

## Problem

$\langle N \rangle$

Describe two different Turing machines,  $M$  and  $N$ , where  $M$  outputs  $\langle N \rangle$ , when started on any input.

## Step-by-step solution

### Step 1 of 2

Given: There are two different Turing machines  $M$  and  $N$ .

We have to prove that, it started with any input  $w$ ,  $M$  outputs  $\langle N \rangle$  and  $N$  outputs  $\langle M \rangle$

Here we need to know the recursion theorem.

**Recursion theorem:** Let  $T$  be a Turing machine that computes a function  $t: \Sigma^* \times \Sigma^* \rightarrow \Sigma^*$ . There is a Turing machine  $R$  that computes a function  $r: \Sigma^* \rightarrow \Sigma^*$ , where for every  $w$ ,

$$r(w) = t(\langle R \rangle, w).$$

To make a Turing machine that can get its own description and computes with it. Machine  $T$  in the statement receives the description of machine as extra input. The recursion theorem produces a new machine  $R$ , which operates exactly  $T$ , does but  $R$ 's description filled automatically.

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### Step 2 of 2

By using Recursion theorem, remember that  $r(w)$  is a Turing machine that prints  $w$  on its tape, and halts.

The Turing machine  $M$  is as follows:

$M =$  " on input  $w$ :

1. No need to note the input
2. Obtain  $\langle M \rangle$  by using the recursion theorem.
3. Compute  $r(\langle M \rangle)$  where  $r$  the computable function is write  $r(\langle M \rangle)$  on the tape.
4. Halt."

• When  $M$  runs, it writes  $N = \langle r(\langle M \rangle) \rangle$  in its tape.

• When  $N$  runs, it writes  $\langle M \rangle$  in its tape automatically.

$\langle N \rangle$  is a member of the turing machine  $M$

$\langle M \rangle$  is a member of the turing machine  $N$

Therefore  $M$  outputs  $\langle N \rangle$  and  $N$  outputs  $\langle M \rangle$ .

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