Alien's Dictionary

A : ["hello", "scaler", "interviewbit"]

B: "adhbcfegskjlponmirqtxwuvzy"

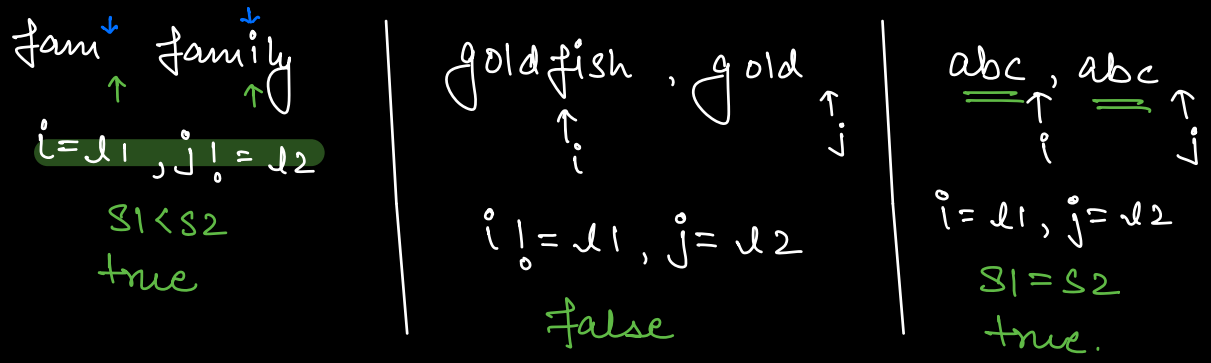
s1 = hello , s2 = hey

s1 < s2

```
bool compare (s1, s2) {
    i=0, j=0
    while( i<s1.size() && j<s2.size()){
        if (s1[i] < s2[j])
            return true;
```

ASCII values are compared.

```
        else if (s1[i] > s2[j])
            return false;
        else {
            i++, j++;
        }
    }
    i == s1.size() ? true : false;
}
```



B: "adhbcjegskjlponmirqtxwuvzy"

a → 1
d → 2
h → 3
b → 4
c → 5
j → 6
e → 7
g → 8
⋮
y → 26

HashMap<Char, Int> map.

```
bool compare (s1, s2) {
    i=0, j=0
    while( i<s1.size() && j<s2.size()){
        if (map[s1[i]] < map[s2[j]])
            return true;
        else if (map[s1[i]] > map[s2[j]])
            return false;
        else {
            i++, j++;
        }
    }
    i == s1.size() ? true : false;
}
```

A : ["hello", "scaler", "interviewbit"] ⇒ N

```
for ( i=0; i<N-1; i++){
    if (!compare(A[i], A[i+1])) {
        return false;
    }
}
return true;
```

TC : O(N·d)
SC : O(1)

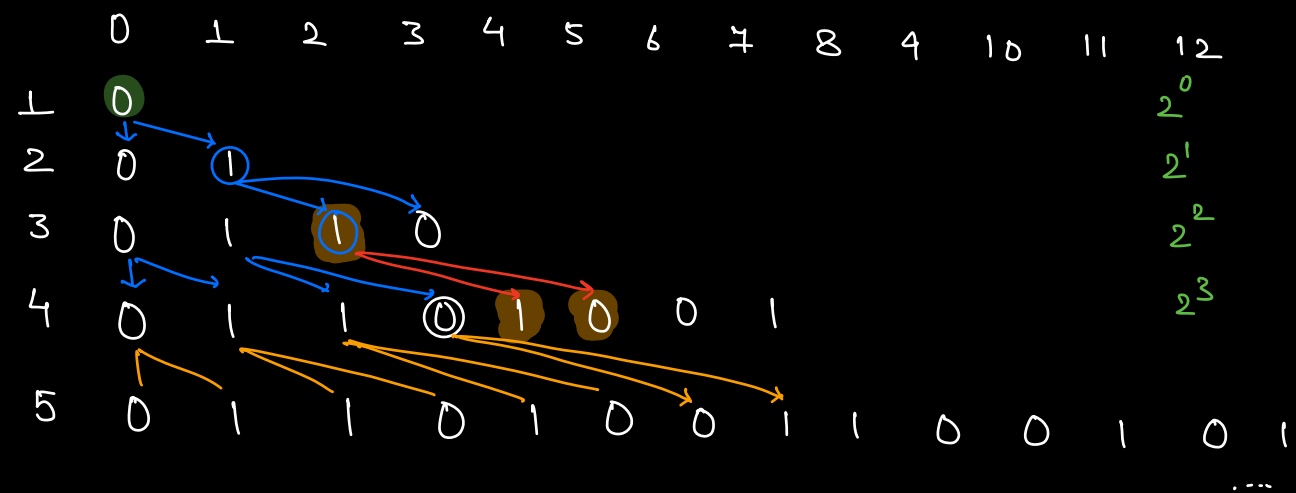
Q. K^{th} Char

→ Each row is generated by replacing all elements from previous row from

$0 \rightarrow \{0\ 1\}$

$1 \rightarrow \{1\ 0\}$

→ Given N & K, find K^{th} character in N^{th} row.



N=5, K=8 ⇒ 1

N=5, K=10 ⇒ 0

N ⇒ $\underbrace{2^0 + 2^1 + 2^2 + \cdots + 2^{N-1}}_{G.P}$

a = 1

r = 2

N = N

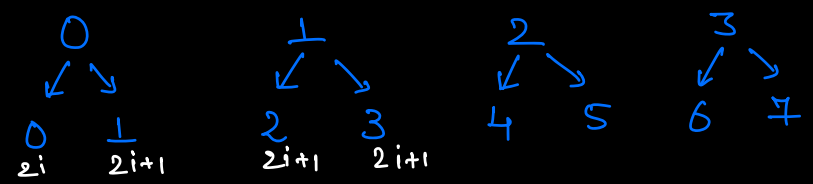
$$Sum = \frac{1 \cdot (1-2^N)}{1-2} = 2^N - 1$$

TC: $O(2^N)$

Constraints.

$N <= 10^5$

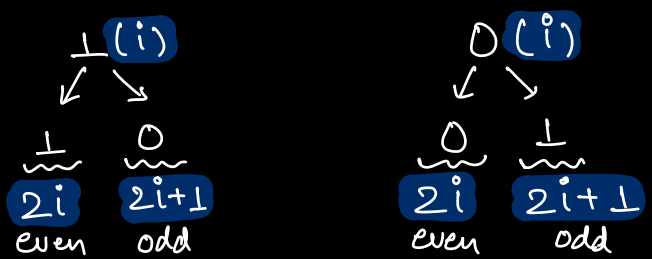
Index :



index i

$2i$ $\quad$ $2i+1$

index $x \Rightarrow$ Parent $\frac{x}{2}$



Element at even index $\Rightarrow$ same as parent

Element at odd index $\Rightarrow$ $\sim$(Parent element)

$N=5, K=8 \Rightarrow 1$

$\hookrightarrow N=4, K=4 \Rightarrow 1$

$\hookrightarrow N=3, K=2 \Rightarrow 1$

$\hookrightarrow N=2, K=1 \Rightarrow \sim 0 = 1$

$\hookrightarrow N=1, K=0$

0

```
int KthChar(N, K) {
    if(K==0) return 0;
    parent = K/2
    par_value = KthChar(N-1, parent);
    if(K%2 == 0) {
        return par_value;
    }
    return 1-par_value;
}
```

Q. Given a Binary Array, Calculate the no. of subarrays whose OR = 0

↳ Bitwise OR of all subarray elements = 0.

```
    0  1  2  3  4
A:  1  0  1  0  0
```

[1-1]
[3-3]
[4-4]
[3-4]
} 4

Observation:

In subarray with bitwise OR = 0, there shouldn't be any 1.

```
    0  1  2  3  4  5  6  7  8  9
A:  1  1  0  0  0  1  0  1  0  0
    X  X                 X
```

3(3+1)/2 = 6

1(1+1)/2 = 1

2(2+1)/2 = 3

⇒ 10

In an array of size N ⇒

# of subarrays = N(N+1)/2

```
ans = 0
c = 0
for( i = 0; i < n; i++ ) {
    if ( a[i] == 0 )
        c++
    else {
        ans += c(c+1)/2
        c = 0
    }
}
return ans;
```

TC: O(N)
SC: O(1)

#

No. of subarrays with Bitwise OR = 1

$$= \frac{n(n+1)}{2} - \text{ans}$$
