

CSE 303 ToC. 115060128.
Assignment: 2.

Problem 1:

- a) Given: $f: A \rightarrow B$ & $g: B \rightarrow C$ are injective.
TPT: $g \circ f: A \rightarrow C$ is also injective.

We will use proof by contradiction:
Suppose $g \circ f: A \rightarrow C$ is not injective
 $\Rightarrow \exists a, a' \in A, g \circ f(a) = g \circ f(a') \quad - (1)$

Since f is injective $\Rightarrow f(a) \neq f(a') \because a \neq a'$
Let $f(a) = b$ & $f(a') = b' \Rightarrow b \neq b' \quad - (2)$

By (1), $g(b) = g(b')$
Since g is injective, $g(b) = g(b') \Rightarrow b = b'$ which contradicts (2).
 $\therefore g \circ f: A \rightarrow C$ is injective. Hence proved

- b) Given $f: A \rightarrow B$ & $g: B \rightarrow C$ are surjective.
TPT: $g \circ f: A \rightarrow C$ is surjective.

We will prove $\forall c \in C, \exists a \in A$ s.t. $g \circ f(a) = c$.
This by definition will prove $g \circ f$ is surjective.
Let c be an arbitrary but fixed element in C .
We have $g(f(a)) = c$.
Since g is surjective $\exists b \in B, g(b) = c$.

Since f is surjective $\exists a \in A, f(a) = b$.

Thus $g(f(a)) = g(b) = c$.

$\Rightarrow g \circ f$ is surjective. Hence proved.

C. Given: $f: A \rightarrow B$ and $g: B \rightarrow C$ are bijective.
 TPT: $g \circ f: A \rightarrow C$ is bijective.

\rightarrow Since f and g are injective.
 $\Rightarrow g \circ f$ is injective Using (a)

Since f and g are surjective
 $\Rightarrow g \circ f$ is surjective Using (b)

Since $g \circ f$ is injective and surjective
 $\Rightarrow g \circ f$ is bijective. Hence proved.

Problem 2:-

a). Let (a, b) and $(b, c) \in R^*$
 we need to prove $(a, c) \in R^*$

$\rightarrow (a, b) \in R^* \Rightarrow \exists n \geq 0$ and sequence $a_0, a_1, \dots, a_n \in A$
 s.t $a_0 = a, a_n = b$ & $(a_k, a_{k+1}) \in R \forall k \in \{0, 1, \dots, n-1\}$ (1)

$(b, c) \in R^* \exists m \geq 0$ and sequence $b_0, b_1, \dots, b_m \in A$ s.t
 $b_0 = b, b_m = c$ & $(b_k, b_{k+1}) \in R \forall k \in \{0, 1, \dots, m-1\}$ (2)

We will prove \exists a sequence, from a to c , i.e
 $(a, c) \in R^*$.

Take $a_0 = a, c_0 = a_0, c_1 = a_1, \dots, c_n = a_n = b$
 $c_{n+1} = b_1, c_{n+2} = b_2, \dots, c_{n+m} = b_m = c$.

~~we~~ Take $l = n + m$.

(2)

We have $(c_k, c_{k+1}) \in R \forall k \in \{0, 1, \dots, l\}$.
Using ① & ②

\therefore By definition $(a, c) \in R^*$
 $\Rightarrow R^*$ is transitive relation.

b). We have $R \subseteq T$.

Let T be any transitive relation on A that contains R .

We need to prove $R^* \subseteq T$.

Let $(a, b) \in R^*$ be arbitrary but fixed elements.
We will show $(a, b) \in T$.

Thus we will get $R^* \subseteq T$

i.e. every element in R^* also belongs to T .

$(a, b) \in R^* \Rightarrow \exists n \geq 0$ & sequence $a_0, a_1, \dots, a_n \in A$
s.t. $a_0 = a, a_n = b$ & $(a_k, a_{k+1}) \in R \forall k \in \{0, 1, \dots, n-1\}$

Since $R \subseteq T; (a_k, a_{k+1}) \in T \forall k \in \{0, 1, \dots, n-1\}$

We have $(a_0, a_1) \in T$ & $(a_1, a_2) \in T \Rightarrow (a_0, a_2) \in T$.
 $(a_0, a_2) \in T$ & $(a_2, a_3) \in T \Rightarrow (a_0, a_3) \in T$.

$(a_0, a_k) \in T$ & $(a_k, a_{k+1}) \in T \Rightarrow (a_0, a_{k+1}) \in T$
 $\Rightarrow (a_0, a_n) \in T \Rightarrow (a, b) \in T$.
 $K \in \{1, 2, \dots, n-1\}$

Thus every element (a, b) in R^* exists in T

$\Rightarrow R^* \subseteq T$.

c) i) We have from b, any Transitive relation on R will contain R^* .

So if we take many transitive relations, their intersection will contain R^* because all of them contain R^* .

Thus $\bigcap T$ will contain R^* .

ii) R^* contains $\bigcap T$.
i.e. $\bigcap T \subseteq R^*$

$\bigcap T$ is the intersection ~~and~~ of all transitive relations on A that contain R .

R^* is one of them.

So by definition if $C = A \cap B$

then $C \subseteq A$

$\Rightarrow \bigcap T \subseteq R^*$.

Problem 3:

a) We will need to prove this in two parts.

Part I: Given $(w, x) \in R^*$ $\Rightarrow \psi_x = \psi_w$

\Rightarrow either $w = x$ which means $\psi_x = \psi_w$ Proved
or $w = uabv$ and $x = ubav$.

Since Parikh vector is counts for occurrences.

If $C = ab$ then $\psi_C = \psi_a + \psi_b$.

(3)

Thus $w = uabv \Rightarrow \psi_w = \psi_u + \psi_a + \psi_b + \psi_v$
 $x = ubav \Rightarrow \psi_x = \psi_u + \psi_b + \psi_a + \psi_v$
 $\Rightarrow \psi_w = \psi_x$ Hence proved.

Part II: Given $\psi_x = \psi_w$.

Because w and x have same occurrences of a_1, a_2, \dots, a_k they are a permutation of same letters appearing same no. of times.

Let y be a sorted element on Σ^* s.t.
 $\psi_y = \psi_x = \psi_w$.
 by sorted we mean lexicographic order

We will check first two elements of x . If they are sorted we will move to 2nd & 3rd element. if they are also sorted we will move ahead. If not, we will swap them. Thus process will be repeated multi fold. to get intermediate x_1, x_2, \dots, x_n . which only one adjacent element swapped bet x_k, x_{k+1} , $k \in \{1, 2, \dots, n\}$
 Hence $\psi_{x_k} = \psi_{x_{k+1}}$, $k \in \{1, 2, \dots, n-1\}$

Since x_n is lexicographically sorted, $x_n = y$.

We can follow same process for w & bring it to y through w_1, w_2, \dots, w_m intermediate strings.

Then we can transform x to w through

$x \rightarrow x_1 \rightarrow x_2 \rightarrow \dots \rightarrow x_n = y = w_m \rightarrow w_{m-1} \rightarrow w_{m-2} \rightarrow \dots \rightarrow w_2 \rightarrow w_1 \rightarrow w$.

Since each of these involves one adjacent element swap all their Parikh vectors are same

And $(x, x_1) \in R^*$
and $(x_1, x_2) \in R^*$

\vdots
 $(w_i, w) \in R^*$

$\Rightarrow (x, w) \in R^*$ by transitivity Rule applied
n+m times.

Thus $(x, w) \in R^*$

$\Rightarrow (w, x) \in R^*$ by taking exact opposite
path.