

Problem

Show that P is closed under union, concatenation, and complement.

Step-by-step solution

Step 1 of 4

The class P is closed under union, concatenation, and complement

Class P: P is the class of languages that are decidable in polynomial time on a deterministic single – tape Turing machine.

$$P = \bigcup_k \text{TIME}(n^k)$$

That is,

Now we have to show that P is closed under union, concatenation and complement.

[Comment](#)

Step 2 of 4

Union:

Assume two languages $P_1 \in P$ and $P_2 \in P$

The Turing machine M that accepts $P_1 \cup P_2$ works as follows:

$M =$ "On input w :

1. Check if $w \in P_1$
2. if not then check if $w \in P_2$
3. Accept w if and only if P_1 or P_2 accepts.
4. If both reject then rejects the input w ".

Since each membership check requires polynomial time the overall time is polynomial.

[Comment](#)

Step 3 of 4

Concatenation:

Assume two languages $P_1 \in P$ and $P_2 \in P$.

The Turing machine M that accepts $P_1 \cdot P_2$ works as follows

$M =$ "On input w of length n :

1. w can be split into two strings in n different ways
2. For each split,
 - (a) Check if the first substring belongs to P_1
 - (b) check if second substring belongs to P_2
3. If any split succeeds, then accept".

Clearly the overall time is polynomial.

[Comments \(1\)](#)

Step 4 of 4

Complement :

Assume a language $P_1 \in P$

The Turing machine M that accepts \bar{P}_1 works as follows

$M =$ "On input w :

1. Check if $w \in P_1$
2. If $w \in P_1$ then reject.
3. If $w \notin P_1$, then accept"

Clearly the overall time is polynomial.

Therefore, the class P is closed under union, concatenation, and complement.

[Comments \(1\)](#)