

(r.v.s sequence limiting r.v.)

seq. of functions

Defn. (convergence in P) :- for any $\varepsilon > 0$, if $Z_n \xrightarrow{P} Z$, then

$$\lim_{n \rightarrow \infty} P(|Z_n - Z| > \varepsilon) = 0.$$

$$\hookrightarrow b_n = \frac{1}{n}, \quad b_n \rightarrow 0.$$

$$\frac{1}{n} < \varepsilon \quad \text{eventually.}$$

$$* \quad \omega : \lim_{n \rightarrow \infty} Z_n(\omega) = Z(\omega)$$

for any $\varepsilon > 0$, $\left\{ \omega : |Z_n(\omega) - Z(\omega)| < \varepsilon \text{ eventually} \right\}$

$\Rightarrow |Z_n - Z| < \varepsilon$ eventually with probability 1.

* For any ε ,

$$\lim_{n \rightarrow \infty} P\left(\bigcap_{m \geq n} |Z_m - Z| < \varepsilon\right) = 1$$

$$|Z_n - Z| < \varepsilon \text{ AND } |Z_{n+1} - Z| < \varepsilon \text{ AND } |Z_{n+2} - Z| < \varepsilon.$$

* Or for any $\varepsilon > 0$;

$$\lim_{n \rightarrow \infty} P\left(\bigcup_{m \geq n} |Z_m - Z| > \varepsilon\right) = 0$$

(Here we have OR, not AND as above. If any ω satisfies, breaks. should never satisfy.)

almost sure implies convergence in probability.

harder.

easier

Defn (convergence in distribution) [weak convergence]

$Z_n \xrightarrow{d} Z$ if $F_{Z_n}(z) \rightarrow F_Z(z)$ at all points where $F_Z(z)$ is continuous.

- * For conv. in probability \rightarrow for a large n , satisfies.
- for almost sure \rightarrow for a large n , and every value of n beyond that.

* fact (The union bound)

Let X_1, X_2, \dots, X_n be any rvs. let $(X_i \in \mathcal{G}_i)$ be events.

Then, $P(\bigcup_{i=1}^n X_i \in \mathcal{G}_i) \leq \sum_{i=1}^n P(X_i \in \mathcal{G}_i)$ (3rd Axiom)

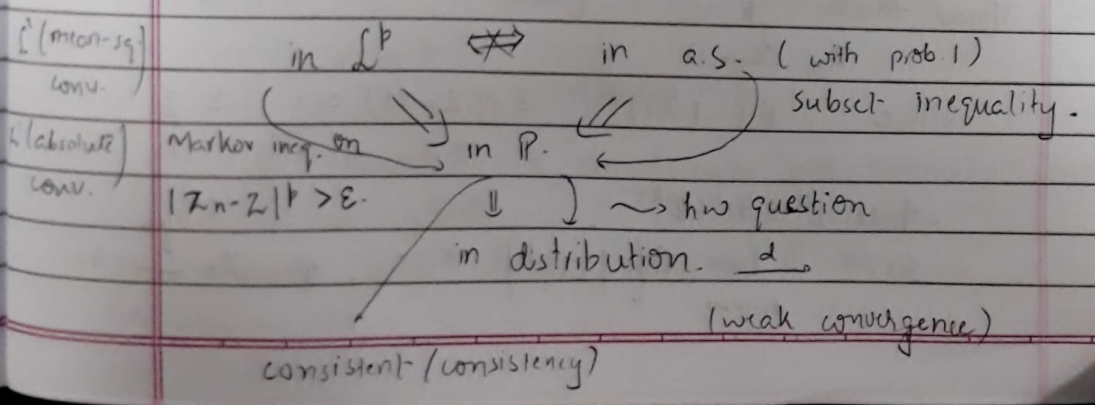
* For a.s. convergence, we want

$\lim_{n \rightarrow \infty} P(\bigcup_{m \geq n} |Z_m - Z| > \epsilon) = 0$ (for any $\epsilon > 0$)

\rightarrow A sufficient condn for almost s. conv. is for any $\epsilon > 0$,

$\lim_{n \rightarrow \infty} \sum_{m=n}^{\infty} P(|Z_m - Z| > \epsilon) = 0$

* Dependencies of convergence.



$$* P\left(\bigcup_{m \geq n} |Z_m - Z| > \varepsilon\right) \geq P(|Z_n - Z| > \varepsilon)$$

Example (counterexample of convergence dependencies)

i) $\xrightarrow{p} \not\Rightarrow \text{a.s.} \quad (p \geq 1).$

Let

- Let $Z = 0$. And. $Z_n = 0$ w.p. $\left(1 - \frac{1}{n}\right)$.

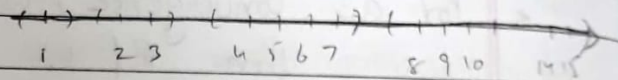
$Z_n = 1$ w.p. $\frac{1}{n}$.

Then, $E(|Z_n - Z|^p) = \frac{1}{n} \rightarrow 0$ as $n \rightarrow \infty$.

Or $Z_n \xrightarrow{p} Z$.

Scrap the above $\Rightarrow \times$.

Example -



4 5 6 7.

$p = 1/4$.

1	0	0	0
0	1	0	0
0	0	1	0
0	0	0	1
Z_4	Z_5	Z_6	Z_7

let $Z_i = 1$, & $Z_j = 0$ for $2^k \leq i, j \leq 2^{k+1} - 1$ in equiprobable fashion. $i \neq j$.

If $2^k \leq i \leq 2^{k+1}$ then

$$P(Z_i = 1) = \frac{1}{2^k}$$

show that $Z_n \xrightarrow{p} 0$

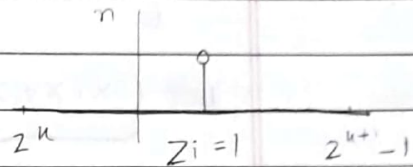
Since $E(|Z_n|^p) = P(Z_n = 1) = \frac{1}{2^k} \leq \frac{2}{n}$

where $2^k \leq n < 2^{k+1}$

So, $\lim_{n \rightarrow \infty} E(|Z_n|^p) = 0$, or $Z_n \xrightarrow{p} 0$.

consequently, $Z_n \xrightarrow{P} 0$

However, $Z_n \xrightarrow{a.s.} 0$ since for any $\varepsilon < 1$,
 $P\left(\bigcup_{m \geq n} |Z_m - 0| > \varepsilon\right) = 1$



Here, not converging a.s., as no matter what n_0 we choose, there will be some Z_i for which $P() \neq 0$. as for a.s., we need convergence for any $n > n_0$. Random sequence should not deviate again, even at a single point.

* in a.s. $\not\Rightarrow$ in L^p

Let $Z_n = 0$ with prob $(1 - \frac{1}{n^2})$
 $= n$ with prob $1/n^2$.

Then $E(|Z_n|^2) = n^2 \cdot 1/n^2 = 1$.
 $E(|Z_n|^3) = n^3 \cdot 1/n^2 = n$.
 $\therefore Z_n \not\xrightarrow{L^p} 0$ for all $p \geq 2$.

However, for any $\varepsilon > 0$
 $P\left(\bigcup_{m \geq n} |Z_m| > \varepsilon\right) \leq \sum_{m=n}^{\infty} P(|Z_m| > \varepsilon)$

$= \sum_{m=n}^{\infty} \frac{1}{m^2} \rightarrow 0$ as $n \rightarrow \infty$. ($Z_n \xrightarrow{a.s.} 0$)

If bounded, & converging in prob, normally converge in L^p as well.
 Also, from above, conv in $P \not\Rightarrow$ in L^p .

✓ Example (data compression example)

W
L

$$L \times W \times 3 \times 8$$

$$1920 \times 1080 \times 3 \times 8 = 6 \text{ MB.}$$

Let $(x_1, x_2, x_3, \dots, x_n)$ be a file.

binary string.

Let $x_i \sim \text{Ber}(p)$, where p is known and x_1, \dots, x_n are iid.

$$\text{Let } p_x(1) = p(1) = p, \text{ \& } p_x(0) = p(0) = (1-p).$$

Consider the function. $\log_2 p(x)$

$$\text{Then } \log_2 p(x) = \log_2(p) \quad ; \text{ w.p. } p$$

$$= \log_2(1-p) \quad ; \text{ w.p. } (1-p)$$

By WLLN;

$$\frac{1}{n} \sum_{i=1}^n \log_2(p(x_i)) \approx E(\log_2 p(x))$$

$$= p \log_2 p + (1-p) \log_2(1-p)$$

$$= \underbrace{-H_2(p)}_{\text{Entropy of } X \sim \text{Ber}(p)} \quad \left[\begin{array}{l} H_2(p), \\ \text{entropy always +ve} \end{array} \right]$$

Z (degenerate r.v.) $\equiv C$ takes it as a $p(c)=1$.

$$\text{where, } H_2(p) = -p \log_2 p - (1-p) \log_2(1-p). \quad (+ve)$$

By WLLN, for any $\epsilon, \delta > 0$,

$$P\left(\left|\frac{1}{n} \sum_{i=1}^n \log_2 p(x_i) + H_2(p)\right| \geq \epsilon\right) \leq \delta.$$

If n is large enough

I.e.,

$$-H_2(p) - \epsilon \leq \frac{1}{n} \sum_{i=1}^n \log_2 p(x_i) \leq -H_2(p) + \epsilon.$$

with prob $> 1-\delta$.

$$\Rightarrow -H_2(p) - \epsilon \leq \frac{1}{n} \log_2 p(x_1, \dots, x_n) \leq -H_2(p) + \epsilon, \text{ w.p. } > 1-\delta$$

$$\text{or, } 2^{-n(H_2(p)+\epsilon)} \leq p(x_1, \dots, x_n) \leq 2^{-n(H_2(p)-\epsilon)} \text{ with prob } > (1-\delta).$$

Define: $A_{n,\epsilon} = \{ \vec{n} : \left| \frac{1}{n} \log p(n_1, \dots, n_n) + H_2(p) \right| \leq \epsilon \}$.

typical set.

\hookrightarrow Then $A_{n,\epsilon}^c$ = atypical set. (Something which not happen normally)
 $H_2(p) = -p \log_2 p - (1-p) \log_2 (1-p)$

$A_{n,\epsilon}$ is a collection of deterministic strings.

Now, WLLN, says that

$$P((x_1, \dots, x_n) \in A_{n,\epsilon}) > 1-\delta.$$

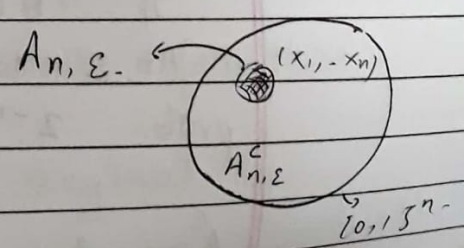
On the typical set,

$$2^{-n(H_2(p)+\epsilon)} \leq p(n_1, \dots, n_n) \leq 2^{-n(H_2(p)-\epsilon)}$$

$$\text{or } p(n_1, \dots, n_n) \approx 2^{-n H_2(p)} \quad (\text{ignoring } \epsilon) \\ \text{for any } (n_1, \dots, n_n) \in A_{n,\epsilon}.$$

$$\hookrightarrow \text{So } |A_{n,\epsilon}| \approx 2^{n H_2(p)} \text{ and } P((x_1, \dots, x_n) \in A_{n,\epsilon}^c) \approx 0.$$

for any A set with $|A|=n$,
 any probability $\rightarrow \frac{1}{n}$ of a particular element.



Given (n_1, \dots, n_n) , check if $(n_1, \dots, n_n) \in A_{n, \epsilon}$.

If no, retain $(1, n_1, \dots, n_n)$ (1, 0 a flag which determines whether in typical set or not)

If yes, retain $(0, n_1, \dots, n_n)$.

$|A_{n, \epsilon}| \approx 2^{nH_2(p)}$ } reinde using $nH_2(p)$ bits. Use lookup table to get a compressed representation for a typical bit string.

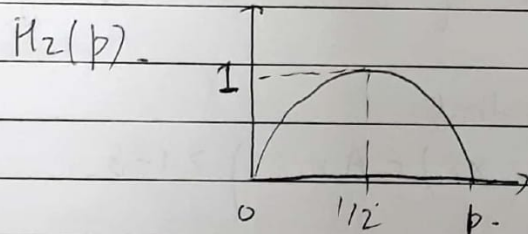
Expected length after compression is

$$\delta(n+1) + (1-\delta)(nH_2(p)+1)$$

$$\approx H_2(p)$$

per bit or per kb compression

0000 0001	
0000 0010	000
0000 0100	L 001
0000 1000	U 010
	T 011
1000 0000	
	111



\therefore farther away from $1/2$, better compression.

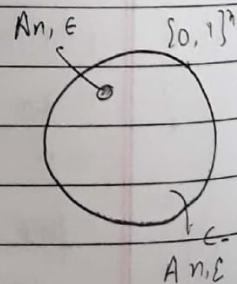
New day

Example. (WLLN and compression).

(x_1, \dots, x_n) ; file.

x_i 's are iid $\text{Ber}(p)$

If n is large, then $\approx 2^{nH_2(p)}$ strings in $A_{n, \epsilon}$ each nearly equiprobable with prob. $2^{-nH_2(p)}$

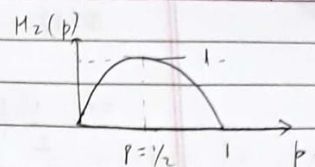


$A_{n, \epsilon}$ has negligible prob. collectively.

Expected length of compressed file $= (1+nH_2(p))(1-\delta) + \delta(n+1)$

1 Expected length $\approx H_2(p)$.

$$H_2(p) = -p \log p - (1-p) \log (1-p)$$



If $p = \frac{1}{2}$, the file is not compressible.

$$p_{x_1, x_2, x_3, \dots, x_n}(x_1, \dots, x_n) = \left(\frac{1}{2}\right)^n = 2^{-n} \quad (\text{or 1 equiprobable, all independent})$$

every string is equiprobable.

and $|\{0,1\}^n| = 2^n$. Entire set $\{0,1\}^n$ is typical. nothing atypical here.

Example: (Cauchy distribution's expectation)

Recall that $X = X^+ - X^-$ and $E(X) = E(X^+) - E(X^-)$ exist if $E(X^+) < \infty$ and $E(X^-) < \infty$.

→ If X is Cauchy distributed then

$$f_X(u) = \frac{1}{\pi} \frac{1}{1+u^2}; \quad u \in \mathbb{R}.$$

↳ $E(X) = \int_{-\infty}^{\infty} \frac{u}{\pi(1+u^2)} du = 0$. However $E(X^+) = \infty$ and $E(X^-) = \infty$.

Let X_1, \dots, X_n be iid Cauchy rvs with pdf $f_X(u)$.

↳ Then $\phi_X(t) = e^{-|t|}$, $t \in \mathbb{R}$.

→ From duality of the FT this result follows:

$$y(t) \xrightarrow{\mathcal{F}} g(\omega) \Leftrightarrow \tilde{y}(t) \xrightarrow{\mathcal{F}} 2\pi g(-\omega)$$

$$f(x) = e^{-x} u(x) \xrightarrow{\mathcal{F}} \frac{1}{1+jt}, \quad f(-x) = e^x u(-x) \xrightarrow{\mathcal{F}} \frac{1}{1-jt}$$

$$\Phi(t) = \int_{-\infty}^{\infty} e^{j t x} e^{-x} u(x) dx.$$

$$\Rightarrow \frac{1}{2} (\Phi(t) + \Phi(-t)) = \frac{1}{1+t^2} \quad \left(\frac{1}{2} \text{ so prob. add up to 1} \right)$$

The characteristic func. should have $\Phi(0) = 1$. Also $\int_{-\infty}^{\infty} f(x) dx = 1$.
check these 2 while applying duality.

$\rightarrow \underbrace{x_1 + \dots + x_n}_n$ has a pdf

$$f_{\underbrace{x_1 + \dots + x_n}_n}(u) = \frac{1}{n} \frac{1}{(1+u^2)^n} ; u \in \mathbb{R}.$$

$$\begin{aligned} \text{Since } \phi_{\underbrace{x_1 + \dots + x_n}_n}(t) &= \phi_{x_1}\left(\frac{t}{n}\right) \phi_{x_2}\left(\frac{t}{n}\right) \dots \phi_{x_n}\left(\frac{t}{n}\right) \\ &= e^{-|t|} , \quad t \in \mathbb{R}. \end{aligned}$$

\therefore averaging out doesn't change the distribution.

$\Phi(t)$ & $f(x)$ are unique pairs. having 1, can get the other deterministically.

If 2 signals have same FT, signals have to be the same.

\rightarrow If $E(X) < \infty$, WLLN tells us that

$$\underbrace{x_1 + \dots + x_n}_n \xrightarrow{P} E(X)$$

\hookrightarrow However $f_{\underbrace{x_1 + \dots + x_n}_n}(u)$ does not show this concentration.

\rightarrow I.e. $E(X) = 0$ is meaningless in the light of WLLN

— x —

Example (convergence of $X_{n:n}$)

let X_1, \dots, X_n be iid $\text{Unif}(0,1)$

let $X_{n:n} = \max\{X_1, \dots, X_n\}$; $X_{1:n} = \min\{X_1, \dots, X_n\}$

and $X_{r:n}$ is the r th largest in the set $\{X_1, \dots, X_n\}$.

order statistics. we will show that

$X_{n:n} \xrightarrow{\text{a.s.}} 1$

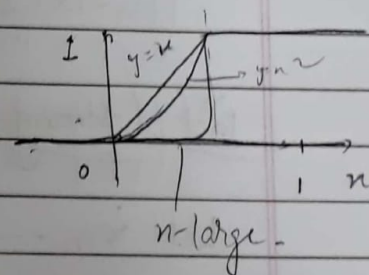
If $n \rightarrow \infty$, $\epsilon \rightarrow 0$

$X_{n:n} \rightarrow 1$

If n is large, the max value better be close to 1.

$$\begin{aligned} \text{Since } F_X(u) &= u; \quad 0 \leq u \leq 1, \\ &= 0, \quad u < 0 \\ &= 1, \quad u \geq 1. \end{aligned}$$

$$\begin{aligned} \text{So } F_{X_{n:n}}(u) &= u^n; \quad 0 \leq u \leq 1, \\ &= 0; \quad u < 0 \\ &= 1; \quad u \geq 1. \end{aligned}$$



We want to show that for any $\epsilon > 0$,

$$\lim_{n \rightarrow \infty} P\left(\bigcup_{m \geq n} |X_{m:n} - 1| > \epsilon\right) = 0.$$

$$\hookrightarrow P(|X_{n:n} - 1| > \epsilon)$$

$$= P(X_{n:n} < 1 - \epsilon)$$

$$= (1 - \epsilon)^n \quad (\rightarrow 0 \text{ as } n \rightarrow \infty)$$

$$\text{i.e., } X_{n:n} \xrightarrow{P} 1$$

\hookrightarrow For a.s. convergence, note that,

For a.s. convergence, note that

$$0 \leq P\left(\bigcup_{m \geq n} |X_{m:m} - 1| > \varepsilon\right) \leq \sum_{m \geq n} P(|X_{m:m} - 1| > \varepsilon) = \sum_{m \geq n} (1 - \varepsilon)^m$$

As $n \rightarrow \infty$,
 \rightarrow so, $X_{n:n} \xrightarrow{\text{a.s.}} 1$

$X_{10000:10000}$

$X_{10000:10000}$

$$1 - X_{n+1:n+1} \leq 1 - X_{n:n}$$

and $X_{n:n} \xrightarrow{P} 1$, it follows that $X_{n:n} \xrightarrow{\text{a.s.}} 1$

$$\rightarrow \text{or } \left\{ \bigcup_{m \geq n} |X_{m:m} - 1| > \varepsilon \right\} = \left\{ |X_{n:n} - 1| > \varepsilon \right\}$$

— x —

as any way it shrinks, the smallest one covers everything.

Fact (Strong law of large numbers)

Let X_1, \dots, X_n be iid r.v. with mean $E(X)$. Then,

$$\frac{X_1 + \dots + X_n}{n} \xrightarrow{\text{a.s.}} E(X)$$

— x —

Fact (Strong law of large numbers (SLLN))

Let X_1, \dots, X_n be iid r.v. with mean $E(X)$. Then

$$\frac{X_1 + \dots + X_n}{n} \xrightarrow{\text{a.s.}} E(X)$$

— x —

Here, we assumed nothing about variance. Also, in WLLN, it was conv. in prob, Here a.s.

Fact (Central Limit Theorem) -

Let X_1, \dots, X_n be iid rvs such that $E(X) = \mu$ and $\text{Var}(X) = \sigma^2$.

Let $V_n := \sum_{i=1}^n \frac{X_i - \mu}{\sqrt{n}\sigma}$. Then

$$V_n \xrightarrow{d} N(0, 1).$$

Proof (sketch):

Tool: Characteristic functions.

Fact: i) If $Z_n \xrightarrow{d} Z$, then $\phi_{Z_n}(t) \rightarrow \phi_Z(t)$ for all $t \in \mathbb{R}$.

ii) If $\phi_{Z_n}(t) \rightarrow \phi_Z(t)$ and $\phi_Z(t)$ is continuous at $t=0$, then $Z_n \xrightarrow{d} Z$ (Levy continuity theorem).

Let $\phi_X(t)$ be the charact. func. of X_1, \dots, X_n .

$$\begin{aligned} \text{Then } \phi_{V_n}(t) &= E(e^{jt V_n}) \\ &= E\left(e^{j t \frac{1}{\sqrt{n}\sigma} \sum_{i=1}^n (X_i - \mu)}\right) \end{aligned}$$

$$= \prod_{i=1}^n \phi_{X-\mu}\left(\frac{t}{\sqrt{n}\sigma}\right) = \left[\phi_{X-\mu}\left(\frac{t}{\sqrt{n}\sigma}\right)\right]^n.$$

Observe that

$$\phi_{X-\mu}(t) = 1 - \frac{\sigma^2}{2} t^2 + o(t^2).$$

$$\phi_X(t) = E(e^{jtx}) = E\left(1 + jtx + \frac{j^2 t^2 x^2}{2!} + \frac{j^3 t^3 x^3}{3!} + \dots\right)$$

$$\lim_{t \rightarrow 0} \frac{o(t^2)}{t^2} = 0, \quad \text{as } E(X-\mu) = 0, \text{ no linear term}$$

defn of.

So, $\phi_{V_n}(t) = \left(1 - \frac{t^2}{2n} + o\left(\frac{t^2}{n}\right) \right)^n$
 $\rightarrow e^{-t^2/2}$ as $n \rightarrow \infty$.

So $V_n \xrightarrow{d} N(0,1)$, $\phi_{N(0,1)}(t)$.

Let $c_n \in \mathbb{C}$ and $c_n \rightarrow c$.

Then, $\left(1 + \frac{c_n}{n} \right)^n \rightarrow e^c$, in \mathbb{C} .

$V_n = \frac{1}{\sqrt{n\sigma}} \sum_{i=1}^n (X_i - \mu)$; $\phi_{V_n}(t) \rightarrow e^{-t^2/2}$ for every $t \in \mathbb{R}$.

So $V_n \xrightarrow{d} Z$ where $Z \sim N(0,1)$.

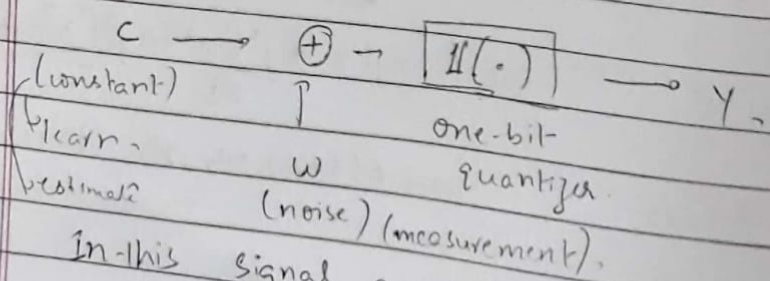
$F_{V_n}(x) \rightarrow F_Z(x)$ at all pts $x \in \mathbb{R}$ where $F_Z(x)$ is continuous.

i.e. $F_{V_n}(x) \rightarrow F_Z(x)$ for all $x \in \mathbb{R}$

(since the latter is continuous)

Example (denoising & quantization)

noise there is averaged out if exposure time is increased
 ISO speed \rightarrow determines the exposure time required to get a neutral image.



In this signal acquisition model,

$$y = \mathbb{1}(c+w \geq 0) = \begin{cases} 1 & \text{if } c+w \geq 0 \\ 0 & \text{if } c+w < 0 \end{cases}$$

$$Y = \left[\frac{1 + \text{sgn}(c + w)}{2} \right] \quad (\text{Observe})$$

↳ If w is unknown, c can't be inferred (accurately) from Y .

↳ Consider 'n' indep. readings

$$Y_i = 1 \quad (c + w_i > 0)$$

where w_1, \dots, w_n are indep. and identically distribution.

↳ Observe that $\{Y_i\}$ are also iid and (Think why?)

$$Y \sim \text{Ber}(1 - F_W(-c))$$

If w is a continuous r.v. $\rightarrow P(w + c > 0)$.

$(1 - F_W(-c))$ is invertible (requires $F_W(u)$ is $\in (0, 1)$ when $u \in \mathbb{R}$)

since the function is monotonic $-f_W(-c) > 0$.

Let w be a sum of n i.i.d. r.v.s

$$Y = [1 + \text{sgn}(c + w)] / 2 \quad (\text{Observe})$$

↳ If w is unknown, c can't be inferred (accurately) from Y .

↳ Consider 'n' indep. readings

$$Y_i = 1 \text{ if } (c + w_i > 0)$$

where w_1, \dots, w_n are indep. and identically distribution.

↳ Observe that $\{Y_i\}$ are also iid and (Think why?)

$$Y \sim \text{Ber}(1 - F_w(-c))$$

If w is a continuous r.v. $\rightarrow P(w + c > 0)$.

$(1 - F_w(-c))$ is invertible (requires $F_w(u)$) in c since the function is monotonic whenever $-f_w(-c) > 0$.

Let w be a symmetric r.v.,
i.e. $f_w(u) = f_w(-u)$.

Then $Y \sim \text{Ber}(F_w(c))$.

↳ Let $S_n = Y_1 + \dots + Y_n$. Then

$$\frac{S_n}{n} \xrightarrow{P} E(Y) = F_w(c) \quad (\text{something with law})$$

$$\text{and } \text{Var}\left(\frac{S_n}{n}\right) = \frac{1}{n} \text{Var}(Y) = \frac{1}{n} F_w(c)(1 - F_w(c)) \leq \frac{1}{4n}$$

$$\rightarrow \text{Let } \hat{c} = F_w^{-1}\left(\frac{S_n}{n}\right)$$

where $F_w(u)$ is known.