

## Step-by-step solution

## Step 1 of 1

 $NTIME(n)_{\ \, \text{is strict subset of}}\ \, PSPACE(n)$ 

At most  $t^{(n)}$  tape cells on each branch can be used by any Turing machine that operates in time  $t^{(n)}$  on each computation branch. So, it can be stated that  $t^{(n)}$  NSPACE $t^{(n)}$ 

- Now, consider the **Savitch's theorem** which says that: "Let  $f: N \to R_{\text{be a function, with}}$   $f(n) \ge n$  then  $\text{NSPACE}(f(n)) \subseteq \text{SPACE}((f(n))^2)_n$ . Therefore according to Savitch's theorem  $\text{NSPACE}(n) \subseteq \text{SPACE}(n^2)_n$ .
- Now, consider the **space hierarchy theorem** which says that "if g is space-constructible (  $1^n \to 1^{g(n)}$  can be computed in space  $O(g(n))_{),}$  f(n) = O(g(n)) then  $SPACE(f(n)) \subsetneq SPACE(g(n))_{.}$

Therefore, according to space hierarchy theorem it can be said that  $SPACE\left(n^{2}\right) \subsetneq SPACE\left(n^{3}\right). \text{ The result follows because } SPACE\left(n^{3}\right) \subseteq PSPACE.$ 

From the above explanation, it can be said that  $\begin{tabular}{l} NTIME(n) \subsetneq PSPACE(n) \ . \label{eq:pspace} \end{tabular}$ 

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