

Problem

$$A \subseteq C \text{ and } B \subseteq \overline{C}.$$

Let A and B be two disjoint languages. Say that language C *separates* A and B if disjoint Turing-recognizable languages that aren't separable by any decidable language.

Describe two

Step-by-step solution

Step 1 of 2

Let A and B are two disjoint languages. Consider that language C , that separates A and B if $A \subseteq C$ and $B \subseteq \overline{C}$. It can be proved as follows:

- It is quite easy to understand this statement. Read the statement carefully things are pretty obvious all you need to understand is $A \subseteq C$ and $B \subseteq \overline{C}$.
- Here, A is a set of C and B is a set of Complement of C . So, it is quite obvious now A and B both are not related to each other anyhow.
- Here, no use of C because even if we use C it won't be able to prove decidability of A on the basis of B and Turing reducibility of A on the basis of B . If it comes to languages those are fully different and belong to different sets then no separators are required. So, here it is worthless to use C as separator.

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Step 2 of 2

Now, consider a Turing machine T that will work as a decider for the language C that separates A and B . **Consider both languages as Regular Expressions that will be decided by M_1 and M_2 .**

1. $S = \langle M, w \rangle$ Where M is a Turing machine.

2. Now, run $\langle M_1, W \rangle$ and $\langle M_2, W \rangle$

3. If M_1 accepts then M rejects and if M_2 accepts M rejects.

- **Remember M will always halt in each situation. Where C decides A or B . Now it is pretty easy to understand the situation.**
- Therefore, it can be said that only the first statement of question is enough to prove the concept "No decidable languages can be used to separate two disjoint Turing recognizable languages".

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