CSE303: Theory of Computation, Fall 2021 Date: 11/08/2021

 $\underset{(\text{ Due: }\underline{11/12/2021}\underline{)}}{\text{Homework }\#\underline{3}\underline{\ }}$

Group Number: __33

Group Members		
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CSE303: Homework #3

Taskit

(1) (a) 5→ a5cc | A

A -> bAccc /E

(b) 5-> XY/UY X -> AD DB

Y->cYIE

U-> aU/E V-> BE | EC

1 D→ aDblE

E-> bEcle

A -> aA | a $B \rightarrow bB \mid b$

C-> cC/c

(c) 5-> asal 1 apr | bad R

P-> aPc| bRc|E

Q->bad|bRc|E

C-> bRelE

(d) $S \rightarrow AB$

A -> a A a | b A b | c B

B-> aB|B|cB|E

(e) $S \rightarrow AB$ $A \rightarrow aABa|bABb|cB$ $B \rightarrow aB|bB|cB|c$ (2) (a) 5→ b5 | aA A-> bA | aB $B \rightarrow aB \mid bB \mid \epsilon$ (b) 5 → aA 1 bC A-> aclbB $B \rightarrow aB | bB | \epsilon$ C-> aclbc (c) 5 → aA | bB | cC A -> aAlbD|cD|E B→ aE/bB/cE/€ C→ aF/bF/cC/E D-> aA | bD | cD E- aE | bB | cE F-> aF | bF | cC (d) 5- aA/bC A-> aB | DD | E B-> a5/bE C-> aG/bC D-> aH/bD/E E-> aF/bE F-> aG/bF/E G-> aHIBGIE H-> aFIBHIE

(e)
$$S \rightarrow 0A|1B|E$$
 $A \rightarrow 0A|1B|E$
 $B \rightarrow 0C|1D$
 $C \rightarrow 0E|1F$
 $D \rightarrow 0G|1A$
 $E \rightarrow 0B|1C$
 $F \rightarrow 0D|1E$
 $G \rightarrow 0F|1G$

(a) (a) $S \rightarrow 0S|1A1S|E$
 $A \rightarrow 0B0A|E$
 $B \rightarrow 1B|E$

(b) $S \rightarrow E|A|ABBC$
 $A \rightarrow AA|a$
 $B \rightarrow BB|B$
 $C \rightarrow aBC|E$

(c) $S \rightarrow a|aa|ab|aba|abb$

(d) $S \rightarrow ABCCCC$
 $A \rightarrow BBBBBBA|E$
 $C \rightarrow B|E$

(4) (6) Suppose L is CFL. Then it must satisfy pumping property. Suppose s = a b a a long to L for all i 70 We will show that $\mu xz \notin L$ for all possible cases. Two cases:

Case L: vxy consists of exactly 1 symbol, (a's or b's)

Subcase 1: vxy consists only of a's.

Let s = $\mu v xyz = a'b^2 p a'$ μxz is not in L

Reason: $\mu xz = a' - (|v| + |y|) |_{z}^{2} p' a' \notin L$ as (|v| + |y|) > 0 μxz has fewer a's than b's.

Subcase 2: vxy consists only of b's Let SE = uvxyz = a'b2'a' uxz is not in L Reason: uxz = a | 2p - (1v1+1y1) a # # L as (1v1+1y1) >0 uxz has fewer b's than a's. Case 2: vxy consists of exactly 2 symbols, (ab's or ba's) Subcase 1: vxy consists of ab's. Let s= uvxyz = a'b' c' uxz is not in, L Reason: uxz = ak, bk2 ak £L where R, + R2 = 3p-(14/4/41)<3p as (14/4/41)>0 Either uxz has fewer b's than a's (when R=P) or uxz has fewer as before b's than often b's. Subcase 2: yxy consists of ba's Similar to subcase 1 Subcase 3: VXY consists of aba's This subcase is impossible as |VXY| > 2p+2>p Contradiction! Hence L is not CFL. (b) Suppose L is CFL. Then it must satisfy pumping property. Suppose we = a b e and p = 25 be the pumping length.

Let $w = u v x y z = a^{-1}ab \in e^{\pm}$ where |vy| = 1 > 1+7870 +7870 and 1/x/y/= s+1<p=23 Then uvicy'z must belong to I for all i>0

Choosing i=t-x+2, we get

uvixyiz = a t-x+2 b \in e

= a t+1 b e t \in t

Reason: uvixyiz has fewer e's than a's

Contradiction! Hence L is not CFL.

Suppose L is CFL. Then, it must satisfy pumping property. Suppose $s=a^p!$ Let s=uvxuz where $|vxy| \le p$ and $|vy| \ge 1$ Then uvxyz must belong to L for all i>0But $uvxyz \ne L$ Reason: Let |vy| = R. Then, $1 \le k \le p$ $uvxyz = a^{p!+k} \ne L$ Because, $p! < p! + k \le p! + p < p! + p! p = p! (1+p) = (p+1)!$ Contradiction! Hence, L is not CFL.