

Problem

A *ladder* is a sequence of strings s_1, s_2, \dots, s_k , wherein every string differs from the preceding one by exactly one character. For example, the following is a ladder of English words, starting with "head" and ending with "free":

head, hear, near, fear, bear, beer, deer, deed, feed, feet, fret, free.

Let $LADDER_{DFA} = \{\langle M, s, t \rangle \mid M \text{ is a DFA and } L(M) \text{ contains a ladder of strings, starting with } s \text{ and ending with } t\}$. Show that $LADDER_{DFA}$ is in PSPACE.

Step-by-step solution

Step 1 of 3

$LADDER_{DFA} \in PSPACE$ $LADDER_{DFA} = \{\langle M, s, t \rangle \mid M \text{ is a DFA and } L(M) \text{ contains a ladder of strings, starting with } s \text{ and ending with } t\}$

The ladder is a sequence of strings where every string differs from preceding one in exactly one character.

To prove: $LADDER_{DFA}$ is in PSPACE

It is known that $PSPACE = NPSPACE$ so, if it is shown that $LADDER_{DFA} \in NPSPACE$ then it implies that $LADDER_{DFA} \in PSPACE$

It is known that NPSPACE is non-deterministic Turing Machine that clears the class of languages which can be decided in polynomial space.

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Step 2 of 3

Follow the following steps:

- If $|s| \neq |t|$, then reject.
- Otherwise, the graph G of size vertices are indexed by strings in $2^{|s|}$ and there is a directed edge from w_1 to w_2 vary in precisely one character and $w_1, w_2 \in L(M)$
- $(M, s, t) \in LADDER_{DFA}$ if and only if there is a path from s to t in G .
- It can be checked in NPSPACE by guessing the path and at each step storing only the name of the current vertex.
- To guess the path at the vertex w_1 non deterministically select a new vertex w_2 that differs from w_1 in precisely one character and verify that M accepts w_2 .
- The machine M always halts by keeping a counter and incrementing it with each guess and rejecting when the counter hits $|\Sigma|^{|s|}$

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Step 3 of 3

Thus, the non-deterministic Turing machine is constructed to decide $LADDER_{DFA}$

Therefore, $LADDER_{DFA} \in NPSPACE$ so, $LADDER_{DFA} \in PSPACE$

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