

## Problem

Show that  $A_{DFA} \in L$ .

## Step-by-step solution

### Step 1 of 1

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**Class  $L$ :** In deterministic Turing machine,  $L$  can be defined as class of languages complexity decidable in logarithmic space i.e.  $L = SPACE(\log n)$

Now we need to prove that  $A_{DFA} \in L$

$$A_{DFA} = \{ \langle B, w \rangle \mid B \text{ is a DFA that accepts input string } w \}$$

To prove  $A_{DFA} \in L$  we need to construct a deterministic Turing machine ( $TM$ ) to decide

$A_{DFA}$  in logarithmic space.

Let  $M$  be the  $TM$  that decides  $A_{DFA}$  in log – space.

The construction of  $TM$   $M$  is as follows:

$M$  = “on input  $\langle B, w \rangle$ , where  $B$  is a DFA and  $w$  is string:

1. Simulate  $B$  on  $w$  by keeping track of  $B$ 's current state and its current head location and updating them appropriately.
2. If the simulation ends in an accept state, accept
3. else if the simulation end in non accepting that, reject”

The SPACE required to carry out this simulation is  $O(\log n)$ , since  $n$  items of values storing its input by  $M$ . Thus we constructed  $TM$   $M$  to decide  $A_{DFA}$  in log – space therefore  $A_{DFA} \in L$ .

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