## Let Extension

Let extension violates the progress rule. In this section, we will provide a counter example.

$$let x': int = Int(1) in y$$

According to (D-Let) rule, we can make a small step, which is

let 
$$x'$$
: int = Int(1) in  $y \mapsto [x' \to Int(1)] y$ 

Observed from the above statement, we know that y is neither **VAL** nor can go any further with a small step (no  $\tau'$  can be the next step of y). Therefore, this example violates the progress theorem.

## **Rec Extension**

Rec extension obeys both the progress rule and the preservation rule. In this section, we will prove this.

• For progress,

What we need to show is, if  $rec(t_0; x. y. t_1)(t)$ :  $\tau$ , then either  $rec(t_0; x. y. t_1)(t)$ :  $\tau$  **VAL** or  $\exists t'$  such that  $rec(t_0; x. y. t_1)(t) \rightarrow t'$ .

1. *t* VAL

By typing inversion for int,  $t = 0 \lor t = n(n \neq 0)$ 

If t = 0, then  $rec(t_0; x. y. t_1)(t) \mapsto t_0$  (D-Rec2), where  $t_0$  is also an int (VAL).

If  $t = n(n \neq 0)$ , then  $rec(t_0; x. y. t_1)(t) \mapsto [x \to n, y \to rec(t_0; x. y. t_1)(n-1)]t_1$  (D-Rec3).

Denote  $[x \to n, y \to rec(t_0; x, y, t_1)(n-1)]t_1$  as t', then we have  $rec(t_0; x, y, t_1)(t) \mapsto t'$ .

2. t is not VAL.

In this case, make a small step according to (D-Rec1)

$$rec(t_0; x. y. t_1)(t) \mapsto rec(t_0; x. y. t_1)(t')$$

Denote  $rec(t_0; x. y. t_1)(t')$  as t'', then we have  $rec(t_0; x. y. t_1)(t) \rightarrow t''$ .

## • For preservation,

What we need to show is, if  $rec(t_0; x.y.t_1)(t)$ :  $\tau$  and  $rec(t_0; x.y.t_1)(t)$ :  $\tau \mapsto t'$ , then t':  $\tau$ .

1. *t* is VAL

By typing inversion for int,  $t = 0 \lor t = n(n \neq 0)$ 

If t = 0, then  $rec(t_0; x, y, t_1)(t) \mapsto t_0$ , where  $t_0: \tau$ , according to D-Rec1.

If  $t = n(n \neq 0)$ , then we can prove by induction.

Base case: the theorem holds for t = 0.

Suppose: the theorem holds for t = n - 1.

We examine the case for t = n.

According to D-Rec3, we have

$$rec(t_0; x. y. t_1)(t) \mapsto [x \to n, y \to rec(t_0; x. y. t_1)(n-1)]t_1$$

where  $rec(t_0; x.y.t_1)(n-1)$ :  $\tau$ . Apply this to T-Rec, then we know  $t_1$ :  $\tau$ . Hence, it holds.

2. t is not VAL

 $t \mapsto t''$ , then by inductive hypothesis,  $rec(t_0; x, y, t_1)(t'')$ :  $\tau$ , according to D-Rec1.