2 - values and types

value = anything that exist, that can be computed, stored, take part in data structure: constants, function return values...

() C types: int, char, long, pointer, array.

Haskell types: bool, int, tuples, records, lists, functions ...

type = set of values, equipped with one or more operations that can be applied uniformly to all these values ex: $\{true, false\} \rightarrow and$, or, not $\times \{13, true, monday\}$ $\{..., 3, -2, 0, 1, 2 ...\} \rightarrow add$, multiply..

primitive types = values that cannot be decomposed into other sub values

5 C: int, float, double, char, long, short, pointers

4 Haskell: bool, int, float, function values

> Python: bool, int, float, str, functions

every programming lang. provides built-in primitive types. some also allow programs to define new primitives

cardinality of a type = the number of distinct values. denoted as # Type $\implies \#$ Bool = 2 # char = 25b # int = 2^{32}

·brilt-in types: can differ with the purpose of the lang. but some types are common but they have different names. ex: Integer -int

*user defined primitive types = enumerated types

List per defined primitive types

Li

type Population is range 0... 1010; (values >> Population = £0,..., 1010}) type Month is (jan, feb, ..., dec);

• discrete ordinal primitive types = datatypes values which have one to one mapping to a range of integers.

- they can be array indices, switch /case labels

- they can be used as for loop variables (in pascal)

L) C: ordinal types

> Pascal, Ada: distinct types

they are different types depending on the composition type.

- word produce conduct, topics, receives, La disjoint union (C-sunion, pascal- variant record, haskell -> lata) H mapping (arrays, functions) bowerset (pascal) set datatype) Grecursive composition (lists, trees, complex duta structures)

· cartesian product = values of several (mostly different) types are grouped into tuples $S = \{a_1b_1c_3\}$ $S \times T = \{(a_1L),(a_1L),(b_1L),(b_1L),(c_1L),(c_1L)\}$

 $\#(S\times T) = \#S, \#T$

multiple 7 (a,b,c) {x,y,z}

homogeneout cartesian product \Rightarrow $S^n = S \times S \times S$ (whose components are both chosen from $S^n = S \times S \times S$) $\Rightarrow 0$ tuple, not empty set (like void) in C

· distoint union= $S = \{ 1, 2, 3 \}$ $S + T = \{ \text{ left } 1, \text{ left } 2, \text{ left } 3, \text{ right } 4 \}$ $T = \{ 3, 4 \}$

#(S+T) = #S + #T

with these tags, we can check wheter the variant was chosen from S or T. and we can all multiple of a value

· mappings = the set of all possible mapping $S = \{a, b\}$ $S \rightarrow T = \{\{a \rightarrow \bot, b \rightarrow 1\}, \{a \rightarrow \bot, b \rightarrow 2\}, \{a \rightarrow \bot, b \rightarrow 3\}, \{a \rightarrow \bot$, ورط, عدم ع , وحط, عمر عدم ع , وعمر عدم ع في الم { 250-d, 50-d, 50-d, 50-b}, { 10-d, 50-b}

 $\#(S \rightarrow T) = \#T^{\#S}$

only integer domain values Stored in the memory *functions*

all types of mappings possible defined by algorithms efficiency, resource usage side effect, output, error, terminate problem

· powerset = the set of all subsets, restricted and special datatype, only in Paxal 5= {1,2,3} P(s)= { {\$, {1,3}, {23, {1,2}, {1,33, {2,33, {1,2,33}}} $\# P(s) = 2^{\# S}$

recursive types = types including themselves in composition in general they defined as: R=...t(..R...R..)

a recursive set equation always has at least solution that is a subset of every other solution

cardinality of a recursive type = infinite Lisever if every individual value of the type is finite) ex: the set of lists is infinitely large, although every individual list in that set is finite

• lists = sequence of values
$$S = Int \times S + \{null\}$$

List $\alpha = \alpha \times (list \alpha) + \{empty\}$

type systems

groups valves into types

type check = checking the types of the operands before operation itself (always)

static typing

operands are checked at compile—time

because each variables and expression

has a fixed type

most of the high lev. lang

clott (strict type checking)

Hastell, ML (type inference)

* More efficient, secure

dynamic typing

operands are checked after they are

(executed)

Lomputed, but before performing the

operation, at run-time

Lower values have fixed types,

but variables and expressions do not have

rython, prolog (interpreted long)

slower (be of multiple check)

and also uses extrem storage

no security

more flexible

type equivalence

T, = T2 how to decide

name equivalence

types should be defined at the same exact place

most lang use

struct a { int int }

struct b { int int }

structural equivalence

types should have same value set (mathematical set equality) * if their structures are same, even if their names are different you can compose them * hard to implement Ly it will give error it you compose or assign than. (because their names) are different

type completeress

first order values

- assignment
- func, parameter
- ' take part in composition
- · return value from a func

functions -> second order values in Pascal, Fortro

pointers in C

first order values in Haskell

C Types:

	Primitive	Array	Struct	Func.
Assignment	\checkmark	×	\checkmark	×
Function parameter	\checkmark	×	\checkmark	X
Function return	\checkmark	X		X
In compositions	\checkmark	(\checkmark)	\checkmark	×
		(')		

Haskell Types:

	Primitive	Array	Struct	Func.
Variable definition	\checkmark	\checkmark	\checkmark	\checkmark
Function parameter	\checkmark	\checkmark	\checkmark	\checkmark
Function return	\checkmark	\checkmark	\checkmark	\checkmark
In compositions	\checkmark	\checkmark	\checkmark	\checkmark

expressions

a segment that is evaluated, and produces a value

literals = return itself 3-3 variable and constant access = 9=3 -> 3 aggregates to construct composite value without decleration/definition x = (12, "ali") variable references = pointers are not references function calls conditional expressions iterative expressions = hastell