order to manage jobs in a distributed computation. This document details the strategy of TreeSearch along with implementation, compilation, and execution details. An example application is also given.

1 Introduction

The computation path of a dynamic search frequently takes the form of a rooted tree. One important property of each node in this tree is that the computation at that node depends only on the previous nodes in the ancestral path leading to the root of the computation. If the search is implemented in the usual way, subtrees operate independently.

For a search of this type, all search nodes at a given depth can be generated by iterating through the search tree, but backtracking once the target depth is reached. Each of the subtrees at this depth can be run independently, and hence it is common to run these jobs concurrently (See [2] Chapter 5 for more information).

The TreeSearch library was built to maximize code reuse for these types of search. It abstracts the structure of the tree and the recursive nature of the search into custom components available for implementation by the user. Then, the ability to generate a list of jobs, run individual jobs, and submit the list of jobs to a cluster are available with minimal extra work.

TreeSearch is intended for execution on a distributed machine using Condor [3], a job scheduler that uses idle nodes of a cluster or network. Condor was chosen as its original development was meant for installation in computer labs and office machines at the University of Wisconsin–Madison to utilize idle computers.

The C++ portion of TreeSearch is independent of Condor. The Python scripts which manage the input and output files as well as modifying the submission script are tied to Condor, but could be adapted for use in other schedulers.

1.1 Acquiring TreeSearch

The latest version of TreeSearch and its documentation is publicly available on GitHub [1] at the address <http://www.github.com/derrickstolee/TreeSearch/>.

2 Strategy

Let us begin by describing the general structure and process of an abstract tree-based search. There is a unique root node at depth zero. Each node in the tree searches in a depth-first, recursive manner. There are

a number of children to select at each node. One may select this child through iteration or selecting via a numerical order. Before searching below the child, a pruning procedure may be called to attempt to rule out the possibility of a solution below that child. Another procedure may be used to find if this node is a solution. Now, the search recurses at this node until its children are exhausted and the search continues back to its parent.

2.1 Subtrees as Jobs

This tree structure allows for search nodes to be described via the list of children taken at each node. Typically, the breadth of the search will be small and these descriptions take very little space. This allows for a method of describing a search node independently of what data is actually stored by the specific search application. Moreover, the application may require visiting the ancestor search nodes in order to have consistent data. With the assumption that each subtree is computationally independent of other subtrees at the same level, one can run each subtree in a different process in order to achieve parallelization. These path descriptions make up the input for the specific processes in this scheme.

Each path to a search node qualifies as a single job, where the goal is to expand the entire search tree below that node. A collection of nodes where no pair are in an ancestor-descendant relationship qualifies as a list of independent jobs. Recognizing that the amount of computation required to expand the subtree at a node is not always a uniform value, TreeSearch allows a maximum amount of time within a given job. In order to recover the state of the search when the execution times out, the concept of *partial jobs* was defined. A partial job describes the path from the root to the current search node. In addition, it describes which node in this path is the original job node. The goal of a partial job is to expand the remaining nodes in the subtree of the job node, without expanding any nodes to the left of the last node in this path. See Figure 1 to an example partial job and its position in the job subtree.

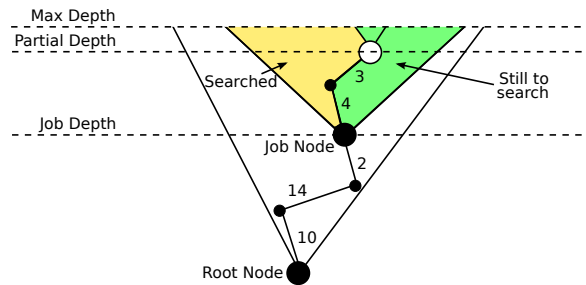


Figure 1: A partial job description.

2.2 Job Descriptions

The descriptions of jobs and partial jobs are described using text files in order to minimize the I/O constraints on the distributed system. The first is the standard job, given by a line starting with the letter J. Following this letter are a sequence of numbers in hexadecimal. The first two should be the same, corresponding to the depth of the node. The remaining numbers correspond to the child values at each depth from the root to the job node.

A partial job is described by the letter P. Here, the format is the same as a standard job except the first number describes the depth of the job node and the second number corresponds to the depth of the current node. For example, the job and partial job given in Figure 1 are described by the strings below:

```
J 3 3 10 14 2
P 3 5 10 14 2 4 3
```

2.3 Customization

The TreeSearch library consists of an iterative implementation of the abstract search. The corresponding actions for a specific application are contacted via extending the `SearchManager` class and implementing certain virtual functions. The list of functions available are given in Table 1.

LONG_T	pushNext()	Deepen the search to the next child of the current node.
LONG_T	pushTo(LONG_T child)	Deepen the search to the specified child of the current node.
LONG_T	pop()	Remove the current node and move up the tree.
int	prune()	Perform a check to see if this node should be pruned.
int	isSolution()	Perform a check to see if a solution exists at this point.
char*	writeSolution()	Create a buffer that contains a description of the solution.
char*	writeStatistics()	Create a buffer that contains custom statistics.

Table 1: List of virtual functions in the `SearchManager` class.

In addition to supplying the logic behind these functions, protected members of the `SearchManager` class can be modified to change the operation of the search. These parameters are listed in Table 2.

Type	Name	Option	Description
int	maxdepth	-m [N]	The maximum depth the search will go. In generate mode, a job will be output with job description given by the current node.
int	killtime	-k [N]	Number of seconds before the search is halted. If the search has not halted naturally, a partial job is output at the current node.
int	maxSolutions	--maxsols [N]	The maximum number of solutions to output. When this number of solutions is reached, a partial job is output and the search halts.
int	maxJobs	--maxjobs [N]	The maximum number of jobs to output (generate mode). When this number of jobs is reached, a partial job is output and the search halts.
bool	haltAtSolutions	--haltatsols [yes/no]	If true, the search will stop deepening if <code>isSolution()</code> signals a solution. If false, the search will continue until specified by <code>prune()</code> or <code>maxdepth</code> .

Table 2: List of members in the `SearchManager` class.

3 Integration with TreeSearch

This section details the specific interfaces for implementation with `TreeSearch`.

It is important to understand the order of events when the search is executing. The search begins when the `doSearch()` method is called. The first call initializes the search, including starting the kill timer. Then, each recursive call expands the current search node at the top of the stack. Figure 3 describes the actions taken by the recursive `doSearch()` method at each search node.

INSERT NEW FLOWCHART: Conceptual order of events in an exhaustive search from root.

Figure 2: The conceptual operation of the `doSearch()` method.

3.1 Virtual Functions

The two most important methods are the `pushNext()` and `pushTo(LONG_T child)` methods. Both deepen the search, manage the stack, and control the job descriptions. Each returns a child description (of type `LONG_T`)

`pushNext()`: Advance the search stack using the next augmentation available at the current node. Return a label of `LONG_T` type which describes the augmentation. Return `-1` if there are no more augmentations to attempt at this stage, signaling a `pop()` operation.

`pushTo(LONG_T child)`: Advance the search stack using the augmentation specified by `child`. If the augmentation fails as specified, return `-1`, and the search will terminate in error. **Note:** If this method is called during the partial portion of a job description, the later augmentations will be called using the `pushNext()` method, so the current augmentation must be stored for later.

`pop()`: Clean up any memory used for the current level and/or revert any data structures to the previous level.

`prune()`: Check if the current node should be pruned (i.e. detect if there are no solutions reachable from the current node by performing the specified augmentation). Return `1` if the node should be pruned, `0` otherwise. A prune signal will be followed by the `pop()` method.

`isSolution()`: Check if there is a solution at this node. If there is a solution, store the solution data and return `1`. The `writeSolution()` method will be called to pass the solution data to output. If the `haltAtSolutions` option is set, a solution will trigger a `pop()` method. Otherwise, the node will be used for augmentations until the maximum depth is reached or the `prune()` method signals a prune for all later nodes.

`writeSolution()`: Return a buffer containing the solution data for the output. This data will be sent to standard out along with the job description of the current node (with an 'S' prefix). If your data is prohibitively large for sending over the network, this job description can be used to generate the current node where the solution data can be recovered. **Note:** The buffer must be allocated using `malloc`, as it will be deallocated using `free`.

`writeStatistics()`: Write a buffer of custom statistics information. Each line must follow the format

T [TYPE] [ID] [VALUE]

where [TYPE] is one of "SUM", "MAX", or "MIN", [ID] is a specified name, and [VALUE] is a number. These statistics are combined by the `compactjobs` script using the [TYPE] specifier (either sum the values or store the maximum or minimum) and the statistics are placed in the `allstats.txt` file. The `combinestats` script converts the `allstats.txt` file into a comma separated value file, grouping by depth over variables with [ID] of the form [SUBID]_AT_[#]. This allows per-depth statistics tracking.

3.2 Helper Methods

The following methods are useful when constructing a `TreeSearch` application.

`importArguments(int argc, char argv)`**: Take the command line arguments and set the standard search options. The options from Table ?? such as `killtime`, `maxJobs`, and `maxdepth` are all set in this way.

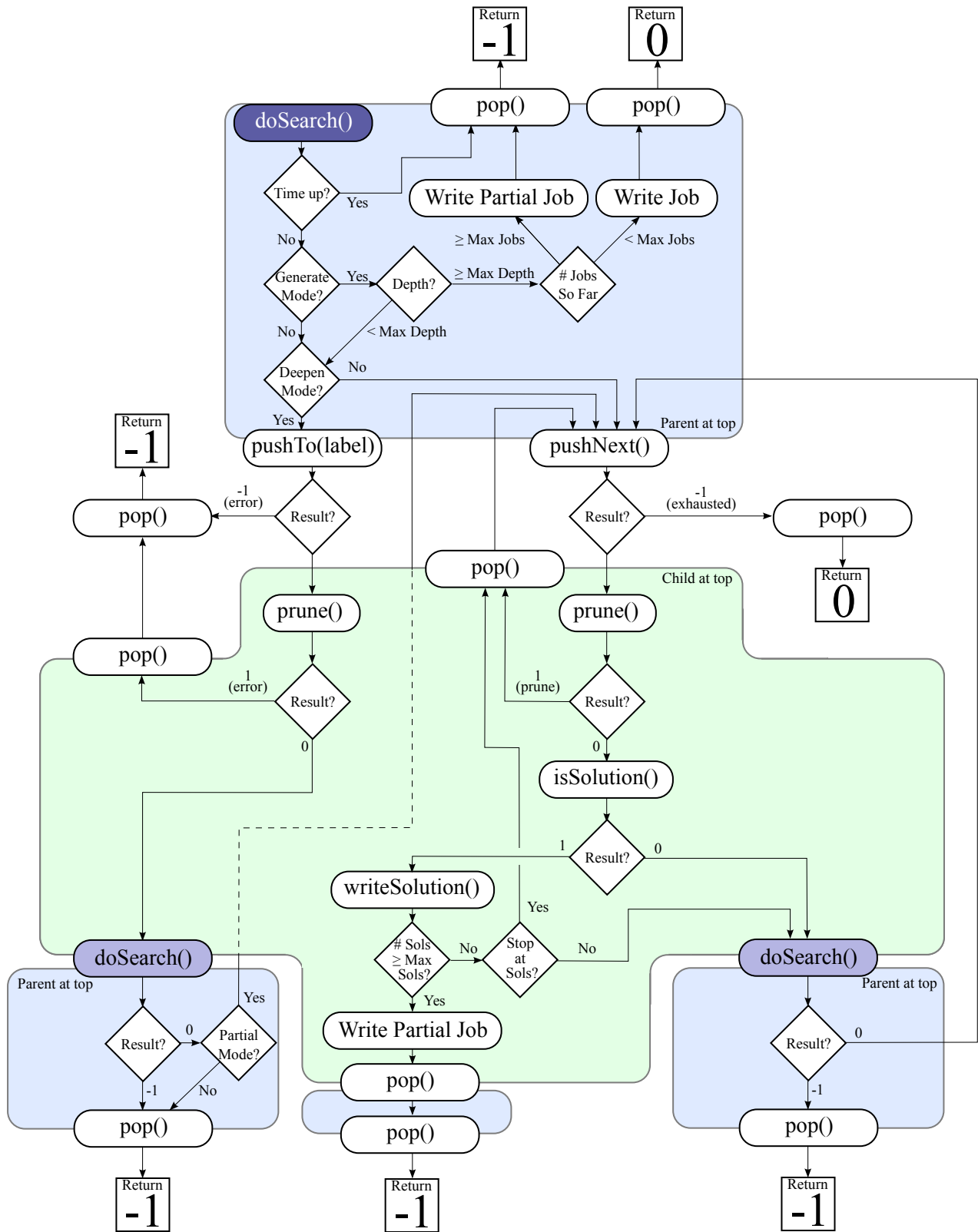


Figure 3: The full operation of the `doSearch()` method.

readJob(FILE* file): Read a job from the given input source, such as standard in. It will read only one line, and prepare the manager for the `doSearch()` method.

doSearch(): This method starts the search on the current job as well as returns the status of the search: 1 if completed with a solution, 0 if completed with no solution, and -1 if halted early due to error or time constraints.

3.3 Compilation

To compile TreeSearch, run `make` in the source directory. This command compiles the object file `SearchManager.o` which must be linked into your executable. Moreover, it compiles the example application presented in Section 5. Your code must reference the header file `SearchManager.hpp` and link the object file `SearchManager.o`.

4 Execution and Job Management

To execute a single process, simply run your executable with the proper arguments. However, to run a distributed job via Condor, a set of scripts were created to manage the input and output files, the Condor submission file, and monitor the progress of the submission during execution.

4.1 Management Scripts

The TreeSearch library works best with independent subtrees and hence does not suffer from scaling issues when the parallelism is increased. However, managing these large lists of jobs requires automation.

4.1.1 Expanding jobs before a run

When the generation step is run, a list of jobs is presented in a single file. Condor requires a separate input and output file for each process. The role of the `expandjobs` script is to split the jobs into individual files and to set up the Condor submission file for the number of jobs that are found.

There are a few customizable options for this script.

- `-f [folder]` – change the folder where the jobs are created. Default is `./`.
- `-m [maximum]` – set the maximum number of jobs allowed. Default is unlimited.
- `-g [groupsize]` – set the number of jobs per process. Default is 1.

Inside the specified folder, the file `condorsubmit.sub` is set with the proper folder and queue size based on the number of jobs found. Any remaining jobs that did not fit within the maximum are held as back jobs. They will be added to the job pool when the jobs are completed.

4.1.2 Collecting data after a run

Once Condor has completed the requested jobs, the output must be collected to discover which jobs completed fully, which are partially complete, and how many solutions have been found. The script `compactjobs` was built for this purpose.

This script takes the output files and reads all new jobs that may have been generated using the staging feature, finds if the input job completed or is partial, and reports on the total number of jobs of each type. Moreover, it will find and store the solutions found, along with the corresponding data.

Finally, it compiles statistics from each run. Using the `writeStatistics` method, the application may report statistics by starting the line with a “T” followed by the type (MAX, MIN, SUM), variable name, and variable value. These are collected using the specified type and compiled with existing statistics from previous batches.

5 Example Application

6 Example Workflow

Create submission template

```
—i
Generate initial jobs
—i
Compact data
—i
Evaluate # of Jobs
—i
Expand with -M/-G
—i
RUN -OR- GENERATE (modify the condorsubmit.sub)
—i
BACK TO COMPACT
DURING A RUN:
check.sh
AFTER A RUN:
stats.sh/stataccept.txt
PREDICTING TIME:
randomjobs.py
```

7 Summary

8 Acknowledgements

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References

- [1] GitHub Inc. Github. <https://github.com/>.
- [2] Petteri Kaski and Patric R. J. Östergård. *Classification Algorithms for Codes and Designs*. Number 15 in Algorithms and Computation in Mathematics. Springer-Verlag, Berlin Heidelberg, 2006.

- [3] Douglas Thain, Todd Tannenbaum, and Miron Livny. Distributed computing in practice: the Condor experience. *Concurrency - Practice and Experience*, 17(2-4):323–356, 2005.