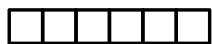


Introduction to Trees

Linear data structures: Array, Stack, Linked List, Queue

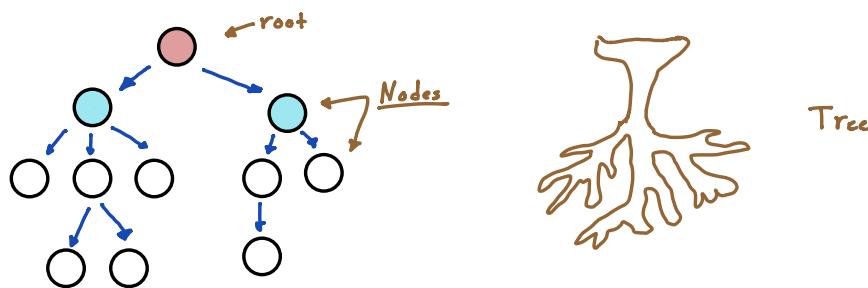


Array 5

How should I decide which data structure to use?

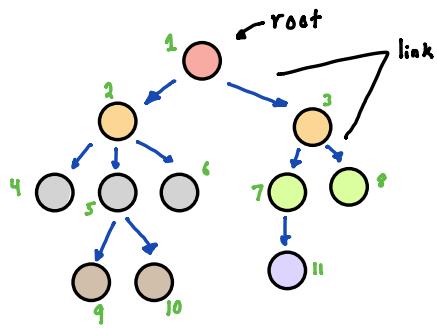
- What needs to be stored?
- Cost of operation
- Memory usage
- Ease of implementation

Tree used for Hierarchical data representation



- Non-linear

- Node contains data and possibly reference to another node



Common Terminology

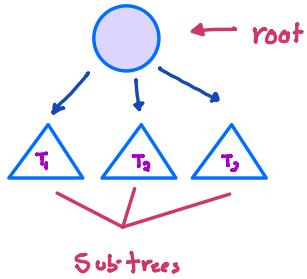
Note: it's don't represent anything in our example

- 2 and 3 linked to Node 1 and our "children" of node 1
- Node 1 is called "Parent" of Nodes 2 and 3
- Children of the same parent are called siblings
- Any node in the tree w/o a child is called a leaf Node
- All nodes with 1 child are called internal nodes
- Parent of a parent is called Grandparent
- When walking in the tree we can only go in one direction
- Ancestor and descendant
- Cousins, uncles

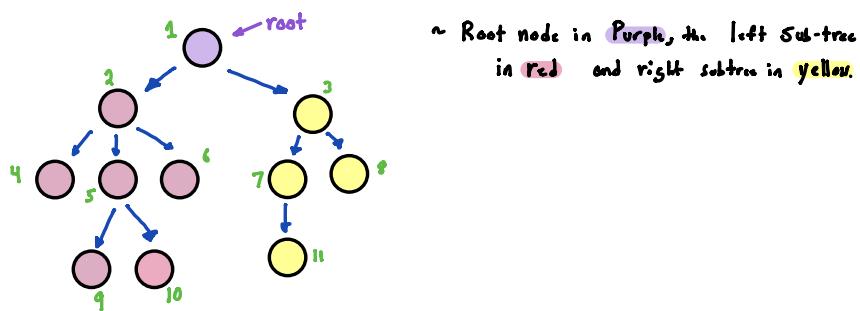
Properties of Trees

- Recursive data structure

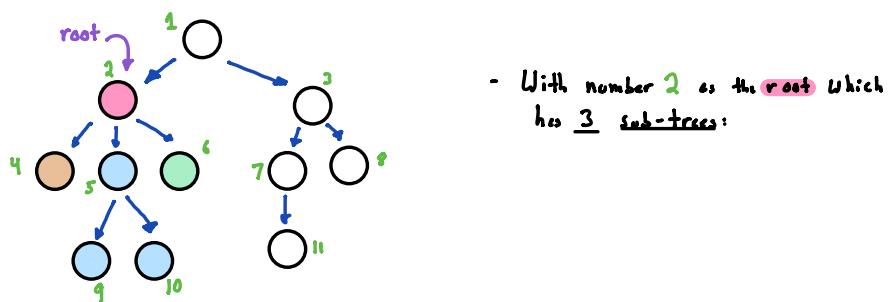
- We can define a tree recursively as a structure that consists of a distinguished node called a root and some sub-trees and the arrangement is such that root of the tree contains link to roots of all the sub-trees (T_1, T_2, T_3) in our diagram.



Example from our diagram above:

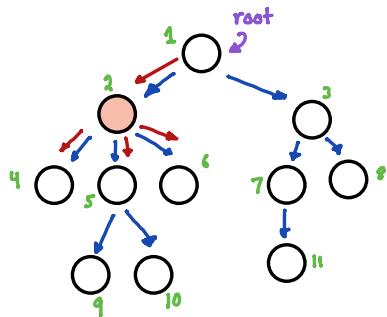


We can also look further at the sub-tree.



Properties of Trees (Continued)

- In a tree with n nodes, there will be exactly $n-1$ links or edges
- All nodes except the root node will have exactly one edge



At node 2, we have one incoming link and 3 outgoing links.

Depth and Height Properties

Depth of x : length of path from root to x

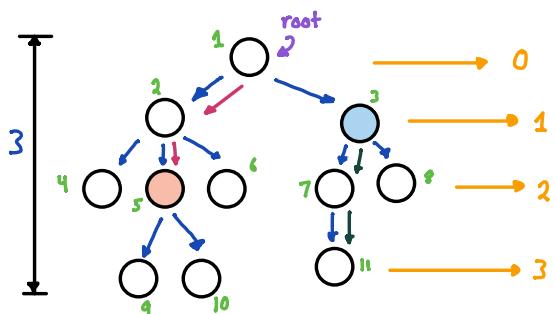
or

Number of edges in path from root to x

Height of x : Number of edges in longest path from x to a leaf.

Height of tree: Height of root node

Example of Depth and Height:



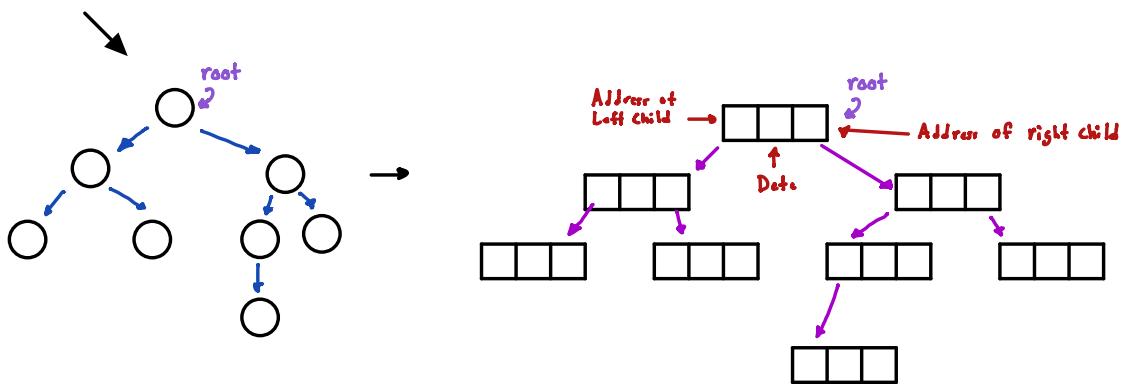
Depths of nodes in Orange

- In node 5 we have 2 edges in the path from root. So the depth is 2.
- In node 3, the longest path from this node to any leaf is 2.

Different types of Trees

Binary Tree:

- a tree in which each node can have at most 2 children
- most famous



Code in Python:

Node Class:

Application of trees

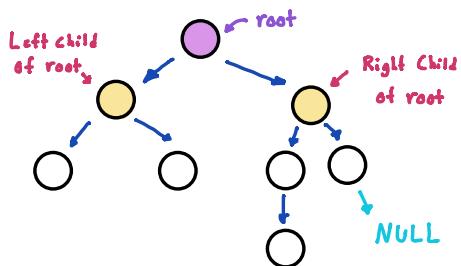
- 1) Storing naturally hierarchical data
→ e.g. file system
- 2) Organizing data for quick search, insertion, deletion
→ e.g. Binary Search trees
- 3) Trie
→ Dictionary
- 4) Network Rooting Algorithm

Class Node:

```
def __init__(self, data):
    self.left = None
    self.right = None
    self.data = data
```

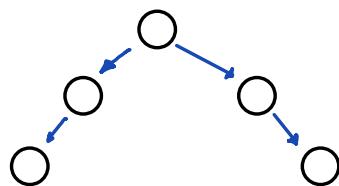
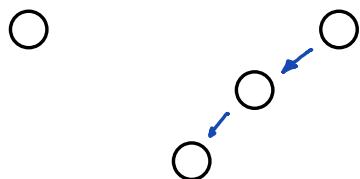
Binary Tree

Properties:



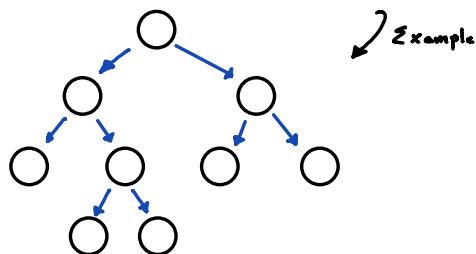
- A node may have both right and left child or a node can have either a right or Left child
- In a program we'll set reference of empty child to node
- For Leaf, we can set both Left and right to NULL

Different types of Binary trees



Strict/Proper binary Tree

↳ each node can have either 2 or 0 children.

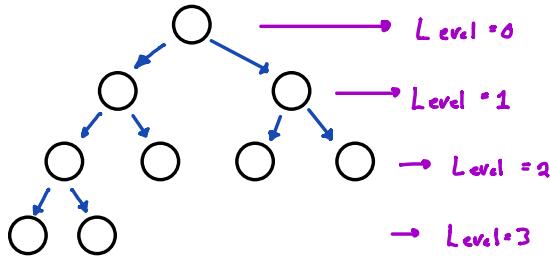


↗ Example

Complete Binary Tree

all levels except possibly the last are completely filled
and all nodes are as left as possible

- max number of nodes at level $i = 2^i$



Perfect Binary Tree

Maximum number of nodes
in a binary tree with height h

$$= 2^0 + 2^1 + \dots + 2^h$$

$$= 2^{h+1} - 1$$

$$= 2^{(\text{no. of levels})} - 1$$

$$n = 2^{h+1} - 1 \quad n = \text{# of nodes}$$

$$\rightarrow 2^{h+1} = (n+1)$$

$$h = \log_2(n+1) - 1$$

$$h = \log_2(15+1) - 1$$

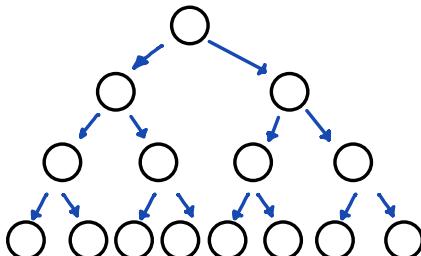
$$h = \log_2 16 - 1$$

$$h = 4 - 1$$

$$h = 3$$

$$2^x = 16$$

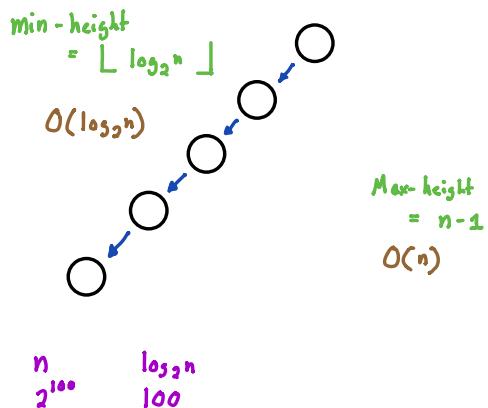
- Solve for x



Height of Complete Binary tree

$$= \lfloor \log_2 n \rfloor \quad \sim \text{floor function}$$

Time Complexity: $O(h)$ \Rightarrow h = height of tree



Balanced Binary Tree

- difference b/w height of Left and right subtree for every node is not more than K (mostly 1)

Height \Rightarrow number of edges in longest path from root to a leaf

Height of an Empty tree = -1

Height of tree with 1 node = 0

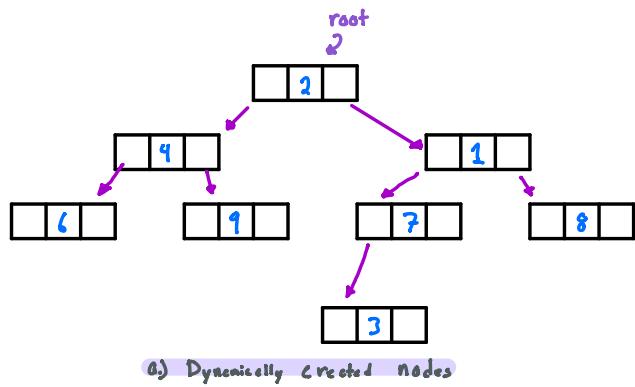
$$\text{diff} = | h_{\text{left}} - h_{\text{right}} |$$

We can implement Binary tree using:

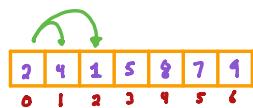
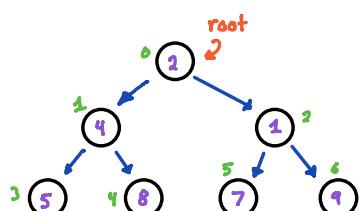
a) Dynamically created nodes

b) arrays

↳ used for heaps



a) Dynamically Created nodes



for node at index i ,
left-child-index = $2i + 1$
right-child-index = $2i + 2$] → in a Complete Binary tree