

## Quick Sort

- Divide and Conquer
- $O(n \log n)$  - Average case running time
- $O(n^2)$  - Worst case running time
- In-place
- Fastest

### High-Level Description of Quick-Sort

The quick-sort algorithm sorts a sequence  $S$  using a simple recursive approach. The main idea is to apply the divide-and-conquer technique, whereby we divide  $S$  into subsequences, recur to sort each subsequence, and then combine the sorted subsequences by a simple concatenation. In particular, the quick-sort algorithm consists of the following three steps (see Figure 12.8):

1. **Divide:** If  $S$  has at least two elements (nothing needs to be done if  $S$  has zero or one element), select a specific element  $x$  from  $S$ , which is called the **pivot**. As is common practice, choose the pivot  $x$  to be the last element in  $S$ . Remove all the elements from  $S$  and put them into three sequences:
  - $L$ , storing the elements in  $S$  less than  $x$
  - $E$ , storing the elements in  $S$  equal to  $x$
  - $G$ , storing the elements in  $S$  greater than  $x$Of course, if the elements of  $S$  are distinct, then  $E$  holds just one element—the pivot itself.
2. **Conquer:** Recursively sort sequences  $L$  and  $G$ .
3. **Combine:** Put back the elements into  $S$  in order by first inserting the elements of  $L$ , then those of  $E$ , and finally those of  $G$ .

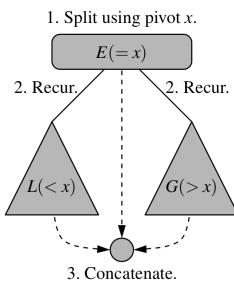
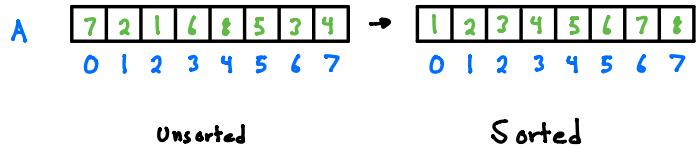
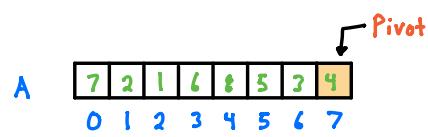


Figure 12.8: A visual schematic of the quick-sort algorithm.

## Quick Sort

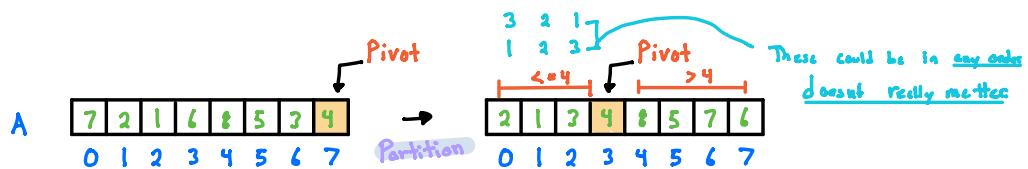


## Quick Sort Overview

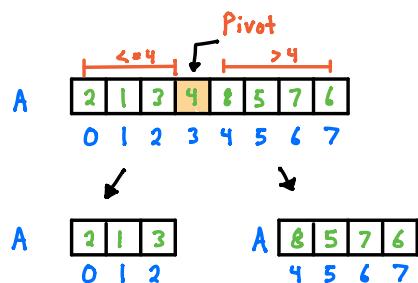


- 1) We first Select one elements in the list (can be any element), we call this selective number, Pivot.

- 2) Now we rearrange the list such that all the elements lesser than the Pivot are towards the left of it and all the elements greater are towards the right of it (Partitioning of a list)



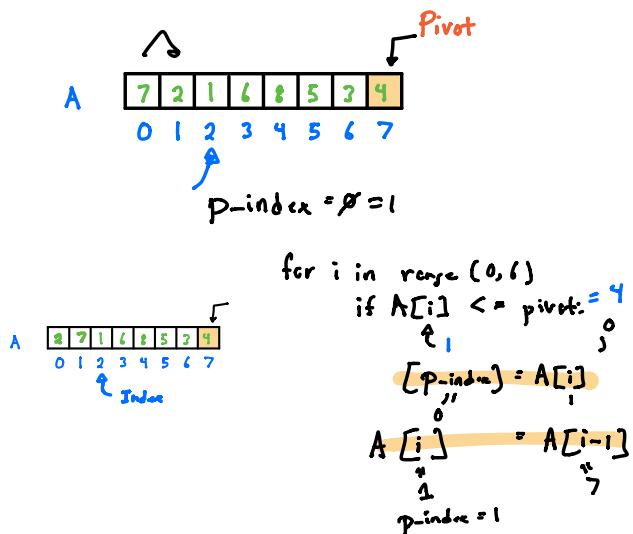
- We can break this problem into two sub-problems, those being sorting the array to the left of the pivot and sorting the array to the right of the pivot.



## Different than Merge Sort

- Unlike merge sort we do not need to create auxiliary arrays entirely. We can work on the same array. We just have to keep track of the start and end index of the segments.

- helper data structure (taking up extra space)



## QuickSort Function Suedo Code

```

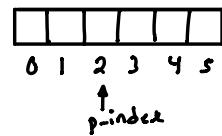
def QuickSort(A, start, end): Index
    if start < end:
        return exit condition
        pIndex = Partition(A, start, end)
        Recursion calls
        QuickSort(A, start, pIndex - 1)
        QuickSort(A, pIndex + 1, end)
    ~ i in range(start_index, end_index)

```

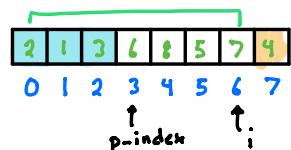
```

def Partition(A, Start, End):
    pivot = A[End] ~ Selecting the Pivot
    Partition_index = start ~ index
    for i in range(Start, End - 1):
        if A[i] <= pivot:
            A[i] = A[p_index]
            A[i-1] = A[i]
        p_index = p_index + 1
    ~ Swap
    A[p_index], A[End] = A[End], A[p_index]
    Return Partition_Index

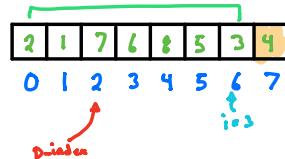
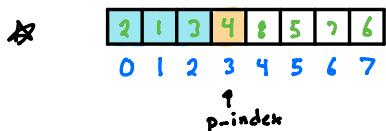
```



Note: At all times, Values to the Left of p-index will be less than the pivot.



- Once we've gone through the range of the array, we will swap at the p-index with our End element, which is our Pivot

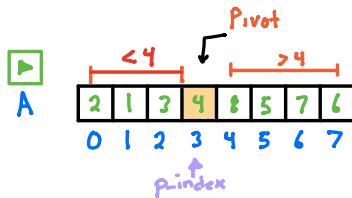


## Start



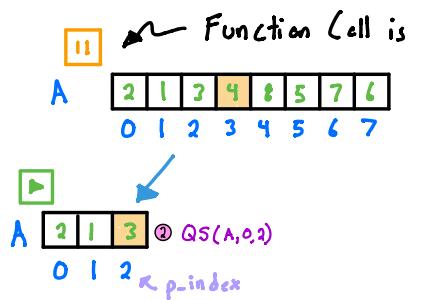
← Unsorted Array

- Since this satisfies our condition  $Start < End$  ( $0 < 7$ ), we will call the Partition function, which will find our Pivot and swap values Left or Right depending on if they're higher or lower than the Pivot point; it then will return a p-index, which will be index 3, for this step.



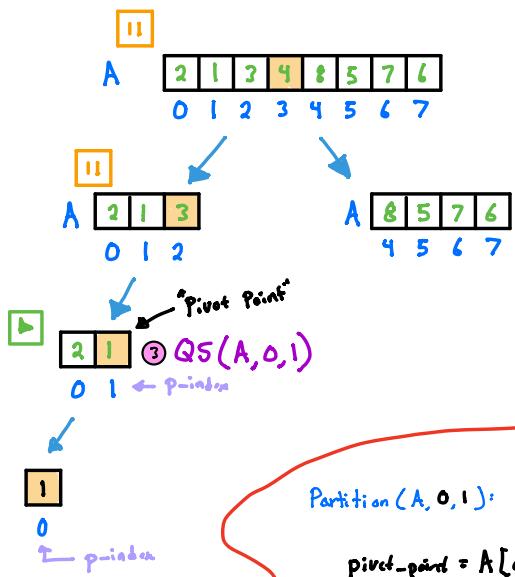
- ← We have 4 as our Pivot  
• The Partition function will rearrange the array to this; it will return the p-index

- With our p-index returned we will call the  $QuickSort(A, Start, pIndex - 1)$ , which will move our "End" index ( $3 - 1 = 2$ ), with our starting being 0, it will still satisfy our condition  $Start < End$  ( $0 < 2$ ), so we call the Partition function it takes in our array, end index value 0 to 2, finds a Pivot point of 3 and reverse things, and return  $pIndex = 2$ .



- machine says it will come back once it's finished.

- ~ All Values to the Left of the Pivot Point (3), were already less than 3.



1  
 0  
 $\leftarrow$  p-index

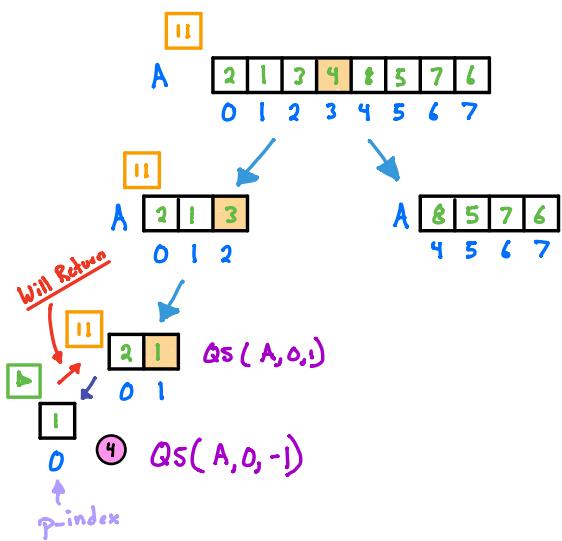
Partition(A, 0, 1):  
 pivot-point = A[end]  $\approx 2$   
 p-index = 0  
 Index 0  
 A[0], A[end] = A[end], [p-index]  
 1                        0  
 So 2 is now in index  
 ①

③ After the p-index is returned `QuickSort(A, Start, pIndex-1)`, is called. Our p-index goes  $(2-1)=1$ . (`Q5(A, 0)`). This still satisfies the condition  $Start < End$  ( $1 < 0$ ). So we call the Partition function again.

def Partition(A, Start, End):

```

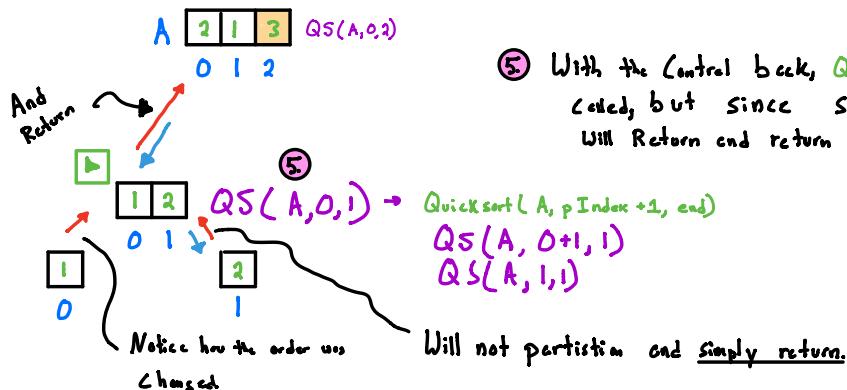
pivot = A[End]           ~ Selecting the Pivot
Partition_index = start   ~ index
for i in range(Start, End-1):
    if A[i] < pivot:
        A[i] = A[Partition_index]
        A[Partition_index] = A[i]
    Partition_index = Partition_index + 1
A[Partition_index], A[End] = A[End], A[Partition_index]
Return Partition_Index
  
```



- ④ We will call the `QuickSort(A, start, pIndex-1)`, with  $QS(A, 0, 0-1) = QS(A, 0, -1)$ , our condition `start > End` is false, so we Will Return, and Control Will return back to `QS(A, 0, 1)`.

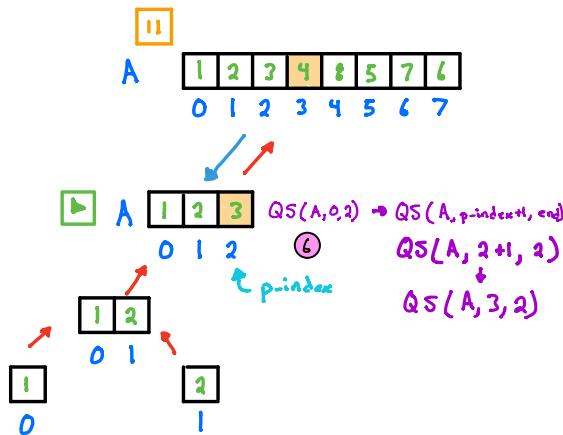
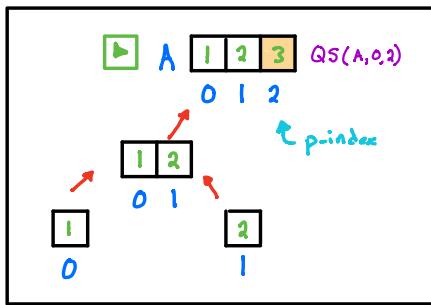
```
def QuickSort(A, start, end):
    Index
    if start < end:
        return
    pIndex = partition(A, start, end)
    Recursive call
    QuickSort(A, start, pIndex-1)
    QuickSort(A, pIndex+1, end)
```

QuickSort(A, 1, 1)  
Returned array

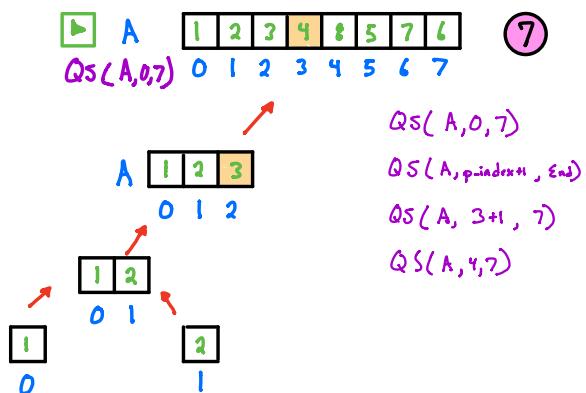


- ⑤ With the Control back, `QuickSort(A, pIndex+1, end)`, is called, but since `start < end` is false, it Will Return and return Control to `QS(A, 0, 2)`.

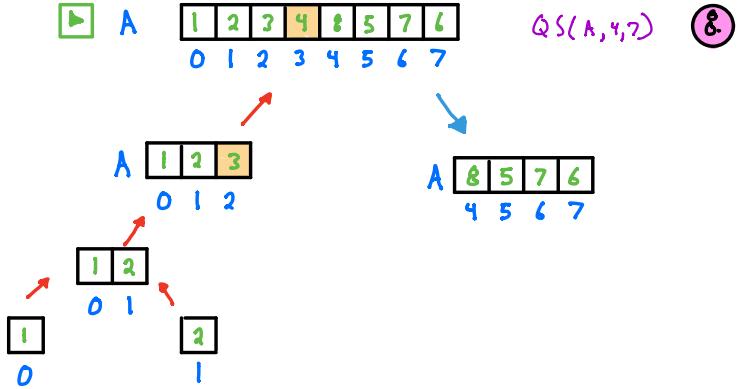
The index order for 1 and 2 were swapped earlier



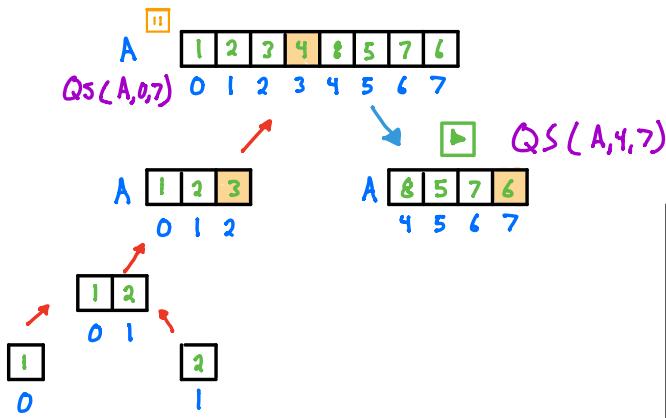
⑥ With Control back here,  
 $\text{Quicksort}(A, \text{pIndex} + 1, \text{end})$  is  
 called and  $\text{QS}(A, 3, 2)$  is  
 invalid, so QS will return  
 $(\rightarrow)$



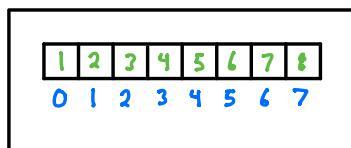
⑦ Our Original array at the top  
 will now resume  $\text{QS}(A, 0, 7)$ . Now  
 it will make a call right of the  
 pivot,  $\text{Quicksort}(A, \text{pIndex} + 1, \text{end})$ ,  
 which will be  $\text{QS}(A, 4, 7)$ .



And here are Starting Index Value is L and end index, so We can now partition the Other side of the array!



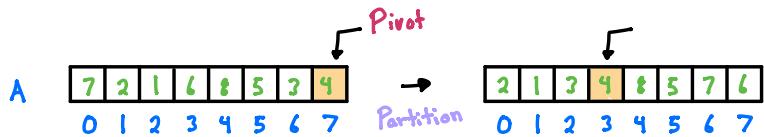
```
def Partition(A, Start, End):
    pivot = A[End] ~ Selecting the Pivot
    Partition_index = start ~ index
    for i in range(Start, End-1):
        if A[i] <= pivot:
            A[i] = A[Partition_index]
            A[Partition_index] = A[i]
            Partition_index += 1
    A[Partition_index], A[End] = A[End], A[Partition_index]
    Return Partition_Index ~ Swap
```



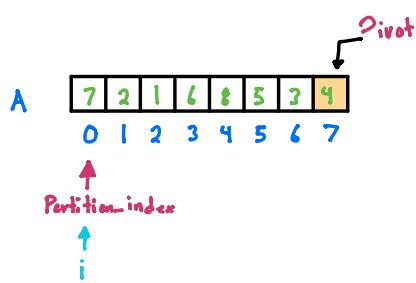
Sorted

```
def QuickSort(A, start, end):  
    if start > end:  
        return  
  
    p_index = Partition(A, start, end)  
  
    # recursive Calls here:  
  
    QuickSort(A, start, p_index-1)  
    QuickSort(A,p_index +1, end)
```

## Partition Logic



- This what we want, to select a pivot and to reArrange a list such that all the elements lesser than the pivot are to the left and all the elements greater than the pivot are to the right of it.



```
def Partition(A ,start, end):  
    pivot = A[end]  
    p_index = start  
  
    for i in range(start,end):  
        if A[i] <= pivot:  
            #swap number at index with the P_index!  
            A[p_index], A[i] = A[i], A[p_index]  
            p_index = p_index +1  
  
    # this is swaping the pivot point between the sorted values  
    A[p_index], A[end] = A[end], A[p_index]  
  
    return p_index
```

**Running Time of Quicktime**

## Actual Code

```
def Partition(A ,start, end):
    pivot = A[end]
    p_index = start

    for i in range(start,end):
        print('Index:', i)

        if A[i] <= pivot:

            #swap number at index with the P_index!
            A[p_index], A[i] = A[i], A[p_index]

            p_index = p_index +1

    # this is swaping the pivot point between the sorted values
    A[p_index], A[end] = A[end], A[p_index]

    return p_index

def QuickSort(A, start, end):

    if start > end:
        return

    p_index = Partition(A, start, end)

    # recursive Calls here:
    QuickSort(A, start, p_index-1)
    QuickSort(A,p_index +1, end)

array = [7,2,1,6,8,5,3,4]
n = len(array)
QuickSort(array,0,n-1)
print(array)
```