## **Problem**

Answer each part for the following context-free grammar G.

$$\begin{array}{l} R \rightarrow XRX \mid S \\ S \rightarrow \mathbf{a}T\mathbf{b} \mid \mathbf{b}T\mathbf{a} \\ T \rightarrow XTX \mid X \mid \varepsilon \\ X \rightarrow \mathbf{a} \mid \mathbf{b} \end{array}$$

- **a.** What are the variables of G?
- **b.** What are the terminals of G?
- **c.** Which is the start variable of *G*?
- **d.** Give three strings in L(G).
- **e.** Give three strings *not* in L(G).
- **f.** True or False:  $T \Rightarrow aba$ .
- **g.** True or False:  $T \stackrel{*}{\Rightarrow}$  aba.
- **h.** True or False:  $T \Rightarrow T$ .

- i. True or False:  $T \stackrel{*}{\Rightarrow} T$ .
- i. True or False:  $XXX \stackrel{*}{\Rightarrow} aba$ .
- **k.** True or False:  $X \stackrel{*}{\Rightarrow} aba$ .
- **l.** True or False:  $T \stackrel{*}{\Rightarrow} XX$ .
- **m.** True or False:  $T \stackrel{*}{\Rightarrow} XXX$ .
- **n.** True or False:  $S \stackrel{*}{\Rightarrow} \varepsilon$ .
- **o.** Give a description in English of L(G).

## Step-by-step solution

Step 1 of 16

Given the context – free grammar G is

$$R \to XRX \mid S$$

$$S \to aTb \mid bTa$$

$$T \to XTX \mid X \mid \varepsilon$$

$$X \to a \mid b$$

Comment

Step 2 of 16

a. The variables of G are R, S, T and X

Variables are the non-terminal symbols that appear in the rules of the grammar.

Comment

<i>Ierminals</i> are	the terminal symbols that appear in the rules of the grammar.
Comments (2)	
	Step 4 of 16
c. The start va	riable of $G$ is $R$ .
$R \to XRX \mid S$	
	is a variable usually occurs on the left – hand side of the topmost rule.
Comments (3)	
	Step 5 of 16
d. Case1: Con	sider the rule $R \to S$
Substitute S w	with rule $S \rightarrow aTb$
$R \rightarrow aTb$	
	ith rule $T \to \varepsilon$ .
$R \to a \in b$	
$R \rightarrow ab$	
Case2: Consid	der the rule $R \to S$
Substitute S w	ith rule $S \rightarrow bTa$
$R \rightarrow bTa$	
Substitute T w	ith rule $T \to \varepsilon$ .
$R \to b \in a$	
$R \rightarrow ba$	
Case3: Consid	der the rule $R \to S$
Substitute S w	ith rule $S \rightarrow aTb$
$R \rightarrow aTb$	
Substitute T w	ith rule $T \to X$ .
$R \rightarrow aXb$	
Substitute X w	ith rule $X \to a$ .
$R \rightarrow aab$	
Therefore, the	3 strings in $L(G)$ are $ab$ , $ba$ and $aab$ .
Comment	
	Step 6 of 16
e. The three s	trings not in $L(G)$ are $aba,b$ and $arepsilon$ . Since these strings cannot be derived from the given grammar $G$ .
Comments (6)	
	Step 7 of 16
. Fole-	
f. False	
$T \Rightarrow aba$ , the	$\epsilon$ string cannot be derived using $G$ .

	<b>Step 8</b> of 16
g. True	
$T \stackrel{*}{\Rightarrow} aba$ , the string can be derived using $G$ .	
Comments (16)	
	Step 9 of 16
h. False	
$T \Rightarrow T$ , the string cannot be derived using $G$ .	
Comment	
	Step 10 of 16
i. True	
$T \stackrel{*}{\Rightarrow} T$ , the string can be derived using $G$ .	
Comments (15)	
	Step 11 of 16
j. True	
$XXX \Rightarrow aba$ , the string can be derived using $G$ .	
Comment	
	Step 12 of 16
k. False	
$\stackrel{*}{X} \Rightarrow aba$ , the string cannot be derived using $G$ .	
Comments (2)	
	Step 13 of 16
I. True	
$T \stackrel{\bullet}{\Rightarrow} XX$ , the string can be derived using $G$ .	
Comments (2)	
**	
	Step 14 of 16
m. Truo	
m. True	
$T \stackrel{*}{\Rightarrow} XXX$ , the string can be derived using $G$ .	
Comments (1)	

n. False		
$S \stackrel{*}{\Rightarrow} \varepsilon$ , the string	cannot be derived using $G_{\cdot}$	
Comment		
	<b>Step 16</b> of 16	
o. $L(G)$ consists	of all strings that are not palindromes and are formed over terminal symbols $\it a$ and $\it b$ .	
Comments (3)		