

Problem

Show that any infinite subset of MIN_{TM} is not Turing-recognizable.

Step-by-step solution

Step 1 of 3

Definition of MIN_{TM} =

If M is a Turing machine, then we say that the length of the description $\langle M \rangle$ of M is the number of symbols in the string description M .

Say that M is minimal if there is no Turing machine equivalent to M that has a shorter description.

$$MIN_{TM} = \{ \langle M \rangle \mid M \text{ is a minimal TM} \}$$

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Step 2 of 3

Now we have to prove that any infinite subset of MIN_{TM} is not Turing – recognizable.

We will prove by taking contradiction.

We assume that there exists A , an infinite subset of MIN_{TM} , such that A is Turing – recognizable.

We know that

“A language is Turing – recognizable if and only if some enumerator enumerates it”.

So, Let E be the enumerator that enumerates A .

By using this E , we construct another TM (Turing – machine) N as follows.

N = “On input w :

1. from recursion theorem. Own description $\langle c \rangle$ is obtained
2. Run the Enumerator E until a machine P is obtained with a longer description than that of N .
3. Simulate P on input w ’.

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Step 3 of 3

As we know that MIN_{TM} is infinite A is infinite subset of MIN_{TM} .

1. When A is infinite, E ’s list must contain a TM with longer description than N ’s description. So obviously N terminates with some $TM P$ which is longer than N . Then N simulates P and so is equivalent to P .
2. It also notify that N is shorter than P . So P cannot be minimal. But P appears on the list that is produced by E .
3. E ’s list must contain a TM with longer description than N ’s description.

From above three conditions we have a contradiction.

Thus our assumption that A is Turing – recognizable is wrong.

Therefore A is not Turing – recognizable.

Thus an infinite subset of MIN_{TM} is not Turing – recognizable.

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