

COMPUTING MACHINERY

CPSC 355

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OPERATIONS

INSTRUCTION	MEANING
MOV X_d, X_n	$X_d = X_n$
ADD $X_d, X_n, X_m,$	$X_d = X_n + X_m$
SUB X_d, X_n, X_m	$X_d = X_n - X_m.$
MADD X_d, X_n, X_m, X_a	$X_d = (X_n \times X_m) + X_a$
MSUB X_d, X_n, X_m, X_a	$X_d = X_a - (X_n \times X_m.)$

BRANCH

Unconditional Branch (immediate)

B label

Branch: unconditionally jumps to pc-relative label.

BL label

Branch and Link: unconditionally jumps to pc-relative label, writing the address of the next sequential instruction to register X30.

```
again: mov x19, 20
mov x20, #10
add x22, x19, x20
b again
```

Conditional Branch

B.cond label

Branch: conditionally jumps to program-relative label if condition is true.

Refer to next slide for various condition codes.

```
mov x19, 1
mov x20, 20
label: add x19, x19, 1
cmp x19, x20
b.le label
```

Condition Codes

Encoding	Name (& alias)	Meaning (integer)	Meaning (floating point)	Flags
0000	EQ	Equal	Equal	Z==1
0001	NE	Not equal	Not equal, or unordered	Z==0
0010	HS (CS)	Unsigned higher or same (Carry set)	Greater than, equal, or unordered	C==1
0011	LO (CC)	Unsigned lower (Carry clear)	Less than	C==0
0100	MI	Minus (negative)	Less than	N==1
0101	PL	Plus (positive or zero)	Greater than, equal, or unordered	N==0
0110	VS	Overflow set	Unordered	V==1
0111	VC	Overflow clear	Ordered	V==0
1000	HI	Unsigned higher	Greater than, or unordered	C==1 && Z==0
1001	LS	Unsigned lower or same	Less than or equal	!(C==1 && Z==0)
1010	GE	Signed greater than or equal	Greater than or equal	N==V
1011	LT	Signed less than	Less than or unordered	N!=V
1100	GT	Signed greater than	Greater than	Z==0 && N==V
1101	LE	Signed less than or equal	Less than, equal, or unordered	!(Z==0 && N==V)

SAMPLE ASSEMBLY LANGUAGE PROGRAM

Code: loopp.s

```
.global main
.balign 4
main: stp x29, x30, [sp, -16]! //push x29 x30 on to stack - fp and lr
     mov x29, sp

     mov x19, 1
     mov x20, 20
     here: add x19, x19, 1
     cmp x19, x20 //compares values inside these registers and sets the flags
     b.le here //if condition is true, branch to 'here'
     █

     ldp x29, x30, [sp], 16 //pop x29, x30 off the stack
     ret

~
~
~
~
~
~
~
-- INSERT --
```

12,1

A11

GDB

SOME COMMANDS

Command	Meaning	Example
<code>r</code>	Run the program	
<code>q</code>	Quit gdb	
<code>c</code>	Continue until next breakpoint or the end of program	
<code>help</code>	Print list of commands	
<code>layout name</code>	Change the layout of windows	<code>layout reg</code>
<code>b label/line number</code>	Set breakpoint	<code>b main</code> <code>b 9</code>
<code>x/fmt addr</code>	Examine memory at the address <code>addr</code> (in hex); <code>fmt</code> - i,x,d,u,s...	<code>x/i 0x440066</code>
<code>p/fmt \$xn</code>	Print the contents of the register	<code>p/d \$x20, p \$x19</code>
<code>display/fmt \$xn</code>	Auto display the contents of the register	<code>display/d \$x23</code>
<code>undisplay</code>	Cancel all display requests	

GDB

```
[raheela.afzal1@csa1:~$ cd Pro*
[raheela.afzal1@csa1:~/Programs$ vi loopp.s
[raheela.afzal1@csa1:~/Programs$ gcc -o loopp loopp.s -g
raheela.afzal1@csa1:~/Programs$ gdb loopp
```

Register group: general					
x0	0x1	1		x1	0xffffffff298 281474976707224
x2	0xffffffff2a8	281474976707240		x3	0x400554 4195668
x4	0xffffffff1b0	281474976706992		x5	0xffffb7ffb2b0 281473768731312
x6	0xffffffff290	281474976707216		x7	0x400010000000400 288231475663340544
x8	0xffffffffffffff	-1		x9	0xff 255
x10	0xffffb7e01cc8	281473766661320		x11	0xffffb7e0eb20 281473766714144

```
B+> 0x400554 <main>      stp    x29, x30, [sp, #-16]!
0x400558 <main+4>      mov     x29, sp
0x40055c <main+8>      mov     x19, #0x1          // #1
0x400560 <main+12>     mov     x20, #0x14       // #20
0x400564 <here>       add     x19, x19, #0x1
0x400568 <here+4>     cmp     x19, x20
```

native process 14951 In: main L3 PC: 0x400554

(gdb) layout reg

(gdb) b main

Breakpoint 1 at 0x400554: file loopp.s, line 3.

(gdb) r

Starting program: /home/grads/raheela.afzal1/Programs/loopp

Breakpoint 1, main () at loopp.s:3

(gdb)

MACRO PREPROCESSORS

Define a piece of text with a macro name: more readable code

```
define(coef, 23)
define(x_reg, x19)
..
add x20, x_reg, coef
```

is expanded to

```
add x20, x19, 23
```

m4 is a macro processor, in the sense that it copies its input to the output, expanding macros as it goes.

Macros are either builtin or user-defined, and can take any number of arguments.

MACRO PREPROCESSORS

When using macros in your program, use the extension .asm followed by these instructions:

```
raheela.afzal1@csa2:~/Programs$ vi ifelse.asm
raheela.afzal1@csa2:~/Programs$ m4 ifelse.asm > ifelse.s
raheela.afzal1@csa2:~/Programs$ gcc -o ifelse ifelse.s -g
raheela.afzal1@csa2:~/Programs$ gdb ifelse
```

IF CONSTRUCTS

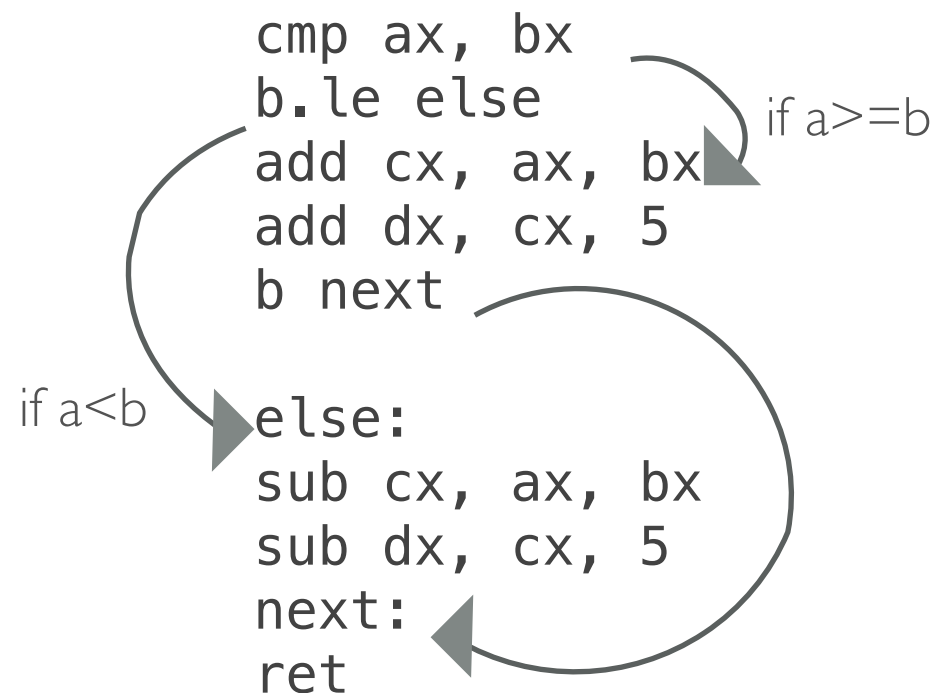
Assembly: ifelse.asm

```
define(ax, x19)
define(bx, x20)
define(cx, x21)
define(dx, x22) mov ax, 20
mov bx, 40
```

```
cmp ax, bx
b.le else      if a >= b
add cx, ax, bx
add dx, cx, 5
b next
```

if a < b

```
else:
sub cx, ax, bx
sub dx, cx, 5
next:
ret
```



C code:

```
int a = 20, b=40;
if(a>b){
    c = a + b;
    d = c + 5;
} else {
    c = a - b;
    d = c - 5;
}
```

Control Flow

LOOP

Assembly:

```
define(x, x19)
test: cmp x,10
      b.ge done
      //instructions
      add x, x, 1
      b test
done: -rest of the program-
```

once $x \geq 10$, exit loop

branch to create a loop

C code:

```
int x;
x=0;
while(x<10){
//instructions

x++;
}
```

References

1. Computer Organization and Design by David A. Patterson & John L. Hennessy (ARM edition)
2. <https://www.gnu.org/software/m4/manual/m4.html>
3. <http://infocenter.arm.com/help/index.jsp?topic=/com.arm.doc.dui0473m/dom1359731152874.html>