Problem

Recall the CFG G_4 that we gave in Example 2.4. For convenience, let's rename its variables with single letters as follows.

$$\begin{split} E &\rightarrow E + T \mid T \\ T &\rightarrow T \times F \mid F \\ F &\rightarrow (E) \mid \mathbf{a} \end{split}$$

Give parse trees and derivations for each string.

- **a.** a
- **b.** a+a
- **c.** a+a+a
- **d.** ((a))

Step-by-step solution

Step 1 of 8

Given Grammar G_4 is

$$E \rightarrow E + T \mid T$$

$$T \to T \times F \mid F$$

$$F \rightarrow (E) \mid a$$

Derivation: The sequence of substitutions to obtain a string is called a *derivation*.

Parse Tree: The pictorial representation of derivation of a string is a parse tree.

Comment

Step 2 of 8

a)

The parse tree to generate string a is as follows:



Comment

Step 3 of 8

$E \Rightarrow I$	
$E \Rightarrow F$	
$E \Rightarrow a$	
Comment	
	01::: 4::10
	Step 4 of 8
b)	
The parse tree to generate string $a + a$ is as follows:	
r F	
F	
a + a	
Comment	
	Step 5 of 8
	Step 3 010
The derivation for the string $a + a$ is as follows:	
$E \Rightarrow E + T$	
$E \Rightarrow T + T$	
$E \Rightarrow F + T$	
$E \Rightarrow a + T$	
$E \Rightarrow a + F$	
$E \Rightarrow a + a$	
Comment	
	Step 6 of 8
	chisp c and
c)	
The parse tree to generate string $a+a+a$ is as follows:	
E	
Ň	
E I	
<i>f</i>	
F	
a + a + a	
The derivation for the string $a+a+a$ is as follows:	
$E \Rightarrow E + T$	
$E \Rightarrow E + T + T$	
$E \Rightarrow T + T + T$	
$E \Rightarrow F + T + T$	
$\nu \rightarrow \nu + \tau + \tau$	
$E \Rightarrow a + T + T$ $E \Rightarrow a + E + T$	
$E \Rightarrow a+T+T$ $E \Rightarrow a+F+T$ $E \Rightarrow a+a+T$	

 $E \Rightarrow a + a + F$ $E \Rightarrow a + a + a$

		Step 7 of 8	
d)			
The parse tree to generate string	((a))is as follows:		
	((~))		
E T F (E) (T) (F)			
Į.			
<u>Í</u>			
(E)			
(<u>†</u>)			
(F)			
((E))			
((<i>T</i>))			
((<i>F</i>))			
((a))			
Comment			
		Step 8 of 8	
The derivation for the string (a)	is as follows:		
$E \Rightarrow T$			
$E \Rightarrow F$			
$E \Rightarrow (E)$			
$E \Rightarrow (T)$ $E \Rightarrow (F)$			
$E \Rightarrow ((E))$			
$E \Rightarrow ((T))$			
$E \Rightarrow ((F))$			
$E \Rightarrow ((a))$			
Comments (2)			