

## Problem

Define the **unique-sat** problem to be

$USAT = \{ \langle \phi \rangle \mid \phi \text{ is a Boolean formula that has a single satisfying assignment} \}$ .

Show that  $USAT \in P^{SAT}$ .

## Step-by-step solution

### Step 1 of 1

#### Given:

$USAT = \{ \langle \phi \rangle \mid \phi \}$ , is basically a Boolean formula which is use for satisfying the single assignment.

#### Proof:

Here, it is to be proved that every Boolean formula which is satisfied has minimum one assignment which is bounded.

For proving this user need to satisfy the assignment of the  $USAT$  formula which is defined syntactically which the help of propositional logic.

$$USAT \in P^{SAT}$$

#### Construction:

Here, user is proving that  $USAT$  should have only single truth value assignment which can be proved with the help propositional logic.

- User need to define a class of bonded truth value.
- It is to be proved that formula  $\phi$  should have minimum one assignment which is bounded. For proving this there is one assignment which is chosen that is  $I_\phi$  which should be syntactically defined.
- Therefore, characterization of the assignment which is uniquely satisfied is  $\phi \in USAT$  and it is proved by the definition which is  $I_\phi$ .
- $I_\phi$ , constraint is use for satisfying all the assignment of  $\phi$  this is because it contains information for the truth value assignment which is use for satisfying the value of  $\phi$ .
- The assignment  $I_\phi$  is syntactically defined with the help of propositional logic for each and every value of the variable  $p$ , substitution of the value  $p$  is the variable free formula.
- Here,  $I_\phi$  is satisfying all the truth value assignment  $\phi$ .
- $I_\phi$ , contains information for the truth value assignment which is use for satisfying the value of  $\phi$  and even it help in determining the other truth value which is use for falsifying the value of  $\phi$ .

It is proved that unique assignments in  $USAT$  are basically bounded and can be syntactically defined with the help of propositional logic.

Hence, it is proved that  $USAT \in P^{SAT}$

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