Problem

Show that if G is a CFG in Chomsky normal form, then for any string w $\sum L(G)$ of length $n \ge 1$, exactly 2n - 1 steps are required for any derivation of w.

Step-by-step solution

Step 1 of 1

Given that G is a CFG in Chomsky Normal Form (CNF).

The length of the string $w \in L(G)$ is $n \ge 1$ for the string w.

It is required to show that exactly 2n-1 steps are required for the derivation of string w.

It can be proved applying the **induction method** by on the string w of length n.

For n = 1: Consider a string "a" of length 1 in Chomsky normal form, so the valid derivation for this will be $S \to a$, where $a \in \sum$ and S is starting symbol.

The number of steps can be obtained as follows:

$$2n-1=2(1)-1$$

= 2-1

Hence it is true that 2n-1 (for n=1) steps are required to derive a string a.

For n = k: Take a string of length $k \ge n$ in Chomsky normal form, so valid derivation for this will take 2k - 1 steps.

The number of steps can be obtained as follows:

$$2n-1=2(k)-1$$
$$=2k-1$$

Assume a string of length at most $k \ge 1$ terminal symbols and it has a string of length

n=k+1 is in Chomsky Normal Form

Since n>1, Consider a language as follows in CNF where derivation starts with start symbols S:

$$S \to BC$$

$$B \to *x$$

$$C \to *y$$

So length of the string starting with start symbol S is |w| = xy where |x| > 0 and |y| > 0.

Using the inductive hypothesis, for the above language in CNF the length of any derivation of string w must be

$$1+(2|x|-1)+(2|y|-1)=2|x|+2|y|+1-1-1=2(|x|+|y|)-1$$

Here n=|x|+|y|.

Since $B \to *x$ has a length of lxl and $C \to *y$ has a length of lyl.

Hence, it is proved that it requires ${\bf 2n\text{-}1}$ steps required for the derivation of string

 $w \in L(G)$ in Chomsky Normal Form(CNF).

Comment