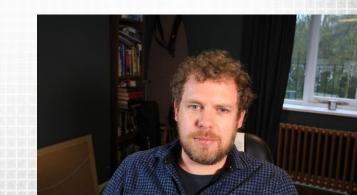
Parallel Computing with GPUs

Warp Level CUDA and Atomics Part 3 - Atomics and Warp Operations



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This Lecture (learning objectives)

- **□**Atomics
 - ☐ Present GPU atomic operations
 - ☐ Demonstrate performance of GPU atomics and locks
- ☐ Warp Operations
 - ☐Present warp shuffle operations for communication of threads within a warp
 - ☐Give examples of warp operations such as sum and warp voting.





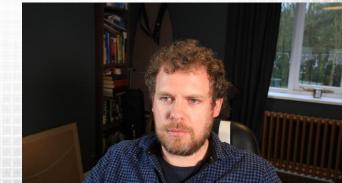
What is wrong with the following

```
__global___ void max_kernel(int *a)
{
    __shared__ int max;
    int my_local = a[threadIdx.x + blockIdx.x*blockDim.x];
    if (my_local > max)
        max = my_local;
}
```



☐ More than one thread may try to modify max at the same time ☐ Race condition

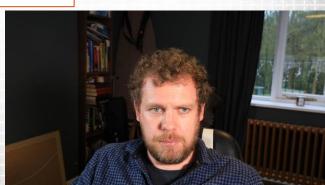
```
__global__ void max_kernel(int *a)
{
    __shared__ int max;
    int my_local = a[threadIdx.x + blockIdx.x*blockDim.x];
    if (my_local > max)
        max = my_local;
}
```



Atomics

□ Atomics are used to ensure correctness when concurrently reading and writing to a memory location (global or shared)

- No race condition
- ☐ Function supported in *most* hardware
 - ☐ Some older generation GPUs lack shared memory and floating point atomic etc.



Atomic Functions and Locks

- ☐ An atomic function
 - ☐ Must guarantee that an operation can complete without interference from any other thread
 - ☐ Does not provide any guarantee of ordering or provide any synchronisation

☐ How can we implement critical sections?

```
__device__ int lock = 0;

__global__ void kernel() {
  bool need_lock = true;
  // get lock
  while (need_lock) {
    if (atomicCAS(&lock, 0, 1) == 0) {
        //critical code section
        atomicExch(&lock, 0);
        need_lock = false;
    }
  }
}
```

```
int atomicCAS(int* address, int compare, int val)
```

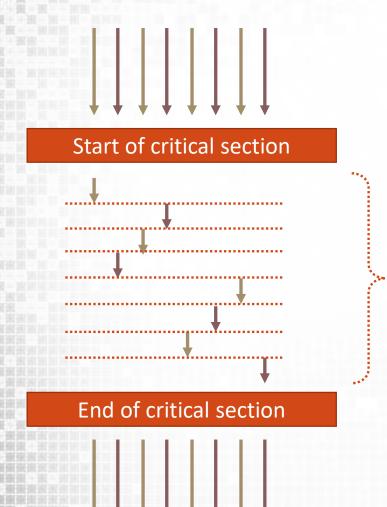
Performs the following in a single atomic transaction (atomic instruction)

```
*address=(*address==compare)? val : *address;
```

Returning the old value at the address



Serialisation



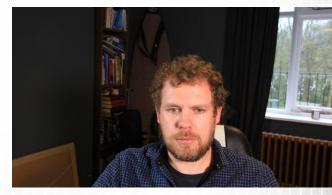
☐What happens to performance when using atomics?

- ☐ In the case of the critical section example
 - ☐ This is serialised for each thread accessing the shared value
- ☐ For the atomic CAS instruction access to the shared lock variable is serialised
 - ☐ This is true of any atomic function or instruction in CUDA



CUDA Atomic Functions / Instructions

- ☐ In addition to atomicCAS the following atomic functions/instructions are available
 - □Addition/subtraction
 - \square E.g. int atomicAdd(int* address, int val) add val to integer at address
 - **□**Exchange
 - ☐ Exchange a value with a new value
 - □Increment/Decrement
 - ☐ Minimum and Maximum
- Variants of atomic functions
 - □64 bit integer and double versions available in Pascal (Compute 6.0)
 - □See docs: https://docs.nvidia.com/cuda/cuda-c-programming-guide/index.html#atomic-functions





Local vs Global Atomics



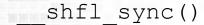
- ☐ Image histogram example
 - ☐ Accumulation of colour values for images
 - ☐ Entropy: measure of the level of disorder (lower entropy == higher contention)
 - https://devblogs.nvidia.com/parallelforall/gpu-pro-tip-fast-histograms-using-shared-atomics-maxwell/

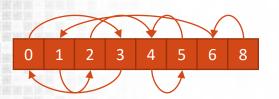


☐ For moving/comparing data between threads in a block to use Shared Memory (SM)	it is possible
☐ For moving/comparing data between threads in a warp (lanes in this context) it is possible to use a warp shuffle (S	
☐ Direct exchange of information between two threads	
☐Can replace atomics	
☐Should <u>never</u> depend on conditional execution!	
☐ Does not require SM	
☐Always faster than SM equivalent	
☐ Implicit synchronisation (no need for syncthreads)	
□EXCEPT on Volta+ hardware (use syncwarp)	
☐ Works by allowing threads to read another threads registers	
https://docs.nvidia.com/cuda/cuda-c-programming-guide/indesshuffle-functions	ex.html#warp-
	The
	The University Of Sheffield.



Shuffle Variants





Shuffled between any two index threads

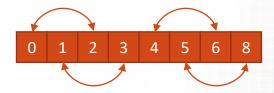
__shfl_up_sync()



Shuffles to nth right neighbour wrapping indices (in this case n=2) __shfl_down_sync()



Shuffles to nth left neighbour wrapping indices (in this case n=2) shfl xor sync()



Butterfly (XOR) exchange shuffle pattern



Shuffle function arguments

☐ Must be a power of 2 and less than or equal to warp size

☐ If smaller than warp size each subsection acts independently (own wrapping)

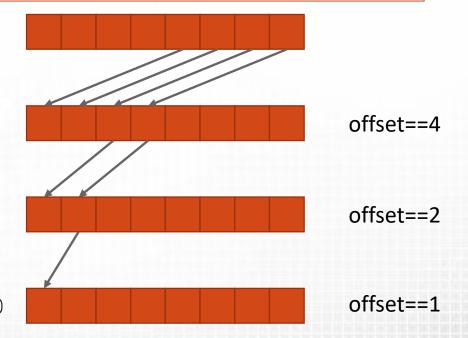
T shfl sync(unsigned mask, T var, int srcLane, int width=warpSize); shfl up sync(unsigned mask, T var, unsigned int delta, int width=warpSize); T shfl down sync (unsigned mask, T var, unsigned int delta, int width=warpSize); delta is the n step used for shuffling T shfl xor sync(unsigned mask, T var, int laneMask, int width=warpSize); □ Source lane determined by bitwise XOR with laneMask ☐ Mask is a bit mask for the warp to indicate which threads participate ☐ T can be int, unsigned int, long, unsigned long, long long, unsigned long long, float or double ☐ Optional width argument

Shuffle Warp Sum Example (down)

```
__global__ void sum_warp_kernel_shfl_down(int *a)
{
  int local_sum = a[threadIdx.x + blockIdx.x*blockDim.x];

  for (int offset = WARP_SIZE / 2; offset>0; offset /= 2)
    local_sum += __shfl_down(local_sum, offset);

  if (threadIdx.x%32 == 0)
    printf("Warp max is %d", local_sum)
}
```



Warp sum in threadIdx.x%32==0

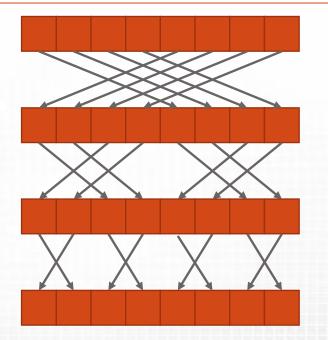


Shuffle Warp Sum Example (xor)

```
__global__ void sum_warp_kernel_shfl_xor(int *a)
{
  int local_sum = a[threadIdx.x + blockIdx.x*blockDim.x];

  for (int mask = WARP_SIZE / 2; mask>0; mask /= 2)
    local_sum += __shfl_xor(local_sum, mask);

  if (threadIdx.x%32 == 0)
    printf("Warp max is %d", local_sum)
}
```



mask==4

mask==2

mask==1



Warp Voting

- ☐ Warp shuffles allow data to be exchanged between threads in a warp
- ☐ Warp voting allows threads to test a condition across all threads in a warp
 - ☐ int all(condition)
 - ☐ True if the condition is met by all threads in the warp
 - ☐ int any(condition)
 - ☐ True is any thread in warp meets condition
 - ☐unsigned int ballot(condition)
 - ☐ Sets the nth bit of the return value based on the nth threads condition value
- □All warp voting functions are single instruction and act as barrier
 - ☐Only active threads participate, does not block like syncth:

Warp Voting Example

```
__global__ void voteAllKernel(unsigned int *input, unsigned int *result)
{
  int i = threadIdx.x + blockIdx.x*blockDim.x;
  int j = i % WARP_SIZE;

  int vote_result = all(input[i]);

  if (j==0)
    result[i / WARP_SIZE] = vote_result;
```

- ☐ For each first thread in the warp calculate if all threads in the warp have true valued input
- ☐ Save the warp vote to a compact array
 - ☐A reduction of factor 32



Global Communication Summary

- ☐ Shared memory is per thread block
- ☐Shuffles and voting for warp level
- ☐Atomics can be used for some global (grid wide) operations
- ☐ What about general global communication?
 - □Not possible within a kernel (*except in Volta not covered*)!
 - Remember a grid may not be entirely in flight on the device
 - ☐ Can be enforced by finishing the kernel

```
step1 <<<grid, blk >>>(input, step1_output);
// step1_output can safely be used as input for step2
step2 <<<grid, blk >>>(step1_output, step2_output);
```



Summary

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Acknowledgements and Further Reading

- Predication: http://docs.nvidia.com/cuda/parallel-thread-execution/index.html#predicated-execution
- □Shuffling: http://on-demand.gputechconf.com/gtc/2013/presentations/S3174-Kepler-Shuffle-Tips-Tricks.pdf
- □Volta: https://devblogs.nvidia.com/cuda-9-features-revealed/

