

Verilog Size Checking

Expressions

- \mathcal{A} is the set of atoms in an expression. In the AST, they correspond to leaves. In System-Verilog, they refer to what the standard calls an “operand” (variables, integers, function calls, slices of a variable, etc.).
- \mathcal{E} is the set of System-Verilog expressions. We have $\mathcal{A} \subset \mathcal{E}$. An expression either contains an expression or is an atom.

Contexts

- Set of contexts identifiers: $\mathcal{C} \simeq \mathbb{N}$.
- A context $q \subseteq \mathcal{Q}$ with:

$$\mathcal{Q} := \underbrace{\{\text{Size } s \mid s \in \mathbb{N}\}}_{\text{Atomic size}} \sqcup \underbrace{\{\text{Id } c \mid c \in \mathcal{C}\}}_{\text{Identity dependencies}} \sqcup \underbrace{\{\text{Add } (c_1, c_2) \mid c_1 \in \mathcal{C}, c_2 \in \mathcal{C}\}}_{\text{Additive dependencies}} \sqcup \underbrace{\{\text{Mul } (s, c) \mid c \in \mathcal{C}, s \in \mathbb{N}\}}_{\text{Multiplicative dependencies}}$$

A context represent all the information needed to compute the size of associated expressions.

- $\Pi : \mathcal{C} \rightarrow 2^{\mathcal{Q}}$ a partial mapping from context identifiers to the smallest possible context set. This means that:

$$\Pi(c) \subset \Pi'(c) \implies \Pi(c) = \Pi'(c)$$

Rules

The statement $\Pi \vdash e : c$ means:

In the context environment Π , all the information needed to compute the size of the expression e is in the context-set identified by c in Π .

We use the following notations:

- Γ compute the size of an atom.
- Φ compute the size of a lvalue.

Base case

$$\frac{s = \Gamma(e) \quad \Pi = \{c \mapsto q\} \quad \text{Size } s \in q \quad e \in \mathcal{A}}{\Pi \vdash e : c}$$

Operators

- $\oplus \in \{+, -, *, /, \%, \&, |, \wedge, \sim, \sim\sim\}$:

$$\frac{\Pi \vdash a : c \quad \Pi \vdash b : c}{\Pi \vdash a \oplus b : c}$$

- $\oplus \in \{+, -, ++, --, \sim\}$:

$$\frac{\Pi \vdash e : c}{\Pi \vdash \oplus e : c}$$

- $\oplus \in \{\$signed, \$unsigned\}$:

$$\frac{\Pi' \vdash e : c' \quad c \notin \text{dom } \Pi' \quad \Pi = \Pi' \sqcup \{c \mapsto q\} \quad \text{Id } c' \in q}{\Pi \vdash \oplus(e) : c}$$

- $\oplus \in \{==, !=, ==?, !=?, ==, !=, >, >=, <, <= \}$:

$$\frac{\Pi' \vdash a : c' \quad \Pi' \vdash b : c' \quad c \notin \text{dom } \Pi' \quad \Pi = \Pi' \sqcup \{c \mapsto q\} \quad \text{Size } 1 \in q}{\Pi \vdash a \oplus b : c}$$

- $\oplus \in \{\&\&, ||, -, <->\}$:

$$\frac{\Pi_1 \vdash a : c_1 \quad \Pi_2 \vdash b : c_2 \quad \text{dom } \Pi_1 \cap \text{dom } \Pi_2 = \emptyset \quad c \notin \text{dom } \Pi_1 \sqcup \text{dom } \Pi_2 \quad \Pi = \Pi_1 \sqcup \Pi_2 \sqcup \{c \mapsto q\} \quad \text{Size } 1 \in q}{\Pi \vdash a \oplus b : c}$$

- $\oplus \in \{\&, \sim\&, |, \sim|, \wedge, \sim\wedge, \wedge\sim, !\}$:

$$\frac{\Pi' \vdash a : c' \quad c \notin \text{dom } \Pi' \quad \Pi = \Pi' \sqcup \{c \mapsto q\} \quad \text{Size } 1 \in q}{\Pi \vdash \oplus a : c}$$

- $\oplus \in \{>>, <<, **, >>>, <<<\}$:

$$\frac{\Pi_1 \vdash a : c \quad \Pi_2 \vdash b : c' \quad \Pi = \Pi_1 \sqcup \Pi_2 \quad \text{dom } \Pi_1 \cap \text{dom } \Pi_2 = \emptyset}{\Pi \vdash a \oplus b : c}$$

- $\oplus \in \{=, +=, -=, *=, /=, \%=, \&=, |=, \wedge=\}$:

$$\frac{s = \Phi(l) \quad \Pi' \vdash b : c' \quad \text{Size } s \in \Pi'(c') \quad c \notin \text{dom } \Pi' \quad \Pi = \Pi' \sqcup \{c \mapsto q\} \quad \text{Size } s \in q}{\Pi \vdash l \oplus b : c}$$

- $\oplus \in \{<<=, >>=, <<<=, >>>=\}$:

$$\frac{s = \Phi(l) \quad \Pi' \vdash b : c' \quad c \notin \text{dom } \Pi' \quad \Pi = \Pi' \sqcup \{c \mapsto q\} \quad \text{Size } s \in q}{\Pi \vdash l \oplus b : c}$$

- If expression:

$$\frac{\Pi_1 \vdash e : c' \quad \Pi_2 \vdash a : c \quad \Pi_2 \vdash b : c \quad \Pi = \Pi_1 \sqcup \Pi_2 \quad \text{dom } \Pi_1 \cap \text{dom } \Pi_2 = \emptyset}{\Pi \vdash e?a:b : c}$$

- Concatenation:

$$\frac{\Pi_1 \vdash e_1 : c_1 \quad \Pi_2 \vdash e_2 : c_2 \quad \text{dom } \Pi_1 \cap \text{dom } \Pi_2 = \emptyset \quad c \notin \text{dom } \Pi_1 \sqcup \text{dom } \Pi_2 \quad \Pi = \Pi_1 \sqcup \Pi_2 \sqcup \{c \mapsto q\} \quad \text{Add } (c_1, c_2) \in q}{\Pi \vdash \{e_1, e_2\} : c}$$

- Replication:

$$\frac{i \in \mathbb{N} \quad \Pi' \vdash e : c' \quad c \notin \text{dom } \Pi' \quad \Pi = \Pi' \sqcup \{c \mapsto q\} \quad \text{Mul } (i, c') \in q}{\Pi \vdash \{i\} e : c}$$

- Inside:

$$\frac{\Pi' \vdash a : c' \quad \Pi' \vdash e_1 : c' \quad \dots \quad \Pi' \vdash e_k : c' \quad c \notin \text{dom } \Pi' \quad \Pi = \Pi' \sqcup \{c \mapsto q\} \quad \text{Size } 1 \in q}{\Pi \vdash a \text{ inside } \{e_1, \dots, e_k\} : c}$$