

# relations

let  $A$  and  $B$  sets. a binary relation from  $a$  to  $b$  is a **subset** of  $A \times B$ . (baginti)  
 ↳ essentially, a binary relation is a set  $R$  of ordered pairs where the first element of each ordered pair comes from  $A$ , and the second element comes from  $B$ .

## example

$A = \{1, 2, 3\}$   $B = \{a, b, c\} \rightarrow R = \{(1, a), (1, b), (2, b), (2, c), (3, a)\}$  is a relation from  $A$  to  $B$ .  
 ↳ it does not have to contain all but it can

$1 R a, 1 R b$

\* a **function** is a relation where each element of  $B$  is mapped to by only one element of  $A$ .

binary relation on a set:

a subset of  $A \times A$ , or a relation from set  $A$  onto itself.

the number of relations from set  $A$  to itself:  $2^{n^2} = 2^{16}$

## properties

(dönüşlü)

**reflexive relations:** if  $(a, a) \in R$  for every  $a \in A$

↳  $R: \{(a, b) \mid a \leq b\} \rightarrow (a, a) \quad a \leq a$  so it is reflexive ✓

$R: \{(a, b) \mid a + b \leq 3\} \rightarrow (a, a) \quad 2a \leq 3$  so it is not reflexive ✗

**symmetric relations:** if  $(b, a) \in R$  whenever  $(a, b) \in R$

↳  $R: \{(a, b) \mid a \leq b\} \rightarrow (a, b) \quad (b, a)$  not always ✗

$R: \{(a, b) \mid a + b \leq 3\}$  ✓

**anti-symmetric relation:** if  $\forall a, b \in A$  if  $(a, b) \in R$  and  $(b, a) \in R$ , then  $a = b$

↳  $R: \{(a, b) \mid a \leq b\} \quad (2, 2) \quad (2, 2) \quad \checkmark$

$R: \{(a, b) \mid a + b \leq 3\} \quad (1, 2) \quad (2, 1) \quad \times$

$R: \{(a, b) \mid a > b\}$  (no counter example) ✓

(geçirli)

**transitive relation:** if whenever  $(a, b) \in R$  and  $(b, c) \in R$ , then  $(a, c) \in R \quad \forall a, b, c \in A$

↳  $R: \{(a, b) \mid a \leq b\} \quad (1, 1) \quad (1, 1) \quad a \leq b \quad b \leq c \rightarrow a \leq c \quad \checkmark$

$R: \{(a, b) \mid a + b \leq 3\} \quad (2, 1) \quad (1, 2) \quad (2, 2) \quad \times$

$R: \{(a, b) \mid a > b\} \quad a > b \quad b > c \quad a > c \quad \checkmark$

## equivalence relations

a relation on a set  $A$  is called an equivalence relation if it is reflexive

a relation on a set  $R$  is called an equivalence relation if it is reflexive, symmetric and transitive.

↳ two elements related by an equivalence relation are said to be equivalent, denoted  $a \sim b$ .

equivalence classes: the set of all elements equivalent to an element  $a$  of  $A$  through the equivalence relation  $R$  is called an equivalence class of  $a$ ,  $[a]_R$

### composing relations

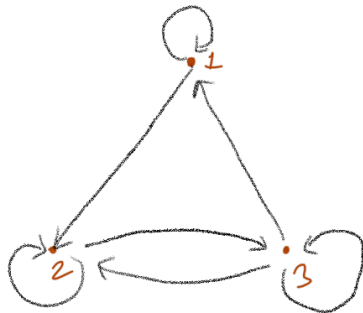
$$(a,b) \in R_1, (b,c) \in R_2 \rightarrow (a,c) \in R_2 \circ R_1$$

ex:  $R_1 = \{(1,3)\}$   $R_2 \circ R_1 = R_2(R_1) = \{(1,4), (1,5)\}$   
 $R_2 = \{(3,4), (3,5)\}$

$$* R = \{(1,1), (2,2), (3,2), (4,3)\}$$

$$R \circ R = \{(1,1), (2,1), (3,1), (4,2)\}$$

### representing relations using digraphs

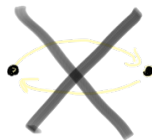


the relation  $R$  on  $\{1, 2, 3\}$  where  
 $R = \{(1,1), (1,2), (2,2), (2,3), (3,1), (3,2), (3,3)\}$

- reflexive ✓ (each node has )
- symmetric ✗ (every node must be connected mutually. )
- antisymmetric ✗ (no allowed)
- transitive ✗ ( $2 \rightarrow 3$   $3 \rightarrow 1$  but no  $2 \rightarrow 1$ )

### partial orders (POSET)

- antisymmetric
- reflexive
- transitive



example:

$$A = \{1, 2, 3, 4\}$$

$$R = \{(1,1), (1,2), (1,3), (1,4), (2,2), (2,3), (2,4), (3,3), (3,4), (4,4)\}$$

is  $(A, R)$  POSET?

↳ reflexive:  $(1,1), (2,2), (3,3), (4,4)$  ✓

↳ antisymmetric:  $(1,2) \checkmark (2,1) \times$   
 $(1,3) \checkmark (3,1) \times$   
 $(1,4) \checkmark (4,1) \times$   
 $(2,4) \checkmark (4,2) \times$   
 $(2,3) \checkmark (3,2) \times$   
 $(3,4) \checkmark (4,3) \times$

$(A, R)$  is POSET

transitive:  $\left. \begin{array}{l} (1,2) (2,3) \rightarrow (1,3) \checkmark \\ (1,2) (2,4) \rightarrow (1,4) \checkmark \\ (1,3) (3,4) \rightarrow (1,4) \checkmark \\ (2,3) (3,4) \rightarrow (2,4) \checkmark \end{array} \right\} \checkmark$

$$A = \{1,2,3\}$$

$$P(A) \subseteq$$

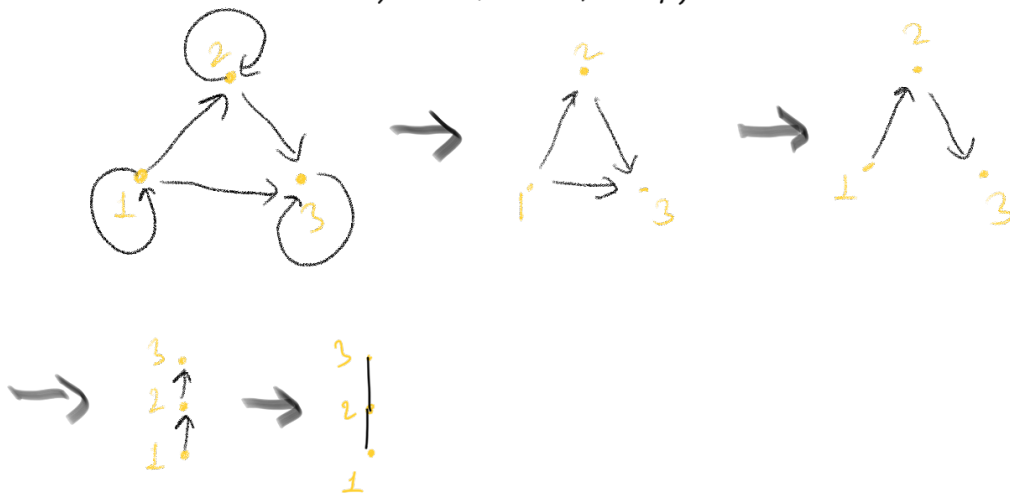
hasse diagram (ordering diagram)

↳ graphical representation of a poset

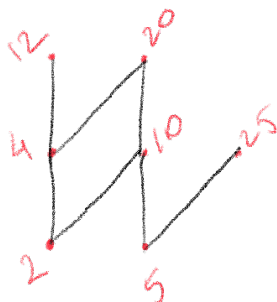
- 1 draw directed graph representation
- 2 remove loops (reflexivity)
- 3 remove the transitive nodes
- 4 arrange each edge so that all arrows point upwards
- 5 remove all arrowheads

example hasse diagram for  $(\{1,2,3\}, \leq)$

$$R = \{(1,1), (1,2), (1,3), (2,2), (2,3), (3,3)\}$$



example  $(\{2,4,5,10,12,20,25\}, |)$



maximal = 12, 20, 25

minimal = 2, 5

greatest = no unique element, no greatest elem

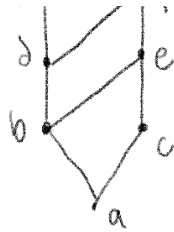
least = no unique element, no least elem

example



maximal = b, c

minimal = a



greatest = no unique element, no greatest element  
least = a

$A = \{a, b, c, d, e, f, g, h, j\}$

$B = \{a, b, c\}$  (a subset)

upper bound ( $B \rightarrow A$ )

GLB (greatest lower bound)

	a	b	c	d	e	f	g	h	j	
a	✓	✓	✓	✓	✓	✓	✓	✓	✓	$a \in J \checkmark$
b	x	✓	x	✓	✓	✓	✓	✓	✓	$b \in J \checkmark$
c	x	x	✓	x	✓	✓	x	✓	✓	

=  $\{e, f, h, j\}$

least bound ( $A \rightarrow B$ ) LUB (least upper bound)

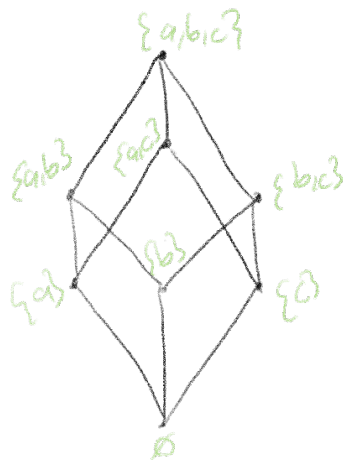
$\{a, b, c\} = \{e\}$  (en aşırsıda hepsini başlayan node)

lower bound ( $A \rightarrow B$ )

	a	b	c	
a	✓	✓	✓	$a \leq a, a \leq b, a \leq c$
b	x	✓	x	
c	x	x	✓	
d	x	x	x	
e	x	x	x	
f	x	x	x	
g	x	x	x	
h	x	x	x	

LB =  $\{a\}$   
GLB =  $\{a\}$

example; the hasse diagram for  $(P(\{a, b, c\}), \subseteq)$



Lattice

every pair of elements has both a least upper bound and a greatest lower bound (partially ordered set)

closure of relations

when relation does not have a certain property, by adding some elements to the relation, making relation to have that property

reflexive closure:

$R = \{(1,1), (1,2), (2,1), (3,2)\} \xrightarrow{\text{add } (2,2), (3,3)} \text{reflexive } (R \text{ on } A = \{1,2,3\})$   
↳ add all pairs in form  $(a,a) \in A$

symmetric closure:

$R = \{(1,1), (1,2), (2,2), (2,3), (3,1), (3,2)\} \xrightarrow{\text{add } (2,1), (1,3)} \text{symmetric}$   
↳ add the inverse of the  $R$

transitive closure:

↳ we cannot produce transitive closure by adding pairs.