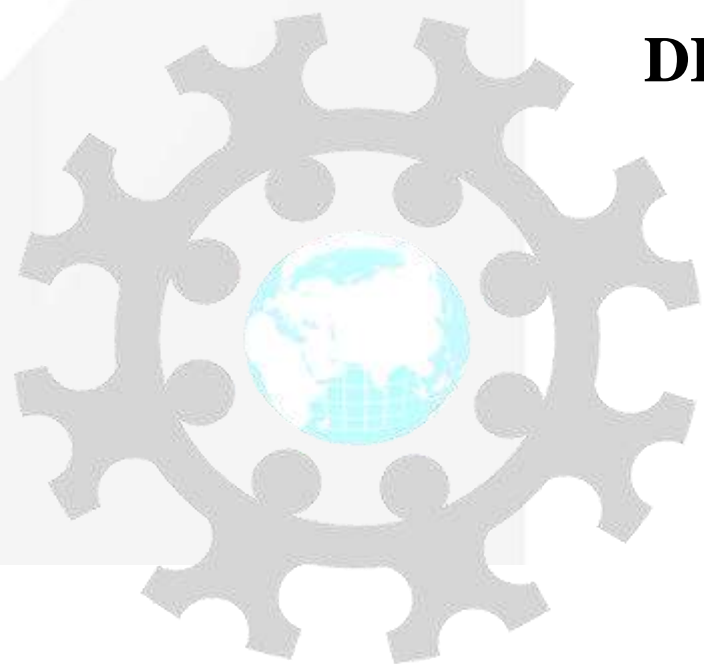


UNIVERSITY INSTITUTE OF COMPUTING

MASTER OF COMPUTER APPLICATIONS

DESIGN AND ANALYSIS OF ALGORITHMS

24CAT611



UNIT-2

DISCOVER . **LEARN** . EMPOWER

Outline

- **Divide and Conquer: General method**
- Binary Search
- Advantages and disadvantages of divide and conquer



Divide & Conquer

Divide and Conquer is an algorithm design paradigm that involves breaking up a larger problem into *non-overlapping* sub-problems, solving each of these sub-problems, and combining the results to solve the original problems. A problem has non-overlapping sub-problems if you can find its solution by solving each sub-problem once.



Divide & Conquer (contd.)

The three main steps in the divide and conquer paradigm are:

- **divide**: involves breaking the problem into smaller, non-overlapping chunks.
- **conquer**: involves solving the sub-problems recursively.
- **combine**: involves combining the solutions of the smaller sub-problems to solve the original problem.

Divide-and-Conquer Technique (cont.)

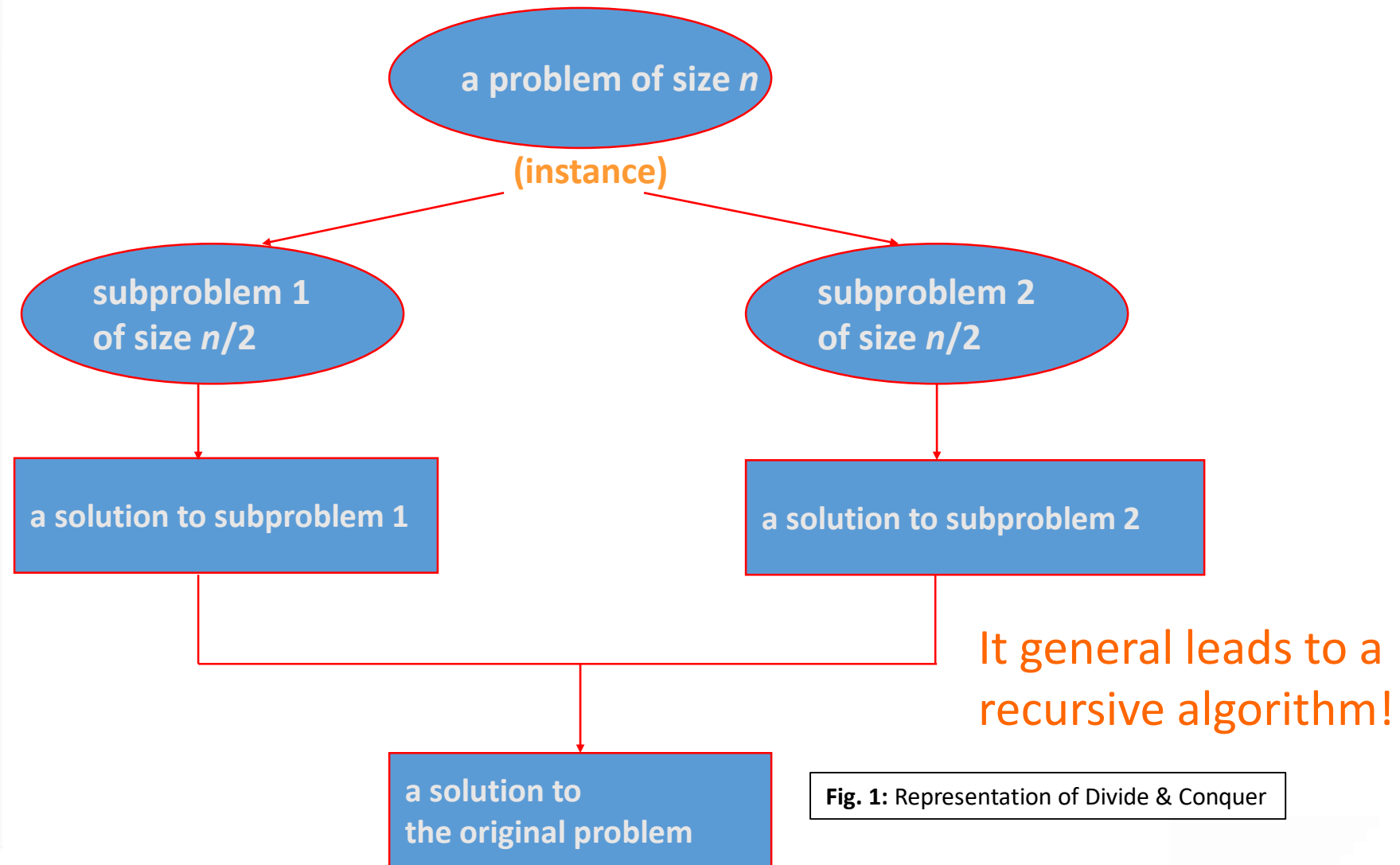


Fig. 1: Representation of Divide & Conquer

Divide-and-Conquer Examples

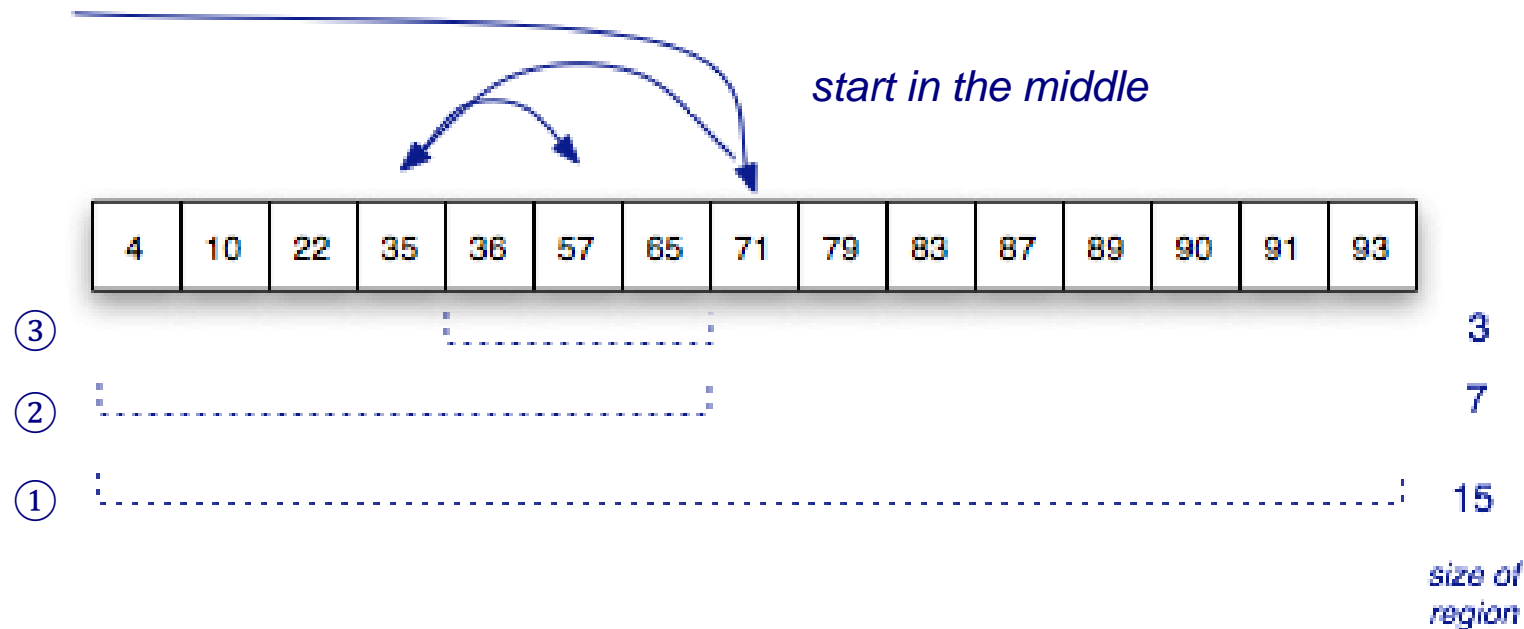
- Sorting: mergesort and quicksort
- Binary tree traversals
- Binary search (?)

Binary Search

- The binary search algorithm uses the divide-and-conquer strategy to search through an array
- The array *must be sorted*
 - the “zeroing in” strategy for looking up a word in the dictionary won’t work if the words are not in alphabetical order
 - binary search will not work unless the array is sorted

Binary Search (contd.)

- To search a list of n items, first look at the item in location $n/2$
 - then search either the region from 0 to $n/2-1$ or the region from $n/2+1$ to $n-1$
- **Example:** searching for 57 in a sorted list of 15 numbers



Binary Search (contd.)

- The algorithm uses two variables to keep track of the boundaries of the region to search

lower

the index *one below* the leftmost item in the region

upper

the index *one above* the rightmost region

4	10	22	35	36	57	65	71	79	83	87	89	90	91	93
---	----	----	----	----	----	----	----	----	----	----	----	----	----	----

initial values when searching an array of n items:

lower = -1

upper = n

Binary Search (contd.)

- The algorithm is based on an iteration (“loop”) that keeps making the region smaller and smaller
 - the initial region is the complete array
 - the next one is either the upper half or lower half
 - the one after that is one quarter, then one eighth, then...

4	10	22	35	36	57	65	71	79	83	87	89	90	91	93
---	----	----	----	----	----	----	----	----	----	----	----	----	----	----

initial values when searching an array of n items:

lower = -1

upper = n

Binary Search (contd.)

- The heart of the algorithm contains these operations:

$\text{mid} = (\text{lower} + \text{upper}) / 2$

return mid if $k == a[\text{mid}]$

upper = mid if $k < a[\text{mid}]$

lower = mid if $k > a[\text{mid}]$

- The first iteration when searching for 57 in a list of size 15:

4	10	22	35	36	57	65	71	79	83	87	89	90	91	93
---	----	----	----	----	----	----	----	----	----	----	----	----	----	----

*

lower = -1

mid = $14 / 2 = 7$

upper for next

upper = 15

iteration: 7

Binary Search (contd.)

- The remaining iterations when searching for 57:

lower = -1
upper = 7
mid = 3
lower = 3

lower = 3
upper = 7
mid = 5
found it!

```
mid = (lower + upper) / 2
return mid if k == a[mid]
upper = mid if k < a[mid]
lower = mid if k > a[mid]
```

[*]

4	10	22	35	36	57	65	71	79	83	87	89	90	91	93
---	----	----	----	----	----	----	----	----	----	----	----	----	----	----

[*]

4	10	22	35	36	57	65	71	79	83	87	89	90	91	93
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This search required only 3 comparisons:

a[7], a[3], a[5]

Binary Search (contd.)

- The number of iterations made by this algorithm when it searches an array of n items is roughly
- To see why, consider the question from the other direction
 - suppose we have an array that starts out with 1 item
 - suppose each step of an iteration doubles the size of the array
 - after n steps we will have 2^n items in the array

The complexity of Binary search algorithm is $O(\log n)$



Divide-and-Conquer - Advantages

- Complexity can be reduced using the concepts of divide and conquer.
- Increase in productivity by allowing multiple programmers to work on different parts of the project independently at the same time.
- Modules can be re-used many times, thus it saves time, reduces complexity and increase reliability.
- Easier to update/fix the program by replacing individual modules rather than larger amount of code.
- Ability to either eliminate or at least reduce the necessity of employing GOTO statement
- Solves difficult problems with less time complexity than its brute-force counterpart.
- Since the sub-problems are independent, they can be computed in parallel

Divide and Conquer - Disadvantages

- Problem decomposition may be very complex and therefore not actually suitable to divide and conquer.
- Recursive nature of the solution may end up duplicating sub-problems, dynamic solutions may be better in some of these cases, like Fibonacci.
- Recursion into small/tiny base cases may lead to huge recursive stacks, and efficiency can be lost by not applying solutions earlier for larger base cases.

Complexity

- The complexity of the divide and conquer algorithm is calculated using the master theorem.

$$T(n) = aT(n/b) + f(n),$$

where,

n = size of input

a = number of subproblems in the recursion

n/b = size of each subproblem. All subproblems are assumed to have the same size.

$f(n)$ = cost of the work done outside the recursive call, which includes the cost of dividing the problem and cost of merging the solutions

Frequently Asked Questions

- What do you understand about binary search?
- Define complexity of binary search.

References

- 1) https://www.tutorialspoint.com/data_structures_algorithms/divide_and_conquer.htm
- 2) **Data Structures and Algorithms made easy** By Narasimha Karumanchi.
- 3) The Algorithm Design Manual, 2nd Edition by Steven S Skiena
- 4) **Fundamentals of Computer Algorithms - Horowitz and Sahani**





THANK YOU