

CS286: Database Systems

1 Lecture 3—9/4/2014

1.1 R*

- Assumptions:
 - There are administrative causes behind distributed data
 - Network: unreliable transport, in-order, packets are intact
 - Independent node failure
 - Slow-ish network
- Research goals:
 - “Site autonomy”: No centralized state or control
 - * Data you touch should determine the sites you talk to
 - * “Distributed system is a system that fails because a machine you’ve never heard of fails”
 - * Load sharing and decentralization
 - * Less communication
 - * Harder to coordinate data consistency
 - * More network connections beyond hub and spoke
 - * Metadata management is harder
 - Location transparency → emulate a centralized DB
 - Don’t assume much about the network or OS
- Highlights:
 - Query optimizer cost modeling
 - Data layouts → horizontal partitioning
 - Replication
 - Distribution
 - Query compilation—unclear as to balance between compilation overhead and work saving
 - Spent a lot of time talking about 2PC → presumed commit

1.2 Gamma

- Assumptions:
 - Fast interconnect—hypercube, more network bandwidth than aggregate disk bandwidth
 - Shared nothing—no disk or memory sharing
- Research goals:
 - Scale

- Highlights:
 - Parallel hybrid-hash join
 - Chained declustering
- Assess:
 - Linear speedup + scale-up
 - Superlinear speedup due to minimized seek count at scale

2 Lecture 4—9/9/2014

- ACID
 - Consistency is not what we typically think
 - Distributed systems: data has a consistent value across sites
 - Databases: data meets contract when transaction completes
- Serializability mathematically gives atomicity and isolation
- Logging gives atomicity and durability
- Ordering:
 - Determines outcome (unless operations are not associative and commutative)
 - Some things are commutable/associable
 - Ordering must be equivalent to some serializable order
 - Implicitly, this provides an API—people don't *need* to reason about concurrency
- What is storage?
 - *Spacial-temporal rendezvous makes everything work!!!!*
- Want to avoid/undo conflicts in space and time
 - *Space*: Shared names
 - *Time*: Ordering
- 2PL: Provides a conflict serialized schedule
 - Ordered by race for locks
 - Ordered by the end of the first phase (“lock point”)
- Multi-version timestamp ordering
 - Every transaction gets a timestamp—this is the only synchronization point
 - For every object:
 - * Writes generate a new version for an object
 - * Reads annotate the version for the object

3 Lecture 5—9/11/2014

- Good graphs:
 - Crossover points
 - Non-monotonicity
 - Good breadth of X
 - Smooth → variance was accounted for
- Infinite resources:
 - Why run infinite resources? Many people assumed infinite resources in their papers.
 - OCC wins because it allows higher parallelism, at the cost of restarting transactions
 - Blocking (2PL) performs well at start, low at the end. Why?
 - * Deadlock starts to cause performance to fail
 - * Lock contention starts to cause transactions to get in each other's way
 - * Locking is a feedback loop—it lengthens transaction time
- Takeaways:
 - MPL is a control variable—choose your infrastructure for your system
- When do we have “infinite” resources?
 - When we have user interaction (Computer \gg human)
 - Vastly overprovisioned compute
 - Work is not going on inside the serving infrastructure (e.g., work is done by clients)

3.1 What happens when you go distributed?

- Why go distributed?
 - Capacity (storage and throughput)
 - Low latency (tolerance)
 - Fault tolerance (durability vs. availability)
- Techniques
 - Sharding—split dataset across many nodes
 - Replication

4 Lecture 6—9/16/2014

- You need replication → resilience to failure
- Tradeoff between replication and performance
- NoSQL:
 - Typically, a key-value store (data/programming model)
 - Typically distributed and sharded/partitioned
 - Usually weaker consistency model
 - No transactions/weak isolation model
 - “Not MySQL” → lots of work at AOL/etc. with MySQL on memcached

- Typically OSS, not enterprise
 - “Scalable”, especially incremental scale → improves organization/administration/ops
 - “Evaporation” of the DBA
- Motivations:
 - Bayou: I want to operate when disconnected
 - Dynamo: Nodes gonna fail
- CAP theorem: if partitions occur, then we can either have consistency or availability
 - Availability: As long as a client can access a server, I can access data (concurrent operations don’t need to communicate)
 - Consistency: “linearizable registers” → if I make a write, you can read my write