Problem

Describe the error in the following fallacious "proof" that $P \neq NP$. Assume that P = NP and obtain a contradiction. If P = NP, then SAT? P and so for some k, SAT? TIME(n^k). Because every language in NP is polynomial time reducible to SAT, you have

$$NP \subseteq TIME(n^k)$$
. Therefore, $P \subseteq TIME(n^k)$.

Step-by-step solution

	Step 1 of 2
Consider the followi	ing deceptive "proof "that $P \neq NP$
1. Suppose that P=N	NP and obtain a contradiction.
2. If P=NP has given	, then for some $k \in N$ and $SAT \in P$, $SAT \in TIME(n^k)$
3. As every language	in $\ NP$ is polynomial time reducible to $SAT_{,then} \ \ NP \subseteq TIME \left(n^{k}\right)$
4. Due to the assump	tion $P=NP$ which is taken above, $P\subseteq TIME(n^k)$
5. Then $TIME(n^k)$	\varsubsetneq $TIMEig(n^{k+1}ig)$ is a contradiction with P \subseteq $TIMEig(n^kig)$
Comment	
	Step 2 of 2
	e used to proof of $P \neq NP$, but it consists some error. Now, consider the step 3, here the time needed for calculating the en into account. So, to make this proof error free, the time needed for calculating the reduction must be taken in to
Comments (1)	