

# Verilog Size Checking

## Expressions

- $\mathcal{A}$  is the set of atoms in an expression. In the AST, they correspond to leaves. In System-Verilog, they refer to what the standard calls an “operand” (variables, integers, function calls, slices of a variable, etc.).
- $\mathcal{R}$  is the set of resizable expressions. It corresponds to the expression whose top level is one of: atom, cast, comparisons, logic operation, reduction, assignments, shift assignments, concatenation, replication and inside operation.
- $\mathcal{E}$  is the set of System-Verilog expressions. We have  $\mathcal{A} \subset \mathcal{E}$ . An expression either contains an expression or is an atom.

## Rules

We use the following notations:

- $\Gamma$  maps atoms to their size,
- $\Phi$  maps lvalues to their size,
- The statement  $e \Rightarrow t$  means “ $e$  has a size  $t$ ”,
- The statement  $e \Leftarrow t$  means “ $e$  can be resized to  $t$ ”.

### Base case

$$\frac{\Gamma(e) = s \quad e \in \mathcal{A}}{e \Rightarrow s} \text{Atom} \Rightarrow$$

### Resize case

$$\frac{e \Rightarrow s \quad s \leq t \quad e \in \mathcal{R}}{e \Leftarrow t} \text{Resize} \Leftarrow$$

## Operators

- $\oplus \in \{+, -, *, /, \%, \&, |, \wedge, \sim, \sim\sim\}$ :

$$\frac{a \Rightarrow t \quad b \Leftarrow t}{a \oplus b \Rightarrow t} \text{LBinOp} \Rightarrow$$

$$\frac{a \Leftarrow t \quad b \Rightarrow t}{a \oplus b \Rightarrow t} \text{RBinOp} \Rightarrow$$

$$\frac{a \Leftarrow t \quad b \Leftarrow t}{a \oplus b \Leftarrow t} \text{BinOp} \Leftarrow$$

- $\oplus \in \{+, -, ++, --, \sim\}$ :

$$\frac{e \Rightarrow t}{\oplus e \Rightarrow t} \text{UnOp} \Rightarrow$$

$$\frac{e \Leftarrow t}{\oplus e \Leftarrow t} \text{UnOp} \Leftarrow$$

- $\oplus \in \{\$signed, \$unsigned\}$ :

$$\frac{e \Rightarrow t}{\oplus(e) \Rightarrow t} \text{Cast} \Rightarrow$$

- $\oplus \in \{==, !=, ==?, !=?, ==, !=, >, >=, <, <=\}$ :

$$\frac{a \Rightarrow t \quad b \Leftarrow t}{a \oplus b \Rightarrow 1} \text{LCmp} \Rightarrow$$

$$\frac{a \Leftarrow t \quad b \Rightarrow t}{a \oplus b \Rightarrow 1} \text{RCmp} \Rightarrow$$

- $\oplus \in \{\&\&, ||, -, <->\}$ :

$$\frac{a \Rightarrow t_a \quad b \Rightarrow t_b}{a \oplus b \Rightarrow 1} \text{Logic} \Rightarrow$$

- $\oplus \in \{\&, \sim\&, |, \sim|, \wedge, \sim\wedge, \wedge\sim, !\}$ :

$$\frac{e \Rightarrow t}{\oplus e \Rightarrow 1} \text{Red} \Rightarrow$$

- $\oplus \in \{>>, <<, **, >>>, <<<\}$ :

$$\frac{a \Rightarrow t \quad b \Rightarrow t_b}{a \oplus b \Rightarrow t} \text{ Shift} \Rightarrow$$

$$\frac{a \Leftarrow t \quad b \Rightarrow t_b}{a \oplus b \Leftarrow t} \text{ Shift} \Leftarrow$$

- $\oplus \in \{=, +=, -=, *=, /=, \%=, \&=, |=, ^=\}$ :

$$\frac{\phi(l) = t \quad e \Leftarrow t}{l \oplus e \Rightarrow t} \text{ LAssign} \Rightarrow$$

$$\frac{\phi(l) = t \quad e \Rightarrow t_e \quad t < t_e}{l \oplus e \Rightarrow t} \text{ RAssign} \Rightarrow$$

- $\oplus \in \{<<=, >>=, <<<=, >>>=\}$ :

$$\frac{\phi(l) = t \quad e \Rightarrow t_e}{l \oplus e \Rightarrow t} \text{ AssignShift} \Rightarrow$$

- If expression:

$$\frac{e \Rightarrow t_e \quad a \Rightarrow t \quad b \Leftarrow t}{e?a:b \Rightarrow t} \text{ LCond} \Rightarrow$$

$$\frac{e \Rightarrow t_e \quad a \Leftarrow t \quad b \Rightarrow t}{e?a:b \Rightarrow t} \text{ RCond} \Rightarrow$$

$$\frac{e \Rightarrow t_e \quad a \Leftarrow t \quad b \Leftarrow t}{e?a:b \Leftarrow t} \text{ Cond} \Leftarrow$$

- Concatenation:

$$\frac{e_1 \Rightarrow t_1 \quad \dots \quad e_k \Rightarrow t_k \quad t = t_1 + \dots + t_k}{\{e_1, \dots, e_k\} \Rightarrow t} \text{ Concat} \Rightarrow$$

- Replication:

$$\frac{i \in \mathbb{N} \quad e \Rightarrow t_e \quad t = i \times t_e}{\{i \ e\} \Rightarrow t} \text{ Repl} \Rightarrow$$

- Inside:

$$\frac{a \Rightarrow t \quad e_1 \Leftarrow t \quad \dots \quad e_k \Leftarrow t}{a \text{ inside } \{e_1, \dots, e_k\} \Rightarrow 1} \text{ LInside} \Rightarrow$$

$$\frac{i \in \{1, \dots, k\} \quad a \Leftarrow t \quad e_1 \Leftarrow t \quad \dots \quad e_{i-1} \Leftarrow t \quad e_i \Rightarrow t \quad e_{i+1} \Leftarrow t \quad \dots \quad e_k \Leftarrow t}{a \text{ inside } \{e_1, \dots, e_k\} \Rightarrow 1} \text{ RInside} \Rightarrow$$