

Problem

Show that the collection of Turing-recognizable languages is closed under the operation of

- Aa. union.
- b. concatenation.
- c. star.
- d. intersection.
- e. homomorphism.

Step-by-step solution

Step 1 of 8

- Suppose, X and Y be two Turing recognizable languages that have Turing machines M_X and M_Y respectively.
- Now the union of these languages is denoted by L_{XY} and the Turing machine recognizing this language is M_{XY} .

"On input w :"

1. Run X and Y alternately on w step by step. If either accepts, *accept*. If both halt and reject, *reject*.

Suppose s be a word from L_{XY} . M_{XY} works for an input string s as shown:

- It executes M_X and M_Y on s individually.
- If at least any one of M_X or M_Y accepts s then M_{XY} also accepts after a finite number of steps and reach to its accepting state.
- If both M_X and M_Y reject and either of them do so by looping then M_{XY} will loop.

[Comment](#)

Step 2 of 8

Hence, it can be said that collection of Turing recognizable languages is closed under union operation.

[Comment](#)

Step 3 of 8

- Suppose, X and Y be two Turing recognizable languages that have Turing machines M_X and M_Y respectively.
- Now the concatenation of these languages is denoted by L_{XY} and the Turing machine recognizing this language is M_{XY} .

Let s be a word from L_{XY} . M_{XY} works for an input string s as shown:

- It divides each string of XY into s_1 and s_2 non-deterministically.
- It runs s_1 to M_X . If M_X halts and rejects, *reject*.
- If runs s_2 to M_Y . If M_Y accepts, *accepts*. If M_Y halts and rejects, *reject*.

[Comments \(2\)](#)

Step 4 of 8

Hence, it can be said that collection of Turing recognizable languages is closed under concatenation operation.

[Comment](#)

Step 5 of 8

• Suppose, X be a Turing recognizable language and the Turing machine is M_X .

• Now X^* is the language obtained from star operation on X .

The Turing machine for this language be M_{X^*} .

M_{X^*} works as follows:

• For an input string s of X , it non-deterministically divided the string into $s_1, s_2 \dots s_n$.

• For each of those divided parts M_X runs, Suppose, M_X all divided parts then s is accepted by M_{X^*} else s is rejected by M_{X^*} .

[Comment](#)

Step 6 of 8

Hence, it can be said that collection of Turing recognizable languages is closed under star operation.

[Comment](#)

Step 7 of 8

• Suppose, X, Y be two Turing recognizable languages that have Turing machines M_X and M_Y respectively.

• Now the intersection of these languages is denoted by L_{XY} and the Turing machine recognizing this language is M_{XY} .

For an input string s from L_{XY} , M_{XY} works for as shown:

• Turing machine M_X runs on s . If it accepts s then M_Y runs on s . Else s is rejected.

• Suppose, M_Y accepts s then it is accepted by the Turing machine otherwise s is rejected.

Hence, it can be said that collection of Turing recognizable languages is closed under intersection operation.

[Comments \(5\)](#)

Step 8 of 8

• Suppose, X a Turing recognizable language that have Turing machine M_X .

• To recognize $h(X)$ the other Turing machine M_Y is simulated in such a way that:

On input s , it will consider all strings w such that $h(w) = s$.

• The $TM M_X$ will execute on input w by going through all strings in w .

If $h(w) = s$ start executing M_X on input w , using merging to interleave with other executions on M_X . *Accept* if any executions accept.

• M_Y will accept s if any of those executions of M_X accepts s . Else s will be rejected.

Hence, it can be said that collection of Turing recognizable languages is closed under homomorphism operation.

[Comment](#)