

# Verilog Size Checking

The statement  $\Gamma \mid n \vdash e : \tau, X$  mean:

In the variable context  $\Gamma$  and the size context  $n$ , the size of the expression  $e$  is the value of the variable  $\tau$ , with  $X$  the size-tagged set of equivalence classes of size variables.

So we have:

- $\Gamma : \mathcal{O} \rightarrow \mathbb{N}$  a mapping from operand to their size, with  $\mathcal{O}$  the set of operands,
- $n \in \mathbb{N} \cup \{*\}$  a size context (with  $*$  for the empty one),
- $e \in \mathcal{E}$ , with  $\mathcal{E}$  the set of expressions,
- $\tau \in \mathcal{V}$ , with  $\mathcal{V}$  the set of size variable,
- $X = (S, f)$  with  $S \in 2^{\mathcal{V}}$  the set of equivalence classes and  $f : S \rightarrow \mathbb{N}$  a size valuation for each equivalence class.

We use the following notations:

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$$\bowtie : \begin{cases} \mathbb{N} \times (\mathbb{N} \cup \{*\}) & \rightarrow & \mathbb{N} \\ (m, *) & \mapsto & m \\ (m, n) & \mapsto & \max(m, n) \end{cases}$$

- $\Delta_X : \mathcal{V} \rightarrow \mathbb{N}$  the function that given a size variable  $x$  returns its size in the environment  $X = (S, f)$ . We have  $\Delta_X(x) = f([x]_S)$  with  $[x]_S$  the equivalence class of  $x$  in  $S$ .
- $\{v := s\}$ : the declaration of a *fresh* size variable  $v$ . This adds a new equivalence class  $\{v\}$  in  $X$  with the size mapping  $\{v\} \mapsto s$ .
- $X/\alpha \sim \beta$  the operation that combines two classes. For the newly created class, the valuation function gives the maximum of the previous classes.

## Base case

$$\frac{\Gamma(x) = m}{\Gamma \mid n \vdash x : v, \{v := m \bowtie n\}}$$

## Operators

$$\oplus \in \{+, -, *, /, \%, \&, |, \wedge, \sim, \sim\sim\}$$

$$\frac{\Gamma \mid n \vdash a : \alpha, A \quad \Gamma \mid n \vdash b : \beta, B}{\Gamma \mid n \vdash a \oplus b : \alpha, A \sqcup B / \alpha \sim \beta}$$

$$\oplus \in \{+, -, ++, --, \sim\}$$

$$\frac{\Gamma \mid n \vdash e : \varepsilon, E}{\Gamma \mid n \vdash \oplus e : \varepsilon, E}$$

$$\oplus \in \{++, --\}$$

$$\frac{\Gamma \mid n \vdash e : \varepsilon, E}{\Gamma \mid n \vdash e \oplus : \varepsilon, E}$$

$$\oplus \in \{\$signed, \$unsigned, signed', unsigned'\}$$

$$\frac{\Gamma \mid * \vdash e : \varepsilon, E}{\Gamma \mid n \vdash \oplus(e) : v, E \sqcup \{v := \Delta_E(\varepsilon) \bowtie n\}}$$

$$\oplus \in \{==, !=, ==?, !=?, ==, !=, >, >=, <, <=\}$$

$$\frac{\Gamma \mid * \vdash a : \alpha, A \quad \Gamma \mid * \vdash b : \beta, B}{\Gamma \mid n \vdash a \oplus b : v, A \sqcup B \sqcup \{v := 1 \bowtie n\}}$$

$$\oplus \in \{\&\&, ||, ->, <->\}$$

$$\frac{\Gamma \mid * \vdash a : \alpha, A \quad \Gamma \mid * \vdash b : \beta, B}{\Gamma \mid n \vdash a \oplus b : v, A \sqcup B \sqcup \{v := 1 \bowtie n\}}$$

$$\oplus \in \{\&, \sim\&, |, \sim|, \wedge, \sim\wedge, \wedge\sim, !\}$$

$$\frac{\Gamma \mid * \vdash e : \varepsilon, E}{\Gamma \mid n \vdash \oplus e : \varepsilon, E \sqcup \{v := 1 \bowtie n\}}$$

$$\oplus \in \{>>, <<, **, >>>, <<<\}$$

$$\frac{\Gamma \mid n \vdash a : \alpha, A \quad \Gamma \mid * \vdash b : \beta, B}{\Gamma \mid n \vdash a \oplus b : \alpha, A \sqcup B}$$

$$\oplus \in \{=, +=, -=, *=, /=, \%=, \&=, |=, \wedge=\}$$

$$\frac{\Gamma \mid * \vdash a : \alpha, A \quad \Gamma \mid \Delta_A(\alpha) \vdash b : \beta, B}{\Gamma \mid n \vdash a \oplus b : v, A \sqcup B \sqcup \{v := \Delta_A(\alpha) \bowtie n\}}$$

$$\oplus \in \{<<=, >>=, <<<=, >>>=\}$$

$$\frac{\Gamma \mid * \vdash a : \alpha, A \quad \Gamma \mid * \vdash b : \beta, B}{\Gamma \mid n \vdash a \oplus b : v, A \sqcup B \sqcup \{v := \Delta_A(\alpha) \bowtie n\}}$$

$$\frac{\Gamma \mid * \vdash e : \varepsilon, E \quad \Gamma \mid n \vdash a : \alpha, A \quad \Gamma \mid n \vdash b : \beta, B}{\Gamma \mid n \vdash e?a:b : \alpha, E \sqcup A \sqcup B / \alpha \sim \beta}$$

$$\frac{\Gamma \mid * \vdash e_1 : \varepsilon_1, E_1 \quad \dots \quad \Gamma \mid * \vdash e_n : \varepsilon_n, E_n}{\Gamma \mid n \vdash \{e_1, \dots, e_n\} : v, E_1 \sqcup \dots \sqcup E_n \sqcup \{v := (\Delta_{E_1}(\varepsilon_1) + \dots + \Delta_{E_n}(\varepsilon_n)) \bowtie n\}}$$

$$\frac{i \in \mathbb{N} \quad \Gamma \mid * \vdash e_1 : \varepsilon_1, E_1 \quad \dots \quad \Gamma \mid * \vdash e_n : \varepsilon_n, E_n}{\Gamma \mid n \vdash \{i\{e_1, \dots, e_n\}\} : v, E_1 \sqcup \dots \sqcup E_n \sqcup \{v := i \times (\Delta_{E_1}(\varepsilon_1) + \dots + \Delta_{E_n}(\varepsilon_n)) \bowtie n\}}$$

$$\frac{\Gamma \mid * \vdash a : \alpha, A \quad \Gamma \mid * \vdash e_1 : \varepsilon_1, E_1 \quad \dots \quad \Gamma \mid * \vdash e_n : \varepsilon_n, E_n}{\Gamma \mid n \vdash a \text{ inside } \{e_1, \dots, e_n\} : v, A \sqcup E_1 \sqcup \dots \sqcup E_n \sqcup \{v := 1 \bowtie n\} / \alpha \sim \varepsilon_1, \dots, \alpha \sim \varepsilon_n}$$