# 嵌入式系統-PA1

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#### Part 1

## [Global Scheduling 10%]

■ Describe how to implement Global scheduling by using pthread. 5%

用 Pthread\_create 建立新的執行緒,建立好的執行緒會以平行的方式執行,而最後再用 Pthread\_join 函數等待執行緒結束,而在設置 Global scheduling 時不必使用 setCore 設定 cpu,因為 setCore 的預設值是-1,而在-1 的情況下 thread 會配給到隨機的 cpu 上。

■ Describe how to observe task migration. 5%

```
#if (PART == 1)
    pthread_mutex_lock(&count_Mutex);
    now = sched_getcpu();
    if(now!= pre){
        now = sched_getcpu();
        std::cout << "The thread " << obj->_ID << " PID : " << obj->PID << " is moved from CPU" << pre << " to CPU" << now << std::endl;
        pre = now;
    }
    pthread_mutex_unlock(&count_Mutex);
#endif</pre>
```

為了觀察 CPU 的變化,在進到 convolution 的 for 迴圈前,令兩個變數 now 跟 pre 並把當前的 core 丟進去,之後在 for 迴圈裡檢查 cpu 是否有變更,有的話就 cout 出來並更新 pre。

### [Partition Scheduling. 5%]

Describe how to implement partition scheduling by using pthread

```
for(int i=0;i<numThread;i++){</pre>
   pthread_create (&threadSet[i]._thread, NULL, threadSet[i].convolution, &threadSet[i]);
for(int i=0;i<numThread;i++){</pre>
   pthread_join (threadSet[i]._thread, NULL);
Thread::setUpCPUAffinityMask (int core_num)
    /*~~~~~~~~Your code(PART1)~~~~~~~*/
  // Pined the thread to core.
    if(core_num == -1){
        return:
      }
    cpu_set_t set;
    CPU_ZERO(&set);
    CPU SET(core num, &set);
    if(sched_setaffinity(0,sizeof(set),&set)==-1)
        std::cout << "WRONG ";
    cur_core = core_num;
// setThreadCore=-1;
    在 partition 時,跟 global 的差異在於,必須將 thread 分配給指定的
```

## [Result. 10%]

Show the scheduling states of tasks. (You have to show the screenshot result of using the input part1 input.txt)

setUpCPUAffinityMask 函式設置 CPU。

cpu,因此這邊使用 setCore 將任務分配到指定的 core 上,之後在利用

## ▲Global Scheduling

## ▲ Partition Scheduling

#### Part 2

[Partition method Implementation. 10%]

Describe how to implement the three different partition methods (First-Fit, Best-Fit, Worst-Fit) in partition scheduling.

```
System::partitionFirstFit ()
    std::cout << "\n============= << std::endl;
#if (PART == 2)
    check->setCheckState(PARTITION_FF);
#endif
    for (int i = 0; i < CORE_NUM; i++)</pre>
       cpuSet[i].emptyCPU();
               ~~~Your code(PART2)~~~~~~*/
    // Implement parititon first-fit and print result.
    for(int i=0;i<numThread;i++){</pre>
        int q=0;
        while(cpuSet[q].utilization()+threadSet[i].utilization()>1){
            q = q+1;
            if(q>CORE_NUM){
                threadSet[i].setCore(-1);
                std::cout<<"Thread-"<<threadSet[i].ID() << " not schedulable"<<std::endl;
                continue:
        if(q<CORE_NUM){</pre>
            std::cout << q<<"_utilization : " << cpuSet[q].utilization() << std::endl;
            cpuSet[q].push_thread(threadSet[i].ID(),threadSet[i].utilization());
std::cout << q<<"_afterutilization : " << cpuSet[q].utilization() << std::endl;</pre>
            threadSet[i].setCore(q);
                        ~END~
    partitionMultiCoreConv();
    cleanMultiResult();
```

First-Fit 的排法要先將前面的 cpu 排滿後才排下一顆,因此這邊用一個 while 判斷是否排滿,排滿的話就排下一顆,而排到超過指定的 CPU 數量後,就判斷這個任務無法排程,並把 core 設成-1。

```
201
       void
202
       System::partitionBestFit ()
203 🖵 {
204
            std::cout << "\n======== Partition Best-Fit Multi Thread Matrix Multiplication======= << std::endl;
205
       #if (PART == 2)
206
            check->setCheckState(PARTITION BF);
207
       #endif
209
            // Implement partition best-fit and print result.
for (int i = 0; i < CORE_NUM; i++)
    cpuSet[i].emptyCPU(); // Reset the CPU set</pre>
210
211
212
213 🖨
            for(int i=0;i<numThread;i++){</pre>
214
215
                 float maxcutilization = -1
                for(int j=0;j<CORE_NUM;j++){</pre>
216 = 217 = 218 = 2
                     if(cpuSet[j].utilization()+threadSet[i].utilization()<=1){</pre>
                          if(maxcutilization<cpuSet[j].utilization()){</pre>
219
                              maxcutilization = cpuSet[j].utilization();
220
221
222
223日
224日
                     else{
                          if(j==(CORE_NUM-1)&&q==-1){
225
                              std::cout<<"Thread-"<< i << " not schedulable"<<std::endl</pre>
226
                              threadSet[i].setCore(-1);
227
228
                          continue
229
230
230 F
                 if(q!=-1){
232
                     cpuSet[q].push_thread(threadSet[i].ID(),threadSet[i].utilization());
233
                     threadSet[i].setCore(q);
234
235
236
                                   ~END~
237
238
239
            partitionMultiCoreConv();
240
            cleanMultiResult();
241
```

Best-Fit 的排法要先把當前利用率最高的先排滿,因此必須在排每個任務時都去比較每個 cpu 的使用率,上圖第 216 行的迴圈就是在比較 cpu 的使用率,而當 cpu 的超過指定的數量時,判斷任務無法排成,並將 core 設成-1。

```
243
244
      System::partitionWorstFit ()
245 □ {
           246
247
248
           check->setCheckState(PARTITION_WF);
249
250
251
                        ~~Your code(PART2)~
           // Implement partition worst-fit and print result.

for (int i = 0; i < CORE_NUM; i++)

cpuSet[i].emptyCPU(); // Reset the CPU set
252
253
254
255 白
           for(int i=0;i<numThread;i++){</pre>
256
               int q=-1;
257
               float mincutilization = 1;
258日
259日
260日
               for(int j=0;j<CORE_NUM;j++){</pre>
                   if(cpuSet[j].utilization()+threadSet[i].utilization()<=1){</pre>
                       if(mincutilization>cpuSet[j].utilization()){
261
                           mincutilization = cpuSet[j].utilization();
262
263
264
265
266 日
                   else{
                       if(j==3&&q==-1){
    std::cout<<"Thread-"<< i << " not schedulable"<<std::endl;</pre>
268
                            threadSet[i].setCore(-1);
269
270
271
                        continue:
272
273 |-
274 |<del>-</del>
               if(q!=-1){
275
                   cpuSet[q].push_thread(threadSet[i].ID(),threadSet[i].utilization());
276
                   threadSet[i].setCore(q);
277
278
279
280
          partitionMultiCoreConv();
281
282
           cleanMultiResult();
```

Worst Fit 的排法是先選擇利用率最小的 cpu 排程,與 Best Fit 相反, 在程式上也差不多,只是將原本找出最大的 cpu 改成找出最小的 cpu,其他就都一樣。

## [Result. 30%]

Show the scheduling states of tasks. (You have to show the screenshot result of using input part2\_input\_5.txt and part2\_input\_10.txt)

```
| Accounts | Accounts
```

▲First-Fit with input\_5.txt

```
| A pact file base | Apply part | Apply part
```

## ▲First-Fit with input\_10.txt

## ▲Best-Fit with input\_5.txt

▲Best-Fit with input\_10.txt

▲Worst-Fit with input\_5.txt

```
Thread 10: 8 PID: 1804 Core: 8 Utilization: 8.888 MatrixSize: 5501
Thread 10: 8 PID: 1804 Core: 8 Utilization: 8.888 MatrixSize: 10: 8 PID: 1805 Core: 1 Utilization: 8.888 MatrixSize: 10: 8 PID: 1805 Core: 1 Utilization: 8.888 MatrixSize: 10: 8 PID: 1805 Core: 8 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 1 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 1 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 1 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 1 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 1 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 2 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 1 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 2 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 2 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.898 MatrixSize: 10: 8 PID: 1805 Core: 3 Utilization: 8.209 MatrixSize: 10: 8 PID: 1805 Core: 4 Utilization: 8.209 MatrixSize: 10: 8 PID: 1805 Core: 5 Utilization: 8.209 MatrixSize: 10: 8 PID: 1805 Core: 5 Utilization: 8.209 MatrixSize: 10: 8 PID: 1805 Core: 5 Utilization: 8.209 MatrixSize: 10: 8 PID: 1805 Core: 5 Utilization: 8.209 MatrixSize: 10: 8 PID: 1805 Core: 5 Utilization: 8.209 MatrixSize: 10: 8 PID: 1805 Core: 5 Utilization: 8.209 MatrixSize: 10: 8 PID: 1805 Core: 5 Utilization: 8.209 MatrixSize: 10: 8 PID: 1805 Core: 5 Utilization: 8.209 MatrixSize: 10: 8 PID: 1805 Core: 8 Utilization: 8.209 MatrixSize: 10: 8 PID: 1805 Core: 8 Utilization: 8.209 Matr
```

▲Worst-Fit with input\_10.txt

## • Part 3

[Scheduler Implementation. 10%]

■ Describe how to implement the scheduler setting in partition scheduling. (FIFO with FF, RR with FF)

```
#if (PART == 3)
    /*~~~~~~~Your code(PART3)~~~~~~*/
    // Set the scheduling policy for thread.

threadSet[i].setSchedulingPolicy(SCHEDULING);
    // ...
    /*~~~~~~~END~~~~~~*/
#endif
```

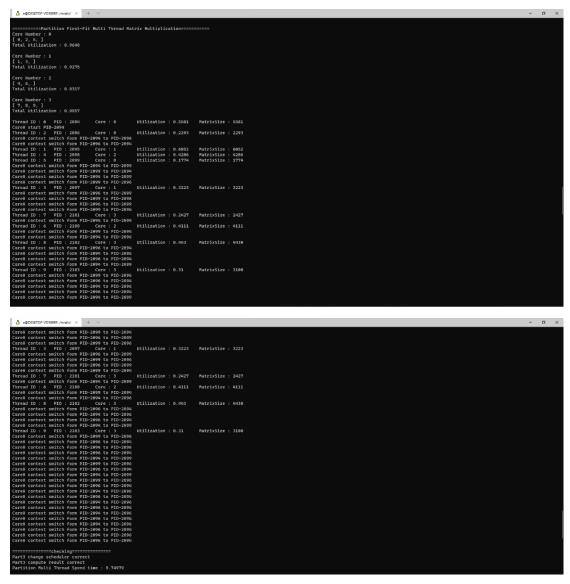
首先在 system.cpp 中根據讀進來的 SCHEDULING 設定 FIFO 或是 Round-Robin

```
void
Thread::setUpScheduler()
    /*~~~~~~~~~Your code(PART3)~~~~~~~~~/
  // Set up the scheduler for current thread
    struct sched_param param;
    int maxpri:
    maxpri = sched_get_priority_max(_schedulingPolicy);
    // std::cout << "setUpScheduler" <<std::endl;</pre>
    if(maxpri == -1){
        std::cout << "!! Failed sched_get_priority_max_ !!" << std::endl;</pre>
    param.sched priority = maxpri;
    PID = syscall (SYS_gettid);
    int ret = sched_setscheduler (PID, _schedulingPolicy, &param);
    if (ret ==-1) {
        std::cout << "!! Failed sched_setscheduler_ !!" << std::endl;</pre>
             ~~~~~~~~END~~~~~~~~~~~~~~~~~~
}
```

接著在 setUpScheduler 裡的 sched\_get\_priority\_max()拿到該排程的最大優先權,再透過 sched\_setscheduler (PID, \_schedulingPolicy, &param);設置排程就完成了。

[Result. 10%]

Show the process execution states of tasks. (You have to show the screen shot result of using input part3\_input.txt)



▲Round-Robin

#### Discussion

Analyze and compare the response time of the program, with single thread and multi-thread using in part 1 and part 2. (Including Single, Global, First-Fit, Best-Fit, Worst-Fit) 5%

首先使用多核的執行時間比起單核快上不少,而在 GLOAB 時我發現 我的電腦在全部的矩陣大小都是一萬時 print 不出 core 切換的資訊, 可能是電腦的核心比較好的原因。

而 First-Fit、Best-Fit 和 Worst-Fit 三個排法中,在輸入為 10 的時候, 最快的是 Worst-Fitst,不過這是因為他有一個 TASK 是給額外的 core 做,因此比較快也是蠻合理的。 ■ Analyze and compare the characteristic of the three different partition methods (First-Fit, Best-Fit, Worst-Fit) 5%

在10個 thread 中,First-Fit 是唯一能夠全部排進去的排法,最滿的利用率是0.9957,最少的則是0.8317,接下來 Best-Fit 的排法沒辦法將最後一個任務排程,最大的利用率是0.8804,最少的則是0.8345,最後的 Worst-Fit 則是無法將倒數第二個任務排程,最大的利用率是0.9152,最少的則是0.6499,First-Fit 的排法跟進來的任務順序有很大的關係,有可能是剛好塞的下,Best-Fit 的排法是先將最大使用率的core 塞滿,理論上感覺可以省下最多的利用率,但從結果來看分配還蠻平均的,我認為主要是任務的利用率都不算小,因此常發生塞不下的情況,最後,Worst-Fit 的排法是先排進利用率最小的core,理論上分配能夠最平均,但結果看來差挺多的,主要是從第七個任務(Thread 6)開始放的任務量差蠻多的,只看前六個任務的話還蠻平均的。

■ Analyze and compare the response time of the program, with two different schedulers. (FIFO with FF, RR with FF) 5%

FIFO 的 context switch 次數明顯比 RR 少非常多,這是因為 FIFO 在執行完當前任務前不會切換給下一個任務,而 RR 則是每個任務會輪流執行,比較執行時間的話,我的結果是 RR 比 FIFO 快 0.5 秒,不過這個差距有可能是受到電腦當時的背景程序影響,否則我認為 context switch 比較少的 FIFO 應該會快一點才對。