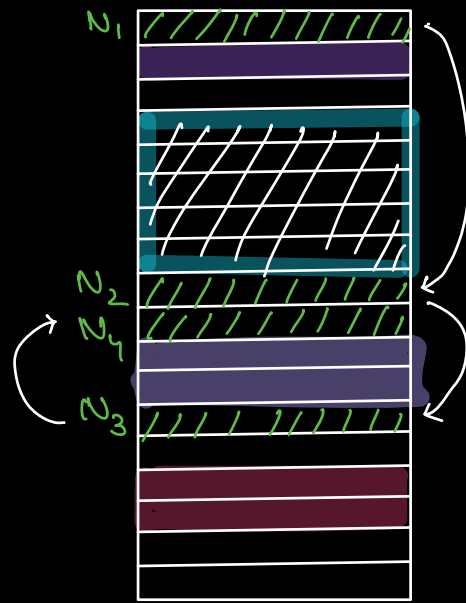


array \Rightarrow Continuous memory locations.

$\Rightarrow O(1)$ random access.

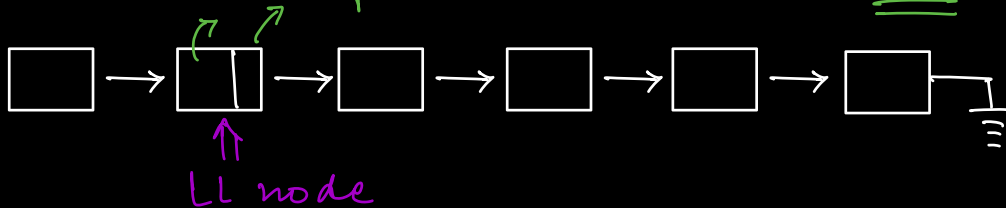


* Worst case TC in dynamic array $\Rightarrow \underline{\underline{O(N)}}$

* Amortized TC of insertion $\Rightarrow O(1)$ in dynamic array.

linked list

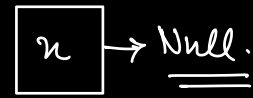
data reference to the next node.



```

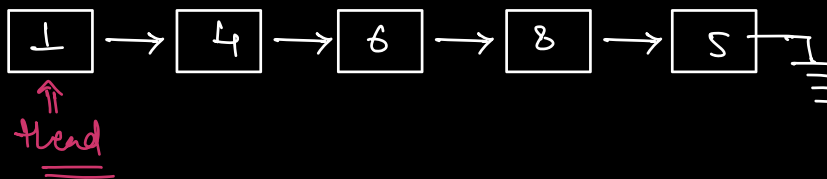
Class Node {
    int data;
    Node next;
    Node (int n) {
        data = n;
        next = Null;
    }
}

```

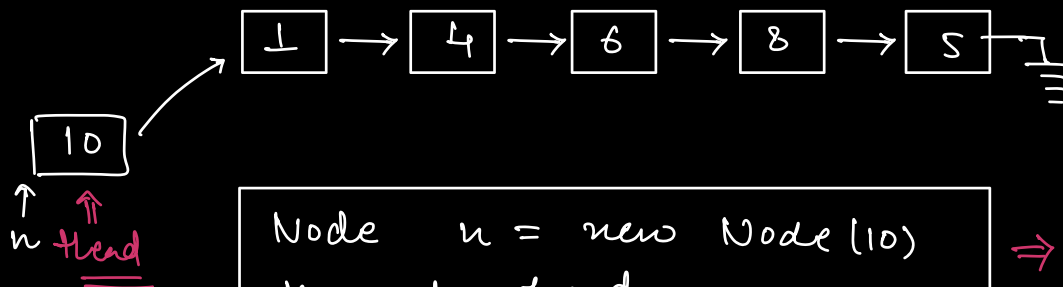


Insertion in L.L

- front
- kth position
- end



① Insert at head.



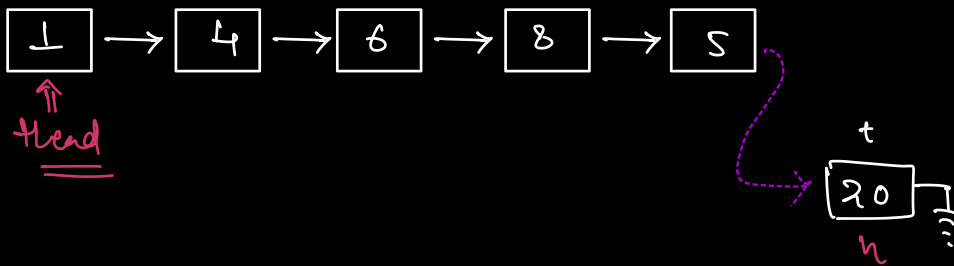
```

Node n = new Node(10)
n.next = head
head = n;

```

⇒ O(1)

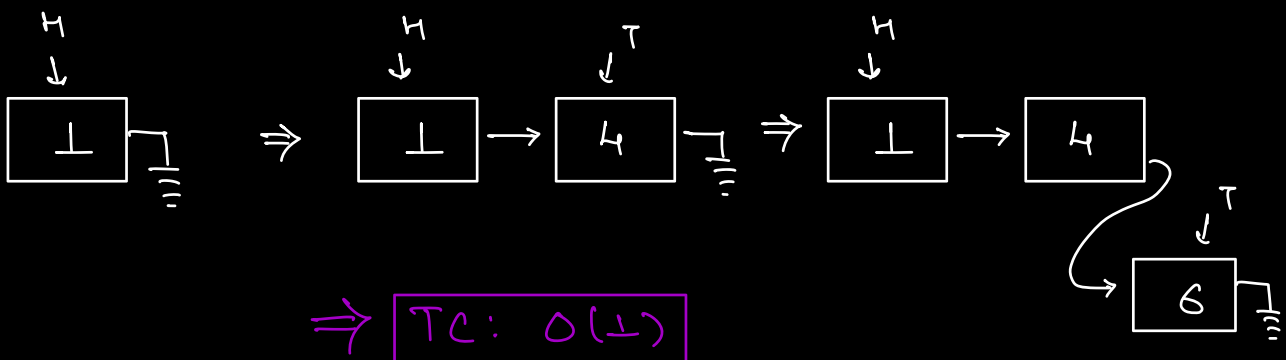
② Insert at End



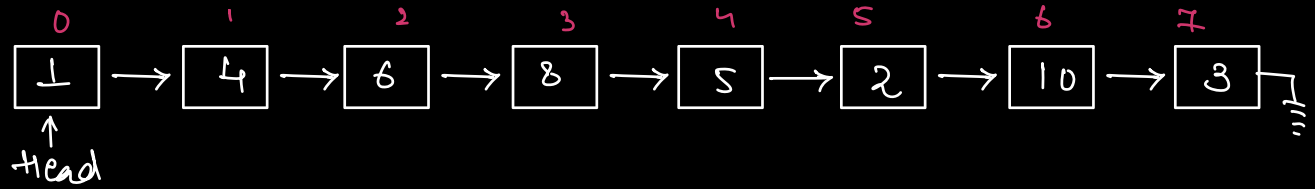
```

Node n = new Node(20);
tail = head;
while (tail != Null) {
    tail = tail.next;
}
tail.next = n;
tail = n
  
```

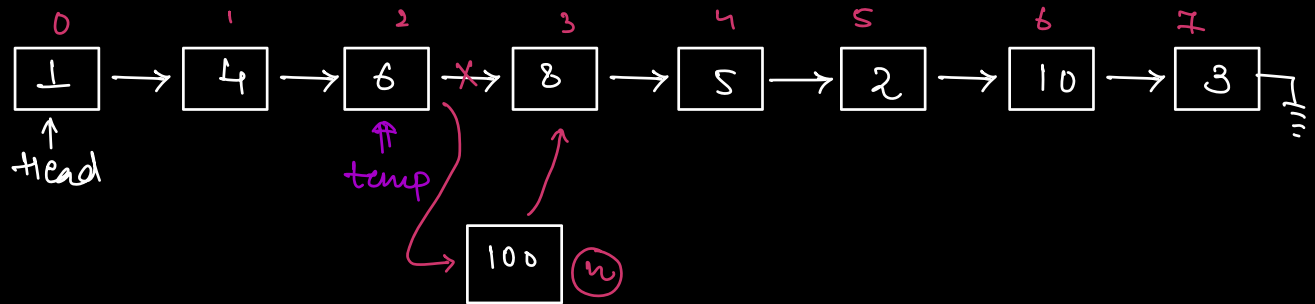
TC: $O(N)$
 \hookrightarrow TC $\Rightarrow O(1)$ { by maintaining the tail reference }



③ Insert at k^{th} position

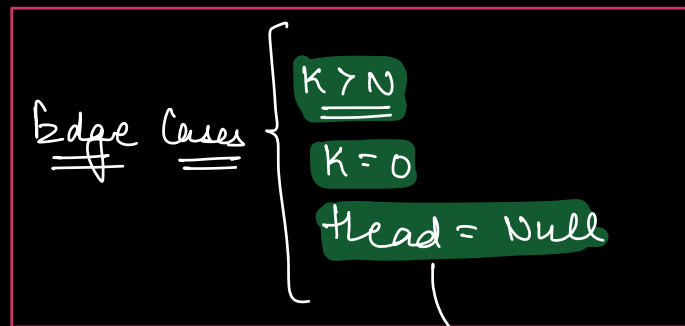


$k = 3$ \Rightarrow data = 100



$n \cdot \text{next} = \text{temp} \cdot \text{next};$

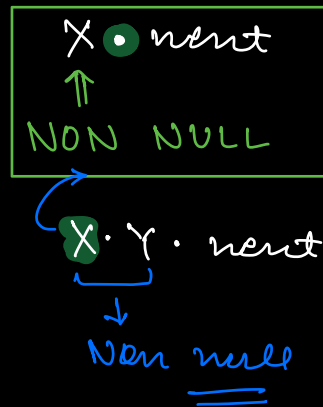
$\text{temp} \cdot \text{next} = n$



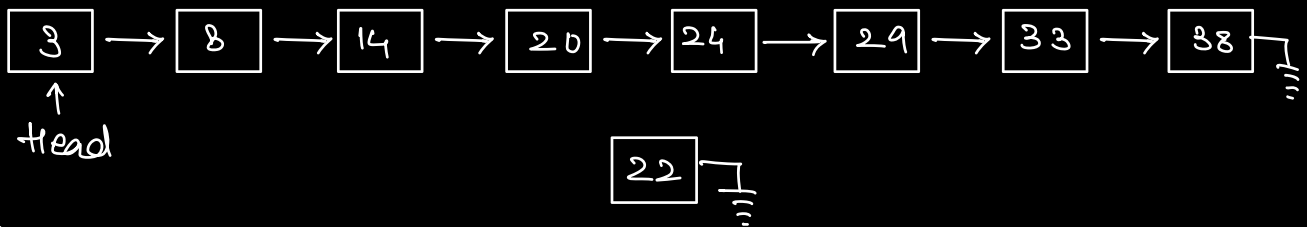
Empty linked list

Edge Cases :

- ① head = null;
- ② N = 1/2/3
- ③ Problem specific



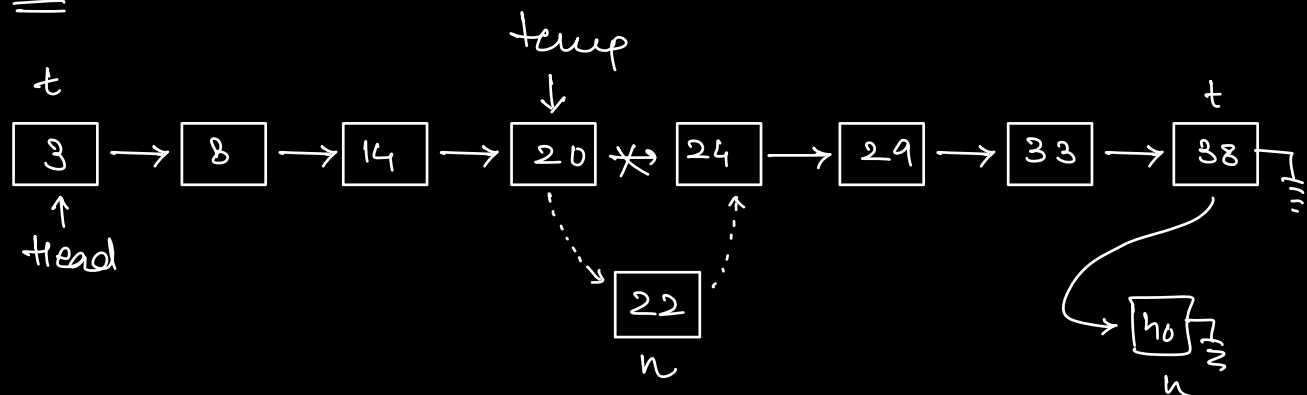
Q. Given a L.L sorted in ascending order
Insert a value at its correct position
in sorted order.



k = 22

⇒ Find the last node with data < k.

k = 22



head of new
LL

Node

insertInSortedOrder(head, k) {

Node n = new Node(k);

if (head == null) { // LL is Empty

return n;

3

if (k <= head.data) {

// insert at front;

n.next = head;

return n;

3

temp = head;

while (temp.next != null &&

temp.next.data < k) {

temp = temp.next;

3

n.next = temp.next;

temp.next = n;

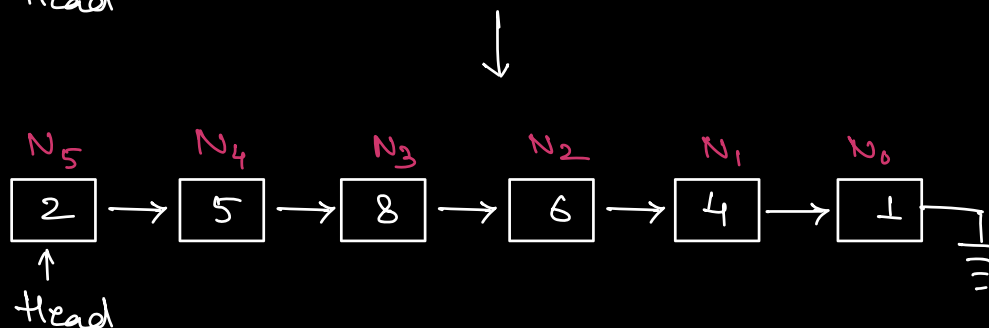
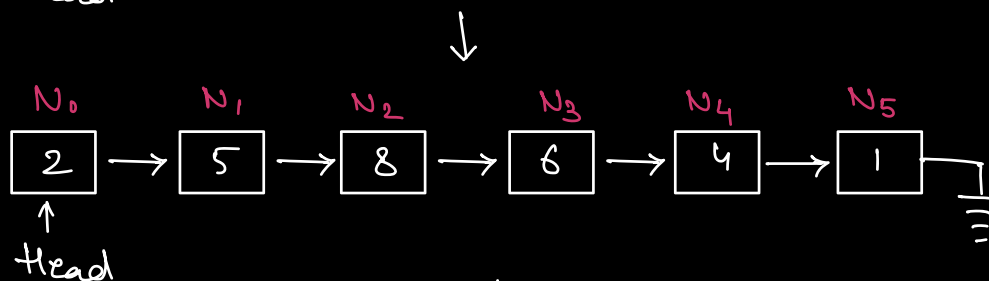
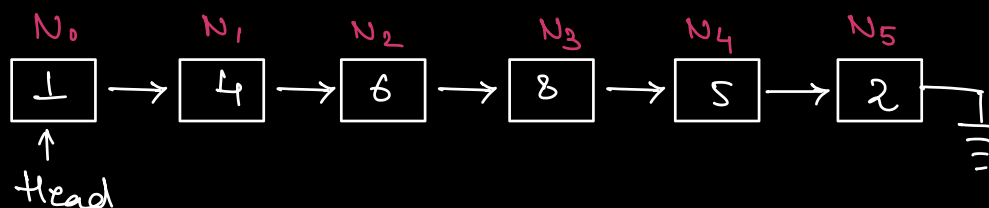
return head;

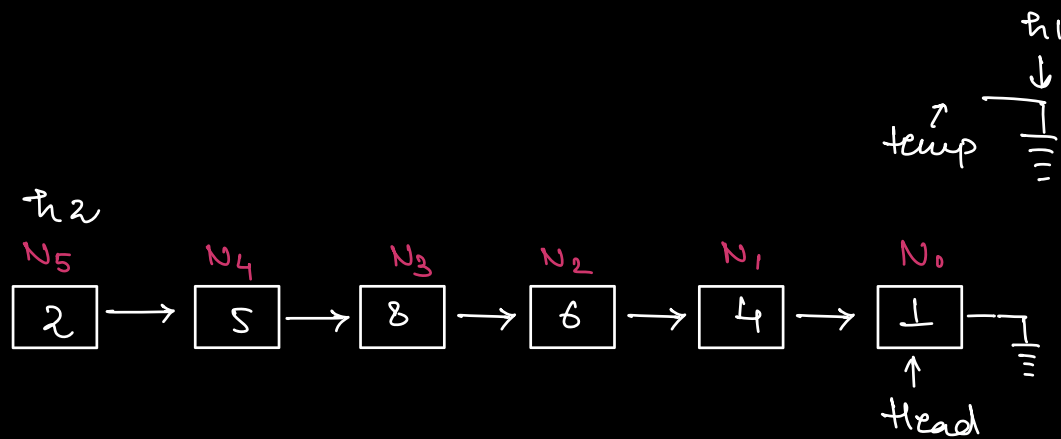
3

TC: O(N) / SC: O(1)

Q. Reverse the given L.L
(Expected SC: O(1))

* Changing the value of nodes isn't allowed.





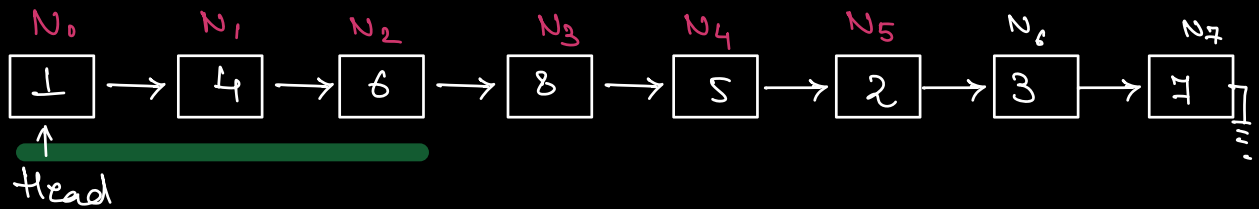
```

Node reverse (head) {
    if (head == Null) return null;
    h1 = head, t = h1, h2 = Null;
    while (h1 != Null) {
        temp = h1;
        h1 = h1.next;
        t.next = h2;
        h2 = temp;
    }
    return h2;
}

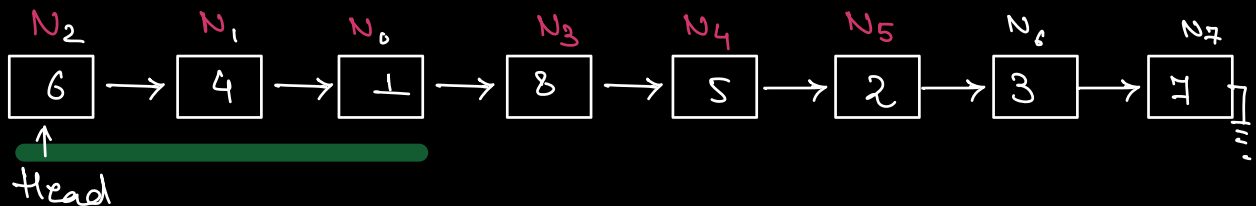
```

TC: $O(N)$
 SC: $O(1)$

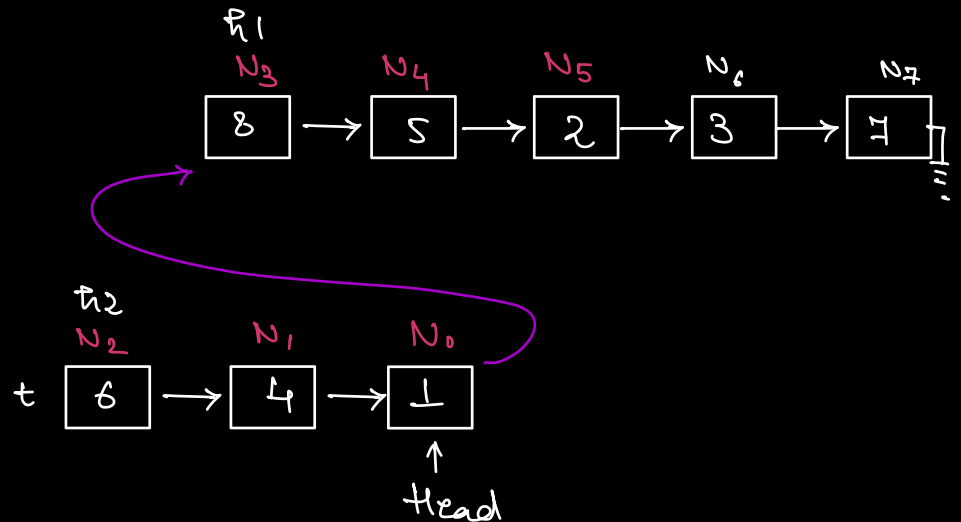
Q. Given a L.L, Reverse the first k nodes of the L.L.



k=3



k=3 2 0



head.next = t1;

```

Node reverseFirstKNodes(head, k) {
    if (k == 0 || head == null) return head;
    h1 = head, t = h1, h2 = null;
    while (k > 0 && h1 != null) {
        temp = h1;
        h1 = h1.next;
        t.next = h2;
        h2 = temp;
        k--;
    }
    head.next = h1;
    return h2;
}

```

3

TC: $O(N)$ / SC: $O(1)$

Q. Reverse L.L in group of k nodes.

Google. Reverse k Groups.

