

## Worksheet No. - 3

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**Branch:** MCA

**Semester:** II

**Subject Name:** DAA LAB

**UID:** 24MCA20454

**Section/Group:** 6 (B)

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**Subject Code:** 24CAP612

**Aim/Overview of the practical:** Find Minimum Cost Spanning Tree of a given undirected graph using Kruskal's algorithm.

**Objective:** To find the Minimum Cost Spanning Tree (MCST) of a given undirected graph using Kruskal's algorithm. The algorithm follows a greedy approach, sorting edges by weight and using the find and union method to avoid cycles. The goal is to understand MST construction and analyze its efficiency.

**Input/Apparatus Used:** IntelliJ Idea as code editor.

### **Algorithms Steps:**

- 1. Sort:** all edges in non-decreasing order based on their weight.
- 2. Initialize:** a parent array where each vertex is its own parent.
- 3. Iterate:** through the sorted edges and check if they form a cycle using the Union-Find data structure.
- 4. Include:** an edge in the MST if it does not create a cycle.
- 5. Repeat:** until we have ' $V-1$ ' edges in the MST (where ' $V$ ' is the number of vertices).

### **Procedure/Algorithm/Code:**

```
package graph;

import java.util.*;

class Edge implements Comparable<Edge> {
    int src, dest, weight;

    public Edge(int src, int dest, int weight) {
```

```
this.src = src;
this.dest = dest;
this.weight = weight;
}

@Override
public int compareTo(Edge other) {
    return this.weight - other.weight;
}
}

class DisjointSet {
    int[] parent, rank;

    public DisjointSet(int n) {
        parent = new int[n];
        rank = new int[n];
        for (int i = 0; i < n; i++) parent[i] = i;
    }

    public int find(int node) {
        if (parent[node] != node) {
            parent[node] = find(parent[node]); // Path compression
        }
        return parent[node];
    }

    public void union(int u, int v) {
        int rootU = find(u);
        int rootV = find(v);
        if (rootU != rootV) {
            if (rank[rootU] > rank[rootV]) {
                parent[rootV] = rootU;
            } else if (rank[rootU] < rank[rootV]) {
                parent[rootU] = rootV;
            } else {
                parent[rootV] = rootU;
                rank[rootU]++;
            }
        }
    }
}
```

```
}  
}
```

```
public class KruskalMST {
```

```
    public static List<Edge> kruskalMST(List<Edge> edges, int V) {  
        List<Edge> mst = new ArrayList<>();  
        Collections.sort(edges); // Sort edges by weight  
        DisjointSet ds = new DisjointSet(V);  
  
        for (Edge edge : edges) {  
            if (ds.find(edge.src) != ds.find(edge.dest)) { // Check if cycle forms  
                mst.add(edge);  
                ds.union(edge.src, edge.dest);  
                if (mst.size() == V - 1) break; // Stop when MST is complete  
            }  
        }  
        return mst;  
    }  
}
```

```
    public static void main(String[] args) {  
        int V = 4; // Number of vertices  
        List<Edge> edges = new ArrayList<>();  
  
        // Adding edges (source, destination, weight)  
        edges.add(new Edge(0, 1, 10));  
        edges.add(new Edge(0, 2, 6));  
        edges.add(new Edge(0, 3, 5));  
        edges.add(new Edge(1, 3, 15));  
        edges.add(new Edge(2, 3, 4));  
  
        List<Edge> mst = kruskalMST(edges, V);  
  
        System.out.println("Graph Edges ");  
        for (Edge edge : edges) {  
            System.out.println(edge.src + " - " + edge.dest + " : " + edge.weight);  
        }  
  
        System.out.println("Minimum Spanning Tree:");
```

```

    for (Edge edge : mst) {
        System.out.println(edge.src + " - " + edge.dest + " : " + edge.weight);
    }
}
}

```

## Output:

```

C:\Users\HELL0\.jdk\openjdk-22.0.2\bin\java.exe "
Graph Edges
2 - 3 : 4
0 - 3 : 5
0 - 2 : 6
0 - 1 : 10
1 - 3 : 15
Minimum Spanning Tree:
2 - 3 : 4
0 - 3 : 5
0 - 1 : 10

Process finished with exit code 0

```

## Learning outcomes (What I have learnt):

1. Sorting helps process edges in optimal order.
2. find() and union() prevent cycles efficiently.
3. Not all edges need to be processed; MST stops early.
4. The final MST is unique for a given set of weights.

## Evaluation Grid:

Sr. No.	Parameters	Marks Obtained	Maximum Marks
1.	Demonstration and Performance		12
2.	Worksheet		8
3.	Viva		10