ESO207 Programming Assignment 1

Devansh Kumar Jha
(200318) and Divyansh Gupta
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1 Problem Statements

Polynomials may be represented as linked lists. Consider a polynomial

$$p(x) = a_1 x^{e_1} + a_2 x^{e_2} + a_3 x^{e_3} + \dots + a_{n-1} x^{e_{n-1}} + a_n x^{e_n}$$

with non-zero terms, p(x) where $0 \le e_1 < e_2 < ... < e_n - 1 < e_n$ are (non-negative) integers. We assume that coefficients a1, . . . , an are non-zero integers. Polynomial p(x) can be represented as a linked list of nodes. Each node has three fields: coefficient, exponent and link to the next node. Let us assume that list is a doubly linked list, with sentinel node, sorted in ascending order of exponents.

(a) (marks 5+15)

Write pseudo-code to add two polynomials p(x) and q(x) in this representation. Your algorithm should take O(n+m) time, where n, m are the number of terms in p(x) and q(x) respectively. Implement your pseudo-code as an actual program.

(b) (marks 10+20)

Write pseudo-code to multiply two polynomials p(x) and q(x) in this representation. Do runtime complexity analysis of your algorithm in terms of n, m, the number of terms in p(x) and q(x) respectively. State this complexity in 'O' notation. Implement your pseudo-code as an actual program.

Note that output list should satisfy all constraints (non-zero coefficients, exponents in strict ascending order etc.) of representation of a polynomial. Make your code non-destructive, that is, it should not modify the lists for p(x) and q(x).

2 Programming Problem 1

2.1 Problem Description

We are required to add the polynomials according to the given input where we are first given two integers n and m which are followed by 2n and 2m integers in the next two lines. The main complexities and boundary cases in the problem were to design a O(m+n) algorithm and to recognize that zero polynomials are the exceptional cases which do not directly fit into the algorithm and so are to be specifically handled. To take care of integer overflow we by default use big integers as our storage containers.

2.2 Data Structure Usage

We will be using a dynamically allocated doubly linked list with a sentinel node for easier and error free implementation.

2.3 Algorithm Explanation

The algorithm is quite simple. We just take two temprorary node pointers which will correspond to the current term of that particular polynomial in the addition process. There are 3 major cases as follows -

• 1st polynomial already traversed

In this case we just add the term of 2nd polynomial to the answer polynomials doubly linked list. Here we don't need to check anything as the question statement says that the coefficients will anyways be non-zero.

• 2nd polynomail already traversed

Similar to the case above with the difference that this time we will be traversing the 1st polynomial and adding new nodes to the resultant polynomial.

• Both polynomials are being traversed

This is the most complex case of the mentioned three with some exceptions to be handled within. While we are simultaneously traversing the 1st and 2nd polynomial the currently checked term might have different variable exponents. In that case we cannot add them and we simply need to add one of them to the list. We chose to add the element with smaller exponent into the list as it ensures the automatic sorting of the resultant polynomial. In case of the exponents being equal we need to confirm that the coefficients don't add to zero. In case they do than we just skip these terms however if not than we add them. So it can easily be seen that in this case the iteration would be a bit more costly as compared 1st and 2nd case iteration.

2.4 Pseudo Code

These are the variables used -

- \bullet \mathbf{HEAD} The first element of addition polynomial.
- \bullet ${\bf HEAD1}$ The first element of first polynomial.
- \bullet $\mathbf{HEAD2}$ The first element of second polynomial.
- \bullet ${\bf SENT}$ Sentinel node for addition polynomial.
- ** Similar goes for TAIL, TAIL1, TAIL2, SENT1, SENT2

Algorithm 1: List Making Algorithm

```
Input: The two linked lists showing p(x) and q(x) and pointers to
              make output polynomial r(x)
   Output: No output
 1 temp1 = HEAD1 and temp2 = HEAD2
 2 /* Temprorary pointers to traverse along the available
       polynomials.
 \mathbf{3} while temp1! = SENT1 or temp2! = SENT2 do
       if temp1 == SENT1 then
           TAIL.insert - node(temp2)
 5
                                              // insert-node() is a function
            which makes a new node and adds to the list
            temp2 \leftarrow (temp2 \rightarrow next)
       else
6
          if temp2 == SENT2 then
 7
              TAIL.insert - node(temp1)
 8
              temp1 \leftarrow (temp1 \rightarrow next)
 9
           else
10
              if temp1 \rightarrow exp == temp2 \rightarrow exp then
11
                  if temp1 \rightarrow cof! = temp2 \rightarrow cof then
12
                      a=make-
13
                          node(temp1 \rightarrow cof + temp2 \rightarrow cof, temp1 \rightarrow cof)
                          /* make-node() function allocates memory for a
                          new node and fills with the parameters given */
                      TAIL.insert - node(a)
14
                  else
15
                  end if
16
                  temp1 \leftarrow (temp1 \rightarrow next)
17
                  temp2 \leftarrow (temp2 \rightarrow next)
18
              else
19
                  if temp1 \rightarrow exp < temp2 \rightarrow exp then
20
                      TAIL.insert - node(temp1)
21
                      temp1 \leftarrow (temp1 \rightarrow next)
22
                  else
23
                      TAIL.insert - node(temp2)
24
                      temp2 \leftarrow (temp2 \rightarrow next)
25
                  end if
26
              end if
27
          end if
28
       end if
29
30 end while
31 return
```

3 Programming Problem 2