# **Problem**

Show that a language is decidable iff some enumerator enumerates the language in the standard string order.

# Step-by-step solution

# Step 1 of 3

To prove this problem:

We need to show equivalence between a Turing machine that decides a language and an enumerator that enumerates it. Hence we have to show the proof in both directions i.e.,

- 1. Language is decidable then enumerator enumerates the language in lexicographic order.
- 2. Enumerator enumerates the language in lexicographic order and then it is decidable.

Comment

#### Step 2 of 3

1. Language is decidable then enumerator enumerates the language in lexicographic order.

#### Proof:

# If direction:

Let us assume that we have a Turing machine M to decide a language L.

Now we can use this M to construct an enumerator E as follows.

We generate the strings in the lexicographic order and input each string into  $\emph{M}$  for  $\emph{L}$ .

If *M* accepts then print that string. Therefore *E* prints all strings of *L* in lexicographic order.

Comment

## **Step 3** of 3

 $2. \ Enumerator \ enumerates \ the \ language \ in \ lexicographic \ order \ and \ then \ it \ is \ decidable.$ 

## Proof:

## And if direction: -

Now we need to consider the other direction.

That is, if we have an enumerator E for a language L, then we can use E to construct a Turing machine M that decides L.

Here we have to consider two cases.

# Case (i): If L is finite:

If L is finite language, then it is decidable, because all finite languages are decidable.

# Case(ii) If L is infinite:-

If L is infinite then a decider for L operates as follows.

- On receiving the input w, the decider enumerates all strings of L in lexicographic order until a string greater than w appears in the lexicographic order.
- This must eventually occur since L is infinite.
- If w has appeared in the enumeration already, then accept.
- If w has not yet appeared in the enumeration then it will never appear and hence we can reject.

So in both cases L is decidable. The theorem proved in both directions.

Therefore,

A language is decidable if and only if some enumerator enumerates the language in lexicographic order.

Comment