

Assignment-3

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1. Minimum spanning tree: Prim's and Kruskal's Algorithm.

(a) Prim's Algorithm

Program:

```
#include <iostream>
#define I 32767      //setting I to Maximum value
using namespace std;

int main()
{
    int cost[8][8] = {{I, I, I, I, I, I, I, I},
                      {I, I, 25, I, I, I, 5, I},
                      {I, 25, I, 12, I, I, I, 10},
                      {I, I, 12, I, 8, I, I, I},
                      {I, I, I, 8, I, 16, I, 14},
                      {I, I, I, I, 16, I, 20, 18},
                      {I, 5, I, I, I, 20, I, I},
                      {I, I, 10, I, 14, 18, I, I}};
    int near[8] = {I, I, I, I, I, I, I, I};
    int t[3][6];
    int i,j,k,u,v,w,n=7,min=I;    //n=7 because we're skipping first row and first column

    /*Finding first min wt edge*/

    for(i=1;i<=n;i++)    //scanning upper triangular matrix for minimum weight edge
    {
        for(j=i;j<=n;j++)
        {
            if(cost[i][j]<min)
            {
                min=cost[i][j];
                u=i;v=j;w=cost[i][j];    // storing co-ordinates of min wt edge in u,v
            }
        }
    }
    near[u]=near[v]=0;    // marking corr near indices to 0 i.e marking edge as visited
    t[0][0]=u;t[1][0]=v;t[2][0]=w;    // storing u,v in t matrix
    for(i=1;i<=n;i++)
    {
        if(near[i]!=0)    // check if node is not visited
        {
            if(cost[i][u]<cost[i][v])    // compare wt of near edges and update them in near array
                near[i]=u;
            else
                near[i]=v;
        }
    }
    /* doing above procedure for all remaining edges */
    for(i=1;i<n-1;i++)
```

```

{
    min=I;
    for(j=1;j<=n;j++)
    {
        if(near[j]!=0 && cost[j][near[j]]<min)
        {
            k=j;
            min=cost[j][near[j]];
        }
    }
    t[0][i]=k;
    t[1][i]=near[k];
    t[2][i]=min;
    near[k]=0;
    for(j=1;j<=n;j++)
    {
        if(near[j]!=0 && cost[j][k]<cost[j][near[j]] )
        {
            near[j]=k;
        }
    }
}
cout << "\n\n\tPrim's Algorithm\n";
cout << "\nOutput is printed in following form: (start vertex, end vertex, weight)" << endl;
cout << "\nMinimum Spanning Tree using Prim's Algorithm:\n";
for (i = 0; i < n - 1; i++)
{
    cout << "(" << t[0][i] << ", " << t[1][i] << ", " << t[2][i] << ")";
    if(i!=n-2)
    {
        cout << " --> ";
    }
}
int t_wt=0;
for(int i=0;i<n;i++)
{
    t_wt+=t[2][i];
}
cout << "\n\nTotal Minimum Weight: " << t_wt << endl;
cout << endl;
}

```

Output:

Prim's Algorithm

Output is printed in following form: (start vertex, end vertex, weight)

Minimum Spanning Tree using Prim's Algorithm:

(1,6,5) --> (5,6,20) --> (4,5,16) --> (3,4,8) --> (2,3,12) --> (7,2,10)

Total Minimum Weight: 71

Complexity:

1. Time complexity of Prim's Algorithm is $O(n^2)$ i.e. $(O(V * E))$
2. If we use Heap to find minimum cost edge then we can reduce time complexity to $O(n \log n)$.

(b) Krushkal's Algorithm**Program:**

```
#include<iostream>
using namespace std;

#define I 32767

int edge[9][3]={ {1,2,15},{1,6,10},{2,5,25},{2,7,14},{3,4,22},
{4,5,52},{4,7,21},{3,6,25},{5,7,34}};

int s[8]={-1,-1,-1,-1,-1,-1,-1,-1};

int t[3][6];

int included[9]={0};

void unionfunc(int u,int v)
{
    if(s[u]<s[v])
    {
        s[u]=s[u]+s[v];
        s[v]=u;
    }
    else
    {
        s[v]=s[u]+s[v];
        s[u]=v;
    }
}

int find(int u)
{
    int x=u,v=0;
```

```

while(s[x]>0)
{
    x=s[x];
}
/*connecting node to head node*/
while(u!=x)
{
    v=s[u];
    s[u]=x;
    u=v;
}
return x;
}
int main()
{
    int u=0,v=0,i,j,k=0,min=I,n=7,e=9;    //e-->number of edges
    i=0;
    while(i<n-1)
    {
        min=I;
        for(j=0;j<e;j++)
        {
            if(included[j]==0 && edge[j][2]<min)
            {
                u=edge[j][0];
                v=edge[j][1];
                min=edge[j][2];
                k=j;
            }
        }
        if(find(u)!=find(v))
        {
            t[0][i]=u;
            t[1][i]=v;
            t[2][i]=min;
            unionfunc(find(u),find(v));
            i++;
        }
        included[k]=1;
    }
    cout << "\n\tKrushkal's Algorithm\n";
    cout << "\nOutput is printed in following form: (start vertex, end vertex, weight)" << endl;
    cout << "\nMinimum Spanning Tree using Krushkal's Algorithm:\n";
    for (i = 0; i < n-1; i++)
    {
        cout << "(" << t[0][i] << ", " << t[1][i] << ", " << t[2][i] << ")";
        if(i!=n-2)
        {
            cout << " --> ";
        }
    }
}

```

```

    }
}
int t_wt=0;
for(int i=0;i<n;i++)
{
    t_wt+=t[2][i];
}
cout << "\n\nTotal Minimum Weight: " << t_wt << endl;
cout << endl;
}

```

Output:

Krushkal's Algorithm

Output is printed in following form: (start vertex, end vertex, weight)

Minimum Spanning Tree using Krushkal's Algorithm:

(1,6,10) --> (2,7,14) --> (1,2,15) --> (4,7,21) --> (3,4,22) --> (2,5,25)

Total Minimum Weight: 107

Complexity:

1. Time complexity of Krushkal's Algorithm is $O(n^2)$ i.e. $(O(V * E))$
2. If we use Heap to find minimum cost edge then we can reduce time complexity to $O(n \log n)$.

2. Shortest path (between source vertex to all other vertices): Dijkstra's Algorithm and Bellman Ford Algorithm

(a) Dijkstra's Algorithm

Program:

```
#include<iostream>
#include<limits.h>
#include<stdlib.h>
#include<stdio.h>
using namespace std;

#define N 9    //N--> number of vertices

// Function to find vertex with minimum distance from set of vertices not yet included
int minDist(int d[], int included[])
{
    int min=INT_MAX,minIndex;
    for(int v=0;v<N;v++)
    {
        if(included[v]==0 && d[v]<=min)
        {
            min=d[v];
            minIndex=v;
        }
    }
    return minIndex;
}

void printSolution(int d[],int source)
{
    cout << "Source: " << source << endl;
    cout << "Vertex \t\tDistance from Source"<<endl;
    for (int i = 0; i < N; i++)
        if(i!=source)
            cout << i << "\t\t" << d[i] << endl;
}

void dijkstrasAlgorithm(int g[N][N],int source)
{
    int d[N],included[N]={0};

    //Initialize all distances as Infinite and included[v]=0
    for(int v=0;v<N;v++)
    {
        d[v]=INT_MAX;
        included[v]=0;
    }
    d[source]=0;
    // Finding shortest path for all vertices
    for(int i=0;i<N-1;i++)
```

```

{
    int min_dist=minDist(d,included);
    included[min_dist]=1;

    for(int v=0;v<N;v++)
    {
        if(included[v]==0 && g[min_dist][v] && d[min_dist]!=INT_MAX && d[min_dist]+
g[min_dist][v] < d[v])
            d[v] = d[min_dist] + g[min_dist][v];
    }
}
printSolution(d,source);
}
int main()
{
    int graph[N][N] = { { 0, 24, 0, 0, 0, 0, 0, 8, 0 },
                        { 24, 0, 18, 0, 0, 0, 0, 11, 0 },
                        { 0, 18, 0, 17, 0, 40, 0, 0, 20 },
                        { 0, 0, 17, 0, 9, 2, 0, 0, 0 },
                        { 0, 0, 0, 9, 0, 42, 0, 0, 0 },
                        { 0, 0, 40, 2, 42, 0, 12, 0, 0 },
                        { 0, 0, 0, 0, 0, 12, 0, 15, 36 },
                        { 8, 11, 0, 0, 0, 0, 15, 0, 27 },
                        { 0, 0, 20, 0, 0, 0, 36, 27, 0 } };
    cout << "\n\tDijkstra's Algorithm Implementation\t\n\n";
    dijkstrasAlgorithm(graph, 0);
}

```

Output:

Dijkstra's Algorithm Implementation

Source: 0

Vertex	Distance from Source
1	19
2	54
3	37
4	46
5	35
6	23
7	8
8	35

Complexity:

The time complexity of Dijkstra's Algorithm is $O(n^2)$.

(b) Bellman-Ford's Algorithm

Program:

```
#include <iostream>
#include <limits.h>
using namespace std;

void printSolution(int d[], int V, int source)
{
    cout << "Source: " << source << endl;
    cout << "Vertex \t\tDistance from Source" << endl;
    for (int i = 0; i < V; i++)
    {
        if (i != source)
        {
            if (d[i] < INT_MAX)
                cout << i << "\t\t" << d[i] << endl;
            else if (d[i] > 32767) //INT_MAX==32767 largest 16 bit integer value
                cout << i << "\t\t" << "INFINITE" << endl;
        }
    }
}

void BellmanFord(int edge[][3], int V, int E, int src)
{
    int dist[V];

    // Step 1: Initialize distances from src to all other vertices as INFINITE
    for (int i = 0; i < V; i++)
        dist[i] = INT_MAX;
    dist[src] = 0;

    /*Step 2: Relax all edges |V| - 1 times. A simple shortest path from src to any other vertex can have
    at-most |V| - 1 edges */
    for (int i = 1; i <= V - 1; i++)
    {
        for (int j = 0; j < E; j++)
        {
            int u = edge[j][0];
            int v = edge[j][1];
            int weight = edge[j][2];
            if (dist[u] != INT_MAX && dist[u] + weight < dist[v])
                dist[v] = dist[u] + weight;
        }
    }

    /*Step 3: check for negative-weight cycles. The above step guarantees shortest distances if graph
    doesn't contain
    negative weight cycle. If we get a shorter path, then there is a cycle.*/
    for (int j = 0; j < E; j++)
    {
```

```

    int u = edge[j][0];
    int v = edge[j][1];
    int weight = edge[j][2];
    if (dist[u] != INT_MAX && dist[u] + weight < dist[v])
    {
        printf("Graph contains negative weight cycle");
        return; // If negative cycle is detected, simply return
    }
}

printSolution(dist, V, src);
return;
}

int main()
{
    int V = 5; // Number of vertices in graph
    int E = 8; // Number of edges in graph
    int edge[][3] = {{0, 1, -1}, {0, 2, 4}, {1, 2, 3}, {1, 3, 2}, {1, 4, 2}, {3, 2, 5}, {3, 1, 1}, {4, 3, -3}};
    cout << "\n\tImplementation of Bellman-Ford's Algorithm\t\n\n";
    BellmanFord(edge, V, E, 3);
    return 0;
}

```

Output:

Implementation of Bellman-Ford's Algorithm

Source: 3

Vertex	Distance from Source
0	INFINITE
1	1
2	4
4	3

Complexity:

For a complete graph, the time complexity of Bellman-Ford Algorithm is $\Theta(n^3)$.

3. Shortest path between any pair of vertices (Floyd-Warshall Algorithm).

Program:

```
#include <bits/stdc++.h>
using namespace std;
#define I INT_MAX

void flyodWarshall(vector<vector<int>> g)
{
    int V = g.size();

    vector<vector<int>> dist = g;

    for (int k = 0; k < V; k++)
        for (int i = 0; i < V; i++)
            for (int j = 0; j < V; j++)
                if (dist[i][k] != I && dist[k][j] != I && dist[i][k] + dist[k][j] < dist[i][j])
                    dist[i][j] = dist[i][k] + dist[k][j];

    cout << endl;
    cout << "The Shortest path between any vertices is:\n";
    for (int i = 0; i < V; i++)
    {
        for (int j = 0; j < V; j++)
        {
            if (dist[i][j] == INT_MAX)
            {
                cout << "INFINITE\t";
                continue;
            }
            cout << dist[i][j] << "\t";
        }
        cout << endl;
    }
}

int main()
{
    int V;
    cout << "\n\tShortest Path between Any Pair of Vertices\n\n";
    cout << "Enter the number of vertex in the g: ";
    cin >> V;
    vector<vector<int>> g(V, vector<int>(V));

    cout << "\nINFINITE: " << I << endl;
    cout << "\nEnter the weights of the each Edges:\n";
    for (int i = 0; i < V; i++)
    {
        for (int j = 0; j < V; j++)
```

```

        {
            cout << ">";
            cin >> g[i][j];
        }
    }

    flyodWarshall(g);
}

```

Output:

Shortest Path between Any Pair of Vertices

Enter the number of vertex in the g: 4

INFINITE: 2147483647

Enter the weights of the each Edges:

```

>0
>3
>2147483647
>7
>8
>0
>2
>2147483647
>5
>2147483647
>0
>1
>2
>2147483647
>2147483647
>0

```

The Shortest path between any vertices is:

0	3	5	6
5	0	2	3
3	6	0	1
2	5	7	0

4. Solve Knapsack problem (for divisible and indivisible objects)

(a) Divisible Objects

Program:

```
#include <bits/stdc++.h>

using namespace std;
struct Item
{
    public:
    int value, weight;
    Item(int value, int weight): value(value),weight(weight){}
};

// comparator used to sort Item according to val/weight ratio
bool cmp(struct Item a, struct Item b)
{
    double r1 = (double)a.value / (double)a.weight;
    double r2 = (double)b.value / (double)b.weight;
    return r1 > r2;
}

void fractionalKnapsack(int W, Item arr[], int n)
{
    cout << "\nWeight Limit: " << W << endl << endl;
    // sorting Item on basis of ratio
    sort(arr, arr + n, cmp);

    // Printing items added to knapsack along with their weights
    for (int i = 0; i < n; i++)
    {
        cout << "Item " << i+1 << "\tValue: " << arr[i].value << "\tWeight: " << arr[i].weight <<
"\tQuantity: "\
<< ((double)arr[i].value / arr[i].weight) << endl;
    }

    int curWeight = 0;          // Current weight in knapsack
    double finalvalue = 0.0;    // Result (value in Knapsack)

    // Looping through all Items
    for (int i = 0; i < n; i++)
    {
        // If adding Item won't overflow, add it completely
        if (curWeight + arr[i].weight <= W)
        {
            curWeight += arr[i].weight;
            finalvalue += arr[i].value;
        }
    }
}
```

```

        // If we can't add current Item, add fractional part of it
        else
        {
            int remain = W - curWeight;
            finalvalue
                += arr[i].value
                * ((double)remain / (double)arr[i].weight);
            break;
        }
    }
    cout << "\nMaximum Value that can be obtained for weight " << W << " is " << finalvalue << endl;
}

int main()
{
    int W = 100; // Weight of knapsack
    Item arr[] = { { 60, 10 }, { 100, 20 }, { 120, 30 }, { 10, 25}, { 55, 22} };

    int n = sizeof(arr) / sizeof(arr[0]);
    cout << "\n\tKnapsack Problem Implementation\n\n";
    cout << "\t\t(Divisible Objects)\n\n";
    fractionalKnapsack(W, arr, n);
    return 0;
}

```

Output:

Knapsack Problem Implementation
(Divisible Objects)

Weight Limit: 100

Item 1	Value: 60	Weight: 10	Quantity: 6
Item 2	Value: 100	Weight: 20	Quantity: 5
Item 3	Value: 120	Weight: 30	Quantity: 4
Item 4	Value: 55	Weight: 22	Quantity: 2.5
Item 5	Value: 10	Weight: 25	Quantity: 0.4

Maximum Value that can be obtained for weight 100 is 342.2

Complexity:

As Sorting takes maximum of time, so the time complexity of Knapsack algorithm for divisible objects is $O(n \log n)$.

(b) Indivisible Objects

Program:

```
#include <iostream>
```

```
using namespace std;
```

```
int knapSackRec(int W, int wt[],int val[], int i,int **dp)
```

```
{
    if (i < 0)
        return 0;
    if (dp[i][W] != -1)
        return dp[i][W];

    if (wt[i] > W)
    {

        dp[i][W] = knapSackRec(W, wt,val, i - 1,dp);
        return dp[i][W];
    }
    else
    {
        dp[i][W] = max(val[i] + knapSackRec(W - wt[i], wt, val, i - 1, dp), knapSackRec(W, wt, val, i - 1, dp));

        return dp[i][W];
    }
}
```

```
int knapSack(int W, int wt[], int val[], int n)
```

```
{
    int **dp;
    dp = new int *[n];
    for (int i = 0; i < n; i++)
        dp[i] = new int[W + 1];

    for (int i = 0; i < n; i++)
        for (int j = 0; j < W + 1; j++)
            dp[i][j] = -1;
    return knapSackRec(W, wt, val, n - 1, dp);
}
```

```
int main()
```

```
{
    int val[] = {10, 20, 30};
    int wt[] = {10, 18, 20};
    int W ;
    int n = sizeof(val) / sizeof(val[0]);
    cout << "\nKnapsack Problem Implementation" <<
        "\n  (Indivisible Objects)\n";
    cout << "\nValues: ";
```

```

for(int i=0;i<n;i++)
{
    cout << " " << val[i];
}
cout << "\nWeights:";
for(int i=0;i<n;i++)
{
    cout << " " << wt[i];
}
cout << "\nEnter Weight limit:\n>";
cin >> W;
cout << "\nWeight limit: " << W << endl;
cout << "Max Value that can be obtained for given weight limit: " << knapSack(W, wt, val, n) <<
endl;
return 0;
}

```

Output:

Knapsack Problem Implementation
(Indivisible Objects)

Values: 10 20 30
Weights: 10 18 20
Enter Weight limit:
>35

Weight limit: 35
Max Value that can be obtained for given weight limit: 40

Complexity:

Time complexity is $O(N*W)$ i.e. $O(n^2)$.

5. Solve the transportation Problem

Program:

```
#include <iostream>
using namespace std;
int main()
{
    int flag = 0, flag1 = 0;
    int s[10], d[10], sn, eop = 1, dm, a[10][10];
    int i, j, sum = 0, min, x[10][10], k, fa, fb;
    /* Getting The Input For the Problem */
    cout << "\n\t Transporation Problem \t\n" << endl;
    cout << "Enter the number of Sources:\n>";
    cin >> sn;
    cout << "\nEnter the number of Destinations\n>";
    cin >> dm;
    cout << "\nEnter the Supply Values:";
    for (i = 0; i < sn; i++)
    {
        cout << "\nSource " << (i+1) << ": ";
        cin >> s[i];
    }
    cout << "\nEnter the Demand Values: ";
    for (j = 0; j < sn; j++)
    {
        cout << "\nDestination " << (j+1) << ": ";
        cin >> d[j];
    }
    cout << "\nEnter costs row wise :\n";
    for (i = 0; i < sn; i++)
    {
        for (j = 0; j < dm; j++)
        {
            cout << ">";
            cin >> a[i][j];
        }
    }
    /* Calculation For the Transportation */
    i = 0;
    j = 0;
    for (i = 0, j = 0; i < sn, j < dm;)
    {
        if (s[i] < d[j]) // Check supply less than demand
        {
            x[i][j] = a[i][j] * s[i]; // Calculate amount * supply
            d[j] = d[j] - s[i]; // Calculate demand - supply
            i++; // Increment i for the deletion of the row or column
        }
        else if (s[i] >= d[j]) //Check the supply greater than equal to demand
```

```

    {
        x[i][j] = a[i][j] * d[j]; // Calculate amount * demand
        s[i] = s[i] - d[j];      // Calculate supply - demand
        j++;                    // Increment j for the deletion of the row or column
    }
}
cout << "\nGiven Cost Matrix is :\n";
for (fa = 0; fa < sn; fa++)
{
    for (fb = 0; fb < dm; fb++)
    {
        cout << a[fa][fb] << "t";
    }
    cout << endl;
}

cout << "\nAllocated Cost Matrix is \n";
for (fa = 0; fa < sn; fa++)
{
    for (fb = 0; fb < dm; fb++)
    {
        if(x[fa][fb]!=0)
        {
            cout << x[fa][fb] << "t";
        }
        else
        {
            cout << "t";
        }
        sum = sum + x[fa][fb];
    }
    cout << endl;
}
cout << "\nTransportation cost: " << sum << endl<<endl;
}

```

Output:

Transporation Problem

Enter the number of Sources:

>4

Enter the number of Destinations

>4

Enter the Supply Values:

Source 1: 10

Source 2: 20

Source 3: 30

Source 4: 40

Enter the Demand Values:

Destination 1: 40

Destination 2: 30

Destination 3: 20

Destination 4: 10

Enter costs row wise :

>1

>2

>3

>7

>9

>3

>4

>6

>8

>1

>5

>4

>6

>9

>8

>4

Given Cost Matrix is :

1 2 3 7

9 3 4 6

8 1 5 4

6 9 8 4

Allocated Cost Matrix is

10

180

80 20

90 160 40

The Transportation cost: 580

===== THE END =====