

Arrays and Matrices

Chapter - 7

2D Arrays

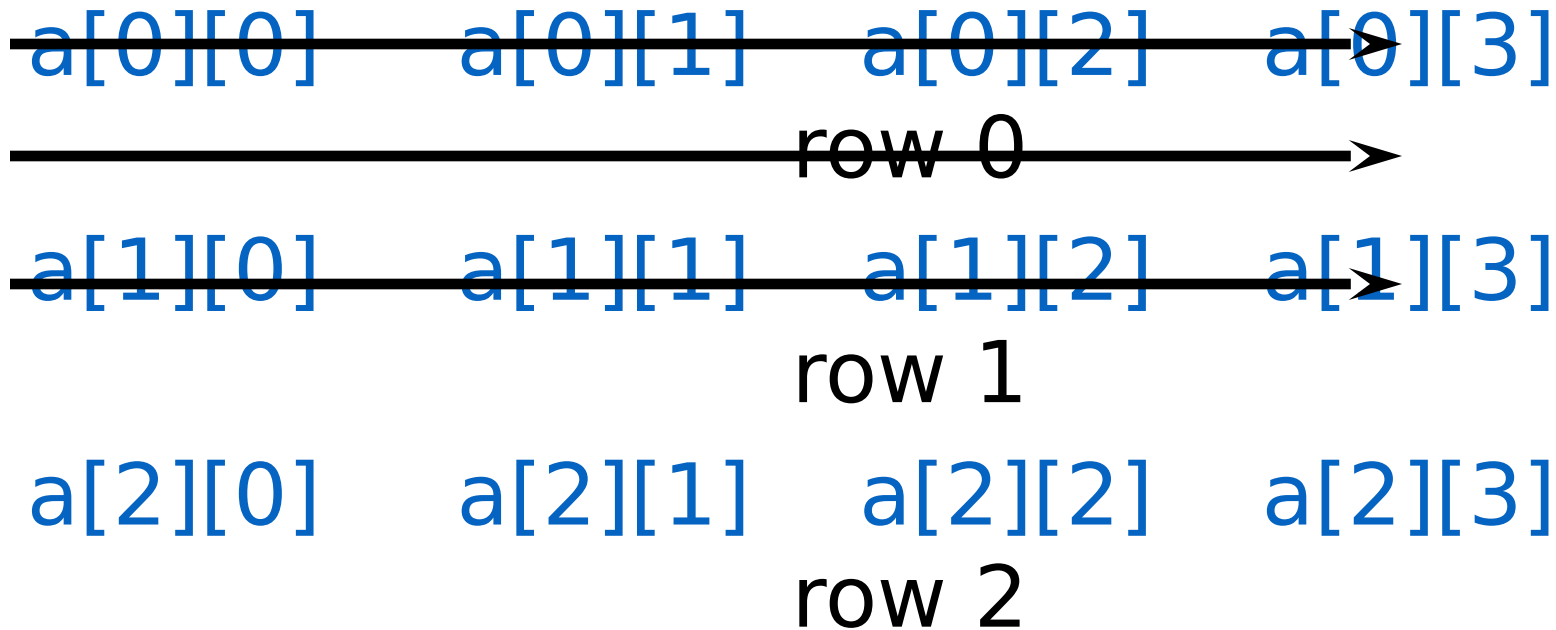
The elements of a 2-dimensional array **a** declared as:

```
int [][]a = new int[3][4];
```

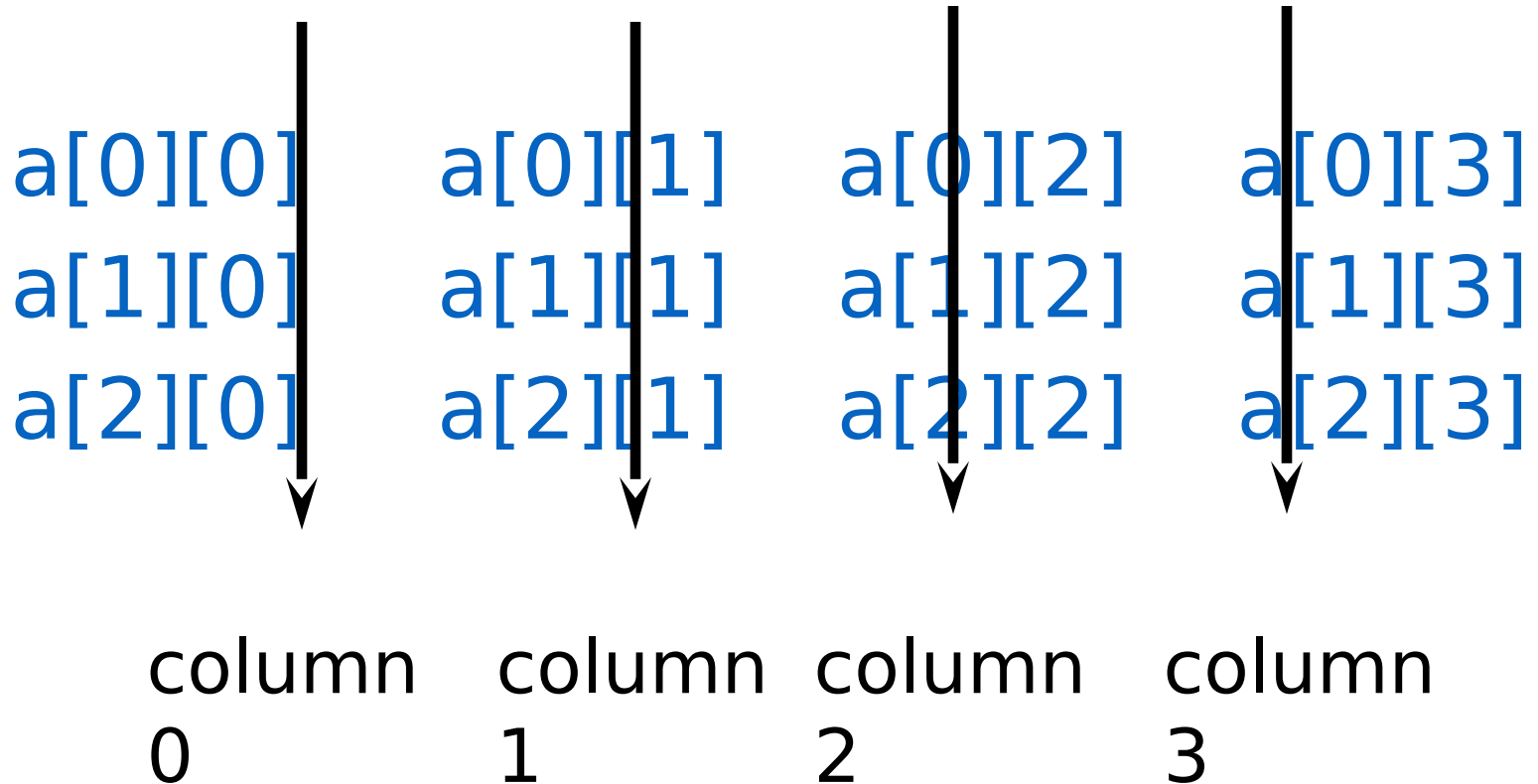
may be shown as a table

a[0][0]	a[0][1]	a[0][2]	a[0][3]
a[1][0]	a[1][1]	a[1][2]	a[1][3]
a[2][0]	a[2][1]	a[2][2]	a[2][3]

Rows Of A 2D Array



Columns Of A 2D Array



2D Array Representation In C++

2-dimensional array `x`

`a, b, c, d`

`e, f, g, h`

`i, j, k, l`

view 2D array as a 1D array of rows

`x = [row0, row1, row 2]`

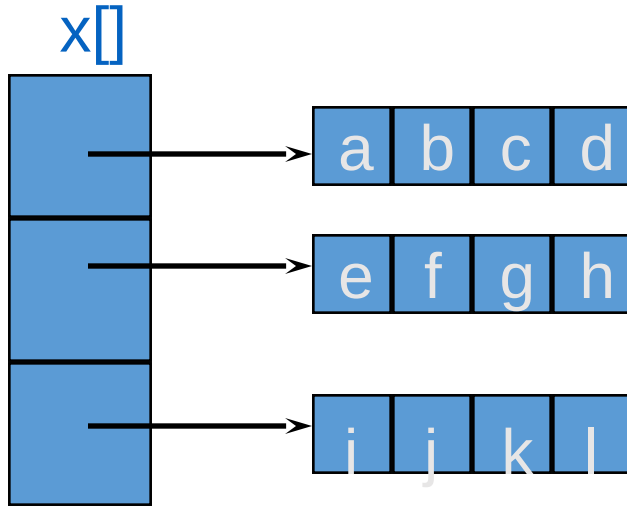
`row 0 = [a,b, c, d]`

`row 1 = [e, f, g, h]`

`row 2 = [i, j, k, l]`

and store as 4 1D arrays

2D Array Representation In C++



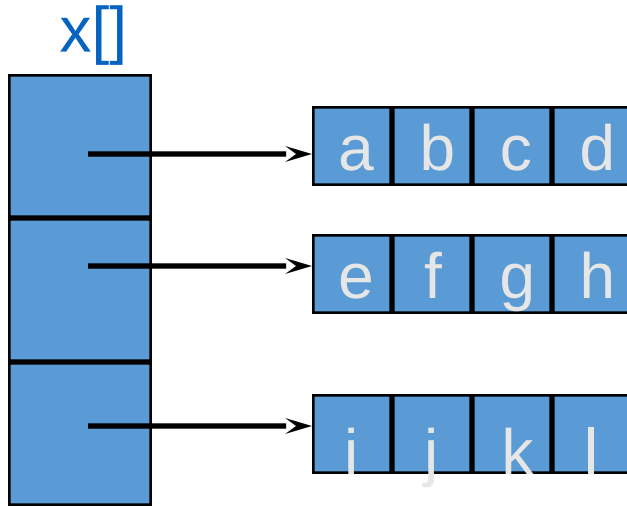
space overhead = overhead for 4 1D arrays

= $4 * 4$ bytes

= 16 bytes

= (number of rows + 1) x 4 bytes

Array Representation In C++



- This representation is called the array-of-arrays representation.
- Requires contiguous memory of size 3, 4, 4, and 4 for the 4 1D arrays.
- 1 memory block of size number of rows and number of rows blocks of size number of columns

Row-Major Mapping

- In general, for 3 dimension,
 $\text{map}(i_1, i_2, i_3) = i_1 u_2 u_3 + i_2 u_3 + i_3$
- Example 3 x 4 array:

a b c d
e f g h
i j k l

- Convert into 1D array y by collecting elements by rows.
- Within a row elements are collected from left to right.
- Rows are collected from top to bottom.
- We get y {a, b, c, d, e, f, g, h, i, j, k, l}



Locating Element $x[i][j]$



- assume x has r rows and c columns
- each row has c elements
- i rows to the left of row i
- so ic elements to the left of $x[i][0]$
- so $x[i][j]$ is mapped to position $ic + j$ of the 1D array

Space Overhead



4 bytes for start of 1D array +
4 bytes for c (number of columns)
= 8 bytes

Disadvantage: Need contiguous memory of size rc

Column-Major Mapping

a b c d
e f g h
i j k l

- Convert into 1D array y by collecting elements by columns.
- Within a column elements are collected from top to bottom.
- Columns are collected from left to right.
- We get $\{a, e, i, b, f, j, c, g, k, d, h, l\}$

Exercise

- Suppose, we want to map the elements of a two dimensional array beginning with bottom row and within a row from left to right order.
 - a) List the indexed of score[3][5] in this order
 - b) Develop a mapping function for the score[u_1 u_2]

(a)

The ordering of the indexes is:

[2][0] [2][1] [2][2] [2][3] [2][4] [1][0] [1][1] [1][2] [1][3] [1][4]

[0][0] [0][1] [0][2] [0][3] [0][4]

(b)

The elements with first index i_1 are preceded by $(u_1 - i_1 - 1)u_2$ elements that have a larger first index. Since the elements with the same first index are stored from right to left,

$$\text{map}(i_1, i_2) = (u_1 - i_1 - 1)u_2 + i_2$$

Matrix

A $m \times n$ matrix is a table with m rows and n columns. m and n are dimensions of the matrix.

a	b	c	d	row 1
e	f	g	h	row 2
i	j	k	l	row 3

- Use notation $x(i,j)$ rather than $x[i][j]$.
- May use a 2D array to represent a matrix.
- Operations performed: Addition, Multiplication, Transpose

Header class Matrix with operator overloading

```
template<class T>
class matrix
{
    friend ostream& operator<<(ostream&, const matrix<T>&);
public:
    matrix(int theRows = 0, int theColumns = 0);
    matrix(const matrix<T>&);
    ~matrix() {delete [] element;}
    int rows() const {return theRows;}
    int columns() const {return theColumns;}
    T& operator()(int i, int j) const;
    matrix<T>& operator=(const matrix<T>&);
    matrix<T> operator+() const; // unary +
    matrix<T> operator+(const matrix<T>&) const;
    matrix<T> operator-() const; // unary minus
    matrix<T> operator-(const matrix<T>&) const;
    matrix<T> operator*(const matrix<T>&) const;
    matrix<T>& operator+=(const T&);
private:
    int theRows,      // number of rows in matrix
        theColumns;  // number of columns in matrix
    T *element;       // element array
};
```

Construct
or and
copy
construct
or

```
template<class T>
matrix<T>::matrix(int theRows, int theColumns)
{
    // matrix constructor.
    // validate theRows and theColumns
    if (theRows < 0 || theColumns < 0)
        throw illegalParameterValue("Rows and columns must be >= 0");
    if ((theRows == 0 || theColumns == 0)
        && (theRows != 0 || theColumns != 0))
        throw illegalParameterValue
            ("Either both or neither rows and columns should be zero");

    // create the matrix
    this->theRows = theRows;
    this->theColumns = theColumns;
    element = new T [theRows * theColumns];
}

template<class T>
matrix<T>::matrix(const matrix<T>& m)
{
    // Copy constructor for matrices.
    // create matrix
    theRows = m.theRows;
    theColumns = m.theColumns;
    element = new T [theRows * theColumns];

    // copy each element of m
    copy(m.element,
        m.element + theRows * theColumns,
        element);
}
```

Overload = operator

```
template<class T>
matrix<T>& matrix<T>::operator=(const matrix<T>& m)
{
    // Assignment. (*this) = m.
    if (this != &m)
    {
        // not copying to self
        delete [] element;
        theRows = m.theRows;
        theColumns = m.theColumns;
        element = new T [theRows * theColumns];
        // copy each element
        copy(m.element,
            m.element + theRows * theColumns,
            element);
    }
    return *this;
}
```

Overload () operator

```
template<class T>
T& matrix<T>::operator()(int i, int j) const
{
    // Return a reference to element (i,j).
    if (i < 1 || i > theRows
        || j < 1 || j > theColumns)
        throw matrixIndexOutOfBounds();
    return element[(i - 1) * theColumns + j - 1];
}
```

Matrix Addition

```
template<class T>
matrix<T> matrix<T>::operator+(const matrix<T>& m) const
{
    // Return w = (*this) + m.
    if (theRows != m.theRows
        || theColumns != m.theColumns)
        throw matrixSizeMismatch();

    // create result matrix w
    matrix<T> w(theRows, theColumns);
    for (int i = 0; i < theRows * theColumns; i++)
        w.element[i] = element[i] + m.element[i];

    return w;
}
```

Matrix multiplication

```
template<class T>
matrix<T> matrix<T>::operator*(const matrix<T>& m) const
{
    // matrix multiply. Return w = (*this) * m.
    if (theColumns != m.theRows)
        throw matrixSizeMismatch();

    matrix<T> w(theRows, m.theColumns); // result matrix

    // define cursors for *this, m, and w
    // and initialize to location of (1,1) element
    int ct = 0, cm = 0, cw = 0;

    // compute w(i,j) for all i and j
    for (int i = 1; i <= theRows; i++)
    {
        // compute row i of result
        for (int j = 1; j <= m.theColumns; j++)
        {
            // compute first term of w(i,j)
            T sum = element[ct] * m.element[cm];

            // add in remaining terms
            for (int k = 2; k <= theColumns; k++)
            {
                ct++; // next term in row i of *this
                cm += m.theColumns; // next in column j of m
                sum += element[ct] * m.element[cm];
            }
            w.element[cw++] = sum; // save w(i,j)

            // reset to start of row and next column
            ct -= theColumns - 1;
            cm = j;
        }

        // reset to start of next row and first column
        ct += theColumns;
        cm = 0;
    }

    return w;
}
```

Special Matrices

- Tridiagonal:

Matrix M is tridiagonal iff $M(i,j) = 0$ for $|i-j| > 1$

- Upper Traingular:

Matrix M is upper triangular iff $M(i,j) = 0$ for $i > j$

- Symmetric:

Matrix M is symmetric iff $M(i,j) = M(j,i)$ for all i and j

Diagonal Matrix

```
template<class T>
class diagonalMatrix
{
    public:
        diagonalMatrix(int theN = 10);
        ~diagonalMatrix() {delete [] element;}
        T get(int, int) const;
        void set(int, int, const T&);
    private:
        int n;           // matrix dimension
        T *element;      // 1D array for diagonal elements
};

template<class T>
diagonalMatrix<T>::diagonalMatrix(int theN)
{// Constructor.
    // validate theN
    if (theN < 1)
        throw illegalParameterValue("Matrix size must be > 0");

    n = theN;
    element = new T [n];
}
```

Get Method for Diagonal Matrix

```
template <class T>
T diagonalMatrix<T>::get(int i, int j) const
{
    // Return (i,j)th element of matrix.
    // validate i and j
    if (i < 1 || j < 1 || i > n || j > n)
        throw matrixIndexOutOfBounds();

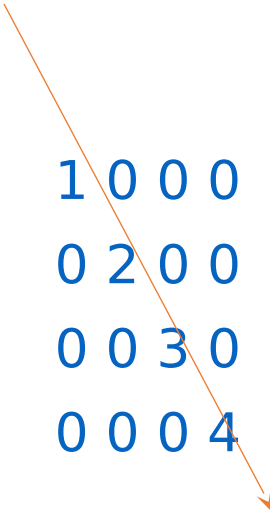
    if (i == j)
        return element[i-1];    // diagonal element
    else
        return 0;                // nondiagonal element
}
```

Set Method for Diagonal Matrix

```
template<class T>
void diagonalMatrix<T>::set(int i, int j, const T& newValue)
{
    // Store newValue as (i,j)th element.
    // validate i and j
    if (i < 1 || j < 1 || i > n || j > n)
        throw matrixIndexOutOfBounds();

    if (i == j)
        // save the diagonal value
        element[i-1] = newValue;
    else
        // nondiagonal value, newValue must be zero
        if (newValue != 0)
            throw illegalParameterValue
                ("nondiagonal elements must be zero");
}
```

Diagonal Matrix



1 0 0 0
0 2 0 0
0 0 3 0
0 0 0 4

- An $n \times n$ matrix in which all nonzero terms are on the diagonal.
- $x(i,j)$ is on diagonal iff $i = j$
- number of diagonal elements in an $n \times n$ matrix is n
- store diagonal only vs n^2 whole

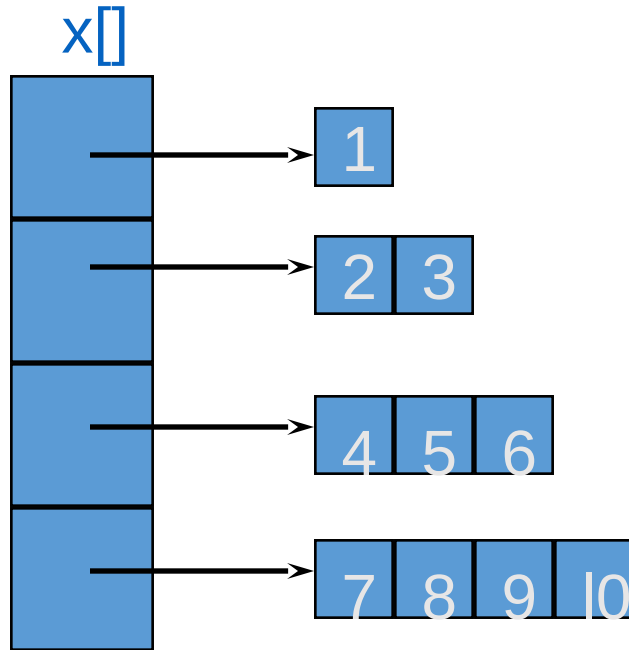
Lower Triangular Matrix

An $n \times n$ matrix in which all nonzero terms are either on or below the diagonal.

1	0	0	0
2	3	0	0
4	5	6	0
7	8	9	10

- $x(i,j)$ is part of lower triangle iff $i \geq j$.
- number of elements in lower triangle is $1 + 2 + \dots + n = n(n+1)/2$.
- store only the lower triangle

Array Of Arrays Representation



Use an irregular 2-D array ... length of rows is not required to be the same.

Creating And Using An Irregular Array

```
// declare a two-dimensional array variable
```

```
// and allocate the desired number of rows
```

```
int ** irregularArray =      int* [numberOfRows];
```

```
// now allocate space for the elements in each  
row
```

```
for (int i = 0; i < numberOfRows; i++)  
    irregularArray[i] = new int [length[i]];
```

```
// use the array like any regular array
```

```
irregularArray[2][3] = 5;
```

```
irregularArray[4][6] = irregularArray[2][3] + 2;
```

Map Lower Triangular Array Into A 1D Array

Use row-major order, but omit terms that are not part of the lower triangle.

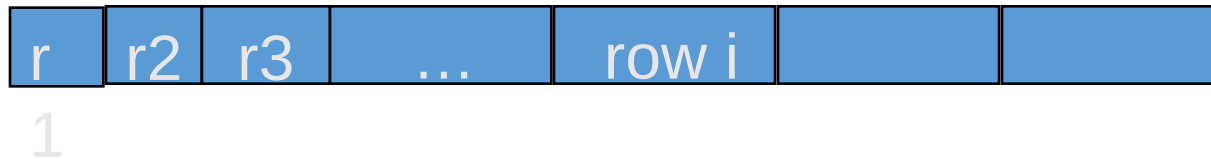
For the matrix

1	0	0	0
2	3	0	0
4	5	6	0
7	8	9	10

we get

1, 2, 3, 4, 5, 6, 7, 8, 9, 10

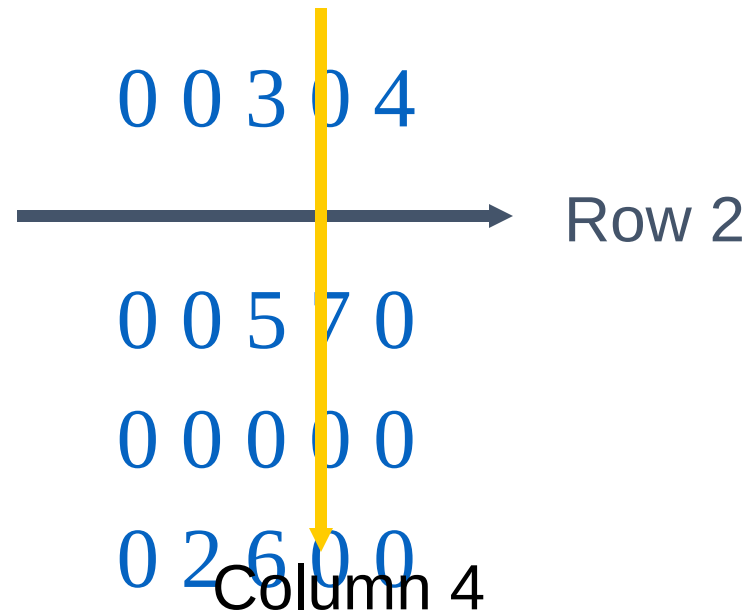
Index Of Element [i][j]



- Order is: row 1, row 2, row 3, ...
- Row i is preceded by rows 1, 2, ..., $i-1$
- Size of row i is i .
- Number of elements that precede row i is
$$1 + 2 + 3 + \dots + i-1 = i(i-1)/2$$
- So element (i, j) is at position $i(i-1)/2 + j$

Sparse Matrices

Matrix  table of values



0	0	3	0	4
0	0	5	7	0
0	0	0	0	0
0	2	6	0	0

Row 2

Column 4

4 x 5 matrix


4 rows

5 columns


20 elements

6 nonzero
elements

Sparse Matrices

Sparse matrix  #nonzero elements/#elements is small.

Examples:

- Diagonal
 - Only elements along diagonal may be nonzero
 - $n \times n$ matrix  ratio is $n/n^2 = 1/n$
- Tridiagonal
 - Only elements on 3 central diagonals may be nonzero
 - Ratio is $(3n-2)/n^2 = 3/n - 2/n^2$

Sparse Matrices

- Lower triangular
- Only elements on or below diagonal may be nonzero
 - Ratio is $n(n+1)/(2n^2) \sim 0.5$

These are structured sparse matrices.
Nonzero elements are in a well-defined portion of the matrix.

Sparse Matrices

An $n \times n$ matrix may be stored as an $n \times n$ array.

This takes $O(n^2)$ space.

The example structured sparse matrices may be mapped into a 1D array so that a mapping function can be used to locate an element quickly; the space required by the 1D array is less than that required by an $n \times n$ array (next lecture).

Representation Of Unstructured Sparse Matrices

Single linear list in row-major order.

scan the nonzero elements of the sparse matrix in row-major order (i.e., scan the rows left to right beginning with row 1 and picking up the nonzero elements)

each nonzero element is represented by a triple

(row, column, value)

the list of triples is stored in a 1D array

Single Linear List Example

0 0 3 0 4

0 0 5 7 0

0 0 0 0 0

0 2 6 0 0

list =

row

column

value

1	1	2	2	4	4
3	5	3	4	2	3
3	4	5	7	2	6

Single Linear List

- Class SparseMatrix

- Array of triples of type MatrixTerm

row, col, value

rows, // number of rows

cols, // number of columns

terms, // number of nonzero
elements

capacity; // size of

- Size of generally not predictable
at time of initialization.

- Start with some default capacity/size (say 10)
 - Increase capacity as needed

Approximate Memory Requirements

500 x 500 matrix with 1994 nonzero elements, 4 bytes per element

2D array $500 \times 500 \times 4 = 1\text{million bytes}$

Class SparseMatrix $3 \times 1994 \times 4 + 4 \times 4$
 $= 23,944 \text{ bytes}$

Array Resizing

```
(newSize < terms)          "Error";  
MatrixTerm *temp =        MatrixTerm[newSize];  
copy(smArray, smArray+terms, temp);  
    [] smArray;  
smArray = temp;  
capacity = newSize;
```

Array Resizing

- To avoid spending too much overall time resizing arrays, we generally set $\text{newSize} = c * \text{oldSize}$, where $c > 0$ is some constant.
- Quite often, we use $c = 2$ (array doubling) or $c = 1.5$.
- Now, we can show that the total time spent in resizing is $O(s)$, where s is the maximum number of elements added to `smArray`.

Matrix Transpose

0 0 3 0 4

0 0 5 7 0

0 0 0 0 0

0 2 6 0 0

0 0 0 0

0 0 0 2

3 5 0 6

0 7 0 0

4 0 0 0

Matrix Transpose

0 0 0 0
0 0 3 0 4
0 0 5 7 0
0 0 0 0 0
0 2 6 0 0

2 3 3 3 4 5
4 1 2 4 2 1
2 3 5 6 7 4

row	1	1	2	2	4	4
column	3	5	3	4	2	3
value	3	4	5	7	2	6

Matrix Transpose

Step 1: #nonzero in each row of transpose.

0 0 3 0 4

0 0 0 0

= #nonzero in each column

of

0 0 0 2

original matrix

0 0 5 7 0

3 5 0 6

= [0, 1, 3, 1, 1]

0 0 0 0 0

0 7 0 0

Step2: Start of each row of transpose

0 2 6 0 0

4 0 0 0

= sum of size of preceding rows of

row 1 1 2 2 4 4

transpose

column 3 5 3 4 2 3

= [0, 0, 1, 4, 5]

value 3 4 5 7 2 6

Step 3: Move elements, left to right, from

original list to transpose list.

Matrix Transpose

Step 1: #nonzero in each row of transpose.

of = #nonzero in each column
original matrix
= [0, 1, 3, 1, 1]

Step2: Start of each row of transpose

= sum of size of preceding
rows of
transpose
= [0, 0, 1, 4, 5]

Step 3: Move elements, left to right,
from

Complexity

m x n original matrix

t nonzero elements

Step 1: $O(n+t)$

Step 2: $O(n)$

Step 3: $O(t)$

Overall $O(n+t)$