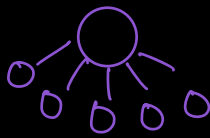


Trees:-

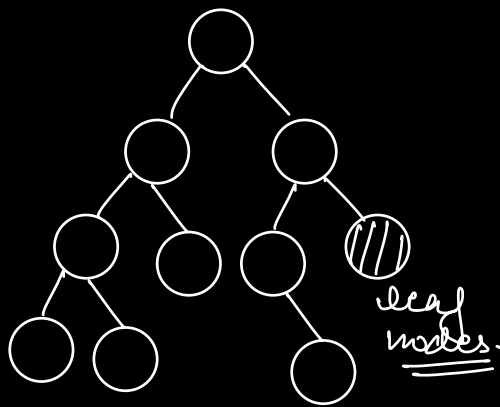
↳ Hierarchical data.

- Directory structure
- XML/HTML
- Organisation structure
- Trie
- B/B+ Trees (DB indexes)
- Segment Trees
- RB/AVL Trees (Self Balancing BST)



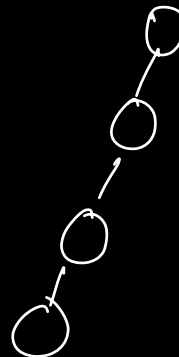
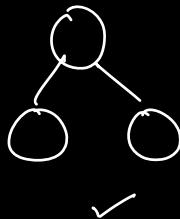
• fb/Google/LinkedIn ...

* Important for Interviews.



• Binary Tree

Tree in which every node can have at max 2 children

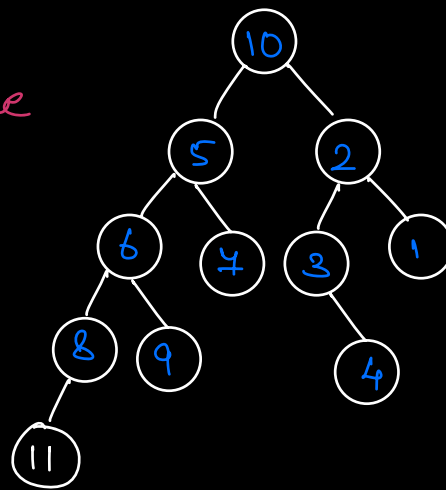


Skewed Tree.

Height of a Node in Binary Tree

Ht of a node:-

→ Distance of the node from the farthest reachable leaf node.
(# of Edges)



$$\text{Ht}(5) = 3$$

$$\text{Ht}(3) = 1$$

$$\text{Ht}(6) = 2$$

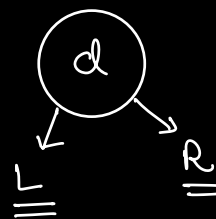
$$\text{Ht}(10) = 4$$

$$\text{Ht}(11) = \text{Ht}(9) = \text{Ht}(4) = \underline{\underline{0}}$$

$$\text{Ht}(\text{NULL}) = -1$$

$$\textcircled{50} \Rightarrow \text{Height} = \underline{\underline{0}}$$

```
class TreeNode {  
    int data;  
    TreeNode left;  
    TreeNode right;  
    TreeNode(int x) {  
        data = x;  
        left = NULL;  
        right = NULL;  
    }  
}
```



3

3

Depth of a node in a Binary Tree.

Distance of node from root node.

$$Dp(6) = 2$$

$$Dp(4) = 3$$

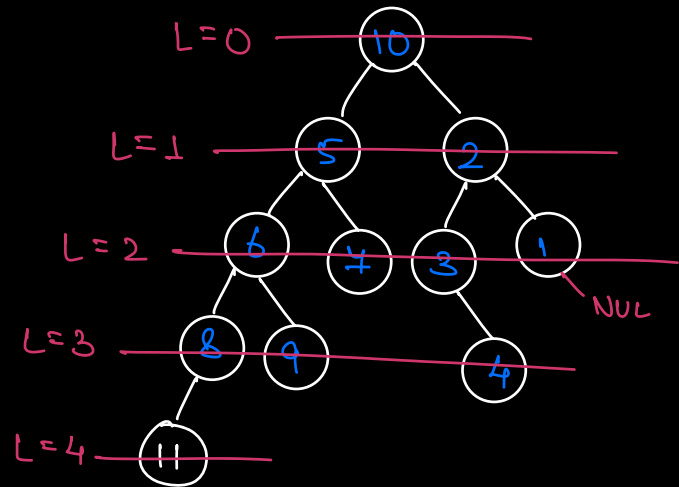
$$Dp(11) = 4$$

$$Dp(3) = 2$$

$$Dp(9) = 3$$

$$Dp(1) = 2$$

$$Dp(10) = \underline{\underline{0}}$$

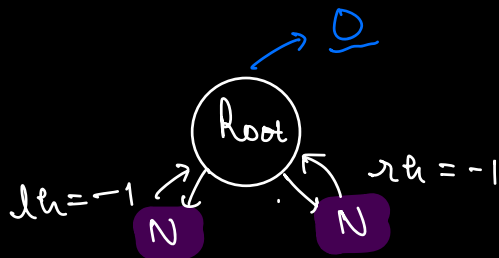


Q. Given a Binary Tree, find its height.

```

int height (root) {
  if (root == NULL)
    return -1;

  int lh = height (root.left); LST
  int rh = height (root.right); RST
  return max (lh, rh) + 1; Root
}
  
```



Post Order traversal

→ Traverse LST before RST

* Root LST RST → PreOrder
* LST Root RST → InOrder
* LST RST Root → PostOrder.

```
Void preOrder (root) {  
    if (root == NULL) return;  
    print (root.data);  
    preOrder (root.left);  
    preOrder (root.right);  
}
```

```
Void inOrder (root) {  
    if (root == NULL) return;  
    inOrder (root.left);  
    print (root.data);  
    inOrder (root.right);  
}
```

```
Void postOrder (root) {  
    if (root == NULL) return;  
    postOrder (root.left);  
    postOrder (root.right);  
    print (root.data);  
}
```

* preOrder

root

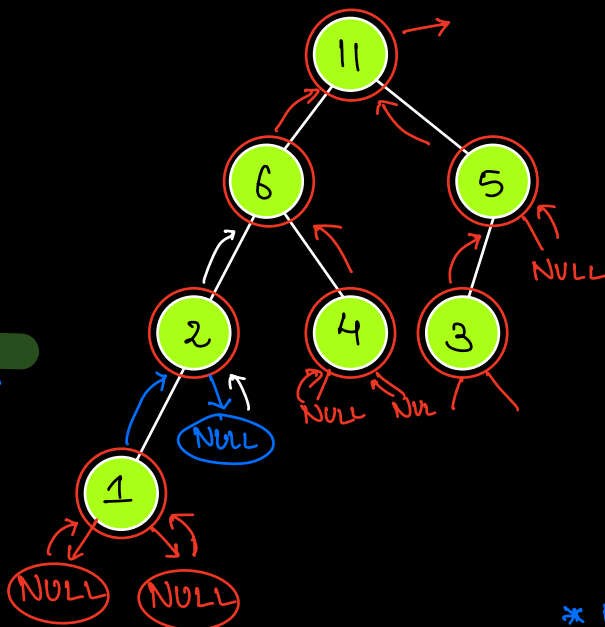
LST

Root

11, 6, 2, 1, 4, 5, 3

↑
Root

└──┬──┬──┬──┬──┬──┬──
LST RST



Pre(11)
Pre(6)
Pre(2)
Pre(1)
Pre(4)
Pre(5)
Pre(3)

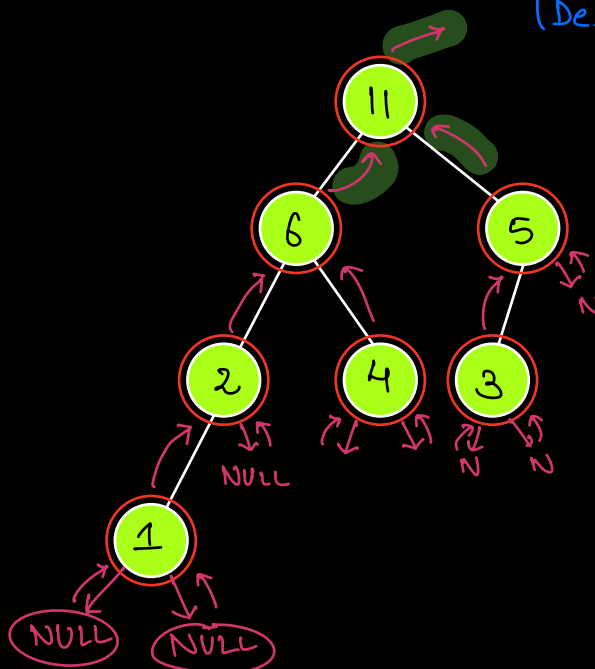
* Post Order

LST

RST

root

1, 2, 4, 6, 3, 5, 11



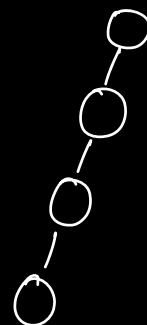
* DFS
(Depth First Search)

TTC: $O(N)$

SC: $O(\text{height})$

↳ $O(N)$

WCS



Skewed

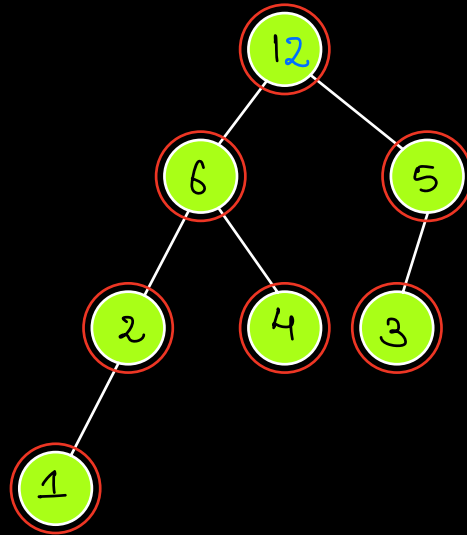
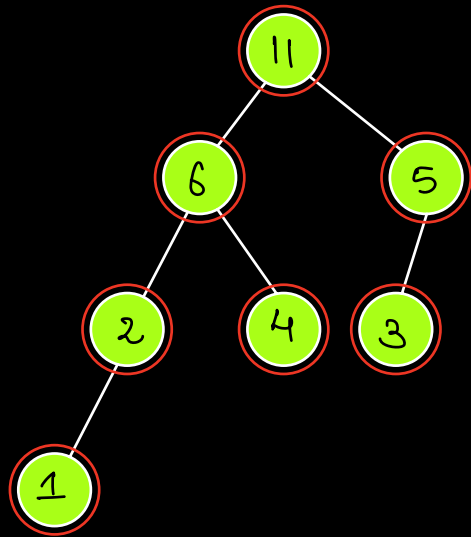
H = N

→ PreOrder :

Root
LST
RST

* Search in a B.T

* Check if 2 trees are identical



→ Fail Fast Approach

* Post Order

└─ Height

└─ No. of nodes in B.T

Q. Given a B.T, search a value k in it

```
bool search (root, k) {  
    if (root == NULL) return false;  
    if (root->data == k)  
        return true;  
    return search (root->left, k)  
        ||  
        search (root->right, k)  
}
```

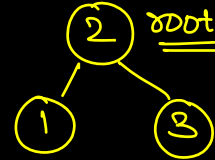
3

A	or	B
↓		↓
T		X

Q.
Amazon
MS

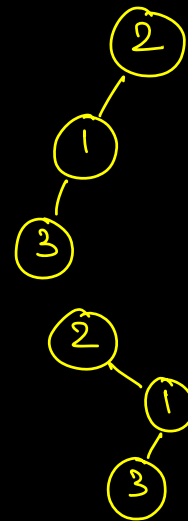
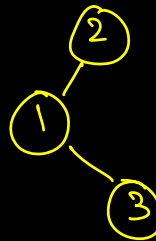
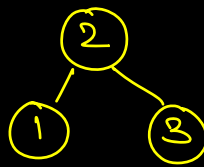
Given the preorder & inorder traversal of a Binary Tree in 2 Arrays. (No duplicates)
Construct the Tree

Pre-Order : 2 1 3
In-Order : 1 2 3

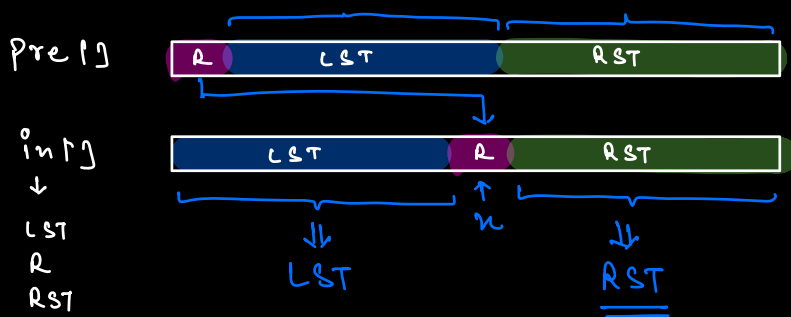


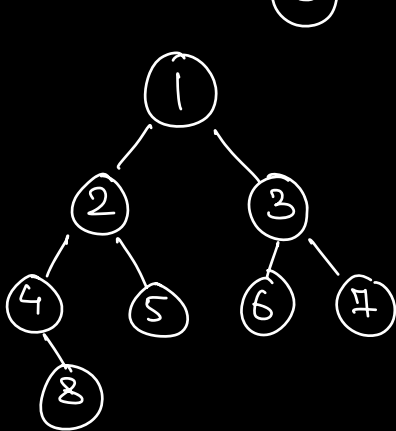
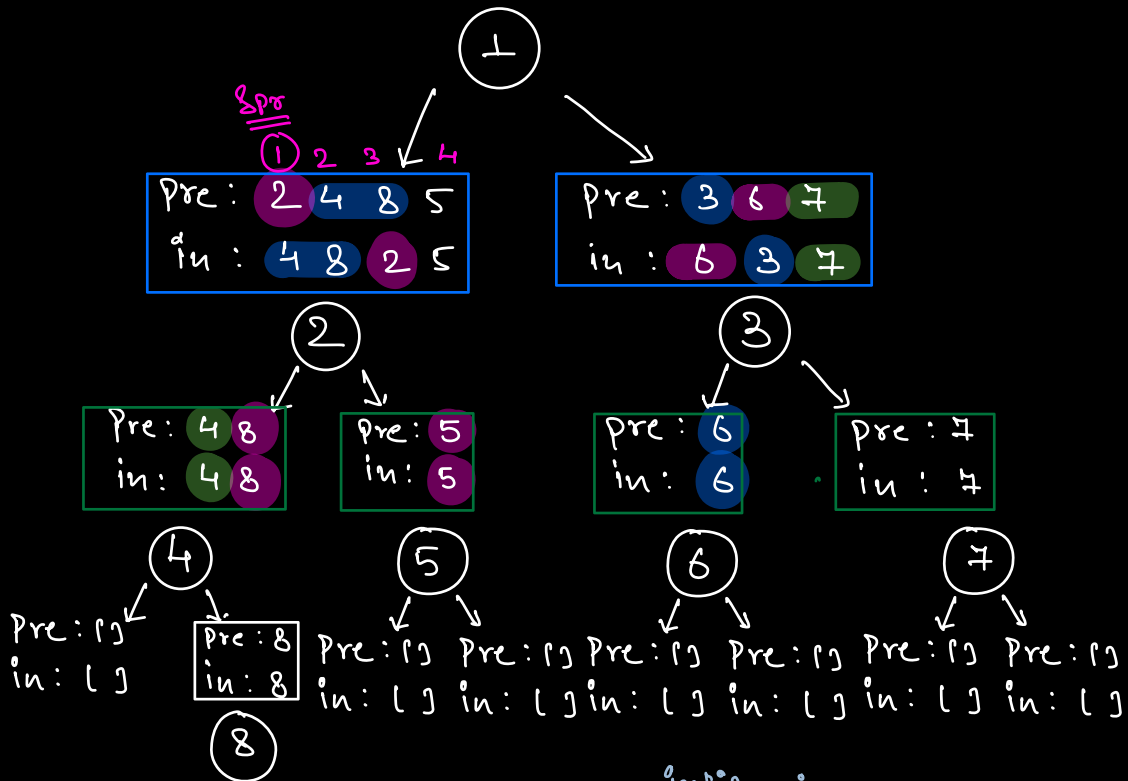
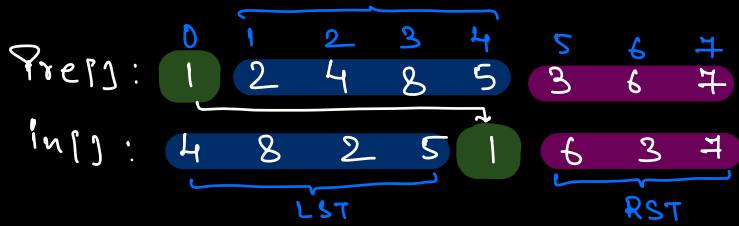
Pre-Order : 2 1 3

→ If only one tree traversal is given then we can't construct a unique tree.



→ First element of pre[] will be root
pre[0] → root.





* HashMap <int, int>

4: 0
 8: 1
 2: 2
 5: 3
 1: 4
 6: 5
 3: 6
 7: 7

```

TreeNode Construct (pre[], in[], sin, ein, spr, epr) {
    // Assumption: Construct(pre, in, sin, ein, spr, epr)
    // returns the root of the tree with inorder
    // from sin to ein in in[] & preorder from spr, epr in pre[].
    if (spr > epr) return NULL;

```

```

    TreeNode root = new TreeNode(pre[spr]);
    // find the index of root in in[].

```

```

    int idn = map.get(pre[spr]);

```

```

    int a = idn - sin; // # of elements in in[] / pre[]
                       // of LST.

```

```

    root.left = Construct(pre, in, sin, idn-1, spr+1, spr+a);

```

```

    root.right = Construct(pre, in, idn+1, ein, spr+a+1, epr);

```

```

    return root;

```

3

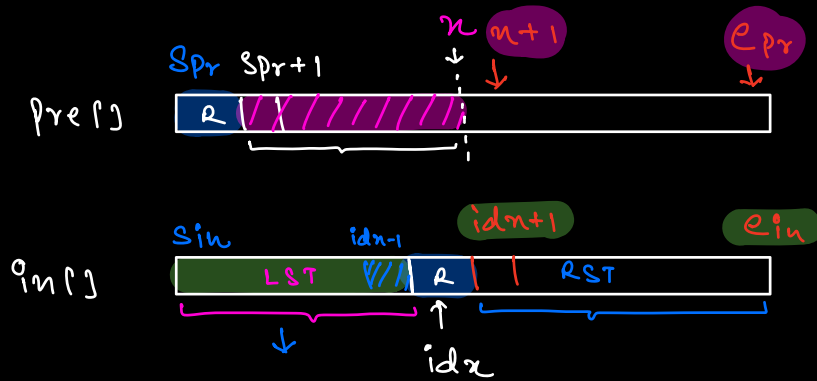
TC: $O(N)$

SC: $O(N)$

↳ Recursion
+
Hash Map.

Without HM

TC: $O(N^2)$



$$[sin, idn-1] \rightarrow idn - sin = a$$

$$[Spr+1, n] \rightarrow idn - sin$$

↓

$$n - (Spr+1) + 1 = idn - sin$$

$$n - Spr - 1 + 1 = idn - sin$$

$$n = Spr + idn - sin$$

$$= \underline{\underline{Spr + a}}$$

* PreOrder[] > Unique Tree ?
 * PostOrder[]

R LST RST

LST RST Root

———— * ————