Dynamic Programming

- 2xn domino tiling
- log outthing
- matrix chaining
- longest common subsequence
- Seam carving
- roller coaster (HWS)
- gery mundering

Domino Tiling (2km)

Tile (n) = number of way to tile a board of size Zxn

Tile(n) = Tile(n-1) + Tile(n-2)

anythine we complete tile(i) we will store in memory



- Requires Optimal Substructure
 - solution to the larger problem contains or solutions to the smaller subproblems versions of this problem

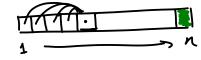
-3 Steps:

- 1. identify recurive structure "what is the last thing we did?"
- 3. Save solutions to any subproblems in memory (memorization)
- 3. Is there a good order to Solve our smaller subproblems? - top-down - bothom-up

Log Cutting the most Cut(i) = kow much money I can make by cutting a by of size i.

$$\frac{\text{cut}(i)}{\text{cut}(i-2)} = \max \begin{cases} \frac{\text{cut}(i-1) + P[1]}{\text{cut}(i-2) + P[2]} \\ \text{cut}(i-3) + P[3] \\ \text{cut}(0) + P(i) \end{cases}$$

anytime we compute cut(i) we will store in memory



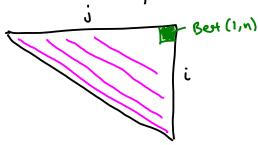
matrix Chaining

Best (i,j) = fewest (best) number of operations to multiply

matrices i ... j

= min (Best (i,k) + Best (k+1,j) + r; rkm Cj)

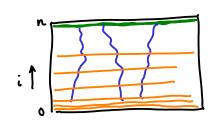
Anytime we compute Best (i,j), Store in memory.



Seam Carving

Seam (i,j) = weight of least weight seam starting at the bottom of the image and ended at pixel [i,j]

Anythine we compute Seam (i,j), we should store in memory.



best seam overall = min (Sean (n, k))

Roller coaster

longest decreasing path - length of longest decreasing path given strings X, 4

LCS(i,j) = length of the longert common subsequence among first i chars of X and j chars of y.

$$CS(i,j) = \begin{cases} 0 & \text{if } i=j=0 \\ CCS(i,j-1)+1 & \text{if } X[i]=\lambda[i] \\ 0 & \text{if } i=j=0 \end{cases}$$

Anytime we compute LCS (iji), store that in memory

