Problem

Define the function majority_n as in Problem 9.24. Show that it may be computed with O(n) size circuits.

Problem 9.24

$$majority_n: \{0,1\}^n \longrightarrow \{0,1\}$$
 as

Define the function

$$majority_n(x_1, \dots, x_n) = \begin{cases} 0 & \sum x_i < n/2; \\ 1 & \sum x_i \ge n/2. \end{cases}$$

Thus, the majority_n function returns the majority vote of the inputs. Show that majority_n can be computed with:

- a. O(n2) size circuits.
- b. O(n log n) size circuits. (Hint: Recursively divide the number of inputs in half and use the result of Problem 9.23.)

Step-by-step solution

Step 1 of 1

Consider the number of inputs taken is n. A **bubble-sort** can be implemented as a circuit. It is used to compare two bits and after comparing, reordering them if necessary is rather easy. The inputs can be called as x_1, x_2 and the outputs can be called as y_1, y_2 . A sub-circuit can be written which accomplishes this as $y_1 = OR(x_1, x_2)$ and $y_2 = AND(x_1, x_2)$. This circuit contains a size of two.

- Now, the action of the bubble-sort algorithm can be mimicked on an array. It can be implemented one step at position to be the n input, n-output subcircuit that passes through all the inputs taken as k and k are unchanged.
- Now, the compare-swap sub-circuit, which is described above, on $< k \ and \ge k+1st$ input can be used to generate the kth and k+1st output. This still has size two. Now, **a pass** can be implemented as the serial concatenation of steps for each of k=1,2,...,n-1, which has a size $\binom{(n-1)*2}{2}$.
- A bubble-sort can be Proceed to implement as the serial concatenation of one passes. Therefore, this gives a size 1(n-1)*2 = O(n)

This O(n) complexity can be achieved only when the **input taken is already in sorted order**. In other word it can be said that, if it shows **a best case** behavior. Therefore, it can be said that $majority_n$ can be computed in O(n) size circuits.

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