

Verilog Size Checking

Expressions

- \mathcal{A} is the set of atoms in an expression. In the AST, they correspond to leaves. In System-Verilog, they refer to what the standard calls an “operand” (variables, integers, function calls, slices of a variable, etc.).
- \mathcal{R} is the set of resizable expressions. It corresponds to the expression whose top level is one of: atom, cast, comparisons, logic operation, reduction, assignments, shift assignments, concatenation, replication and inside operation.
- \mathcal{E} is the set of System-Verilog expressions. We have $\mathcal{A} \subset \mathcal{E}$. An expression either contains an expression or is an atom.

Rules

We use the following notations:

- Γ maps atoms to their size,
- Φ maps lvalues to their size,
- The statement $e \Rightarrow t$ means “ e has a size t ”,
- The statement $e \Leftarrow t$ means “ e can be resized to t ”.

Base case

$$\frac{\Gamma(e) = s \quad e \in \mathcal{A}}{e \Rightarrow s} \text{Atom} \Rightarrow$$

Resize case

$$\frac{e \Rightarrow s \quad s \leq t \quad e \in \mathcal{R}}{e \Leftarrow t} \text{Resize} \Leftarrow$$

Operators

- $\oplus \in \{+, -, *, /, \%, \&, |, \wedge, \sim, \sim\sim\}$:

$$\frac{a \Rightarrow t \quad b \Leftarrow t}{a \oplus b \Rightarrow t} \text{LBinOp} \Rightarrow$$

$$\frac{a \Leftarrow t \quad b \Rightarrow t}{a \oplus b \Rightarrow t} \text{RBinOp} \Rightarrow$$

$$\frac{a \Leftarrow t \quad b \Leftarrow t}{a \oplus b \Leftarrow t} \text{BinOp} \Leftarrow$$

- $\oplus \in \{+, -, ++, --, \sim\}$:

$$\frac{e \Rightarrow t}{\oplus e \Rightarrow t} \text{UnOp} \Rightarrow$$

$$\frac{e \Leftarrow t}{\oplus e \Leftarrow t} \text{UnOp} \Leftarrow$$

- $\oplus \in \{\$signed, \$unsigned\}$:

$$\frac{e \Rightarrow t}{\oplus(e) \Rightarrow t} \text{Cast} \Rightarrow$$

- $\oplus \in \{==, !=, ==?, !=?, ==, !=, >, >=, <, <= \}$:

$$\frac{a \Rightarrow t \quad b \Leftarrow t}{a \oplus b \Rightarrow 1} \text{LCmp} \Rightarrow$$

$$\frac{a \Leftarrow t \quad b \Rightarrow t}{a \oplus b \Rightarrow 1} \text{RCmp} \Rightarrow$$

- $\oplus \in \{\&\&, ||, -, <->\}$:

$$\frac{a \Rightarrow t_a \quad b \Rightarrow t_b}{a \oplus b \Rightarrow 1} \text{Logic} \Rightarrow$$

- $\oplus \in \{\&, \sim\&, |, \sim|, \wedge, \sim\wedge, \wedge\sim, !\}$:

$$\frac{e \Rightarrow t}{\oplus e \Rightarrow 1} \text{Red} \Rightarrow$$

- $\oplus \in \{>, <, **, >>, <<<\}$:

$$\frac{a \Rightarrow t \quad b \Rightarrow t_b}{a \oplus b \Rightarrow t} \text{ Shift} \Rightarrow$$

$$\frac{a \Leftarrow t \quad b \Rightarrow t_b}{a \oplus b \Leftarrow t} \text{ Shift} \Leftarrow$$

- $\oplus \in \{=, +=, -=, *=, /=, \%=, \&=, |=, ^=\}$:

$$\frac{\phi(l) = t \quad e \Leftarrow t}{l \oplus e \Rightarrow t} \text{ LAssign} \Rightarrow$$

$$\frac{\phi(l) = t \quad e \Rightarrow t_e \quad t < t_e}{l \oplus e \Rightarrow t} \text{ RAssign} \Rightarrow$$

- $\oplus \in \{<=<, >>=, <<<=, >>>=\}$:

$$\frac{\phi(l) = t \quad e \Rightarrow t_e}{l \oplus e \Rightarrow t} \text{ AssignShift} \Rightarrow$$

- If expression:

$$\frac{e \Rightarrow t_e \quad a \Rightarrow t \quad b \Leftarrow t}{e?a:b \Rightarrow t} \text{ LCond} \Rightarrow$$

$$\frac{e \Rightarrow t_e \quad a \Leftarrow t \quad b \Rightarrow t}{e?a:b \Rightarrow t} \text{ RCond} \Rightarrow$$

$$\frac{e \Rightarrow t_e \quad a \Leftarrow t \quad b \Leftarrow t}{e?a:b \Leftarrow t} \text{ Cond} \Leftarrow$$

- Concatenation:

$$\frac{e_1 \Rightarrow t_1 \quad \dots \quad e_k \Rightarrow t_k \quad t = t_1 + \dots + t_k}{\{e_1, \dots, e_k\} \Rightarrow t} \text{ Concat} \Rightarrow$$

- Replication:

$$\frac{i \in \mathbb{N} \quad e \Rightarrow t_e \quad t = i \times t_e}{\{i \ e\} \Rightarrow t} \text{ Repl} \Rightarrow$$