PROGRAMMING AND DATA STRUCTURES

BINARY TREES (HEAP)

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SPRING 2023

OUTLINE

- Characteristics of the Heap
- Operations on the Heap
- → Implementation of the Heap class

STUDENT LEARNING OUTCOMES

At the end of this chapter, you should be able to:

- Describe the properties of the Heap
- Trace operations on the Heap
- Implement the Heap generic data structure
- Use the Heap data structure
- Evaluate the complexity of the operations on the Heap

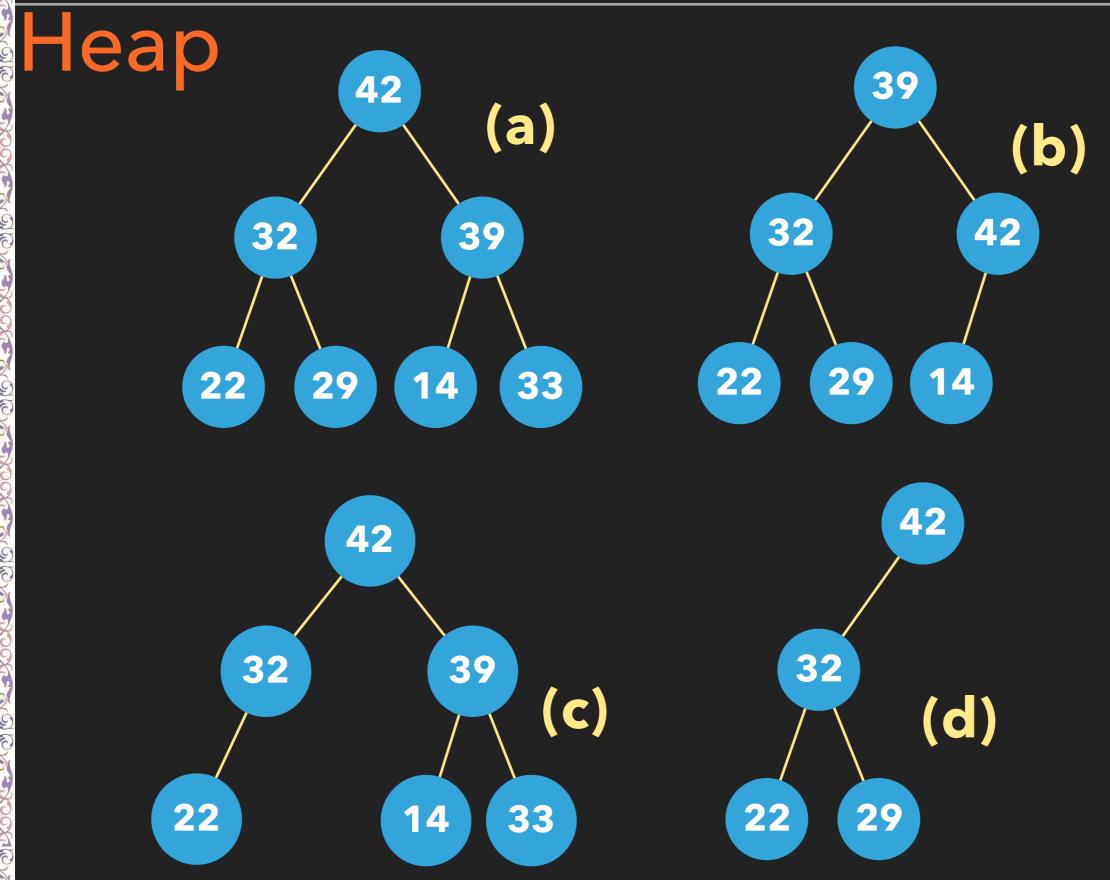
Heap

- Special binary tree used to order data (not to retrieve data)
- Used for efficient sorting (heap sort)
- Used to implement the priority queue

Heap

- Properties of the heap
 - ◆ Property 1: Complete binary tree All the levels are filled except the last level All leaves on the last level are placed leftmost
 - Property 2: every node is greater than or equal to any of its children (Max Heap) [Min Heap: less than or equal]

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Heap

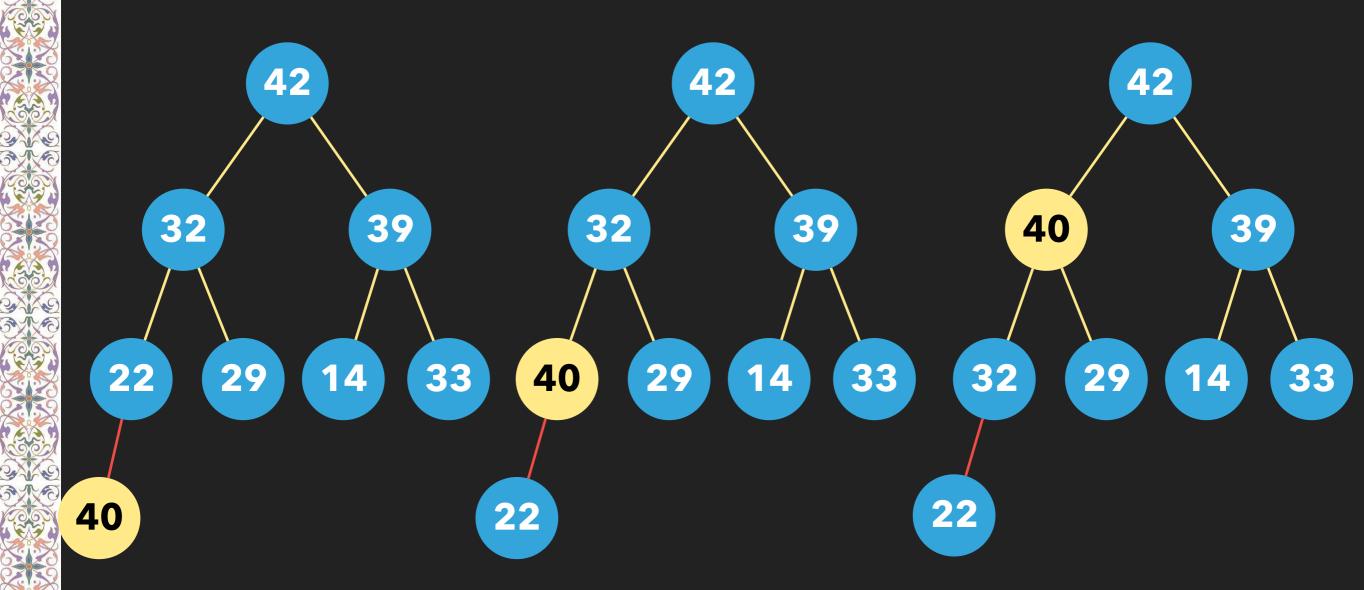
Two main operations on the Heap

Adding a new node while keeping the heap properties

Removing a node while keeping the heap properties

Heap (add)

 \rightarrow Adding a new node to the heap (40)



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End

Heap (add)

Adding a new node to the heap

```
Algorithm add

Add the new node at the end of the heap

Current node = added node

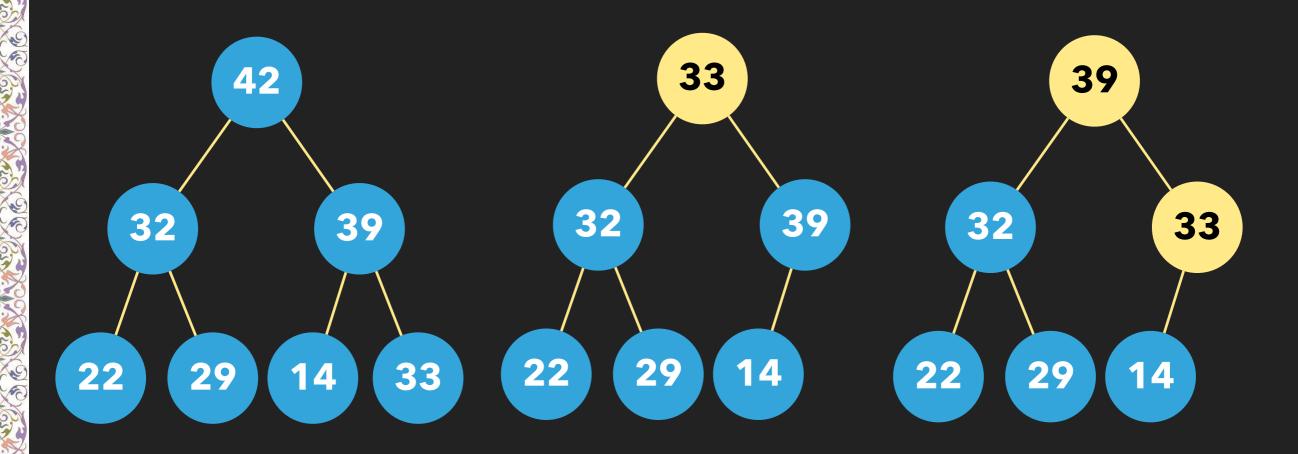
While (current node > its parent)

Swap current node with its parent

Current node becomes the parent
```

Heap (remove)

- ightharpoonup Removing a node from the heap (42)
- Always the root



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Heap (remove)

Removing a node from the heap (root)

```
Algorithm remove
```

```
Copy the value of the last node to the root
```

```
Current node = root
```

```
While (current node < its children)
```

```
Swap current node with the largest of its children
```

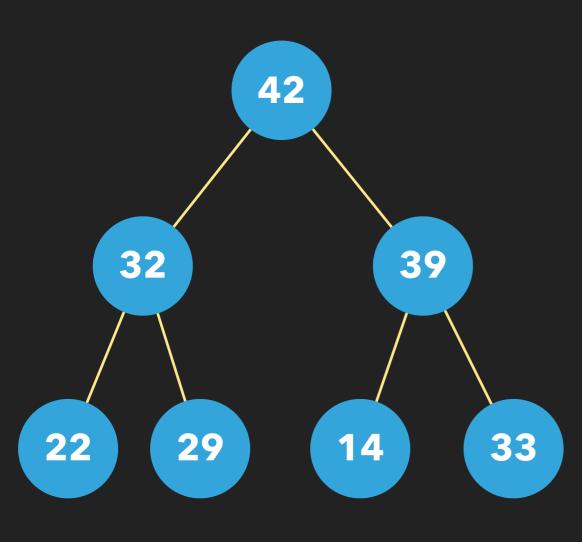
Current node becomes the largest child

End

Heap implementation

ArrayList to store the heap nodes

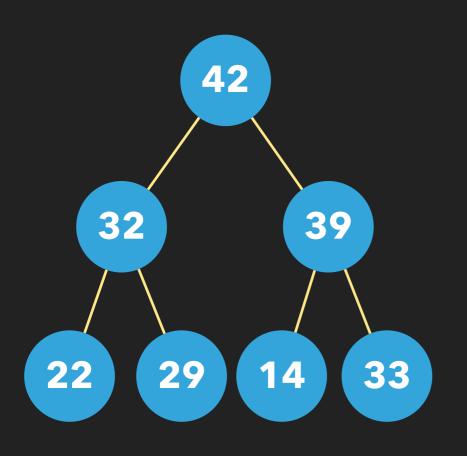
Easy access to children and parent

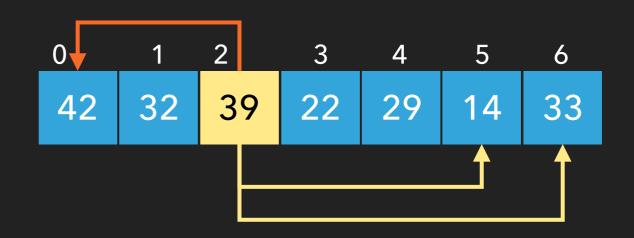


ArrayList with the nodes of the heap

0	1	2	3	4	5	6
42	32	39	22	29	14	33

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```
IndexOf(Parent) = (IndexOf(current) - 1) / 2
```

```
IndexOf(Left child) = 2 * IndexOf(current) + 1
IndexOf(Right child) = 2 * IndexOf(current) + 2
```

"interface" Collection<E> "interface" Queue<E> PriorityQueue<E> (MinHeap)

```
Heap<E extends Comparable<E>>
-list: ArrayList<E>
+Heap()
+add(E): void
+remove(): E
+contains(E): boolean
+size(): int
+isEmpty(): boolean
+clear(): void
+toString(): String
```

Heap.java

```
public class Heap<E extends Comparable<E>>> {
   private ArrayList<E> list;
   public Heap(){
        list = new ArrayList<>();
   public int size(){
        return list.size();
   public boolean isEmpty(){
        return list.isEmpty();
   public void clear(){
        list.clear();
   public String toString(){
        return list.toString();
```

```
public boolean contains(E value) {
    for(int i=0; i<list.size(); i++) {
        if(list.get(i).equals(value))
            return true;
    }
    return false;
}</pre>
```

Heap.java

```
public void add(E value) {
        list.add(value);
        int currentIndex = list.size()-1;
        while(currentIndex > 0) {
            int parentIndex = (currentIndex-1)/2;
            E current = list.get(currentIndex);
            E parent = list.get(parentIndex);
            if(current.compareTo(parent) > 0) {
                list.set(currentIndex, parent);
                list.set(parentIndex, current);
            else
                break;
            currentIndex = parentIndex;
```

Heap.java

```
public E remove() {
        if(list.size() == 0) return null;
        E removedItem = list.get(0);
        list.set(0, list.get(list.size()-1));
        list.remove(list.size()-1);
        int currentIndex = 0;
        while (currentIndex < list.size()) {</pre>
            int left = 2 * currentIndex + 1;
            int right = 2 * currentIndex + 2;
            if (left >= list.size())
                break;
            int maxIndex = left;
            E max = list.get(maxIndex);
            if (right < list.size())</pre>
            if(max.compareTo(list.get(right)) < 0)</pre>
                maxIndex = right;
            E current = list.get(currentIndex);
            max = list.get(maxIndex);
            if(current.compareTo(max) < 0){</pre>
                list.set(maxIndex, current);
                list.set(currentIndex, max);
                currentIndex = maxIndex;
            else
                break;
        return removedItem;
```

Test.java

```
public class Test {
   public static void main(String[] args) {
        Heap<String> heap = new Heap<>();
        heap.add("Kiwi");
        heap.add("Strawberry");
        heap.add("Apple");
        heap.add("Banana");
        heap.add("Orange");
        heap.add("Lemon");
        heap.add("Watermelon");
        System.out.println("Heap: " + heap.toString());
        System.out.println("Removed: " + heap.remove());
        System.out.println("Heap: " + heap.toString());
        System.out.println("Heap contains Pear?: " +
                            heap.contains("Pear"));
```

Performance of the Heap operations

Method	Complexity		
Heap()	0(1)		
size()	0(1)		
clear()	0(1)		
isEmpty()	0(1)		
add(E)	O(log n)		
remove(E)	O(log n)		
contains(E)	O(n)		
toString()	O(n)		

Summary

- Heap Special binary tree
- Operations: Add and Remove mainly
- Implementation Using an ArrayList
- Performance of the operations on the Heap (logarithmic complexity for add and remove)
- Heap is a balanced binary tree always (height = log(number of nodes))