

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)}$$

We can't  
use limits!

$$\begin{aligned} f &\leq g & \times \\ f &\geq g & \wedge \\ f &= g & \times \end{aligned}$$

$$f(n) = 1000n$$

$$g(n) = n^2$$

$g$  is "worse"  
than  $f$ .

For algorithm  $A$  and input  $x$

$t_A^*(x) :=$  # of steps taken  
by  $A$  on i/p  $x$ .

worst-case time of  $A$

$$t_A(n) = \max_{x: |x|=n} t_A^*(x)$$

$t_A: \mathbb{N} \rightarrow \mathbb{N}$  (No of definition  
of continuity!)

1) Consider behavior only on large inputs.

We say  $f \leq g$  if there is some  $n_0 \in \mathbb{N}$  s.t. for all  $n \geq n_0$

$$f(n) \leq g(n)$$

Note:  $1000n \leq n^2$  why?

Pick  $n_0 = 1000$   
and satisfy  
the definition.

Not  
Enough!

For any constant  $k$

$$kn \leq n^2 \quad (\text{Pick } n_0 = k)$$

Also, observe  $\lim_{n \rightarrow \infty} \frac{kn}{n^2} = 0$

2) Constant factors do not matter

$$2n^2 \leq n^2 \checkmark$$

Compare  $2n^2$  and  $n^2$   $n^2 \leq 2n^2$

$$f(n) = \begin{cases} n^2, & \text{if } n \text{ odd} \\ 2n^2, & n \text{ even} \end{cases}$$

$$g(n) = \begin{cases} 2n^2, & n \text{ odd} \\ n^2, & n \text{ even} \end{cases}$$

$$f(n) = n^2$$

$$g(n) = 2n^2$$

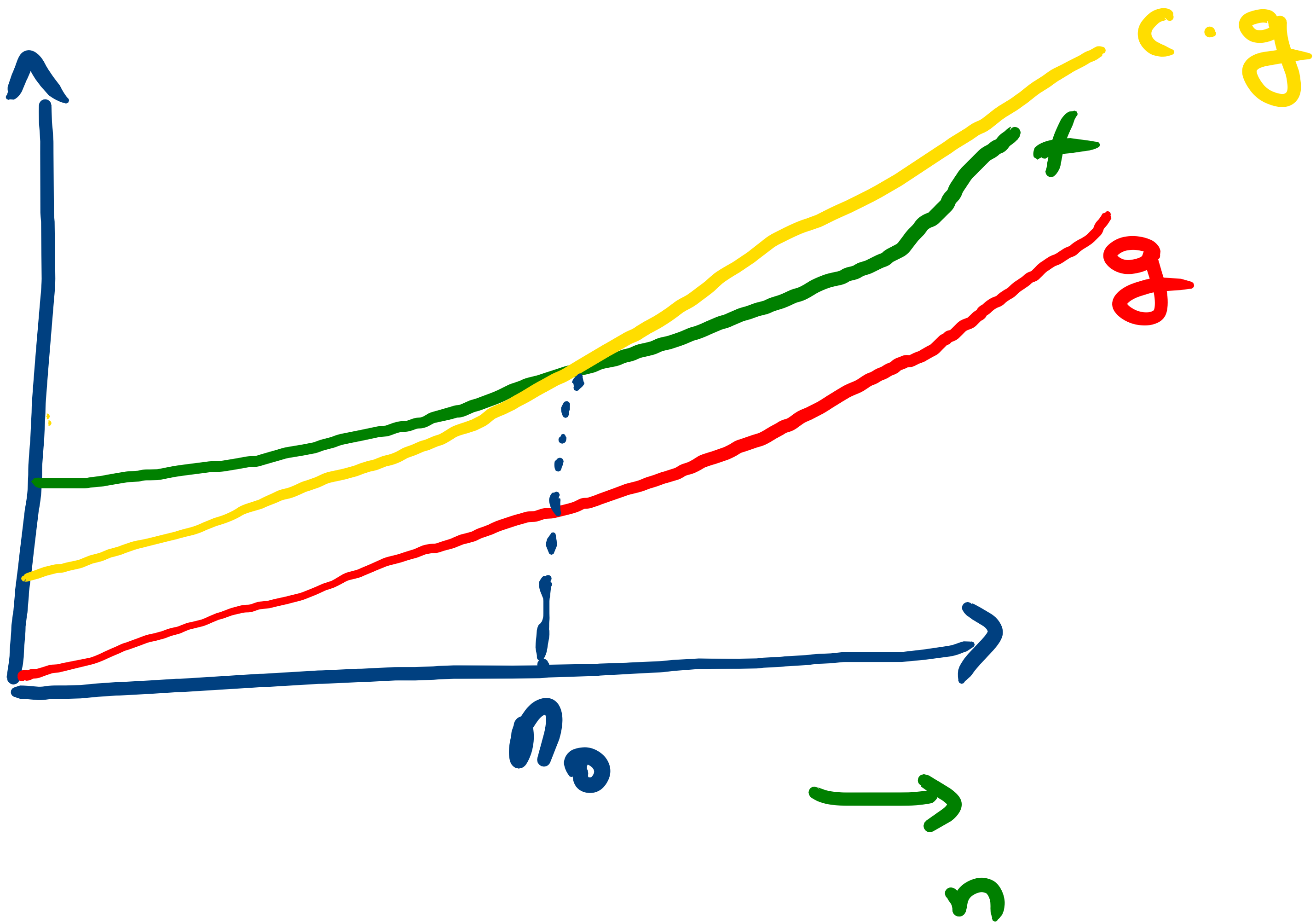
$$\lim_{n \rightarrow \infty} \frac{f(2n)}{f(n)} = 4$$

$$\lim_{n \rightarrow \infty} \frac{g(2n)}{g(n)} = 4$$

$f \leq g$  if there exists  
 $n_0 \in \mathbb{N}$ ,  $c > 0$  s.t. for all

$$n \geq n_0 \quad f(n) \leq c \cdot g(n)$$

$O(g) = \{ f \mid f \leq g \text{ as above} \}$   
"big-oh" of  $g$





For any fn.  $g$ .  $O(g)$  is  
a set of functions.

$$f \in O(g) \sim f = O(g)$$

$$f \notin O(g) \quad f \neq O(g)$$

$$O(n^2) \subseteq O(n^3) \subseteq O(2^n)$$

$$f(n) = 2^n, \quad n \leq 1000$$

$$2^{1000} + n^2, \quad n > 1000$$

$$n_0 = 1001, \quad c = 2^{1000} \left\{ \begin{array}{l} f \leq c \cdot n^2 \\ \text{for } n \geq n_0 \end{array} \right\}$$

To Prove  $f \notin O(n)$   
You have to show for all  
constants  $c > 0$  there are  
infinitely many  $n'$  s.t.  
 $f(n') > c \cdot n'$

"big-omega" of  $g$

$$\Omega(g) = \left\{ f \mid \begin{array}{l} \text{there are constants } n_0, c > 0 \\ \text{s.t. } (c \cdot g(n) \leq f(n)) \\ \text{for all } n \geq n_0 \end{array} \right\}$$

Observe  $\frac{n^2}{1000} = \Omega(n) / \frac{n^2}{1000} = \Omega(n)$

$$f(n) = \begin{cases} n, & n \text{ even} \\ n^2, & n \text{ odd} \end{cases}$$

$$g(n) = \begin{cases} n^2, & n \text{ even} \\ n, & n \text{ odd} \end{cases}$$

