Problem

Show that if P = NP, a polynomial time algorithm exists that takes an undirected graph as input and finds a largest clique contained in that graph. (See the note in Problem 7.38.)

Step-by-step solution

Step 1 of 2

Class - P: P is a class of Languages that are decidable in polynomial time on a deterministic single tape Turing machine.

 $\textbf{Class-NP}: \ NP \text{ is a class of Languages that are decidable in polynomial time on a nondeterministic Turing machine.}$

Clique: A clique in an undirected graph is a sub graph, wherein every two nodes are connected by an edge.

- If P = NP then we have to show that "a polynomial time algorithm exists that takes an undirected graph as input and finds the largest clique in the graph.
- A k clique is a clique that have k-nodes.
- ${f \cdot}$ We know that "clique is in $N\!P$ ".
- So if P = NP then clique is in P.

Therefore, if P = NP then clique is recognizable in polynomial time.

Comment

Step 2 of 2

The following algorithm will find the largest clique in the graph:

- 1. Let n be the no. of nodes in the given graph G.
- i be the variable which runs from 1 ton.
- 2. Using the polynomial time algorithm form clique, check whether there exist a clique of size i.
- 3. Output the Large $\,i$ for which a clique exists.

To find the maximum clique, we start with i, the maximum clique size.

Remove one node and see if there is still a clique of size i.

If not, restore that node and remove another node.

If so, respect the process until we are left with a graph of $\,i\,$ nodes, which must be a clique.

• This algorithm will take almost "n trials to find which node to remove and at most "n nodes to be removed.

Then the total running time is polynomial.

Comment