

### Problem

The pumping lemma says that every regular language has a pumping length  $p$ , such that every string in the language can be pumped if it has length  $p$  or more. If  $p$  is a pumping length for language  $A$ , so is any length  $p' \geq p$ . The **minimum pumping length** for  $A$  is the smallest  $p$  that is a pumping length for  $A$ . For example, if  $A = 01^*$ , the minimum pumping length is 2. The reason is that the string  $s = 0$  is in  $A$  and has length 1 yet  $s$  cannot be pumped; but any string in  $A$  of length 2 or more contains a 1 and hence can be pumped by dividing it so that  $x = 0$ ,  $y = 1$ , and  $z$  is the rest. For each of the following languages, give the minimum pumping length and justify your answer.

<sup>A</sup>a.  $0001^*$

<sup>A</sup>b.  $0^*1^*$

c.  $001 \cup 0^*1^*$

<sup>A</sup>d.  $0^*1^+0^+1^* \cup 10^*1$

e.  $(01)^*$

f.  $\epsilon$

g.  $1^*01^*01^*$

h.  $10(11^*0)^*0$

i.  $1011$

j.  $\Sigma^*$

### Step-by-step solution

#### Step 1 of 10

a.

The pumping length could not be 3 as 000 being in the language it cannot be pumped. Consider the string length be 4 or more and divide in  $xyz$  as  $x$  being 000,  $y$  being the first 1 and  $z$  being everything after then it satisfies every condition of pumping lemma.

Hence the minimum pumping length is 4.

[Comment](#)

#### Step 2 of 10

b.

String  $\epsilon$  is in the language but it could not be pumped so the pumping length could not be 0. According to pumping lemma's three condition arises which are as follows:

- If user divide the string in  $xyz$  as  $x$  is  $\epsilon$
- $y$  is first symbol (0|1)
- $z$  being the everything after then it holds.

Hence the minimum pumping length is 1.

[Comment](#)

#### Step 3 of 10

c.

The string  $001$  is in the language but if it is generated by  $001$  then it cannot be pumped. If string  $s$  is larger than 3 and in the language, then it is generated by  $0^*1^*$ . Dividing the string according to Pumping Lemma's condition into a string  $xyz$  where  $x$  is  $\epsilon$ ,  $y$  be the first symbol and  $z$  be the remaining, user can pump the string.

**Hence the minimum pumping length is 4.**

[Comment](#)

#### Step 4 of 10

d.

The string  $11$  is in the language but it cannot be pumped so the length is not 2. Let  $s$  be the string in the language of length at least 3. If  $s$  is generated by  $0^*1^*0^*1^*$  it can be divided in  $xyz$  as  $x$  is  $\epsilon$ ,  $y$  is the first string and  $z$  is everything else so it can be pumped.

Again if  $s$  is generated by  $10^*1$  then also user can write it as  $xyz$  where  $x$  is  $1$ ,  $y$  is  $0$ , and  $z$  is the remainder so it could be pumped.

**Hence the minimum pumping length is 3.**

[Comment](#)

#### Step 5 of 10

e.

Let  $s$  be a string in the language.

Now  $s$  could be  $\epsilon$  but it cannot be pumped so the length is not 0. Next  $s$  could be  $01$  which if divide in  $xyz$  as  $x$  is empty string  $\epsilon$ ,  $y$  is  $01$ , and  $z$  is everything after then it satisfies the three conditions of pumping lemma.

Pumping length is not 1 because since there is no string of length 1 in the language.

**Hence the minimum pumping length is 2.**

[Comment](#)

#### Step 6 of 10

f.

Let  $s$  be a string in the language then  $s$  is  $\epsilon$  and according to pumping lemma it cannot be pumped. As per the pumping lemma, pumping length should greater then equal to 1.

**Hence the minimum pumping length is 0.**

[Comments \(1\)](#)

#### Step 7 of 10

g.

The minimum pumping length of the language could not be 2 as  $00$  being in the language it could not be pumped. Let  $s$  be the string of length at least 3 in the language so minimally  $s$  could be  $100$  or  $010$  or  $001$ .

Now dividing  $s$  in  $xyz$  in all the three cases user get, for  $100$   $x$  is  $\epsilon$ ,  $y$  is  $1$ , and  $z$  is the remainder, for  $010$   $x$  is  $0$ ,  $y$  is  $1$ , and  $z$  is the remainder and for  $001$   $x$  is  $00$ ,  $y$  is  $1$ , and  $z$  is the remainder. All satisfies Pumping Lemma's three conditions.

**Hence the minimum pumping length is 3.**

[Comment](#)

#### Step 8 of 10

h.

The minimum length string in the given language is  $100$  but it cannot be pumped. Next minimum length string is  $10100$ . If divide it according to the pumping lemma in  $xyz$  then  $x$  be  $10$ ,  $y$  be  $10$ , and  $z$  be the rest of the string then  $y$  can be pumped.

Hence the minimum pumping length is 4.

[Comment](#)

Step 9 of 10

i.

If  $s$  be a string in the language then  $s$  is  $1011$ . If set  $p = 4$ , then claim  $s$  is pumpable (which it is not, as it is the only string in the language). This should be 5.

Hence then the minimum pumping length is 5.

[Comment](#)

Step 10 of 10

j.

Say  $s$  be a string in the language  $\Sigma^*$ . According to pumping lemma if divide  $s$  in  $xyz$  then  $x$  be the empty string,  $y$  is  $(\varepsilon|0|1)$  and  $z$  is empty string. Now  $\varepsilon$  could not be pumped.

Hence the minimum pumping length is 1.

[Comment](#)