$$G = (V, \Sigma, R, \langle \text{STMT} \rangle)_{\text{be the following grammar.}}$$

$$\langle \text{STMT} \rangle \rightarrow \langle \text{ASSIGN} \rangle \mid \langle \text{IF-THEN} \rangle \mid \langle \text{IF-THEN-ELSE} \rangle$$

$$\langle \text{IF-THEN} \rangle \rightarrow \text{if condition then } \langle \text{STMT} \rangle$$

$$\langle \text{IF-THEN-ELSE} \rangle \rightarrow \text{if condition then } \langle \text{STMT} \rangle \text{ else } \langle \text{STMT} \rangle$$

$$\langle \text{ASSIGN} \rangle \rightarrow \text{a:=1}$$

$$\Sigma = \{\text{if, condition, then, else, a:=1} \}$$

$$V = \{\langle \text{STMT} \rangle, \langle \text{IF-THEN} \rangle, \langle \text{IF-THEN-ELSE} \rangle, \langle \text{ASSIGN} \rangle \}$$

G is a natural-looking grammar for a fragment of a programming language, but G is ambiguous.

- a. Show that G is ambiguous.
- b. Give a new unambiguous grammar for the same language.

Step-by-step solution

Step 1 of 3

Ambiguous and unambiguous grammar Consider $G = (V, \Sigma, R, \langle STMT \rangle)$ To show that G is ambiguous It is required to find out two leftmost derivations: (Stmt) $\Rightarrow \langle \text{if-then} \rangle$ ⇒ if condition then ⟨Stmt⟩ ⇒ if condition then ⟨if-then-else⟩ ⇒ if condition then if condition then ⟨Stmt⟩ else ⟨Stmt⟩ ⇒ if condition then if condition then ⟨Assign⟩ else ⟨Stmt⟩ \Rightarrow if condition then if condition then a:=1 else $\langle Stmt \rangle$ \Rightarrow if condition then if condition then a:=1 else $\langle Assign \rangle$ ⇒ if condition then if condition then a:=1 else a:=1 Now the second derivation will be: (Stmt) ⇒ (if-then-else) ⇒ if condition then ⟨Stmt⟩ else ⟨Stmt⟩ ⇒ if condition then ⟨if-then⟩ else ⟨Stmt⟩ \Rightarrow if condition then if condition then $\langle Stmt \rangle$ else $\langle Stmt \rangle$

- ⇒ if condition then if condition then ⟨Assign⟩ else ⟨Stmt⟩
- \Rightarrow if condition then if condition then a:=1 else $\langle Stmt \rangle$
- ⇒ if condition then if condition then a:=1 else ⟨Assign⟩
- ⇒ if condition then if condition then a:=1 else a:=1

Step 2 of 3

In both cases when it takes "in-then" or "if-then else" result is same. Hence both have same left derivation and it is ambiguous.

Comment

Step 3 of 3

For making unambiguous grammar, it is required to make correct interpretation of the above two. For that when "if-then-else" is derived, it should not allow then part for deriving "if-then". So introduce a new variable $\langle Stmt \rangle$ the new grammar is:

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\begin{split} &\langle Stmt \rangle \! \Rightarrow \! \langle if\text{-then} \rangle | \; \langle Assign \rangle \; | \; \langle if\text{-then-else} \rangle \\ &\langle if\text{-then-else} \rangle \Rightarrow \; if \; condition \; then \; \langle Stmt1 \rangle \; else \; \langle Stmt \rangle \\ &\langle Stmt1 \rangle \! \Rightarrow \! \langle if\text{-then-else} \rangle \; | \; \langle Assign \rangle \\ &\langle if\text{-then} \rangle \! \Rightarrow \! if \; condition \; then \; \langle Stmt \rangle \\ &\langle Assign \rangle \! \Rightarrow \! a \! : \! = \! 1 \end{split}
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Comments (2)