## **Problem**

Show that if  $NP = P^{SAT}$ , then NP = coNP.

## Step-by-step solution

## Step 1 of 1

If  $NP = P^{SAT}$  is assumed then, there is only need to show that NP = coNP. It can be conclude directly from the assumption  $NP = P^{SAT}$ , that  $P^{SAT} \subseteq NP$ . As  $coNP \subseteq P^{SAT}$  is already known, which result in  $coNP \subseteq NP$ .

• Now it is known that P is closed under the complement operation, so is  $P^{SMT}$ , because it can be just swap the reject and accept states. It can be concluded that:

$$L \in P^{SAT} \Longrightarrow \overline{L} \in P^{SAT}$$

- The given statement can be managed by using the prediction  $NP = P^{SAT}$ , because if this would be the case any language in  $P^{SAT}$  would be in NP and vice versa.
- For which NP also has to be closed under the complement operation  $L \in NP \Rightarrow \overline{L} \in NP$ , that is just the same as  $L \in NP \Rightarrow L \in coNP$  from the definition of coNP.
- In other words, from the above explanation it can be said that  $NP \subseteq coNP$  . Hence NP = coNP

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