Divide-and-Conquer: Master Theorem

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Algorithmic Design and Techniques
Algorithms and Data Structures

Outline

$$T(n) = T\left(\frac{n}{2}\right) + O(1)$$

$$T(n) = T\left(\frac{n}{2}\right) + O(1)$$

$$\Big\downarrow$$

$$T(n) = O(\log n)$$

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n)$$

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n)$$

$$\Big\downarrow$$

$$T(n) = O(n^2)$$

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$

$$\downarrow$$

$$T(n) = O(n^{\log_2 3})$$

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

$$\downarrow$$

$$T(n) = O(n \log n)$$

Master Theorem

Theorem

If $T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$

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If $T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil \right) + O(n^d)$ (for constants $a > 0, b > 1, d \geq 0$), then:

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$$T(n) = \begin{cases} \\ \\ \end{cases}$$

Master Theorem

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$$T(n) = \begin{cases} O(n^d) & \text{if } d > \log_b a \\ & \\ & \end{cases}$$

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$$T(n) = \begin{cases} O(n^d) & \text{if } d > \log_b a \\ O(n^d \log n) & \text{if } d = \log_b a \\ & \end{cases}$$

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If $T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$ (for constants $a > 0, b > 1, d \geq 0$), then:

$$T(n) = \begin{cases} O(n^d) & \text{if } d > \log_b a \\ O(n^d \log n) & \text{if } d = \log_b a \\ O(n^{\log_b a}) & \text{if } d < \log_b a \end{cases}$$

Master Theorem Example 1

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n)$$

Master Theorem Example 1

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n)$$

$$a = 4$$

Master Theorem Example 1

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n)$$

$$a = 4$$

$$b = 2$$

Master Theorem Example 1

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n^{1})$$

$$a = 4$$

$$b = 2$$

$$d = 1$$

Master Theorem Example 1

$$
\begin{aligned}
T(n) &= 4T\left(\frac{n}{2}\right) + O(n) \\
a &= 4 \\
b &= 2 \\
d &= 1
\end{aligned}
$$

Since $d < \log_b a$, $T(n) = O(n^{\log_b a}) = O(n^2)$

Master Theorem Example 2

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$

Master Theorem Example 2

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$

$$a = 3$$

Master Theorem Example 2

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n)$$
$$a = 3$$
$$b = 2$$

Master Theorem Example 2

$$T(n) = 3T\left(\frac{n}{2}\right) + O(n^{1})$$

$$a = 3$$

$$b = 2$$

$$d = 1$$

Master Theorem Example 2

$$
\begin{aligned}
T(n) &= 3T\left(\frac{n}{2}\right) + O(n) \\
a &= 3 \\
b &= 2 \\
d &= 1
\end{aligned}
$$

Since $d < \log_b a$,
$T(n) = O(n^{\log_b a}) = O(n^{\log_2 3})$

Master Theorem Example 3

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

Master Theorem Example 3

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

$$a = 2$$

Master Theorem Example 3

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n)$$

$$a = 2$$

$$b = 2$$

Master Theorem Example 3

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^1)$$

$$a = 2$$

$$b = 2$$

$$d = 1$$

Master Theorem Example 3

$$
\begin{aligned}
T(n) &= 2T\left(\frac{n}{2}\right) + O(n) \\
a &= 2 \\
b &= 2 \\
d &= 1
\end{aligned}
$$

Since $d = \log_b a$,
$T(n) = O(n^d \log n) = O(n \log n)$

Master Theorem Example 4

$$T(n) = T\left(\frac{n}{2}\right) + O(1)$$

Master Theorem Example 4

$$T(n) = 1\,T\left(\frac{n}{2}\right) + O(1)$$

$$a = 1$$

Master Theorem Example 4

$$T(n) = T\left(\frac{n}{2}\right) + O(1)$$

$$a = 1$$

$$b = 2$$

Master Theorem Example 4

$$T(n) = T\left(\frac{n}{2}\right) + O(n^{0})$$

$$a = 1$$

$$b = 2$$

$$d = 0$$

Master Theorem Example 4

$$T(n) = T\left(\frac{n}{2}\right) + O(1)$$
$$a = 1$$
$$b = 2$$
$$d = 0$$

Since $d = \log_b a$, $T(n) = O(n^d \log n) = O(n^0 \log n) = O(\log n)$

Master Theorem Example 5

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^2)$$

Master Theorem Example 5

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^2)$$

$$a = 2$$

Master Theorem Example 5

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^2)$$

$$a = 2$$

$$b = 2$$

Master Theorem Example 5

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^2)$$

$$a = 2$$

$$b = 2$$

$$d = 2$$

Master Theorem Example 5

$$T(n) = 2T\left(\frac{n}{2}\right) + O(n^2)$$
$$a = 2$$
$$b = 2$$
$$d = 2$$

Since $d > \log_b a$, $T(n) = O(n^d) = O(n^2)$

Outline

Master Theorem

Theorem

If $T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$ (for constants $a > 0, b > 1, d \geq 0$), then:

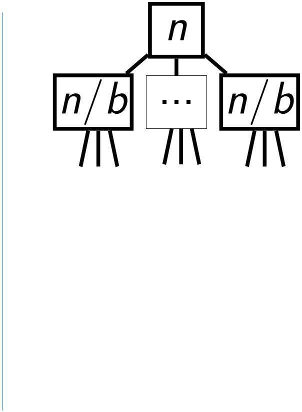
$$T(n) = \begin{cases} O(n^d) & \text{if } d > \log_b a \\ O(n^d \log n) & \text{if } d = \log_b a \\ O(n^{\log_b a}) & \text{if } d < \log_b a \end{cases}$$

$$T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$$

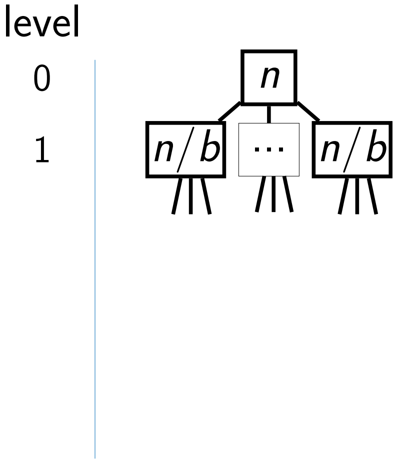
$$\boxed{n}$$

$$T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$$

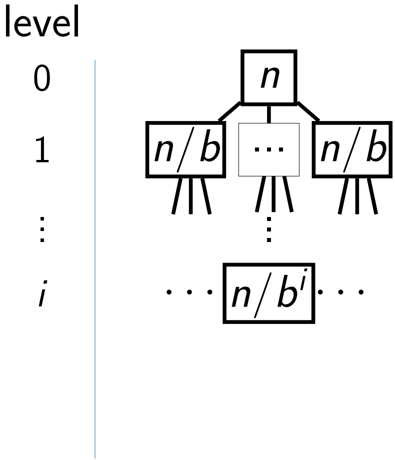
level



$$T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$$



$$T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil \right) + O(n^d)$$



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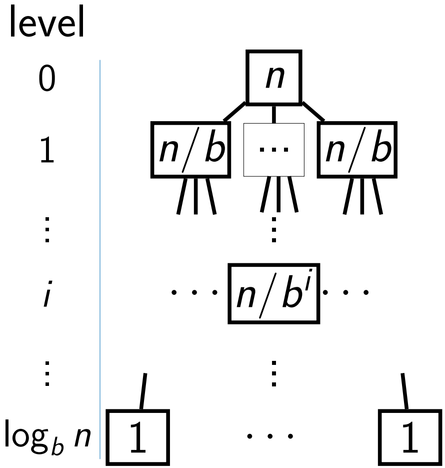

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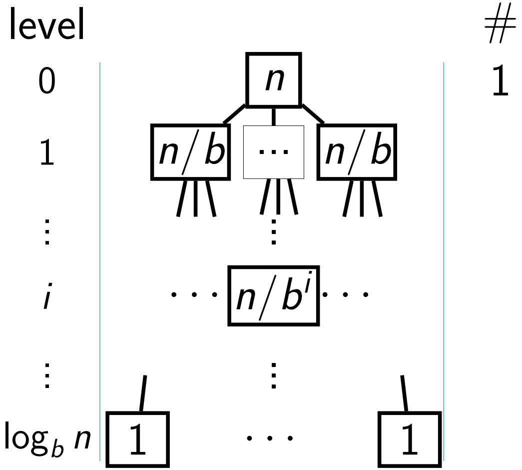

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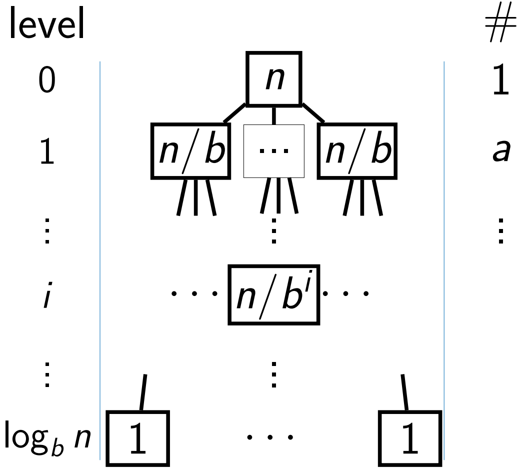

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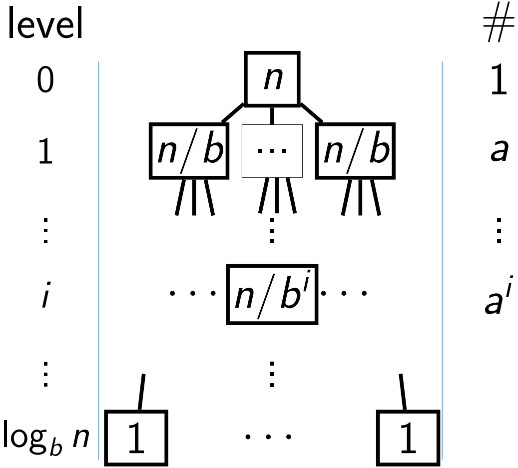

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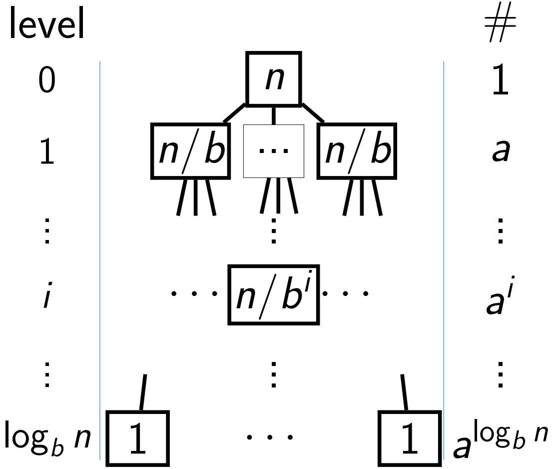

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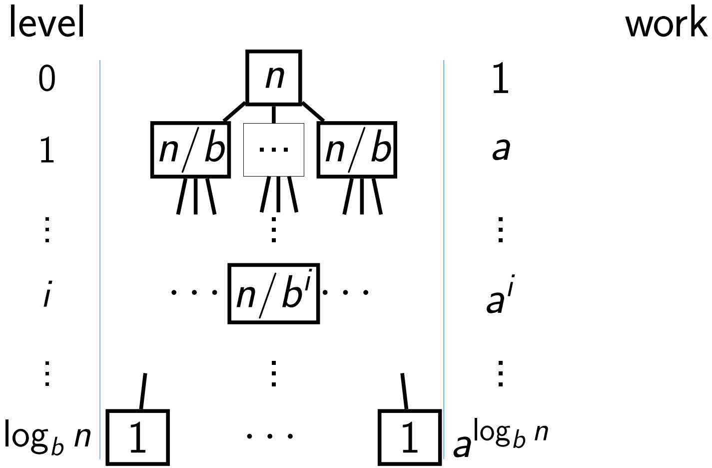

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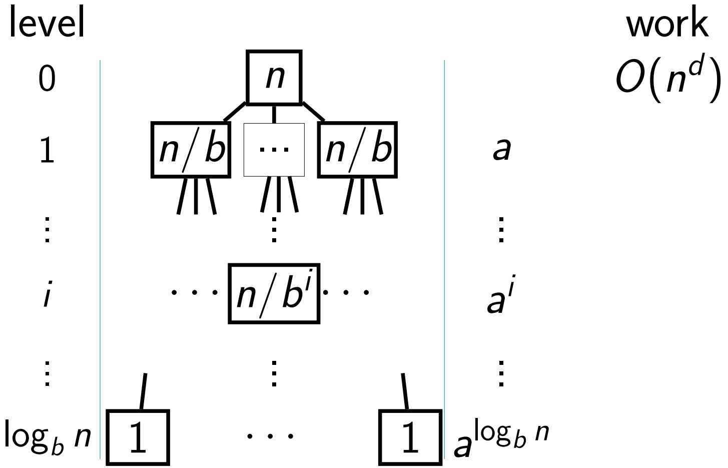

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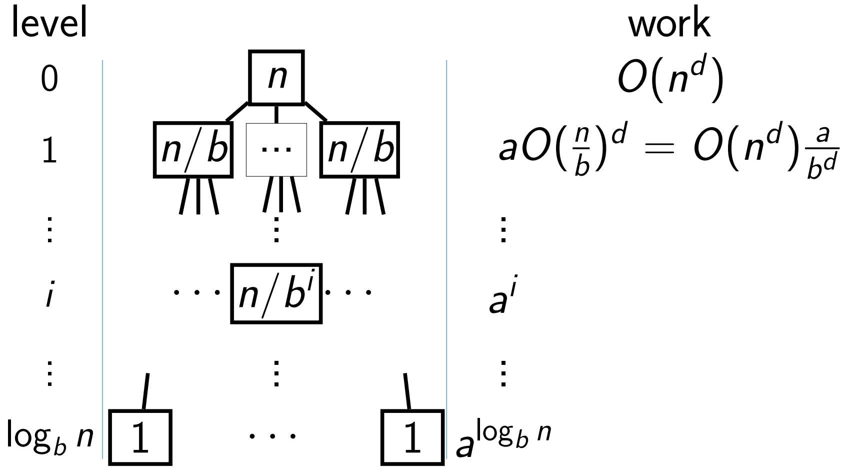

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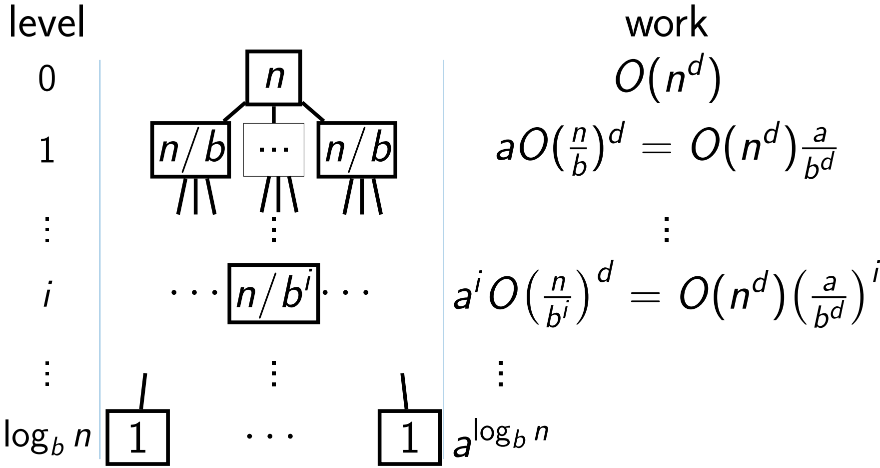

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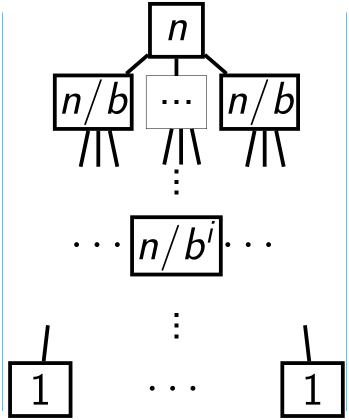

level	work
0	$O(n^d)$
1	$aO(\frac{n}{b})^d = O(n^d)\frac{a}{b^d}$
$\vdots$	$\vdots$
i	$a^i O\left(\frac{n}{b^i}\right)^d = O(n^d)\left(\frac{a}{b^d}\right)^i$
$\vdots$	$\vdots$
$\log_b n$	$a^{\log_b n} = O(n^{\log_b a})$

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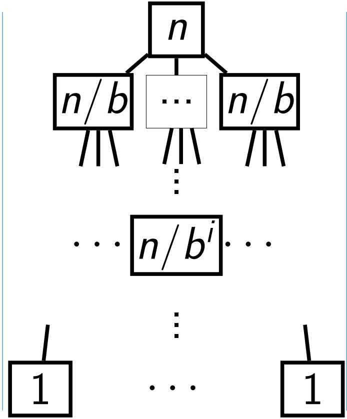
$\vdots$	$\vdots$
$\log_b n$	$a^{\log_b n} = O(n^{\log_b a})$



Total: $\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i$

Geometric Series

For $r \neq 1$:

$$a + ar + ar^2 + ar^3 + \cdots + ar^{n-1}$$

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$$
\begin{aligned}
& a + ar + ar^2 + ar^3 + \cdots + ar^{n-1} \\
=& a\frac{1 - r^n}{1 - r}
\end{aligned}
$$

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$$
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& a + ar + ar^2 + ar^3 + \cdots + ar^{n-1} \\
=& a\frac{1 - r^n}{1 - r} \\
=& \begin{cases} \\ \end{cases}
\end{aligned}
$$

Geometric Series

For $r \neq 1$:

$$\begin{aligned} & a + ar + ar^2 + ar^3 + \cdots + ar^{n-1} \\ =& a\frac{1 - r^n}{1 - r} \\ =& \begin{cases} O(a) & \text{if } r < 1 \\ & \end{cases} \end{aligned}$$

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$$\begin{aligned} & a + ar + ar^2 + ar^3 + \cdots + ar^{n-1} \\ = & a\frac{1 - r^n}{1 - r} \\ = & \begin{cases} O(a) & \text{if } r < 1 \\ O(ar^{n-1}) & \text{if } r > 1 \end{cases} \end{aligned}$$

Case 1: $\frac{a}{b^d} < 1$ $(d > log_b a)$

$$\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i$$

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$$\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i$$

$$= O(n^d)$$

Case 2: $\frac{a}{b^d} = 1$ ($d = log_b a$)

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$$\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i$$

$$= \sum_{i=0}^{\log_b n} O(n^d)$$

Case 2: $\frac{a}{b^d} = 1$ ($d = log_b a$)

$$
\begin{aligned}
&\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i \\
=&\sum_{i=0}^{\log_b n} O(n^d) \\
=&(1 + \log_b n)O(n^d)
\end{aligned}
$$

Case 2: $\frac{a}{b^d} = 1$ $(d = log_b a)$

$$
\begin{aligned}
&\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i \\
=&\sum_{i=0}^{\log_b n} O(n^d) \\
=&(1 + \log_b n)O(n^d) \\
=&O(n^d \log n)
\end{aligned}
$$

Case 3: $\frac{a}{b^d} > 1$ $(d < log_b a)$

$$\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i$$

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$$= O\left(O(n^d)\left(\frac{a}{b^d}\right)^{\log_b n} \right)$$

Case 3: $\frac{a}{b^d} > 1$ $(d < log_b a)$

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$$= O\left(O(n^d)\left(\frac{a}{b^d}\right)^{\log_b n}\right)$$

$$= O\left(O(n^d)\frac{a^{\log_b n}}{b^{d \log_b n}}\right)$$

Case 3: $\frac{a}{b^d} > 1$ $(d < log_b a)$

$$\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i$$

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$$= O\left(O(n^d)\frac{a^{\log_b n}}{b^{d \log_b n}}\right)$$

$$= O\left(O(n^d)\frac{n^{\log_b a}}{n^d}\right)$$

Case 3: $\frac{a}{b^d} > 1$ ($d < log_b a$)

$$
\begin{aligned}
&\sum_{i=0}^{\log_b n} O(n^d)\left(\frac{a}{b^d}\right)^i \\
=&O\left(O(n^d)\left(\frac{a}{b^d}\right)^{\log_b n}\right) \\
=&O\left(O(n^d)\frac{a^{\log_b n}}{b^{d\log_b n}}\right) \\
=&O\left(O(n^d)\frac{n^{\log_b a}}{n^d}\right) \\
=&O(n^{\log_b a})
\end{aligned}
$$

Summary

Master theorem is a shortcut:

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Theorem

If $T(n) = aT\left(\left\lceil \frac{n}{b} \right\rceil\right) + O(n^d)$ (for constants $a > 0, b > 1, d \geq 0$), then:

$$T(n) = \begin{cases} O(n^d) & \text{if } d > \log_b a \\ O(n^d \log n) & \text{if } d = \log_b a \\ O(n^{\log_b a}) & \text{if } d < \log_b a \end{cases}$$