## **Problem**

Consider the problem of determining whether a Turing machine M on an input we ver attempts to move its head left when its head is on the left-most tape cell. Formulate this problem as a language and show that it is undecidable.

## Step-by-step solution

## Step 1 of 2

Consider the problem of determining whether a Turing machine M on an input w ever attempts to move its head left when its head is on the left-most tape cell. This problem is formulated as a language:

 $L = \{ \langle M, w \rangle | M \text{ attempts to move its head left when its head is on the leftmost tape cell} \}$ 

Comment

## Step 2 of 2

Assume that the language L is decidable and  $\hat{M}$  be a TM that decides the language L. Construct a TM, A that decides the halting problem.

A = "on input < M, w>:

- Construct a TM, A', from A. The TM A' moves w one tape cell to the right and marks the leftmost cell with #.
- 2. Run the TM A' on  $\langle M, w \rangle$ .
- If A' encounters # then A' moves to the right side and simulates M reaching the leftmost tape cell.
- 4. If *M* halts and accepts on *w* then *A'* simulates to move its head left when its head is on the leftmost tape cell."

Now, TM A runs  $\hat{M}$  on the input  $\stackrel{<A, \ w>}{}$ . If  $\hat{M}$  accepts, A accepts. Otherwise, A rejects. It is assumed that  $\hat{M}$  be a TM that decides the language L. If M halts and accepts on w, then only A' moves its head left when its head is on the left-most tape cell. If A decides the halting problem, then halting problem is decidable. Thus, the halting problem is undecidable. It is a contradiction.

Therefore, the language L is undecidable.

Comments (1)