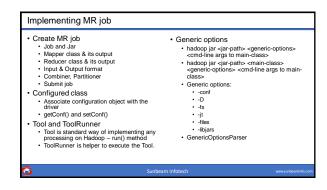


Implementing MR job

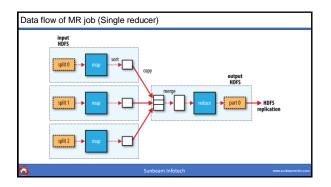
Implement Mapper class
Handle individual record
class MyMapper extends Mapper-Keyln, Valueln, KeyOut, ValueOut> {...}
Override map() method
Input from InputFormat record by record and key-value pair output to merge stage
Implement Reducer class
Perform aggregation on set of values (corresponding to each key)
class MyReducer extends Reducer-Keyln, ValueIn, KeyOut, ValueOut> {...}
Override reduce() method
Input from merge stage in key-values pair and key-value pair output to OutputFormat
Hadoop Writable
Like java wrapper classes, but optimized for serialization over the network.
IntWritable, ByteWritable, ShortWritable, LongWritable, DoubleWritable, BooleanWritable, Text
ArrayWritable, MapWritable, NullWritable

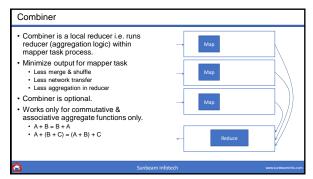


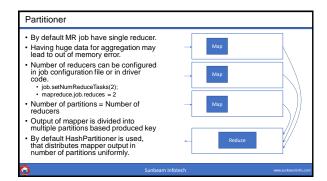
Executing MR

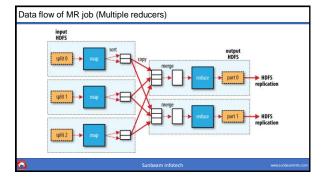
• hadoop jar command
• hadoop jar sjar-paths <generic-options> <cmd-line args to main-class>
• hadoop jar sjar-paths <main-class> <generic-options> <cmd-line args to main-class>
• MR job configurations
• fs.defaultFS = hdfs://namenode:9000/
• mapreduce.framework.name = yarn
• yarn.resourcemanager.ddress = resourcemanager:8032
• Understanding MR summary
• Number of mapper & reducer tasks.
• Number of input & output records for mapper
• Number of input & output records for reducer
• Custom Job counters
• MR log files review
• SHADOOP_HOME/logs

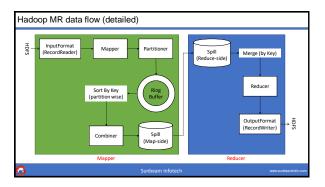
Input/Output Format and Input Splits InputFormat – how to read data · Output format - how to write data · FileInputFormat, TextInputFormat · FileOutputFormat, TextOutputFormat KeyValueTextInputFormat, NLineInputFormat DBOutputFormat • RecordWriter - Write individual record DBInputFormat • Number of output files = Number of RecordReader - Logical division of reducers record · Output written on HDFS (replicated) · Number of mappers = Number of input Number of input splits ≈ Number of HDFS blocks

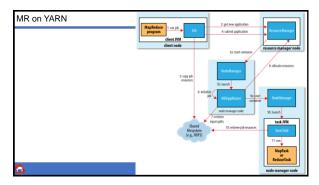


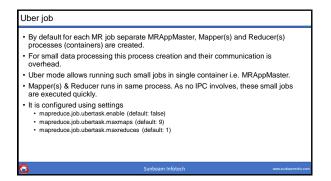


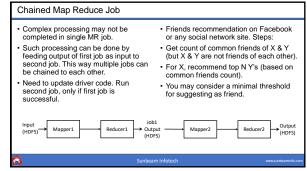


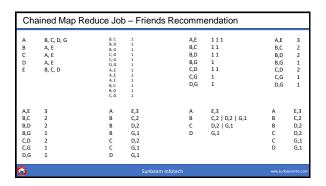


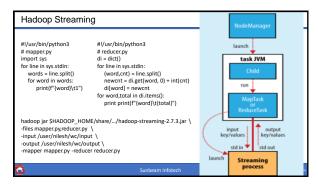


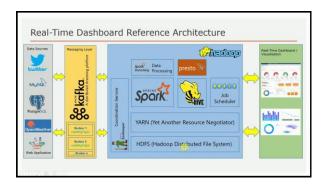




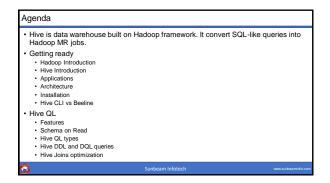


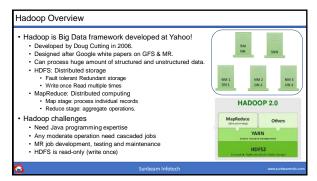


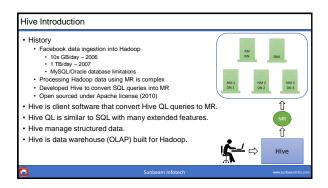


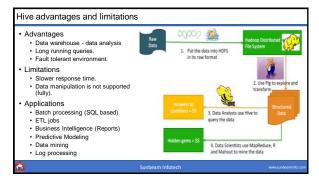


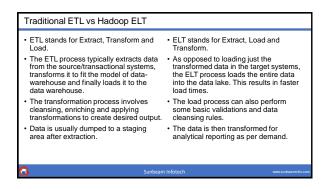


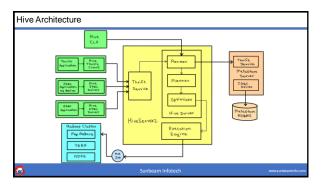


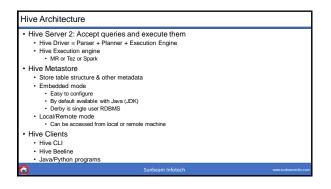


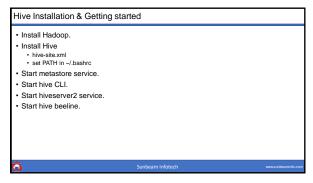


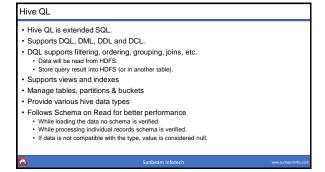


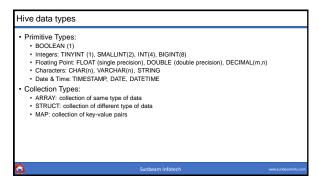


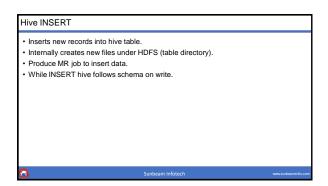


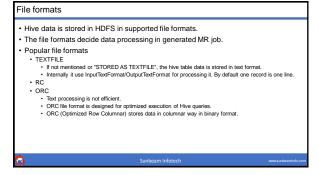


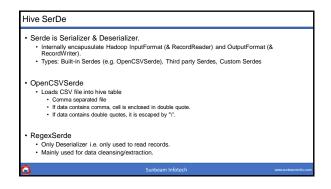


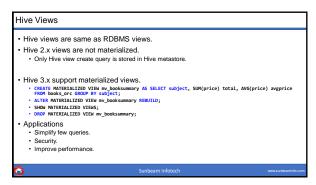


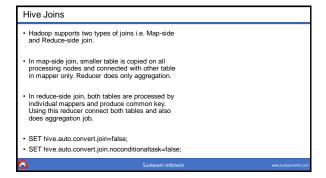


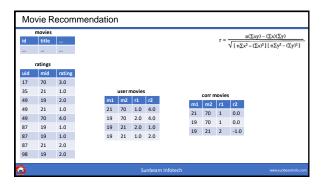


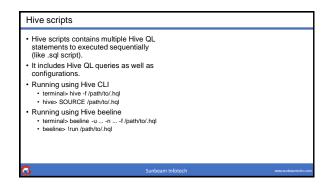


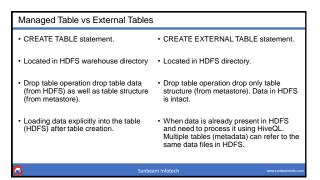


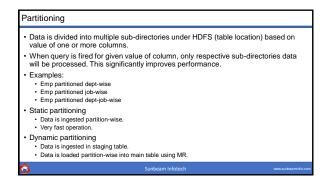


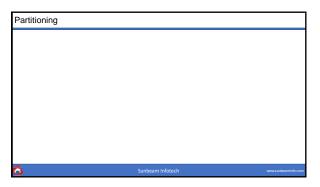












Bucketing

Data in bucketed tables is divided into multiple files.

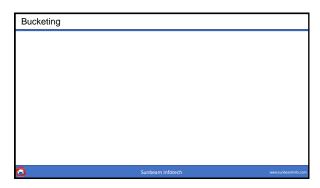
When data is processed using MR job, number of reducers will be same as number of buckets.

To insert data into bucketed table, it must be uploaded via staging table.

Usually buckets are created on unique column(s) to uniformly divide data across multiple reducers.

It provides better sampling and speedup map side joins.

It is mandatory for DML operations.



Hive have many built-in functions.

• Hive have many built-in functions.

• Single Row Functions

• Row → Function → Row

• e.g. LENGTH(), CONCAT(), ROUND(), ...

• Group/Aggregate Functions

• Rows → Function → Row

• e.g. SUM(), AVC(), COUNT(), ...

• Table generation Functions

• Row → Function → Rows

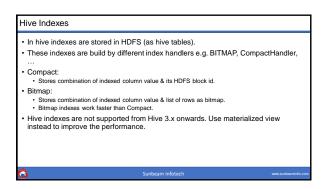
• e.g. EXPLODE(), ...

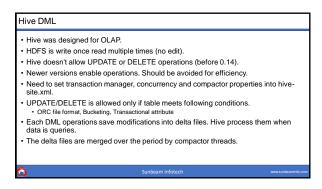
• Hive function help is available in Hive documentation:

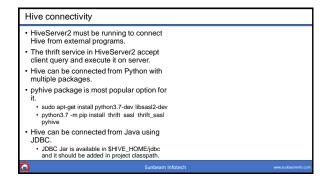
• https://cwiki.apache.org/confluence/display/
Hive/LanguageManual+UDE

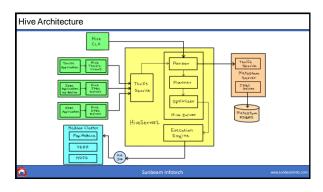
• Hive UDF can be written in Java/Python.

- Similar to RDBMS index.
- To speed up SELECT queries (searching & grouping).
- Indexes internally store addresses of records for given column values.
- Creating index is time-taking job (for huge data). If indexing is done under load, then clients query performance is too low.
- In Hive indexes are created, but deferred for build (using ALTER statement).
- CREATE INDEX query doesn't create index, rather keep ready for building later.
- Index building should be triggered explicitly, when server is less loaded.



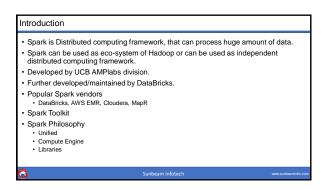


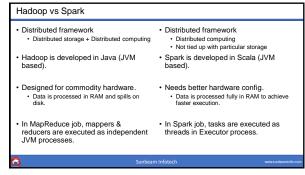




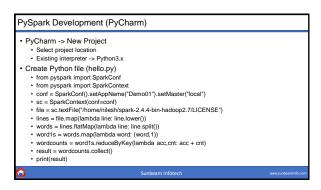


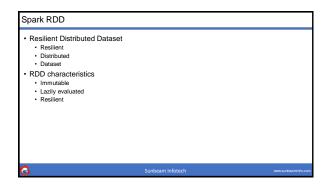


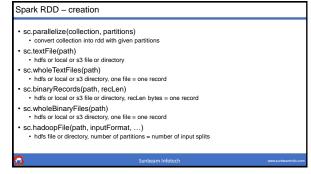


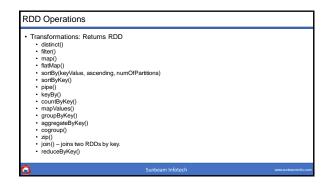


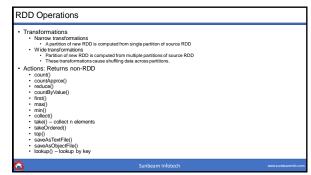


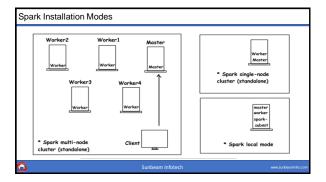


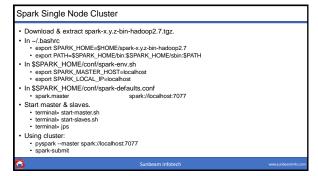


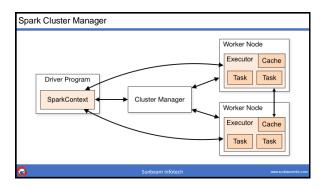


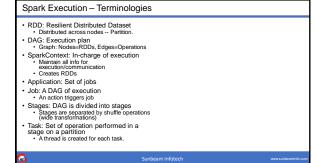


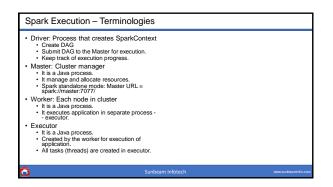


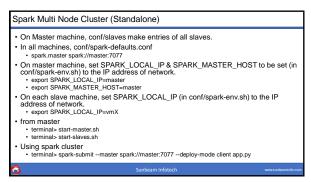


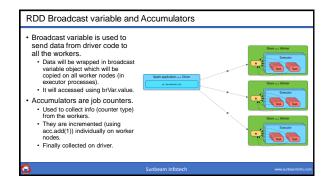


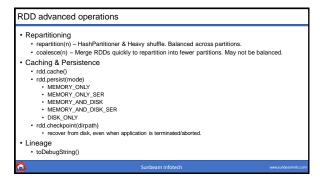




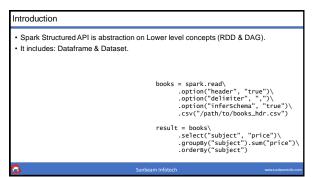


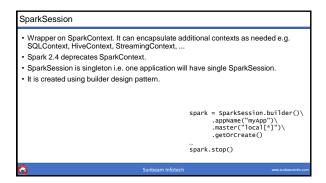


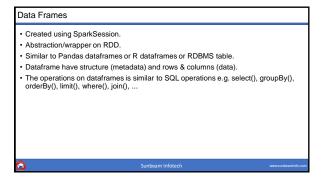


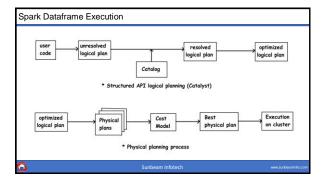


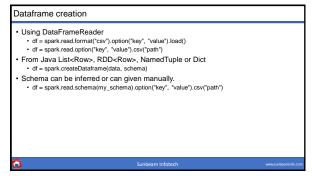


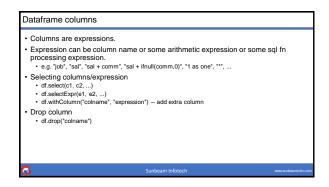


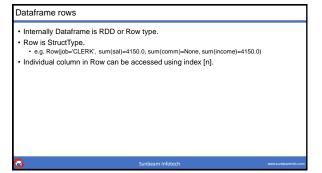


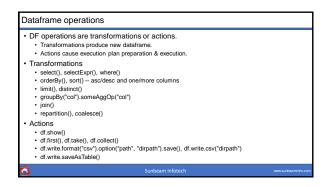


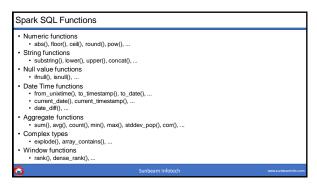


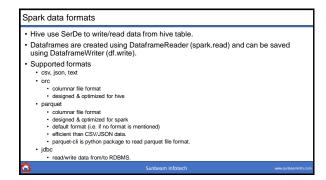




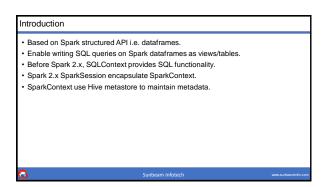


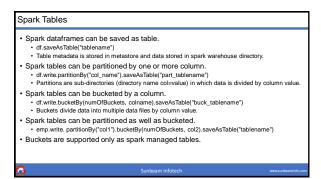


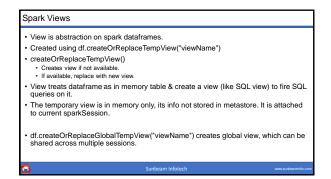




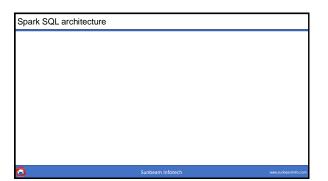


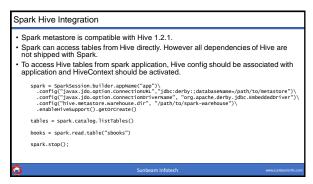




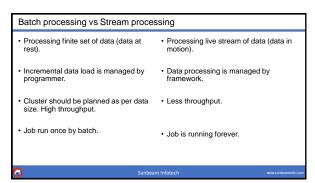


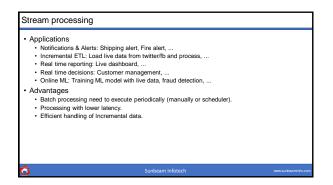


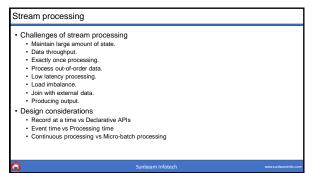


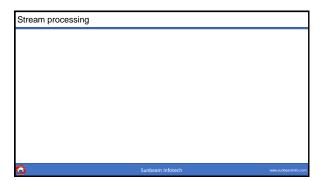


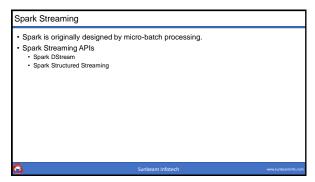


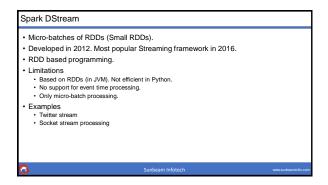


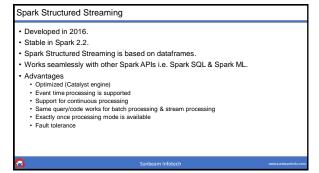


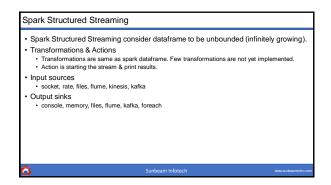


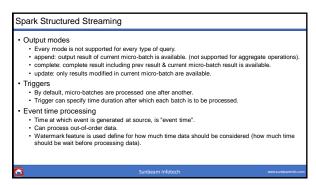


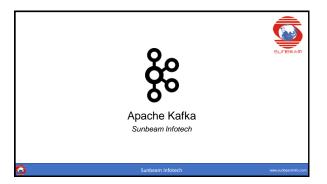


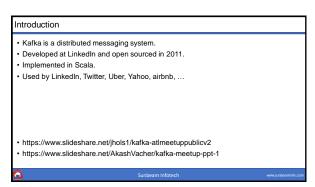


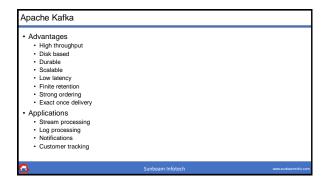


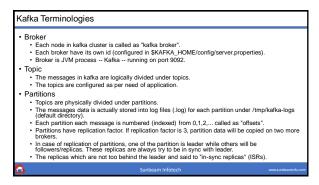


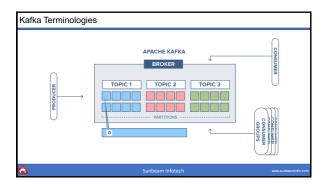


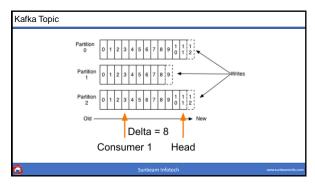










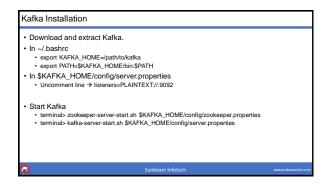


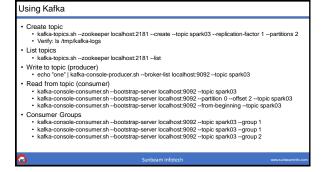
Kafka Terminologies

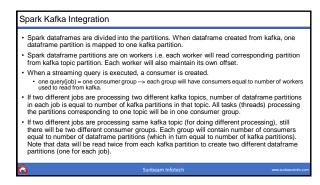
Producer
Client process of kafka which send messages to the broker, is said to be producer.
These client processes can be implemented in Java, Python, C/C++, ...
Producer publish data to the topics, which is sent to partitions, so that partitions are load balanced. It may use round-robin algorithm or may do key-based division.
Producer always write data to the leader of the partition.

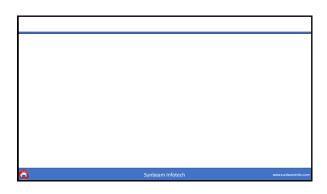
Consumer
Client process of kafka which receive messages from the broker, is said to be consumer.
These client processes can be implemented in Java, Python, C/C++, ...
Clients subscribe to the kafka topics.
It will read the data from any of the replica, which have data required for that client.
The data in a partition is guaranteed to be received in sequence (of writing). However data in different partitions is not guaranteed to read in sequence of writing.

• Consumer groups
• Each topic can have one or more consumers — divided into one or more consumer groups.
• Consumer group – kafka ensure than only one consumer from a consumer group will get a message.
• if there are three consumer groups as follows:
• (a) – C1a, C1b, C1c
• (a) – C2a, C2b
• (3a) – C3a, C3b, C3c, C3d
• (if message "m" of is received into kafka, it will be received by exactly one consumer from each group.
• Possible organization of consumers in a group:
• examplet:
• (a) – C1
• (a) – C2
• (a) – C3
• (b) – C3
• (b) – C3
• (c) –

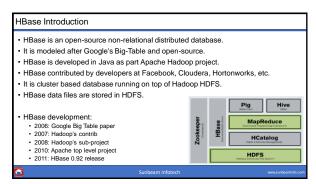


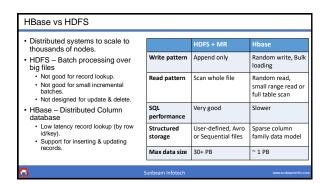


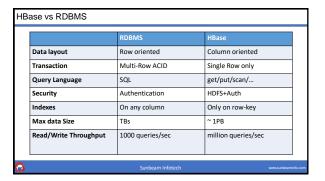


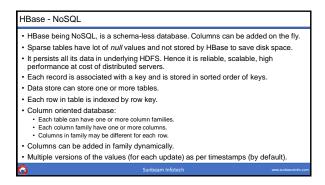


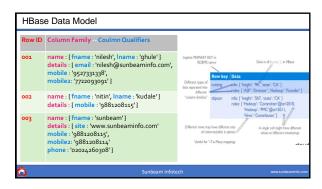


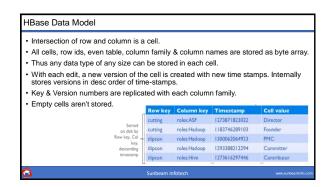


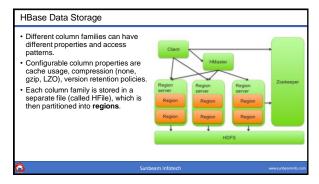


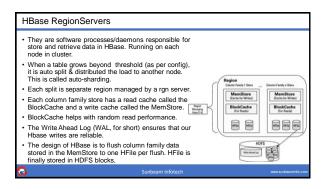


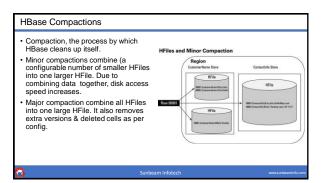


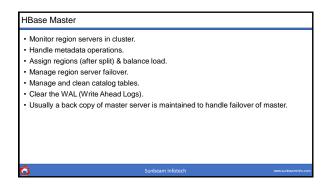


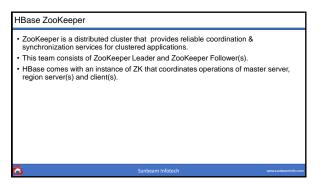


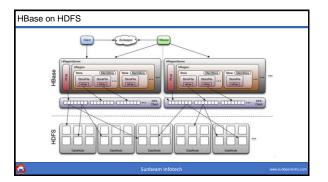


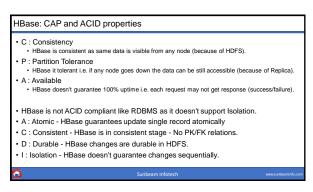


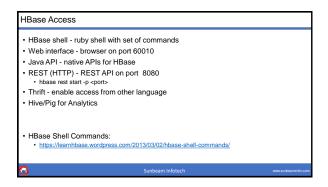


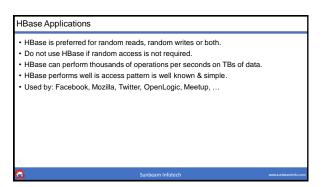






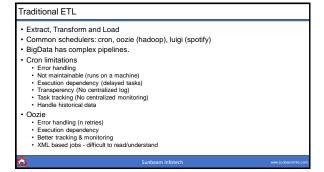


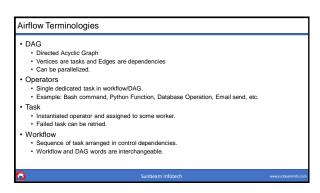


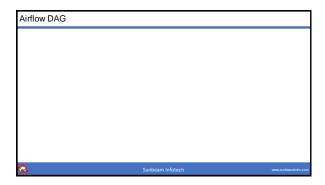


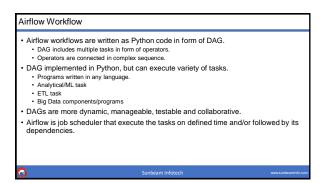


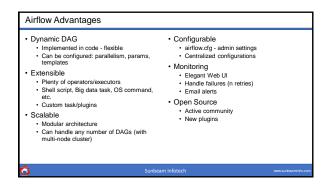


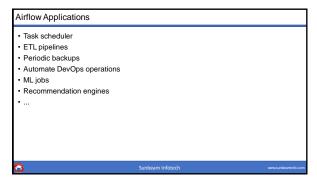


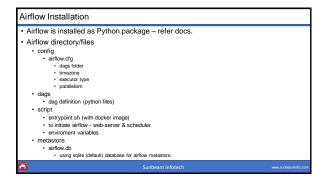


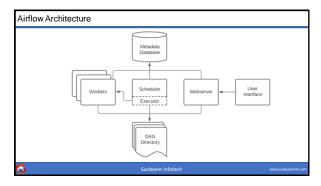


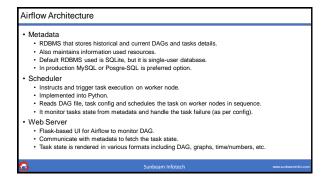


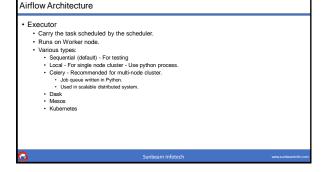


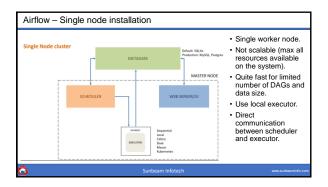


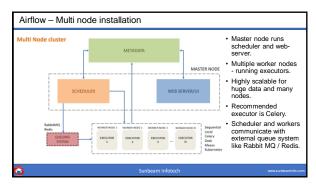


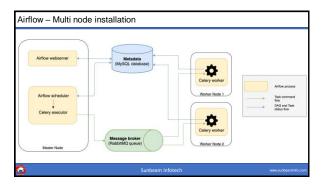












Airflow Working

Scheduler periodically pings for DAG folder and communicate with metastore.

If any DAG is available for execution, scheduler starts a DAG run for it. DAG run is an object representing an instantiation of DAG.

Scheduler update DAG state as "Running".

For each task, task object is instantiated. The task state is updated as "Scheduled".

Scheduler assigns priority to each task in DAG (as per config) and push them into queuing system. The task state is updated to "Queued".

Worker pull the task from queue, set its state as "Running" and start executing it.

Upon completion of each task, task state is updated as succeed or failed.

When all tasks in DAG run are executed, status of DAG run is updated as succeed or failed.

Airflow web-server periodically get the data from meta-store and render it for users.

If any new DAG is found it DAGs folder, scheduler begin its execution (as above).

