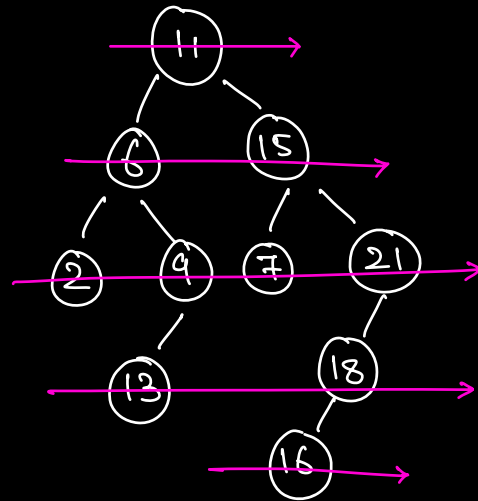


* Different type of Tree traversals & Views:

Q.1 Given a Binary Tree, Print the level order traversal.

11, 6, 15, 2, 9, 7, 21, 13, 18, 16



①

Queue: Queue <Node> q;



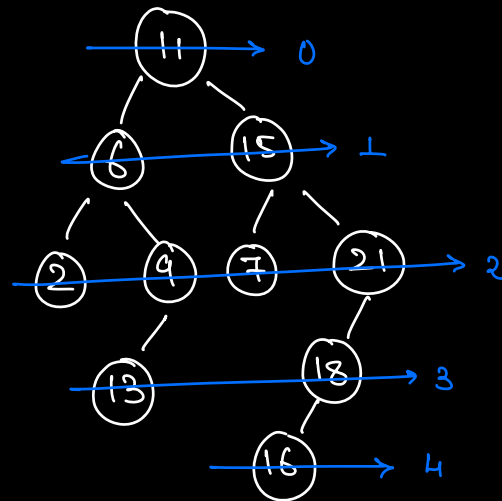
11, 6, 15, 2, 9, 7, 21, 13, 18, 16

② [[11],
 [6, 15],
 [2, 9, 7, 21]
 [13, 18]
 [16]]

list < list < int > > ans;

* Give the output in the form of list of list of int.

* All the nodes of a level will be in separate list.



① Along with the nodes we can maintain the level of the Node in the Queue.

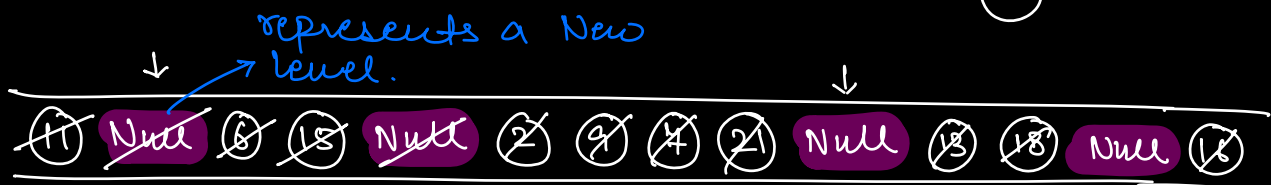
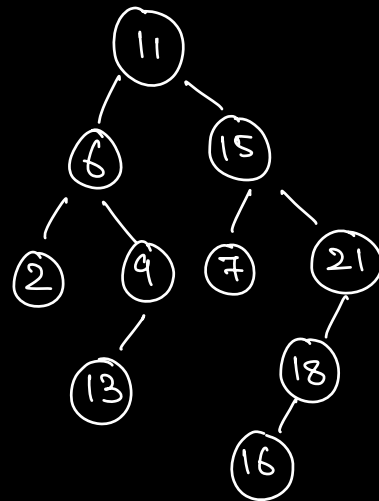
~~<11, 0>~~ <6, 1> <15, 1> ~~<2, 2>~~ ~~<9, 2>~~ <7, 2> <21, 2> <13, 3> ...

Queue < pair < Node, int > >

[11]
 [6, 15]
 [2, 9,

① Queue <Node> q;

* Add a marker/delimiter Node after every level.

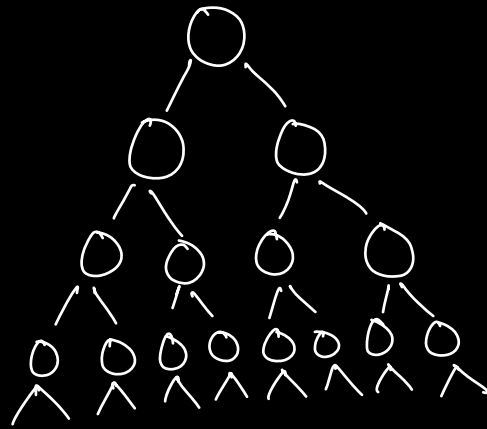


[[11]
 [6, 15]
 [2, 9, 4, 21]
 [13, 18]
 [16]]
 ==

TC: $O(N)$

SC: Queue size

Man. no. of nodes in
 the Queue.

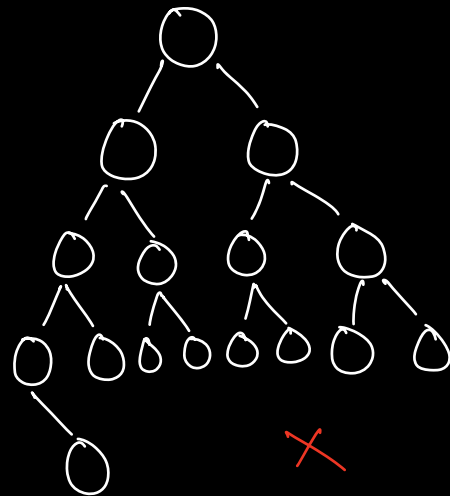
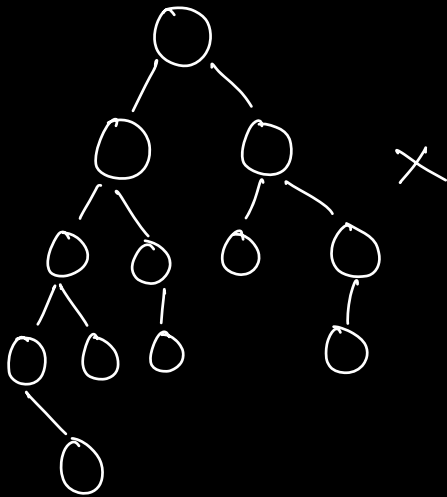


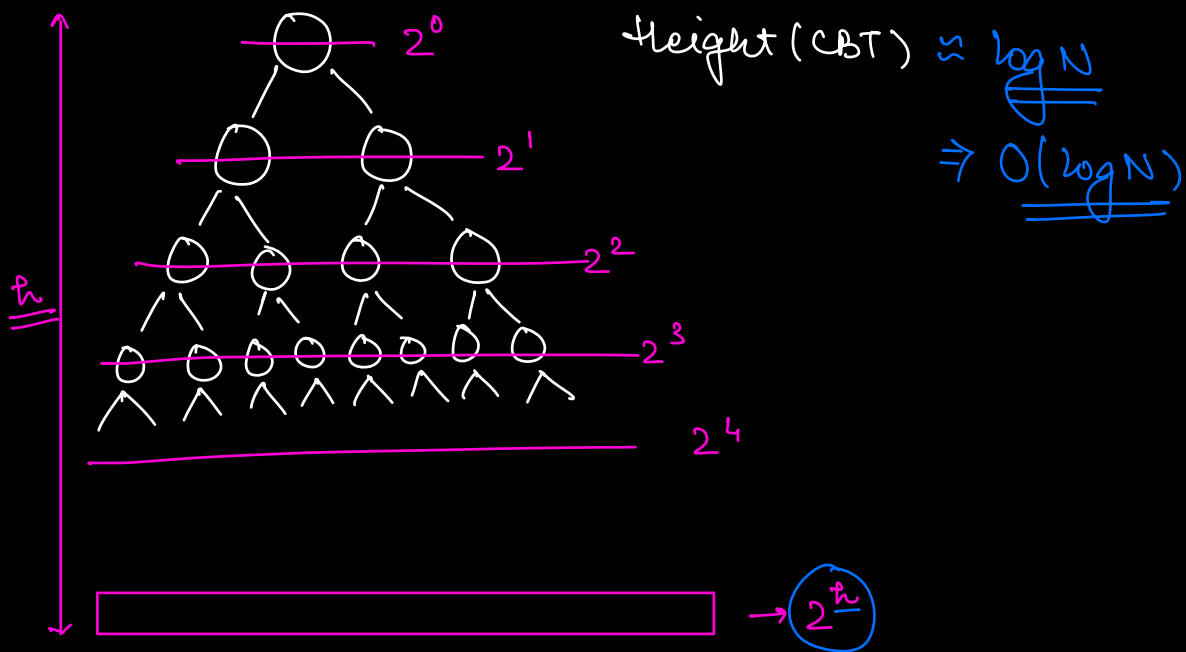
* Complete Binary Tree (CBT)

A B.T where all the levels are completely filled

↳ except possibly the last level.

→ Nodes in the last level are left aligned.





$$\# \text{ of nodes} = 2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^h$$

(Geometric
Progression)

$$\text{Sum} = \frac{a(1-r^n)}{(1-r)}$$

$$a = 2^0 = 1$$

$$r = 2$$

$$\# \text{ of terms} = \underline{\underline{h+1}}$$

$$\# \text{ of nodes} = \frac{2^0(1-2^{h+1})}{1-2} = 2^{h+1} - 1$$

$$\# \text{ of nodes} = 2^{h+1} - 1 = \underline{\underline{N}}$$

$$2^{h+1} = N+1$$

$$\log_2(2^{h+1}) = \log_2(N+1)$$

$$h+1 = \log(N+1)$$

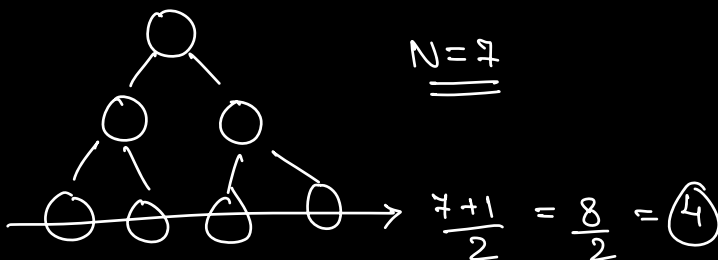
$$h = \log(N+1) - 1 \approx \underline{\underline{\log N}}$$

$$\# \text{ of nodes in last level} = 2^{(\log(N+1) - 1)}$$

$$a^{\log_a x} = x$$

$$= \frac{2^{\log_2(N+1)}}{2}$$

$$= \underline{\underline{\frac{N+1}{2}}}$$

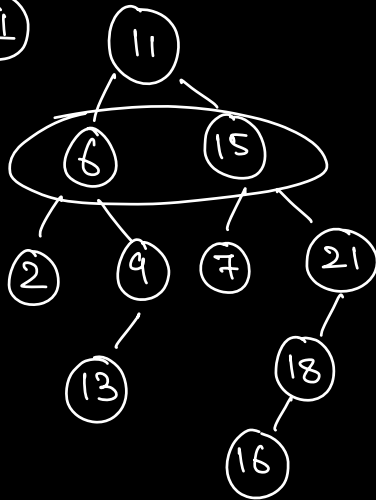


SC: Max size of the Queue

$$\Rightarrow \frac{(N+1)}{2}$$

$$\Rightarrow \underline{\underline{O(N)}}$$

III



of dequeue
operations.

Size of
Queue

~~11~~ 6 15

1

1

~~6~~ ~~15~~ 2 9 7 21

2

2

~~2~~ ~~9~~ ~~7~~ ~~21~~ 13 18

4

4

~~13~~ ~~18~~ 16

2

2

~~16~~

1

No. of dequeue operations at each level \equiv
Size of the Queue.

```

* list<list<int>> levelOrder(root) {
    if (root == NULL)
        return empty list;
    Queue<TreeNode> q;
    list<list<int>> ans;
    q.enqueue(root);
    while (q.size() > 0) {
        list<int> level;
        size = q.size();
        for (i = 0; i < size; i++) {
            TreeNode temp = q.front();
            q.dequeue();
            level.add(temp.data);
            if (!temp.left)
                q.enqueue(temp.left);
            if (!temp.right)
                q.enqueue(temp.right);
        }
        ans.add(level);
    }
    return ans;
}

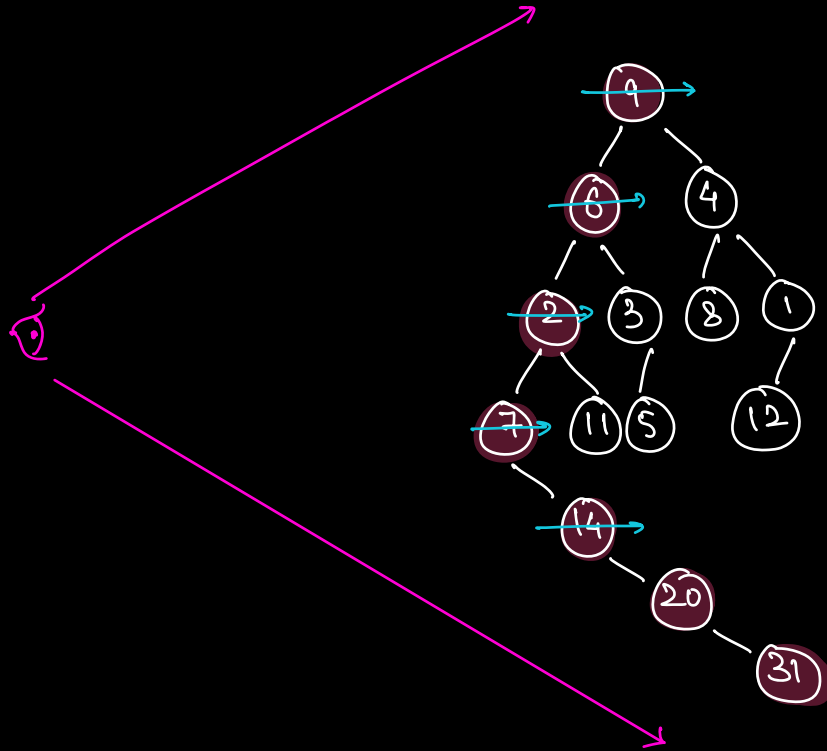
```

TC: $O(N)$

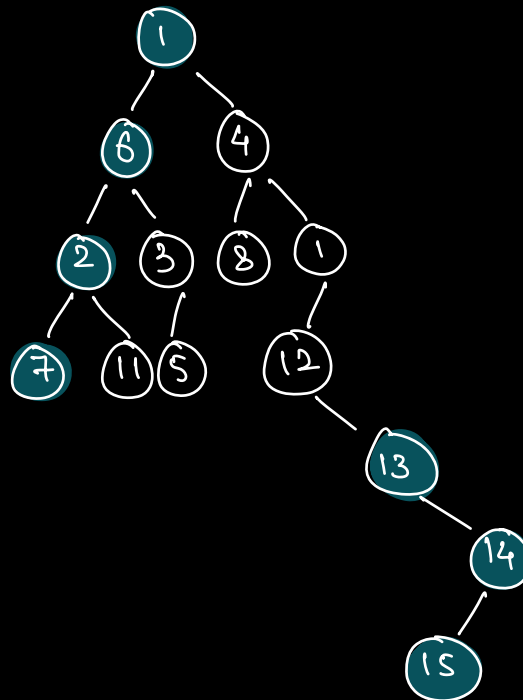
SC: $O(N)$

Q-2 Given a B.T, print its Left View.

Amazon
MS
Adobe.

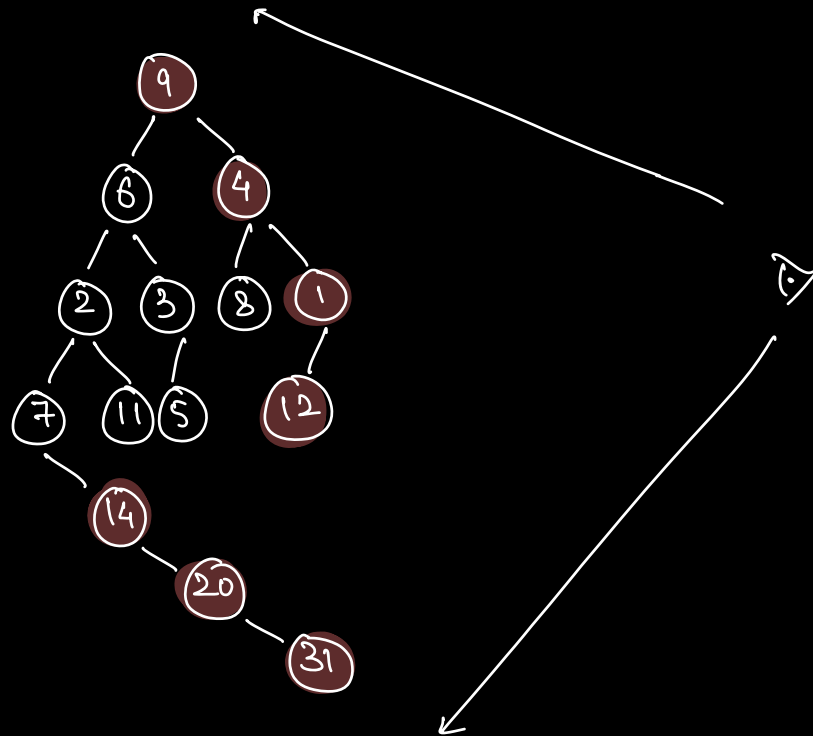


Left View : first Node of every level.



Q-2 Given a B.T, print its right View.

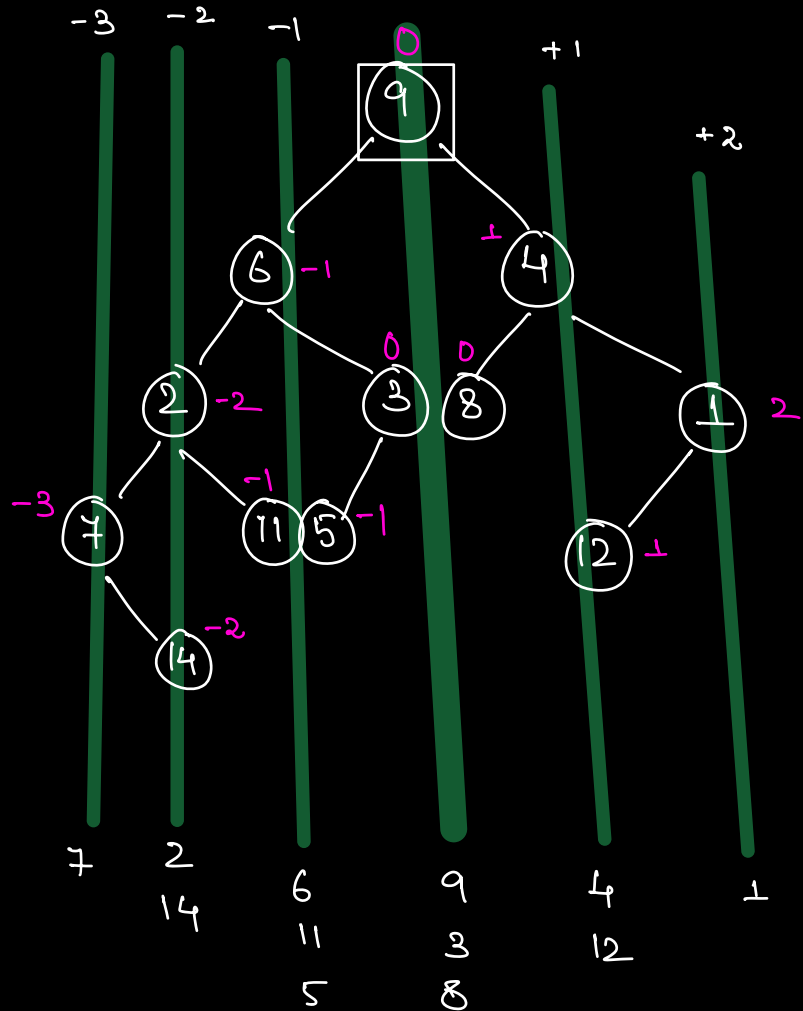
Amazon
MS
Adobe.



Right View : Last node at each level.

Q. Vertical Order Traversal

[[7],
 [2,14],
 [6,11,5],
 [9,3,8],
 [4,12],
 [1]]
 ==



-3 → 7

-2 → 2, 14

-1 → 6, 11, 5

0 → 9, 3, 8

⋮

→ <K, V>

→ HashMap<int, list<int>>

HashMap<int, List<TreeNode/Int>>

↑
distance (d)
from the
root.

↑
list of nodes
which are (d) distance
apart from root.

→ Map: Global

preOrder (root, dist) {

if (root == null)

return;

if (!map.containsKey(dist)) {

map.insert(dist, new ArrayList<int>());

3

map.get(dist).add(root.data);

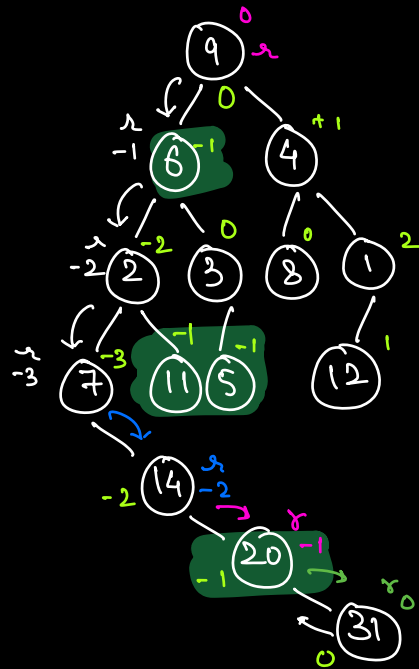
preOrder (root.left, dist+1);

preOrder (root.right, dist+1);

3

Hash Map

dist: List of int

$$0 : [9, 31]$$
$$-1 : [6, 20,$$
$$-2 : [2, 14,$$
$$-3 : 17,$$


- * PreOrder
- * Inorder
- * Postorder

} Dfs.

⇒ Queue : Level Order Traversal ✓

Queue $\langle \text{pair} \langle \text{TreeNode}, \text{int} \rangle \rangle$
↑↑
dist

```

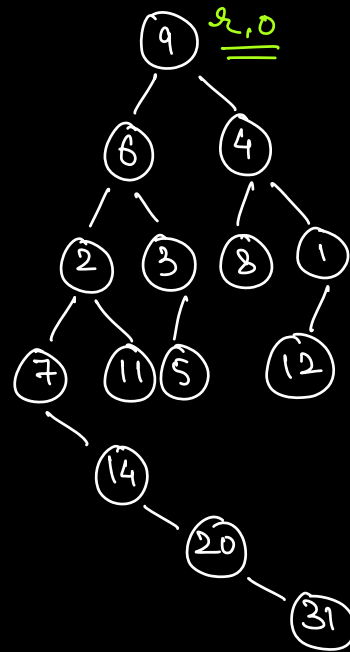
Class TreeInfo {
    Tree NODE node;
    int dist;
}

```

3

Queue (TreeInfo)

Queue < TreeInfo > q;



~~{9, 0}~~ ~~{6, -1}~~ ~~{4, 1}~~ ~~{2, -2}~~ ~~{3, 0}~~ ~~{8, 0}~~ ~~{1, 1}~~
~~{7, -3}~~ ~~{11, -1}~~ ~~{5, -1}~~ {12, 1} {14, -2}

HashMap

dist : List of int

0 : [9, 3, 8,

-1 : [6, 11, 5,

1 : [4,

-2 : [2,

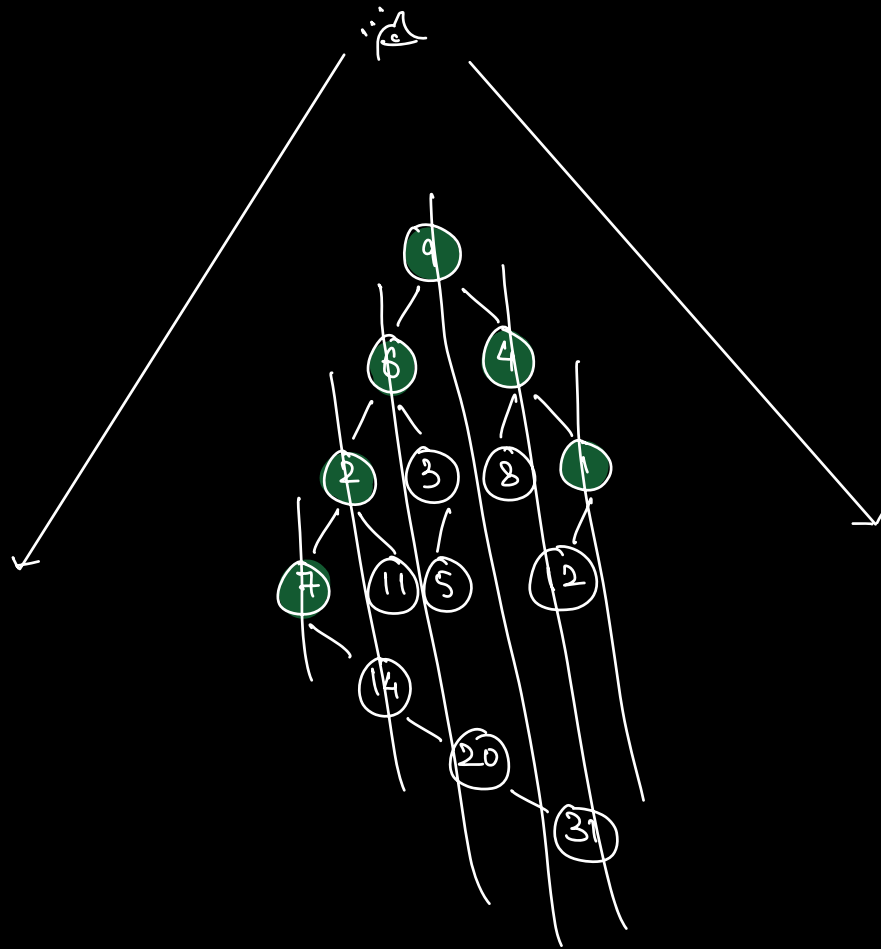
2 : [1,

-3 : [7,

min_dist → -3

max_dist → +3

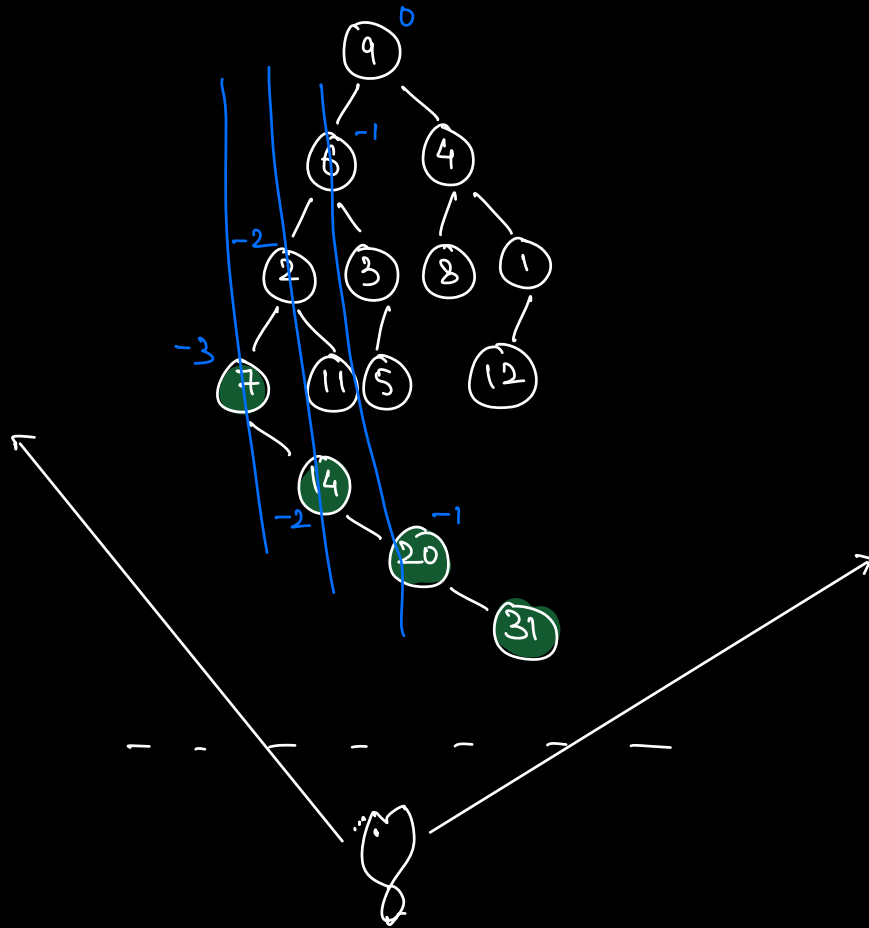
Q. Top View of Binary Tree.



⇒ First node in the list against every key (dist) in the Map.

Instead of list create single entry key->value pair in above code

Q. Bottom View of Binary Tree.



Bottom View : last node in the list against every key in the Map.

_____ * _____