CS 228 : Logic in Computer Science

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- Resolvent not unique : $\{A_1, A_3, \neg A_3, A_4\}$ is also a resolvent.

3 rules in Resolution

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- Let F be a formula in CNF. Let R be a resolvent of two clauses of F. Then $F \vdash R$ (Prove!)

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Show that resolution can be used to determine whether any given formula is satisfiable.

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- $Res^0(F) =$clauses in F, and there are finitely many clauses that can be derived from F.
- Hence, there is some m such that $Res^m(F) = Res^{m+1}(F)$. Denote it by $Res^*(F)$.

Example

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- $Res^2(F) = Res^1(F) \cup \{A_1, A_2, \neg A_3\} \cup \{A_1, A_3, \neg A_2\}$

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- Since $\emptyset \notin Res^0(F)$ ($\emptyset$ is not a clause), there is an $m > 0$ such that $\emptyset \notin Res^m(F)$ and $\emptyset \in Res^{m+1}(F)$.
- Then $\{A\}, \{\neg A\} \in Res^m(F)$. By the rules of resolution, we have $F \vdash A, \neg A$, and thus $F \vdash \bot$. Hence, F is unsatisfiable.

Resolution

Prove the converse: If F is unsatisfiable, then $\emptyset \in \mathit{Res}^*(F)$.

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- If $F = \{\{p\}\}$ or $F = \{\{\neg p\}\}$, F is satisfiable.
- Hence, $F = \{\{p\}, \{\neg p\}\}$. Clearly, $\emptyset \in Res^1(F)$.

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  - Let G_0 be the conjunction of all C_i in F such that $\neg p_{n+1} \notin C_i$.
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- Let $F_0 = \{C_i - \{p_{n+1}\} \mid C_i \in G_0\}$
- Let $F_1 = \{C_i - \{\neg p_{n+1}\} \mid C_i \in G_1\}$

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Let $F = \{\{p_1, p_3\}, \{p_2\}, \{\neg p_1, \neg p_2, p_3\}, \{\neg p_2, \neg p_3\}\}$ and $n = 2$.

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- $F_0 = \{\{p_1\}, \{p_2\}, \{\neg p_1, \neg p_2\}\}$ and $F_1 = \{\{p_2\}, \{\neg p_2\}\}$
- If p_{n+1} is assigned *false* in F, then F is equivalent to F_0

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- Hence $F \equiv F_0 \vee F_1$.

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- If p_{n+1} is assigned *false* in F, then F is equivalent to F_0
- If p_{n+1} is assigned *true* in F, then F is equivalent to F_1
- Hence $F \equiv F_0 \vee F_1$.
- As F is unsatisfiable, F_0 and F_1 are both unsatisfiable.

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- By induction hypothesis, $\emptyset \in \mathit{Res}^*(F_0)$ and $\emptyset \in \mathit{Res}^*(F_1)$.

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- By induction hypothesis, $\emptyset \in Res^*(F_0)$ and $\emptyset \in Res^*(F_1)$.
- Hence, $\emptyset \in Res^*(G_0)$ or $\{p_{n+1}\} \in Res^*(G_0)$, and $\emptyset \in Res^*(G_1)$ or $\{\neg p_{n+1}\} \in Res^*(G_1)$.

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- If $\emptyset \in \mathit{Res}^*(G_0)$ or $\emptyset \in \mathit{Res}^*(G_1)$, then $\emptyset \in \mathit{Res}^*(F)$.
- Else, $\{p_{n+1}\} \in \mathit{Res}^*(G_0)$ and $\{\neg p_{n+1}\} \in \mathit{Res}^*(G_1)$.

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- Hence, $\emptyset \in Res^*(G_0)$ or $\{p_{n+1}\} \in Res^*(G_0)$, and $\emptyset \in Res^*(G_1)$ or $\{\neg p_{n+1}\} \in Res^*(G_1)$.
- If $\emptyset \in Res^*(G_0)$ or $\emptyset \in Res^*(G_1)$, then $\emptyset \in Res^*(F)$.
- Else, $\{p_{n+1}\} \in Res^*(G_0)$ and $\{\neg p_{n+1}\} \in Res^*(G_1)$.
- Hence $\emptyset \in Res^*(F)$.

Resolution Summary

Given a formula ψ, convert it into CNF, say ζ. ψ is satisfiable iff $\emptyset \notin \mathit{Res}^*(\zeta)$.

- If ψ is unsat, we might get $\emptyset$ before reaching $\mathit{Res}^*(\zeta)$.
- If ψ is sat, then truth tables might be faster : stop when some row evaluates to 1.

Propositional Logic : Summary

- Syntax, Semantics
- Encoding problems into logic
- Sound and Complete Proof Engine
- Semantic/Provable equivalence of formulae
- Normal forms, satisfiability, hardness
- Resolution for SAT checking

Moving On

Propositional Logic

SAT solvers, heuristics, competitions for SAT solvers and so on. In many cases, parts of a complex problem reduced to SAT solving.

What we propose to do now

Move on to other logics, and their applications in CS.