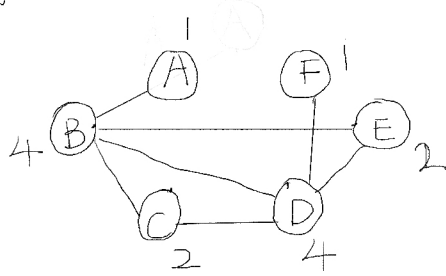


Problem 1

1.1 The vertices = A B C D E F

The edges = AB, BC, BD, BE, CD, DE, DF



1.2. 6 vertices = A B C D E F

1.3. 7 edges

1.4 It is connected.

1.5 It is connected. So there is only one component in the graph.

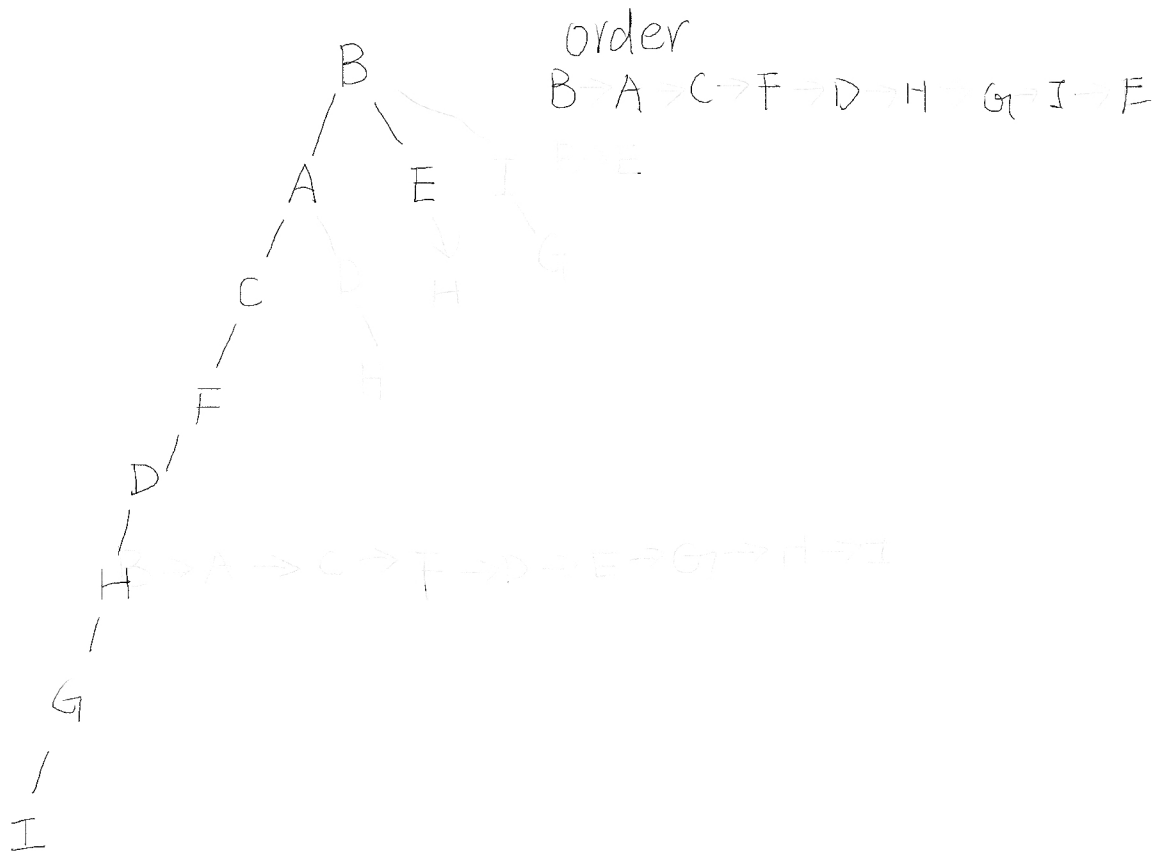
1.6 2, A and F

1.7 Yes. $A \rightarrow B \rightarrow E \rightarrow D \rightarrow B \rightarrow C \rightarrow D \rightarrow F$

1.8 No. Not all vertices have even degrees.

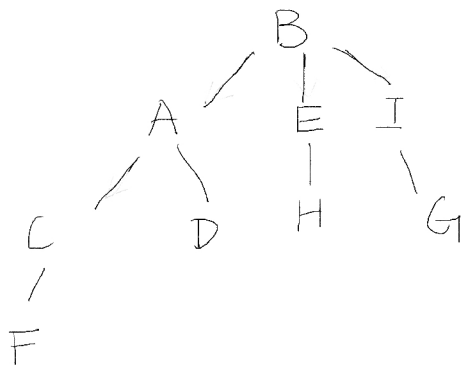
Problem 2

2.1



2.2 $B \rightarrow A \rightarrow C \rightarrow F$

2.3



order

$B \rightarrow A \rightarrow E \rightarrow I \rightarrow C \rightarrow D \rightarrow H \rightarrow G \rightarrow F$

B

2.4 $B \rightarrow A \rightarrow C \rightarrow F$

2.5 All edges = $AC = 10$

$IB = 11$

$DF = 13$

$AP = 16$

$EH = 17$

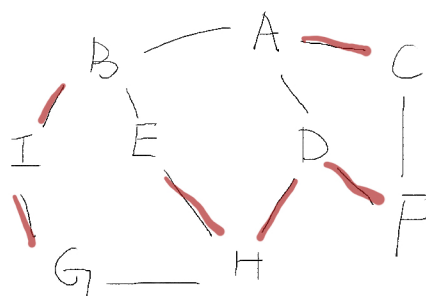
$IG = 19$

$DH = 20$

$$10 + 11 + 13 + 17 + 19 + 20 = 90$$

The lowest cost = 90.

It can be called Minimum spanning tree.



2.6

		L(A)	L(C)	L(D)	L(E)	L(F)	L(G)	L(H)	L(I)
Start	B	∞	∞	∞	∞	∞	∞	∞	∞
BI	I	2 B	∞	∞	23B	∞	∞	∞	(1 B)
BA	A	(2 B)	∞	∞	23B	∞	30I	∞	(1 B)
BAC	C	(2 B)	(3 A)	3 A	23B	∞	30I	∞	(1 B)
BAD	D	(2 B)	(3 A)	(3 A)	23B	53C	30I	∞	(1 B)
BADH	H	(2 B)	(3 A)	(3 A)	23B	50C	30I	5 D	1 B

Problem 3

3.1 There are V vertices in a complete graph.

You can start from every vertex.

It's a full permutation.

And the same circuit will be added twice.

Thus, $\frac{(V-1)!}{2}$

3.2 $\sum \text{degree}(V) = 2|E| = 4V$

$$E = 2V$$

3.3 The complete graph = $\frac{V(V-1)}{2}$

Thus, the answer is $\frac{V(V-1)}{2} - E$

Problem 4

4.1 There's only one edge.

$$\text{Thus } |E| = p \times 1 = p$$

4.2 First find all possible edges.

For complete graph (any pair of vertices in the graph), there are $\frac{n(n-1)}{2}$ edges.

Thus, the expectation is $\frac{n(n-1)}{2} p$.

Problem 5

5.1

$$5 \times 9 = 45$$

45 is an odd number.

$\sum(\text{degree}) = 2|E|$ is an even number.

So it is impossible.

5.2.

There are 12 people.

If all degrees are different. (0-11)

But if one vertex's degree is 0,

it is impossible there is one vertex's degree is 11.

But there are 12 people (vertices)

So there exist two vertices have same degree.

5.3.

According to handshake lemma, the number of vertices of odd degree must be even.

So if there is one odd-degree vertex, in the component, there must be another odd-degree vertex.

And a component is connected.

Thus any vertex with odd degree there exists a path from it to another vertex with odd degree.