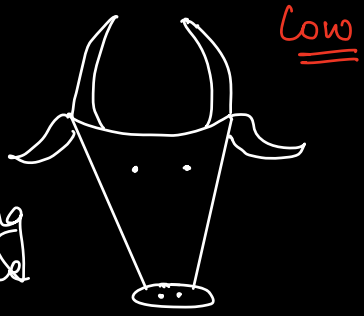
Q.1 Aggressive Cows.

Google



Given :

1) A sorted array +ve integers having the positions of the rooms where we can keep the cows. (size = N)

2) K ⇒ No. of cows.

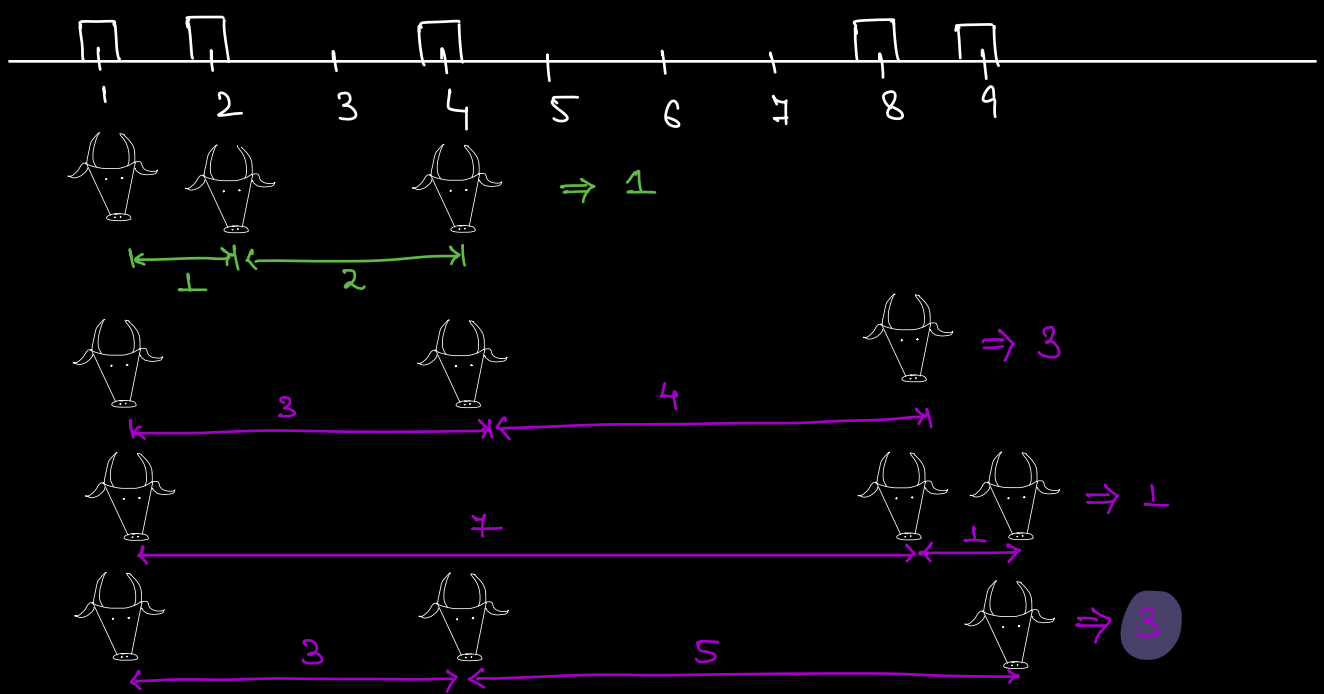
Cows are aggressive.

distance b/w any two cows.

Return the MAX value of MIN distance possible b/w any two cows.

A : 0 1, 1 2, 2 4, 3 8, 4 9 } Room positions.

K = 3



⇒ Place (K) cows in (N) rooms such that the min distance b/w any 2 closest cows is MAXIMUM

Brute Force

Try out all the possible ways of placing K cows in N rooms.

↳ N_{C_K} $\left(\frac{N!}{K!(N-K)!}\right)$

⇒ Iterate over all N_{C_K} possibilities & keep updating the distance b/w any 2 closest cows.

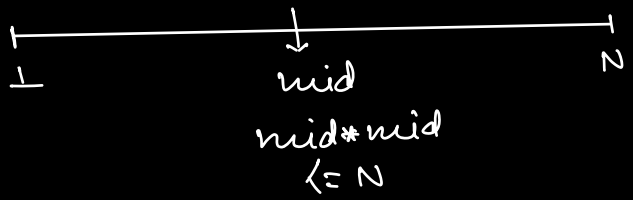
$$N_{C_K} * N \simeq O(N!)$$

sqrt(N)

i = 1

1 x 1 <= N ?
2 x 2 <= N ?
3 x 3 <= N ?
:
n x n <= N

} $\sqrt{N}$



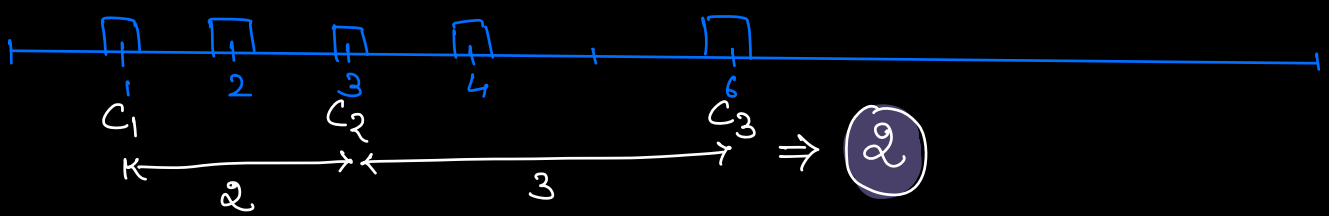
⇒ max value s.t
n * n <= N.

Target ⇒ dist b/w 2 closest cows.

ans_{max}	$A[N-1] - A[0]$
ans_{min}	1

} $Rt[1, ans_{max}]$

A: { 1 (0) 2 (1) 3 (2) 4 (3) 6 (4) } K = 3.



$ans \in [1, 5]$

d	check(A, K, d)
5	x
4	x
3	x
2	✓

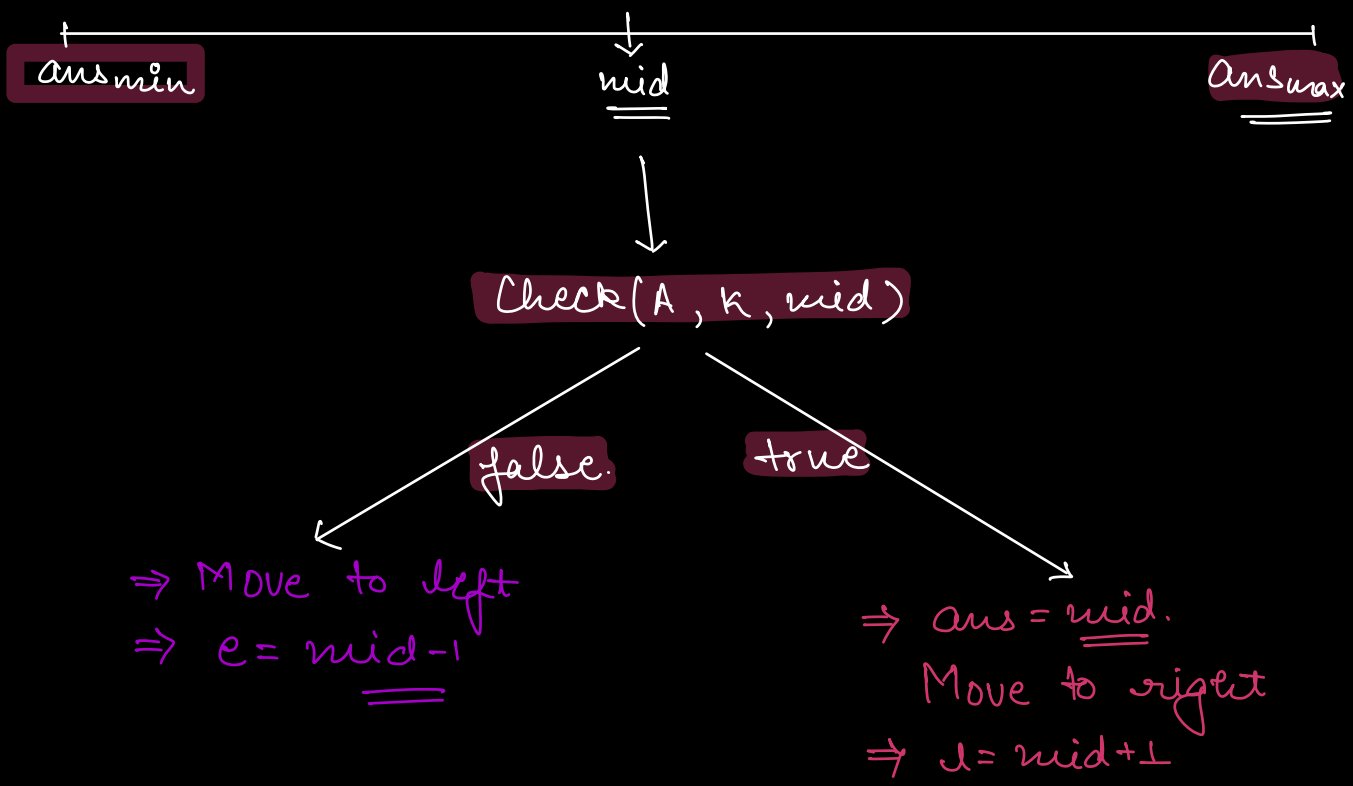
Range (R) →

```
for( d = ans_max ; d >= 1 ; d--) {
    // Check if we can place K cows
    // maintaining a min dist of d.
    if(check(A, K, d)) {
        return d;
    }
}
```

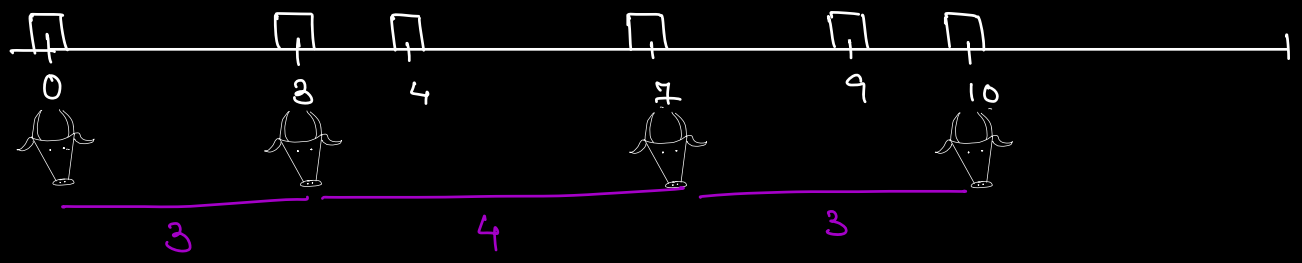
O(N)

TC : O(R · N)

SC : O(1)


ansmin
mid
ansmax
Check(A, K, mid)
false.
true
⇒ Move to left
⇒ e = mid - 1
⇒ ans = mid.
Move to right
⇒ s = mid + 1

A: { 0, 3, 4, 7, 9, 10 } K=4



ans ∈ [1, 10]

d	Check(A, K, d)
10	X
9	X
8	X
7	X
6	X
5	X
4	X
3	✓

```
bool check(A, d, K) {
    // returns true if it is possible to place
    // K cows maintaining the dist. b|w
    // any 2 cows >= d.
    int prevPos = A[0]
    int cowsPlaced = 1;
    for(i=1; i<N; i++){
        if( A[i] - prevPos >= d ){
            cowsPlaced++;
            prevPos = A[i]
        }
        if(cowPlaced == K){
            return true;
        }
    }
    return false;
}
```

$\Rightarrow \log R * N$

TC: $O(N \log R)$

SC: $O(1)$

Q.2 Google

Given **N tasks** & **K** workers.

- Array of size N.
- A[i] is the time required to complete i^{th} task.

1) One task can only be performed by one worker. (Workers can't share a task)
2) A worker can only perform tasks which are contigous to each other.
3) All workers can do their tasks parallely.

* Find the minimum amount of time required to complete all the tasks by the worker team.

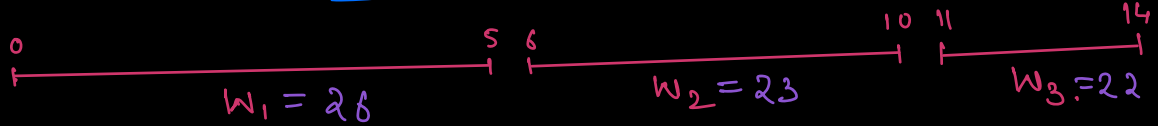
size → N = 15, K = 3

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
A:	3	5	1	7	8	2	5	3	10	1	4	7	5	4	6

$w_1 = 31$ $w_2 = 25$ $w_3 = 15$

$t = \max(w_1, w_2, w_3)$

$= 31$



$t = 26.$

$t = 25$ ✓

Target ⇒ Min time to complete all the tasks.

Range ⇒ [max(A[i]), sum(A[])]

- ans_{min} ⇒ max(A[i]) {N workers}
- ans_{max} ⇒ sum(A[]) {1 worker}

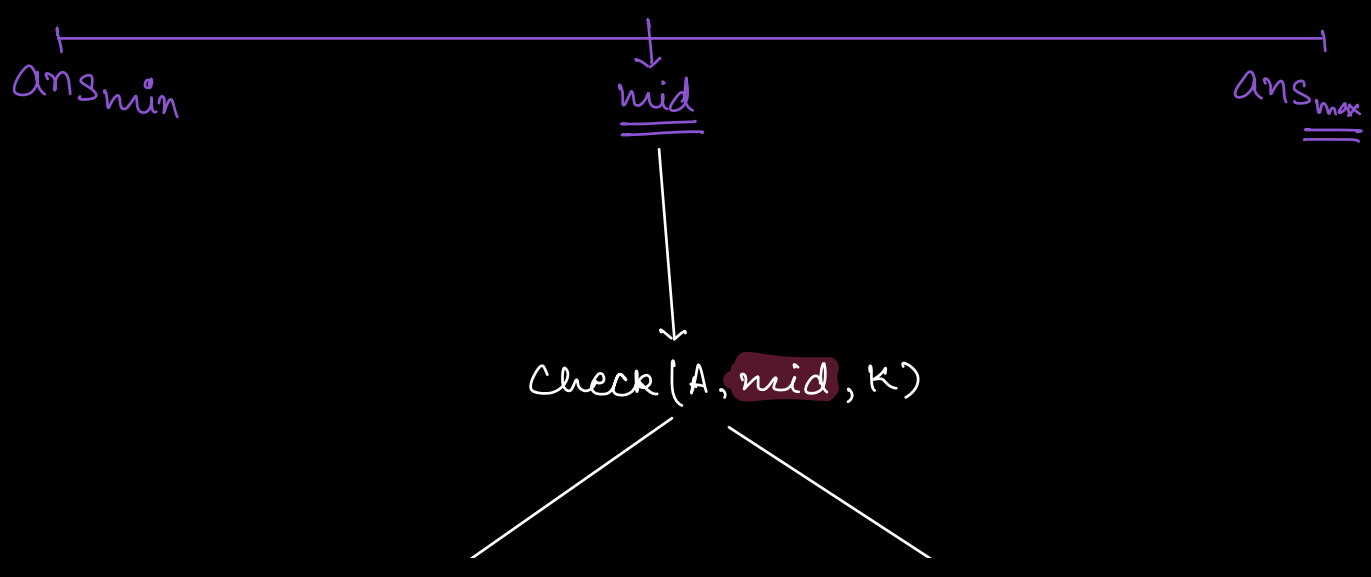
⇒ Linearly

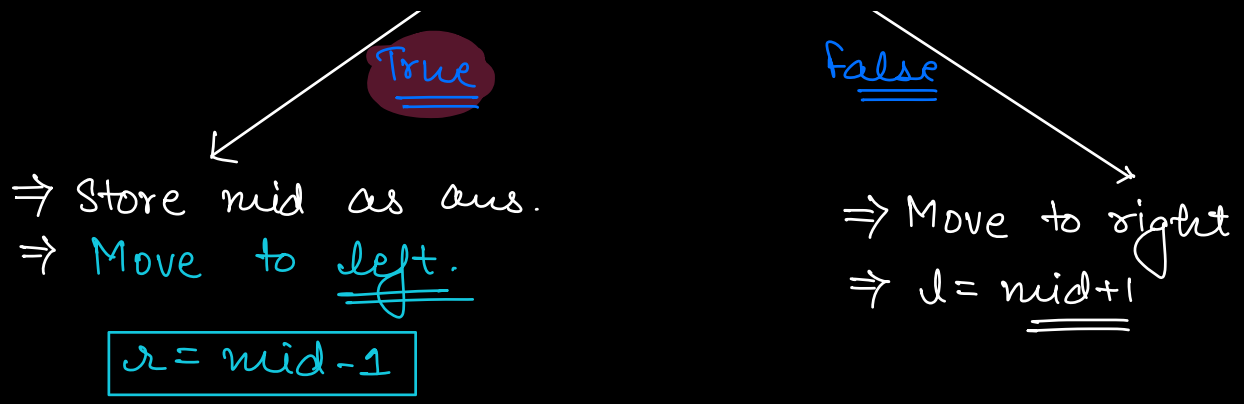
```
for(t = ansmin; t <= ansmax; t++){
    if(check(A, t, K)){
        return t;
    }
}
```

O(N)

Check if it is possible to complete N tasks in t time using K workers.

⇒ TC: O(R·N)





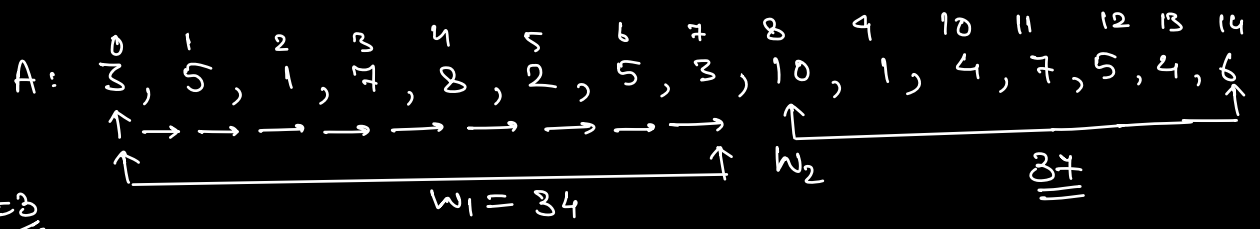

```
l = ans_min
r = ans_max        } → log R
while( l <= r ) {
    m = (l+r)/2
    if( check(A, mid, K) )
        ==========
    }
    else {
        ==========
    }
}
```

TC: $O(\log R \cdot N)$

$R \equiv [ans_{min}, ans_{max}]$

$R \equiv ans_{max} - ans_{min} + 1$

ans_{max} ↑ sum(A) ans_{min} ↑ max(A[])

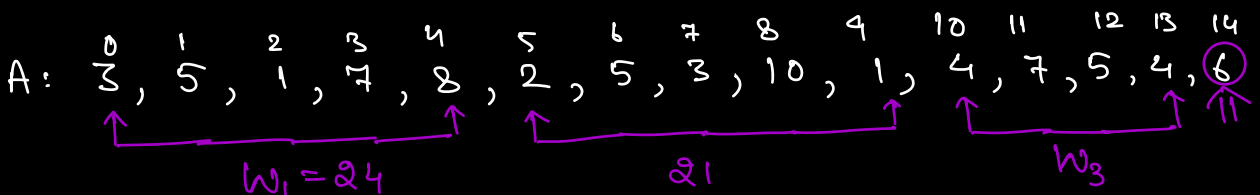


K=3

ansmin ⇒ 10
ansmax ⇒ 71
} mid = 40

ans = 40
⇒ [10, 39]

mid = 24 ×



⇒ [25, 39]

———— * ————

$a = (\log)^{b} \times c$