Operating System

Ch12: I/O systems

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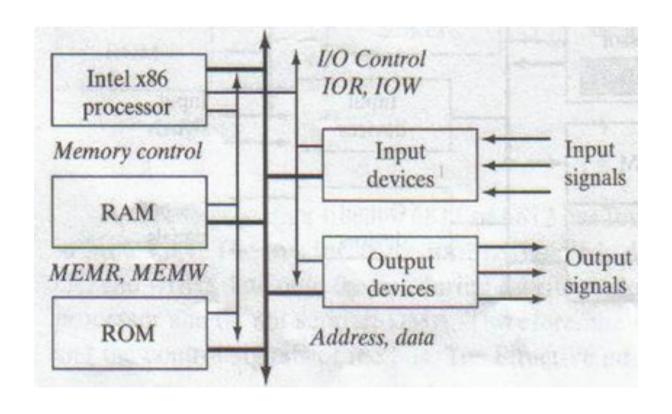
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Revisited: Computer System Operation

- I/O operation
 - ✓ I/O request via I/O instruction
 - Direct I/O vs. Memory-mapped I/O
 - ➤ Communicates with registers in I/O controller
 - ➤ Typically, IR (Instruction Register) & DR (Data Register)
 - √ I/O method
 - Programmed I/O
 - > Interrupt
 - DMA (Direct Memory Access)

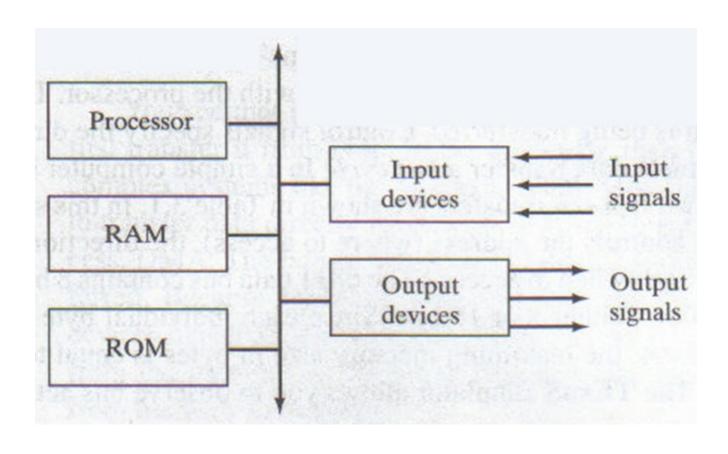
Direct I/O

- Separate addresses: memory vs. I/O
- Separate instructions
 - ✓ MEMR / MEMW (Load / Store)
 - ✓ IOR / IOW (IN / OUT)

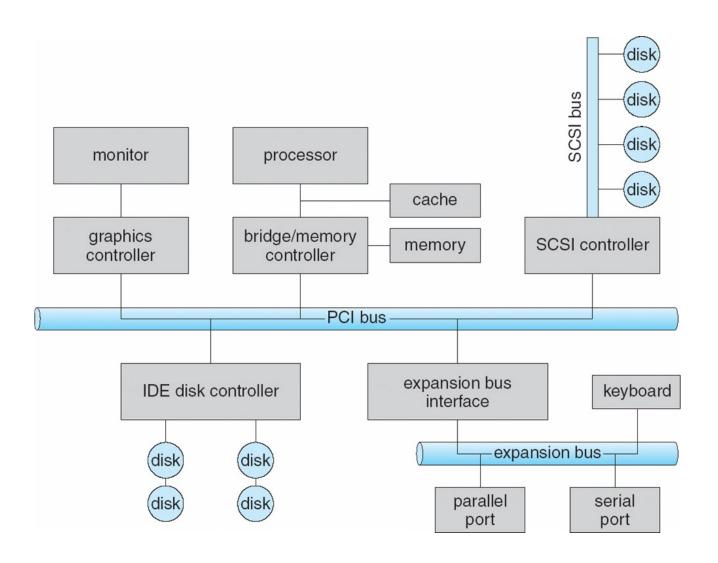


Memory-Mapped I/O

- Integrated addresses: memory & I/O
- Integrated instructions
 - ✓ MEMR / MEMW (Load / Store)



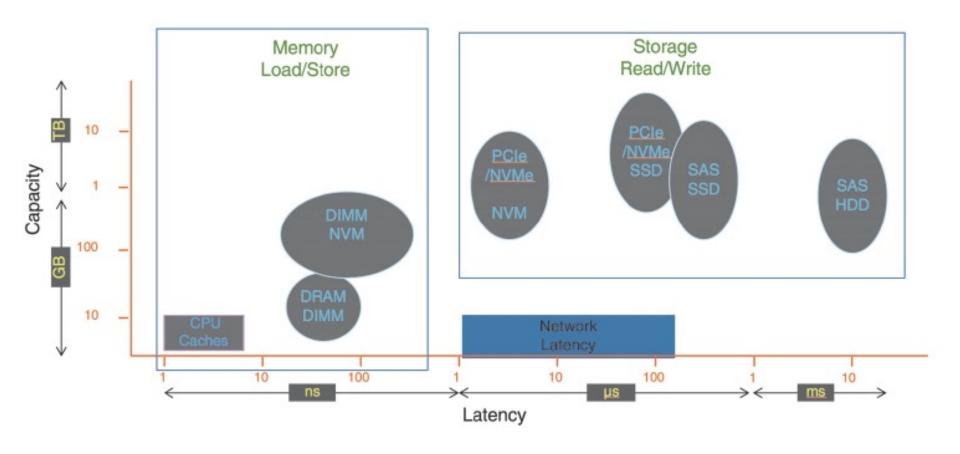
A Typical PC Bus Structure



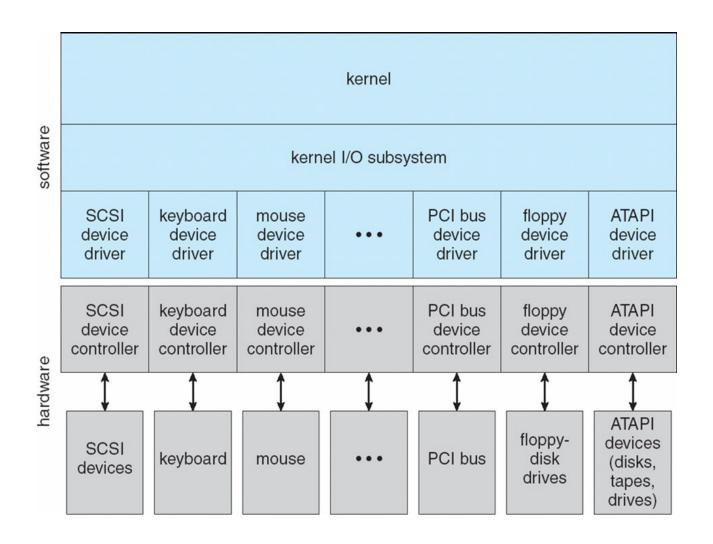
Device Controller

- Device controller (or host adapter)
 - ✓ I/O devices have components:
 - > Mechanical component
 - > Electronic component
 - ✓ The electronic component is the device controller
 - May be able to handle multiple devices
 - ✓ Controller's tasks
 - Convert serial bit stream to block of bytes
 - Perform error correction as necessary
 - Make available to main memory

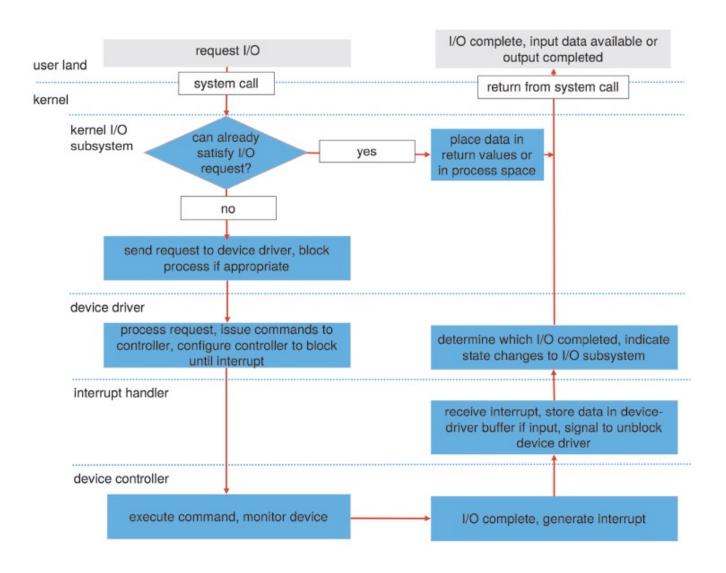
I/O Performance of Storage and Network Latency



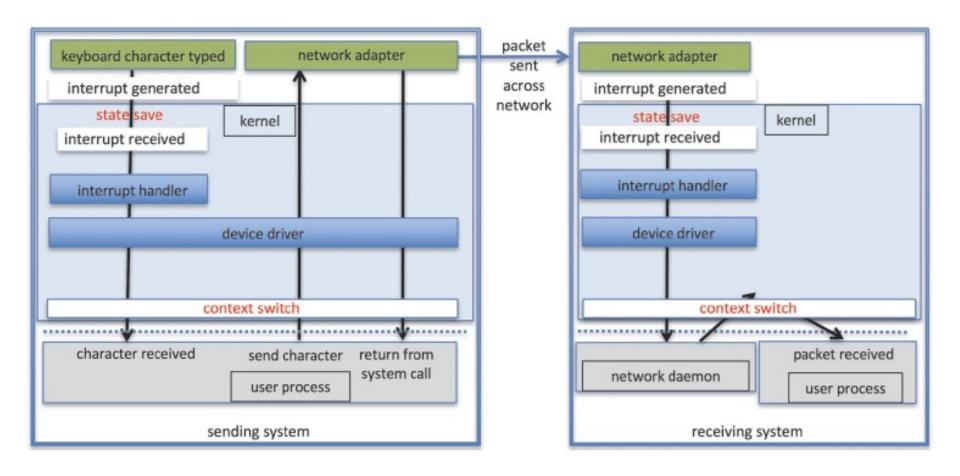
Kernel I/O Structure



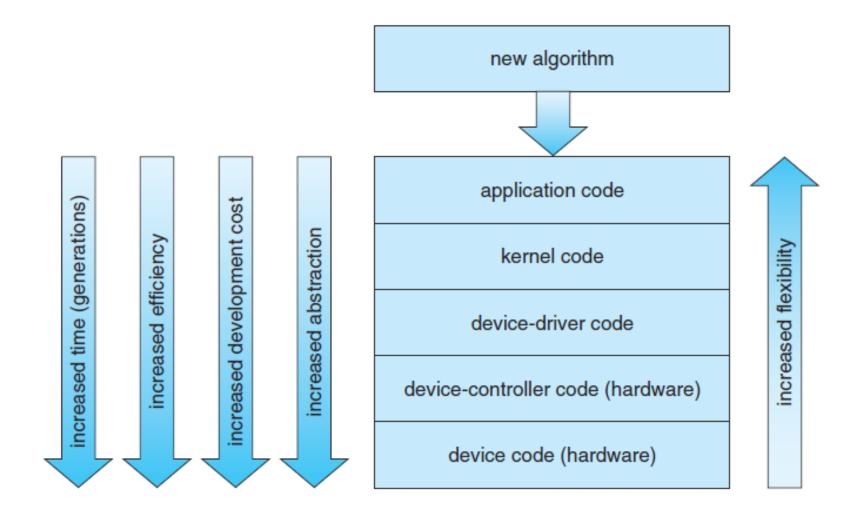
Life Cycle of An I/O Request



Intercomputer Communications



Device-Functionality Progression

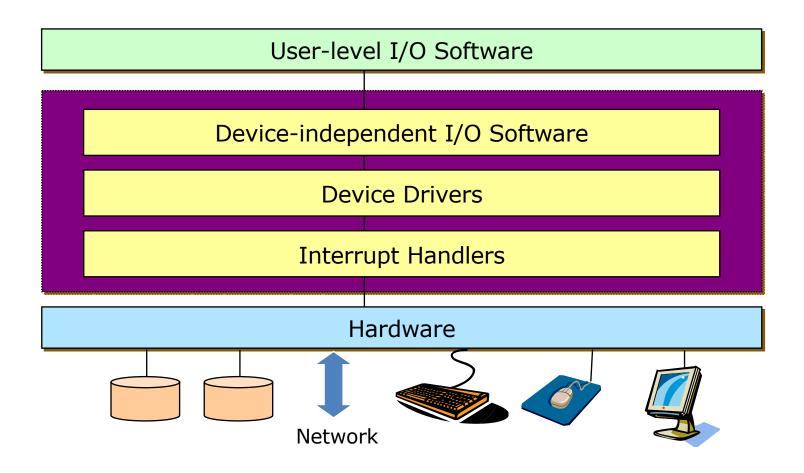


Goals of I/O Software

Goals

- ✓ Device independence
 - Programs can access any I/O device without specifying device in advance
- ✓ Uniform naming
 - Name of a file or device should simply be a string or an integer
- ✓ Error handling
 - > Handle as close to the hardware as possible
- ✓ Synchronous vs. asynchronous
 - blocked transfers vs. interrupt-driven
- ✓ Buffering
 - > Data coming off a device cannot be stored in final destination
- ✓ Sharable vs. dedicated devices
 - Disks vs. tape drives
 - Unsharable devices introduce problems such as deadlocks

I/O Software Layers



Interrupt Handlers

Handling interrupts

Critical actions

- : Acknowledge an interrupt to the PIC.
- : Reprogram the PIC or the device controller.
- : Update data structures accessed by both the device and the processor.

Reenable interrupts

Noncritical actions

: Update data structures that are accessed only by the processor.

(e.g., reading the scan code from the keyboard)

Return from interrupts

Noncritical deferred actions

- : Actions may be delayed.
- : Copy buffer contents into the address space of some process (e.g., sending the keyboard line buffer to the terminal handler process).

Bottom half (Linux)

Device Drivers

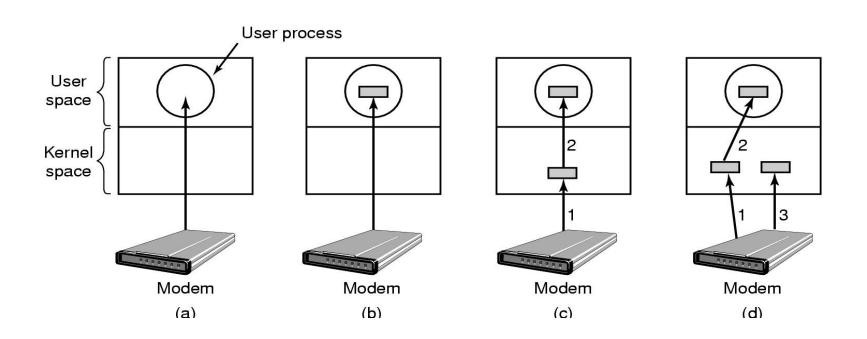
Device drivers

- ✓ Device-specific code to control each I/O device interacting with deviceindependent I/O software and interrupt handlers
- ✓ Requires to define a well-defined model and a standard interface of how they interact with the rest of the OS
- ✓ Implementing device drivers:
 - > Statically linked with the kernel
 - Selectively loaded into the system during boot time
 - Dynamically loaded into the system during execution (especially for hot pluggable devices)

- Uniform interfacing for device drivers
 - ✓ In Unix, devices are modeled as special files
 - They are accessed through the use of system calls such as open(), read(), write(), close(), ioctl(), etc.
 - > A file name is associated with each device
 - ✓ Major device number locates the appropriate driver
 - Minor device number (stored in i-node) is passed as a parameter to the driver in order to specify the unit to be read or written
 - ✓ The usual protection rules for files also apply to I/O devices.

Buffering

- √ (a) Unbuffered
- √ (b) Buffered in user space
- √ (c) Buffered in the kernel space
- √ (d) Double buffering in the kernel



Error reporting

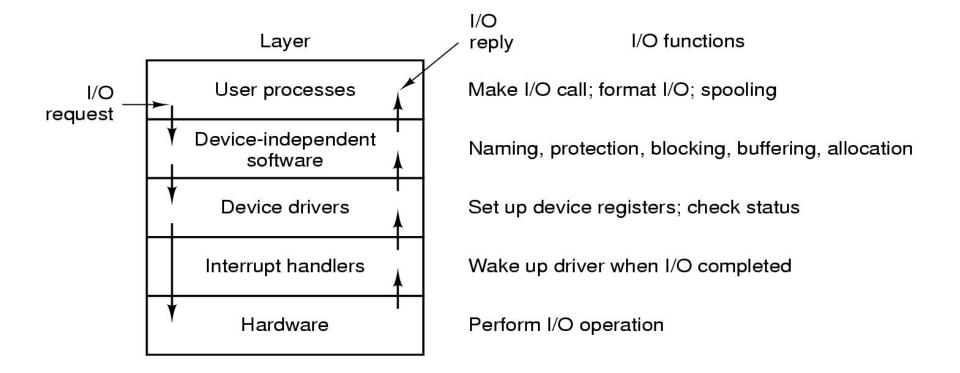
- ✓ Many errors are device-specific and must be handled by the appropriate driver, but the framework for error handling is device independent
- ✓ Programming errors vs. actual I/O errors
- √ Handling errors
 - > Returning the system call with an error code
 - > Retrying a certain number of times
 - > Ignoring the error
 - Killing the calling process
 - > Terminating the system

- Allocating and releasing dedicated devices
 - ✓ Some devices cannot be shared
 - (1) Require processes to perform open()'s on the special files for devices directly
 - > The process retries if open() fails
 - (2) Have special mechanisms for requesting and releasing dedicated devices
 - > An attempt to acquire a device that is not available blocks the caller
- Device-independent block size
 - ✓ Treat several sectors as a single logical block
 - ✓ The higher layers only deal with abstract devices that all use the same block size.

User-Space I/O Software

- Provided as a library
 - √ Standard I/O library in C
 - > fopen() vs. open()
- Spooling
 - ✓ A way of dealing with dedicated I/O devices in a multiprogramming system.
 - ✓ Implemented by a daemon and a spooling directory
 - ✓ Printers, network file transfers, USENET news, mails, etc.

I/O Systems Layers



Thank You! Q&A