

5 - abstraction and parameter passing

abstraction = make a program reusable by enclosing it in a body, hiding the details, and defining a mechanism to access it.

↳ separating the usage and implementation of program segments

ex: programming languages over machine language

☆ declare once, use many times

☆ code reusability

abstraction types

function and procedure abstraction =

- functions are abstractions over expressions

- procedures are abstractions over commands

↳ no value but contains executable statements as detail (void functions of C) ^{exists because of their side effects}

selector abstraction = `[]` operator selects elements of an array

generic abstraction = abstraction over declaration

- same declaration pattern applied to different data types (template functions in C++)

`template <class T>`: type of T declared at compile time

↳ for every different T (int, float...) it declares different functions

↳ it is different from polymorphism (done at run time) (faster)

So it is more efficient since everything done at compile time

iterator abstraction = iteration over a user defined data structure

↳ in C++, `begin()` and `end()` functions of linked list.

entity

expression

command

selector

declaration

command block

abstraction

function

procedure

selector function

generic

iterator

parameters =

formal parameters = a variable and its type as they appear in the prototype of the function

actual parameters = the variable or expression corresponding to a formal parameter that appears in the function in calling environment

modes: **in** = passes info from caller to callee

out = callee writes values in caller

in/out = caller tells callee value of variable, which may be updated by callee

parameter passing mechanisms

call mechanisms: call by value, call by reference

copy mechanism: assignment based

↳ C only allows copy-in mechanism. → is called as **pass by value**.

binding mechanism: based on binding of the formal parameter variable/identifier to actual parameter value/identifier

↳ **constant binding**: formal parameter is constant during the function (in Haskell)

↳ **variable binding**: formal parameter variable is bound to the actual parameter variable. same memory area is shared by two variable references → also known as **pass by reference**

pass by name: substitution based

textually substitute the argument in a procedure call for the corresponding parameter in the body of the procedure

↳ the argument expression is re-evaluated each time the formal parameter is passed (**normal evaluation**)

↳ the procedure can change the values of variables used in the argument expression and hence change the expression's value (in imperative languages not in Haskell)

evaluation order

normal order evaluation: actual parameter is re-evaluated each time it is used

↳ has an advantage ex: $x \ y = \text{if } x > 0 \text{ then } y \text{ else } x$

if this is not true y is not evaluated (y can be like $\frac{0}{5}$)

↳ some parameters evaluated once or more, or not evaluated at all

↳ side effects are repeated

eager evaluation: actual parameters are evaluated first, then passed

↳ faster than normal evaluation

↳ evaluates parameters exactly once

church-rosser property = if a function does not have side effect, normal order and eager evaluation produce same result (order of the evaluation does not matter)
(or others)

↳ then more efficient solution: **lazy evaluation**

lazy evaluation: faster, computes the same result

↳ evaluates the parameter first time that is needed (does not reevaluate)

↳ so, evaluates parameters at most once

ex: Haskell, infinite values

corresponding principle = if for each form of declaration there exists a corresponding parameter mechanism, the PL satisfies this principle

C/C++ ✓ pascal X