

## Problem

Let

$$f(x) = \begin{cases} 3x + 1 & \text{for odd } x \\ x/2 & \text{for even } x \end{cases}$$

for any natural number  $x$ . If you start with an integer  $x$  and iterate  $f$ , you obtain a sequence,  $x, f(x), f(f(x)), \dots$ . Stop if you ever hit 1. For example, if  $x = 17$ , you get the sequence 17, 52, 26, 13, 40, 20, 10, 5, 16, 8, 4, 2, 1. Extensive computer tests have shown that every starting point between 1 and a large positive integer gives a sequence that ends in 1. But the question of whether all positive starting points end up at 1 is unsolved; it is called the  $3x + 1$  problem.

Suppose that  $A_{TM}$  were decidable by a TM  $H$ . Use  $H$  to describe a TM that is guaranteed to state the answer to the  $3x + 1$  problem.

## Step-by-step solution

### Step 1 of 1

The solution can be built in stages. Later stages will use Turing machines that are constructed in earlier stages. The first step is to construct a Turing machine  $M_{\text{query}}$  that takes an input  $x$  and accepts if iterating  $f$  starting from  $x$  eventually yields 1 and loops forever otherwise.

$M_{\text{query}}$ . On input  $\langle x \rangle$ :

1. If  $x = 1$ , then accept.
2. If  $x$  is odd, update  $x \leftarrow 3x + 1$ . If  $x$  is even, update  $x \leftarrow \frac{x}{2}$ .
3. Go to step 1.

Our next step is to construct a Turing machine  $M_{\text{loop}}$  which iterates over all positive integers, looking for a counter-example to the  $3x + 1$  conjecture.

That is,  $M_{\text{loop}}$  searches for some  $x$  such that iterating  $f$  starting from  $x$  never reaches 1 and accepts if it finds such an  $x$ . To do this, it seems tempting to have  $M_{\text{loop}}$  simulate  $M_{\text{query}}$  first on  $x = 1$ , next on  $x = 2$ , and so forth. Whenever iterating  $f$  on  $x$  yields 1, the simulation of  $M_{\text{query}}$  will eventually end and  $M_{\text{loop}}$  would then proceed to the next number  $x + 1$ . But what if  $M_{\text{loop}}$  actually finds a counter-example  $x$ ? In this case, the simulation of  $M_{\text{query}}$  on  $x$  will never terminate, and  $M_{\text{loop}}$  will be in the unfortunate situation that it has found what it is looking for, but it doesn't know it has found it!

To get around this, use  $H$ . Instead of simulating  $M_{\text{query}}$  on  $x$ , it has  $M_{\text{loop}}$  check whether or not  $M_{\text{query}}$  would accept  $x$  by passing  $\langle M_{\text{query}}, x \rangle$ , to  $H$ .

$M_{\text{loop}}$ . On input  $\langle w \rangle$ :

1. Ignore the input  $w$ .
2. For each natural number  $y = 1, 2, 3, \dots$ :
3. Run  $H$  on  $\langle M_{\text{query}}, y \rangle$ .
4. If  $H$  rejects, then accept. Otherwise, continue the loop.

Finally, in order to solve the  $3x + 1$  problem, it is required to know whether or not  $M_{\text{loop}}$  finds a counter-example. Again, it might be tempted to simulate  $M_{\text{loop}}$  and see if it ever finds a counter-example and accepts. The problem is that there may not be any counter-example, in which case  $M_{\text{loop}}$  will loop forever and our simulation will not terminate. The trick is to use  $H$  again to see if  $M_{\text{loop}}$  finds a counter example.

$M_{3x+1}$ . On input  $\langle w \rangle$ :

1. Ignore the input  $w$ .
2. Run  $H$  on  $\langle M_{\text{loop}}, \epsilon \rangle$ .
3. If  $H$  accepts, then print "There is a counter-example to the  $3x + 1$  conjecture."

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