

Practice Midterm Exam (Solutions)

(9:45 AM – 11:00 AM : 75 Minutes)

- This exam will account for 30% of your overall grade.
- There are six (6) questions, worth 75 points in total. Please answer all of them.
- This is a *closed book, closed notes* exam. *No cheat sheets* are allowed.
- You are allowed to *use scratch papers* for your calculations.
- You are *not allowed to use your own calculator*. A scientific calculator will be available inside the Respondus Lockdown Browser.

GOOD LUCK!

Question	Parts	Points
1. DFA Construction	(i), (ii)	10 + 10 = 20
2. DFA Composition	–	10
3. Regular Expressions	(i), (ii)	5 + 5 = 10
4. NFA to DFA	–	15
5. Non-regularity	–	15
6. Context-free Grammar	–	5
Total		75

QUESTION 1. [20 Points] DFA Construction. Write down a DFA in the 5-tuple form to accept each of the following two regular languages.

Assume that $\Sigma = \{a, b\}$.

Your answers do not need to include DFA diagrams (though you may draw them on your scratch papers if you like).

(a) [10 Points] $L = \{w \mid aaba \text{ is a substring of } w\}$

(b) [10 Points] $L = \{w \mid w \text{ does not contain } ab\}$

SOLUTION.

(i) 5-tuple, $M = \{Q, \Sigma, q, F, \delta\}$, where,

Set of states, $Q = \{q_0, q_1, q_2, q_3, q_4\}$

Set of symbols, $\Sigma = \{a, b\}$

Start state, $q = q_0$

Set of accept states, $F = \{q_4\}$

Transition function, δ :

state	a	b
q_0	q_1	q_0
q_1	q_2	q_0
q_2	q_2	q_3
q_3	q_4	q_0
q_4	q_4	q_4

(ii) 5-tuple, $M = \{Q, \Sigma, q, F, \delta\}$, where,

Set of states, $Q = \{q_0, q_1, q_2\}$

Set of symbols, $\Sigma = \{a, b\}$

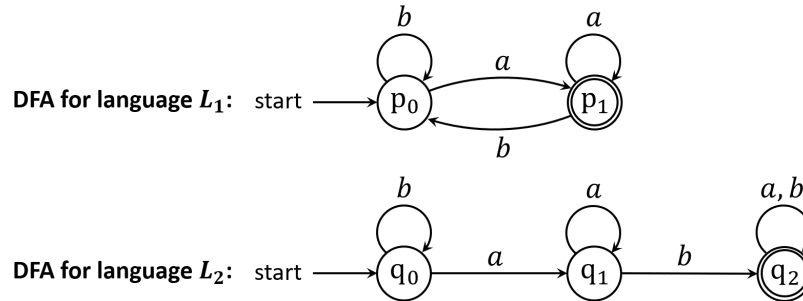
Start state, $q = q_0$

Set of accept states, $F = \{q_0, q_2\}$

Transition function, δ :

state	a	b
q_0	q_1	q_0
q_1	q_1	q_2
q_2	q_2	q_2

QUESTION 2. [10 Points] DFA Composition. Consider the following two DFAs.



Write down a DFA in the 5-tuple form that accepts the language $L_1 \cup L_2$.

Your answer does not need to include the DFA diagram (though you may draw it on your scratch papers if you like).

SOLUTION.

5-tuple, $M = \{Q, \Sigma, q, F, \delta\}$, where,

Set of states, $Q = \{ (p_0, q_0), (p_0, q_1), (p_0, q_2), (p_1, q_0), (p_1, q_1), (p_1, q_2) \}$

Set of symbols, $\Sigma = \{ a, b \}$

Start state, $q = (p_0, q_0)$

Set of accept states, $F = \{ (p_0, q_2), (p_1, q_0), (p_1, q_1), (p_1, q_2) \}$

Transition function, delta:

state	a	b
(p0, q0)	(p1, q1)	(p0, q0)
(p0, q1)	(p1, q1)	(p0, q2)
(p0, q2)	(p1, q2)	(p0, q2)
(p1, q0)	(p1, q1)	(p0, q0)
(p1, q1)	(p1, q1)	(p0, q2)
(p1, q2)	(p1, q2)	(p0, q2)

QUESTION 3. [10 Points] Regular Expressions. Write down regular expressions for the following two languages.

(a) [5 Points] $L = \{w \mid n_a(w) = 2 \text{ or } n_b(w) = 2\}, \Sigma = \{a, b, c\}$

(b) [5 Points] $L = \{w \mid \text{binary number } w \text{ is not divisible by } 8\}, \Sigma = \{0, 1\}$

SOLUTION.

In the solutions below the letter ‘U’ represents the union operation.

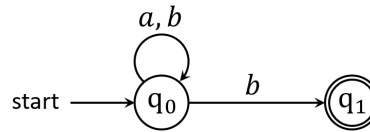
(i) Let $A = a \cup c$
 $B = b \cup c$

Regular expression = $B^*aB^*aB^* \cup A^*bA^*bA^*$

(ii) Let $S = \text{Sigma}$

Regular expression = $1 \cup 01 \cup 1S \cup S^*(SS1 \cup S1S \cup 1SS)$

QUESTION 4. [15 Points] NFA to DFA. Consider the following NFA.



Convert this NFA into a DFA and write it down in the 5-tuple form.

Your answer does not need to include the DFA diagram (though you may draw it on your scratch papers if you like).

SOLUTION.

5-tuple, $M = \{Q, \Sigma, q, F, \delta\}$, where,

Set of states, $Q = \{ q_0, (q_0, q_1) \}$

Set of symbols, $\Sigma = \{ a, b \}$

Start state, $q = q_0$

Set of accept states, $F = \{ (q_0, q_1) \}$

Transition function, δ :

state	a	b
q_0	q_0	(q_0, q_1)
(q_0, q_1)	q_0	(q_0, q_1)

QUESTION 5. [15 Points] Non-regularity. Use the pumping lemma to prove that the following language is not regular.

$$L = \{w \mid w = a^m b^n, 0 \leq m < n\}, \quad \Sigma = \{a, b\}$$

SOLUTION.

Suppose L is regular. Then it must satisfy the pumping property.

Let s be the number of states in any finite state machine that accepts L .

Let $w = a^s b^{(s+1)} = xyz$, where, $x = a^p$, $y = a^q$, and $z = a^r b^{(s+1)}$,

such that $|xy| \leq s$, $|y| \geq 1$, and $p + q + r = s$.

Then $xy^i z$ must belong to L for all $i \geq 0$.

$$\begin{aligned} \text{Now } xyyz &= a^p a^q a^q a^r b^{(s+1)} \\ &= a^{(p+q+r+q)} b^{(s+1)} \\ &= a^{(s+q)} b^{(s+1)}. \end{aligned}$$

But since $|y| = q \geq 1$, we have $s + q \geq s + 1$.

So, $xyyz = a^{(s+q)} b^{(s+1)}$ is not in L , which is a contradiction!

Hence, L is not regular.

QUESTION 6. [5 Points] Context-free Grammar. Write down a context-free grammar to accept the following language:

$$L = \{w \mid w = a^{n+1}b^{n+2}, n \geq 0\}, \quad \Sigma = \{a, b\}$$

SOLUTION.

S \rightarrow aSb

S \rightarrow abb