

Verilog Size Checking

Path

Definition of paths and their relation to sub-expressions.

Expressions

- \mathcal{A} is the set of atoms in an expression. In the AST, they correspond to leaves. In System-Verilog, they refer to what the standard calls an “operand” (variables, integers, function calls, slices of a variable, etc.).
- \mathcal{R} is the set of resizable expressions. It corresponds to the expression whose top level is one of: atom, cast, comparisons, logic operation, reduction, assignments, shift assignments, concatenation, replication and inside operation.
- \mathcal{E} is the set of System-Verilog expressions. We have $\mathcal{A} \subset \mathcal{E}$. An expression either contains an expression or is an atom.

Rules

We use the following notations:

- Γ maps atoms to their size,
- Φ maps lvalues to their size,
- The statement $e \Rightarrow t \dashv f$ means “ e has size t , with f mapping each sub-expressions of e to their size”,
- The statement $e \Leftarrow t \dashv f$ means “ e can be resized to t , with f mapping each sub-expressions of e to their size”.

Function combinator

$$\begin{aligned} \text{Binary}(t, f, g) : \begin{cases} [] & \mapsto t \\ \text{Left} :: p & \mapsto f(p) \\ \text{Right} :: p & \mapsto g(p) \end{cases} & \quad \text{Unary}(f, t) : \begin{cases} [] & \mapsto t \\ \text{Arg} :: p & \mapsto f(p) \end{cases} \\ \text{Narry}(t, f_1, \dots, f_k) : \begin{cases} [] & \mapsto t \\ \text{Args}(i) :: p & \mapsto f_i(p) \end{cases} & \quad \text{Ternary}(t, f, g, h) : \begin{cases} [] & \mapsto t \\ \text{Arg} :: p & \mapsto f(p) \\ \text{Left} :: p & \mapsto g(p) \\ \text{Right} :: p & \mapsto h(p) \end{cases} \end{aligned}$$

Base case

$$\frac{\Gamma(e) = s \quad e \in \mathcal{A}}{e \Rightarrow s \dashv \{ [] \mapsto s \}} \text{Atom} \Rightarrow$$

Resize case

$$\frac{e \Rightarrow s \dashv f \quad s \leq t \quad e \in \mathcal{R}}{e \Leftarrow t \dashv f [[] \mapsto t]} \text{Resize} \Leftarrow$$

Operators

- $\oplus \in \{+, -, *, /, \%, \&, |, \wedge, \sim, \sim\sim\}$:

$$\frac{a \Rightarrow t \dashv f_a \quad b \Leftarrow t \dashv f_b}{a \oplus b \Rightarrow t \dashv \text{Binary}(t, f_a, f_b)} \text{LBinOp} \Rightarrow \quad \frac{a \Leftarrow t \dashv f_a \quad b \Rightarrow t \dashv f_b}{a \oplus b \Rightarrow t \dashv \text{Binary}(t, f_a, f_b)} \text{RBinOp} \Rightarrow$$

$$\frac{a \Leftarrow t \dashv f_a \quad b \Leftarrow t \dashv f_b}{a \oplus b \Leftarrow t \dashv \text{Binary}(t, f_a, f_b)} \text{BinOp} \Leftarrow$$

- $\oplus \in \{+, -, ++, --, \sim\}$:

$$\frac{e \Rightarrow t \dashv f}{\oplus e \Rightarrow t \dashv \text{Unary}(t, f)} \text{UnOp} \Rightarrow \quad \frac{e \Leftarrow t \dashv f}{\oplus e \Leftarrow t \dashv \text{Unary}(t, f)} \text{UnOp} \Leftarrow$$

- $\oplus \in \{\text{\$signed}, \text{\$unsigned}\}$:

$$\frac{e \Rightarrow t \dashv f}{\oplus(e) \Rightarrow t \dashv \text{Unary}(t, f)} \text{Cast} \Rightarrow$$

- $\oplus \in \{==, !=, ==?, !=?, ==, !=, >, >=, <, <= \}$:

$$\frac{a \Rightarrow t \dashv f_a \quad b \Leftarrow t \dashv f_b}{a \oplus b \Rightarrow 1 \dashv \text{Binary}(1, f_a, f_b)} \text{LCmp} \Rightarrow \quad \frac{a \Leftarrow t \dashv f_a \quad b \Rightarrow t \dashv f_b}{a \oplus b \Rightarrow 1 \dashv \text{Binary}(1, f_a, f_b)} \text{RCmp} \Rightarrow$$

- $\oplus \in \{\&\&, ||, -, <->\}$:

$$\frac{a \Rightarrow t_a \dashv f_a \quad b \Rightarrow t_b \dashv f_b}{a \oplus b \Rightarrow 1 \dashv \text{Binary}(1, f_a, f_b)} \text{Logic} \Rightarrow$$

- $\oplus \in \{\&, \sim\&, |, \sim|, \wedge, \sim\wedge, \hat{\sim}, !\}$:

$$\frac{e \Rightarrow t \dashv f}{\oplus e \Rightarrow 1 \dashv \text{Unary}(1, f)} \text{Red} \Rightarrow$$

- $\oplus \in \{>>, <<, **, >>>, <<<\}$:

$$\frac{a \Rightarrow t \dashv f_a \quad b \Rightarrow t_b \dashv f_b}{a \oplus b \Rightarrow t \dashv \text{Binary}(t, f_a, f_b)} \text{Shift} \Rightarrow \quad \frac{a \Leftarrow t \dashv f_a \quad b \Rightarrow t_b \dashv f_b}{a \oplus b \Leftarrow t \dashv \text{Binary}(t, f_a, f_b)} \text{Shift} \Leftarrow$$

- $\oplus \in \{=, +=, -=, *=, /=, \%, \&=, |=, \wedge=\}$:

$$\frac{\phi(l) = t \quad e \Leftarrow t \dashv f}{l \oplus e \Rightarrow t \dashv \text{Unary}(t, f)} \text{LAssign} \Rightarrow \quad \frac{\phi(l) = t \quad e \Rightarrow t_e \dashv f \quad t < t_e}{l \oplus e \Rightarrow t \dashv \text{Unary}(t, f)} \text{RAssign} \Rightarrow$$

- $\oplus \in \{<=<, >>=, <<=<, >>=<\}$:

$$\frac{\phi(l) = t \quad e \Rightarrow t_e \dashv f}{l \oplus e \Rightarrow t \dashv \text{Unary}(t, f)} \text{AssignShift} \Rightarrow$$

- If expression:

$$\frac{e \Rightarrow t_e \dashv f_e \quad a \Rightarrow t \dashv f_a \quad b \Leftarrow t \dashv f_b}{e?a:b \Rightarrow t \dashv \text{Ternary}(t, f_e, f_a, f_b)} \text{LCond} \Rightarrow \quad \frac{e \Rightarrow t_e \dashv f_e \quad a \Leftarrow t \dashv f_a \quad b \Rightarrow t \dashv f_b}{e?a:b \Rightarrow t \dashv \text{Ternary}(t, f_e, f_a, f_b)} \text{RCond} \Rightarrow$$

$$\frac{e \Rightarrow t_e \dashv f_e \quad a \Leftarrow t \dashv f_a \quad b \Leftarrow t \dashv f_b}{e?a:b \Leftarrow t \dashv \text{Ternary}(t, f_e, f_a, f_b)} \text{Cond} \Leftarrow$$

- Concatenation:

$$\frac{e_1 \Rightarrow t_1 \dashv f_1 \quad \dots \quad e_k \Rightarrow t_k \dashv f_k \quad t = t_1 + \dots + t_k}{\{e_1, \dots, e_k\} \Rightarrow t \dashv \text{Narry}(t, f_1, \dots, f_k)} \text{Concat} \Rightarrow$$

- Replication:

$$\frac{i \in \mathbb{N} \quad e \Rightarrow t_e \dashv f \quad t = i \times t_e}{\{i \ e\} \Rightarrow t \dashv \text{Unary}(t, f)} \text{Repl} \Rightarrow$$