Problem

Consider the following two-person version of the language *PUZZLE* that was described in Problem 7.28. Each player starts with an ordered stack of puzzle cards. The players take turns placing the cards in order in the box and may choose which side faces up. Player I wins if all hole positions are blocked in the final stack, and Player II wins if some hole position remains unblocked. Show that the problem of determining which player has a winning strategy for a given starting configuration of the cards is PSPACE-complete.

Step-by-step solution

Step 1 of 4

Consider the two-people version of the language PUZZLE. In this, every player uses an ordered puzzle card stack to start.

- The cards are placed in each turn in order, by the player, in the box and can select which part faces up. Player I wins if every hole place are closed in the resultant stack and player I wins if some position of the holes remains in unblocked state.
- Now, it can be shows that "the problem of determining which player has a winning strategy for a given starting configuration of the cards is *PSPACE* complete".

Comment

Step 2 of 4

Consider the formula φ , which is defined as:

$$\varphi = \exists p_1 \forall p_2 \exists p_3 \left[\left(p_1 \vee p_2 \right) \wedge \left(p_2 \vee p_3 \right) \wedge \left(\overline{p_2} \vee \overline{p_3} \right) \right]$$

In the above given formula for φ , player I picks the value of P_1 and then player I picks the value of P_2 and finally player I picks the value of P_3 .

- Now, assign true with the value 1 and false with the value 0. Suppose that player I picks p_1 =1, then the player I picks p_2 =0, and finally player I picks p_3 =1.
- · Now by using these values in the sub-formula which is described above as

$$(p_1 \lor p_2) \land (p_2 \lor p_3) \land (\overline{p_2} \lor \overline{p_3})$$

is1, therefore the player I has won the game. Then, it can be said that "player I has a strategy of winning for the given puzzle".

Comment

Step 3 of 4

Now, a slightly change in formula, which is given above, will change the strategy of winning of the player.

· Consider the modified formula:

$$(p_1 \lor p_2) \land (p_2 \lor p_3) \land (p_2 \lor \overline{p_3})$$

The above formula is for winning strategy for the player II.

Comment

Step 4 of 4

Consider the PUZZLE, which is described above, is PSPACE-complete. It is because, it is same as TQBF. The PUZZLE problem is PSPACE-complete can be proved by using the fact of PUZZLE=TQBF.

• Consider the formula $\varphi = \exists p_1 \forall p_2 \exists p_3 ... [\psi]$ became TRUE when setting of p_1 exists in such a way that, for every setting of p_2 , a setting of p_3 exist is such a way and so on. Where, p_3 sows TRUE value under the setting of the taken variable.

- ${\mbox{\footnote{h}}}$ The same process will be applied for the winning strategy for the player $I\!I$.
- If the formula φ has of the form $\forall p_1, p_2, p_3 \exists p_4, p_5 \forall p_6 [\psi]$, player I would make the first three move in e PUZZLE, to assign values to p_1, p_2 and p_3
- \cdot Then, player II makes two moves to assign p_4 and p_5 . Finally, the player I assigned a value p_6 .

Therefore, $\varphi \in TQBF$ exactly when $\varphi \in PUZZLE$. Hence it can be said that "the problem of determining which player has a winning strategy for a given starting configuration of the cards is PSPACE -complete".

Comment