CSE-5311 Algorithms

* Big-O notation

$T(n) \le C \times F(n)$

$T(n)$: Real Function, $F(n)$: O function

1) Find constant C
2) Find the smallest n (n_0)

* Big-Θ (theta) Notation
* Big-Ω (omega) Notation

* Array Operations
  - lookup
  - append
  - insert
  - delete
  - extend
  - slice

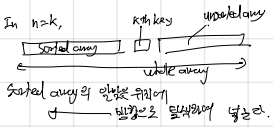
* Dynamic Array Operations
  - lookup
  - append
  - insert
  - delete

* Searching Algorithms

1. Linear Search
2. Binary Search

* Sorting Algorithms

1. Insertion Sort - n^2
2. Merge Sort - $n \cdot \log n$
3. Selection Sort - n^2
4. Bubble Sort - n^2
5. Quick Sort - n^2
   Three-way Quick Sort - n^2
   Median-of-Three Quick Sort - n^2



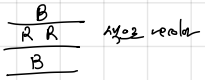
* Quick sort의 개선된 pivot

* k-th smallest selection
1. Sort the array - $n \cdot \log n$
2. Find the smallest in each iteration - $O(n \cdot k)$
3. Quick Select Algorithm - $O(n^2)$
4. Median of Medians - $O(n)$

* Inversion Count Problem
1. Brute Force - $O(n^2)$
2. Merge Sort - $O(n \cdot \log n)$
3. Bubble Sort - $O(n^2)$
4. Quick Sort - $O(n \cdot \log n)$

* Tree (Graph) Traversal
1) Breadth First Traversal
2) Depth First Traversal
   - Pre-order Traversal — Graph default.
   - In-order Traversal
   - Post-order Traversal

* Tree
1) Binary Tree
2) Complete Binary Tree - Heap
   - $i \to \lfloor \frac{i}{2} \rfloor, 2i, 2i+1$
3) Disjoint Set
4) Binary Search Tree
   - AVL Tree
   - Red Black Tree — LR, RR 은 rotation 횟수

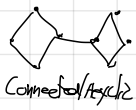


xyz 순 realm

* Splay Tree

* Graph Theory
- Distance
- Eccentricity
- Degree (of Vertex)

- Simple Graph
  - Connected / Disconnected Graph
  - Complete Graph
  - Cycle / Acyclic Graph
  - (Complete) Bipartite Graph
  - Complement Graph



- Spanning Tree (Minimum Spanning Tree)

- Cut Vatex,edge, and Set

Algorithm	Best Case	Average Case	Worst Case
Insertion Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Merge Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
Bubble Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Quick Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$
Three-way Quick Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$
Median-of-Three Quick Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$
Quick Select	$O(n \log n)$	$O(n \log n)$	$O(n^2)$
Median of Medians	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$

Algorithms

1.
```
def search(arr, x):
    for i, element in enumerate(arr):
        if element == x:
            return i
    return -1
```

2. 

$32 = 2^5 \rightarrow k=5$

$\log_2 2^{k+1} = 6$

* The last element which binary search needs to find.

3. A [1 ... n]

a. A[n] — O(1)

b. add all values while iterating to the array and divide n. O(n)

c. n : odd

$\frac{n+1}{2}$ → O(1)

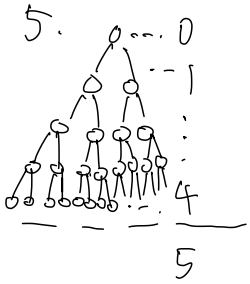
4. (a) : $O(n^3)$ (b) : $O(n^2)$ (c) : $O(\log n)$

(d) : $O(\log^2 n)$ (e) : $O(2^n)$ (f) : $O(n^5)$

$x = \log n$

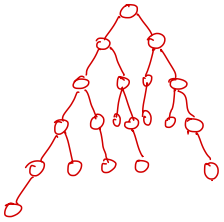
$\log x + x^2$

(g) : $O(n^2 \cdot \log n)$

5. 

$2^5 - 1 = 31$

$2^6 - 1 = 63$

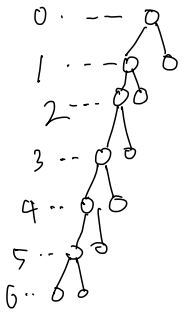


MIN: 양쪽으로 크게 뻗어 그려고,

나머지는 AVL 균형에 맞게 채우기

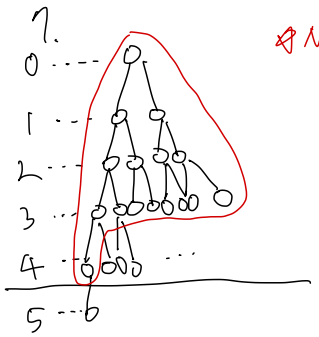
6. * Full Tree: Every node except leaf has two children w/o complete graph condition

* Perfect tree : All internal nodes have two children and all leaves are in the same level. (→ complete & full tree)



⇒ MIN: 13

Max: $2^7 - 1 = 127$ (perfect tree)

7. 

* Non-leaf node

MIN: $2^4 - 1 + 1 = 16$

Max: $2^5 - 1 = 31$

8. * Depth of Root = 0

9. J, Y, W, Z, P, A, X, C

10.

Best	Worst
logn	n
logn	logn
logn	logn

11.

Full : a, c, f

Complete : a, f

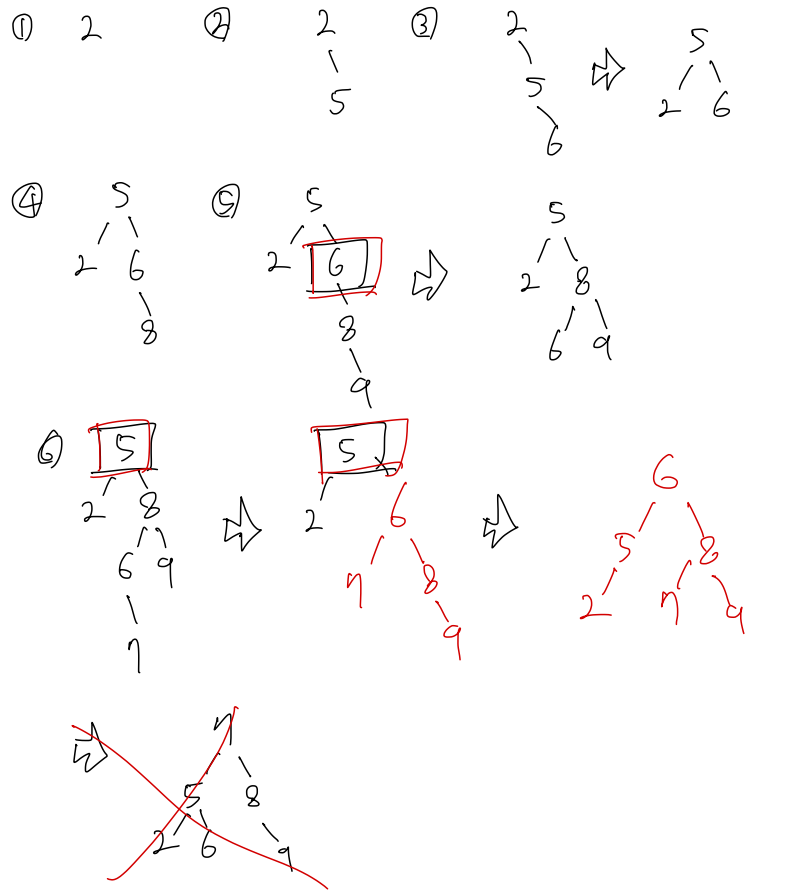
AVL balanced : a, c, e, f

Perfect : a,

12.

	sorted array	reversed ..
quick	n^2	n^2
Heap	$n \cdot \log n$	$n \cdot \log n$
Insertion	n	n^2
selection	n^2	n^2

13.



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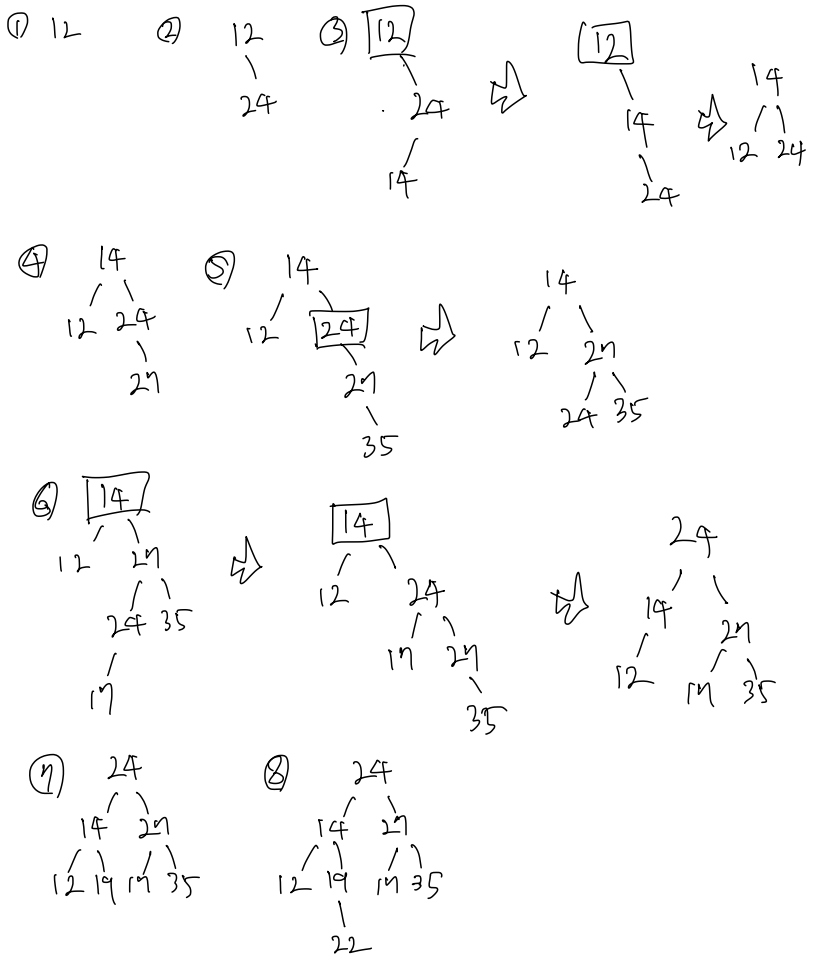

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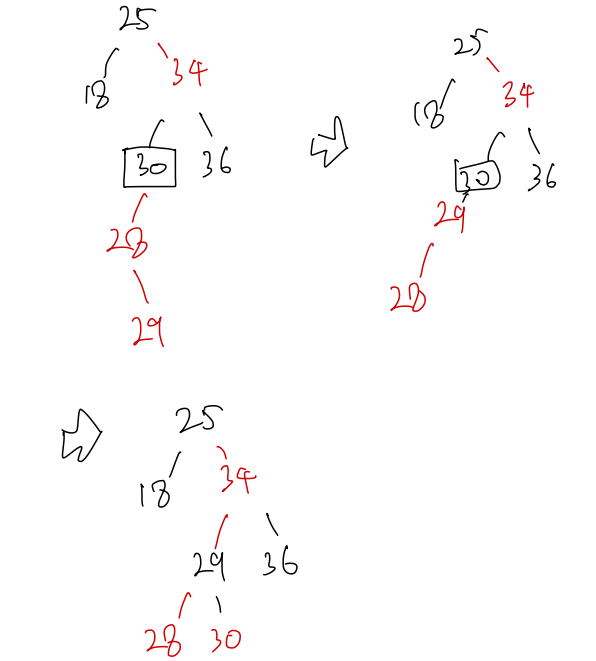

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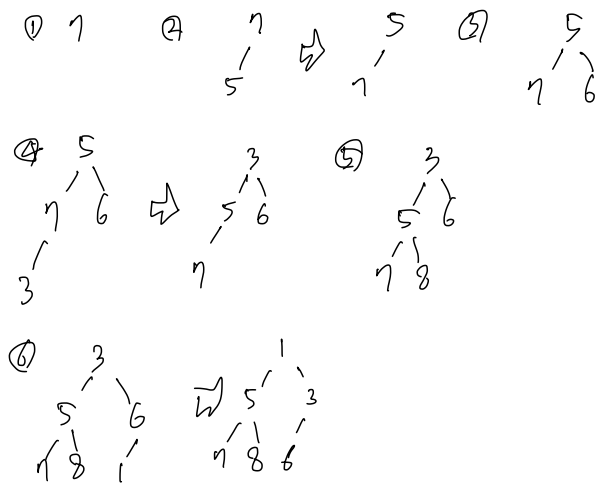

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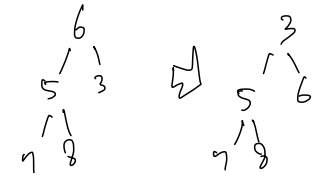

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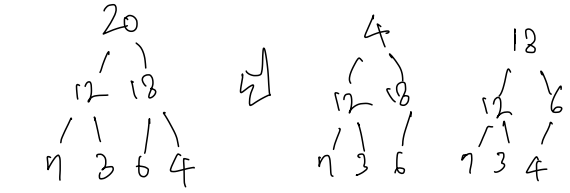


18 12 6 7 3 4

19. Y X H G T C A F B Q R

20.

Y X (P) G T (H) A F B Q R (C)

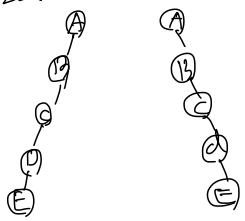
21.

H, ~~K~~ ≤ ? ≤ ~~X~~ N

H, I, J, K, L, M, N, ~~O, P, Q, R, S, T, U~~

* 기존 heap에서의 서열 관계도 고려해야 할 듯

22.



23.

* Sorted Set 2개 합치기 — 두 set에서 하나씩 pop 하면서 비교

: O(2n) = O(n)

24.

* All possible spanning trees of graph G

have the same number of edges and vertices.

(n vertices and n-1 edges)

```
missing_number(array):
    n = 10
    total = n × (n + 1)/2

    sum = 0
    for i in array:
        sum+ = i

    missing_number = total − sum
    return missing_number

remove_duplicates(array):
    n = len(array)
    present = [False] * n
    unique = []

    for num in array :
        if not present[num] :
            present[num] = True
            unique.append(num)

    return unique
```

```
find_pairs(array, target):
    array.sort()

    left = 0
    right = len(array) − 1
    pairs = []

    while left < right :
        sum = array[left] + array[right]

        if sum == target :
            pairs.append((array[left], array[right]))
            left+ = 1
            right− = 1
        elif sum < target :
            left+ = 1
        else:
            right− = 1

    return pairs
```

25.

N=1 : 1

N=3 : 2

N=7 : 4

N=15 : 8

$2^{h+1}-1 = N$ 2^h 개

$2^{h+1} = N+1$

$2^h = \left\lceil \frac{N+1}{2} \right\rceil$개

26.

1, 3, 5, 7, 9

1 2 3 4 5

~~2^k~~ -1 : # of nodes

k : # of leaf nodes.

27.

24 <

~~68~~

~~62~~

~~80~~

54

28.

* The maximum difference in Leaves of AVL

: $h - \frac{h}{2} = \frac{h}{2}$

where $h = \log n$

29.