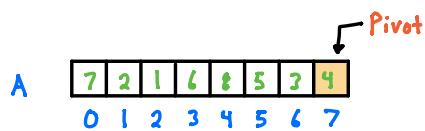


## Quick Sort

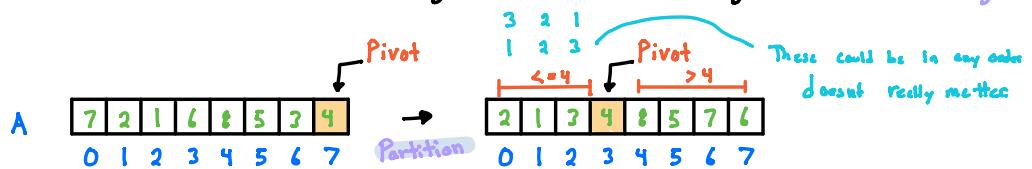
- Divide and Conquer
- $O(n \log n)$  - Average case running time
- $O(n^2)$  - Worst case running time
- In-place

## Quick Sort Overview

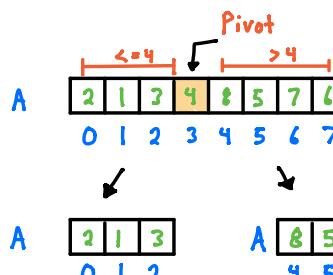


1) We first Select one elements in the list (Can be any element), We Call this selective number, Pivot.

- 2) Now we rearrange the list such that all the elements lesser than the Pivot are towards the left of it and all the elements greater are towards the right of it (Partitioning of a list)



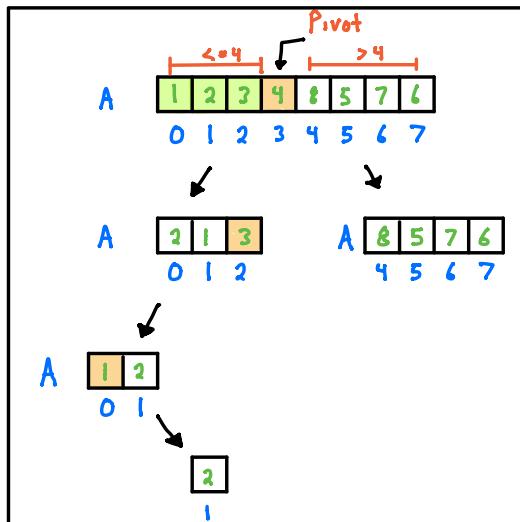
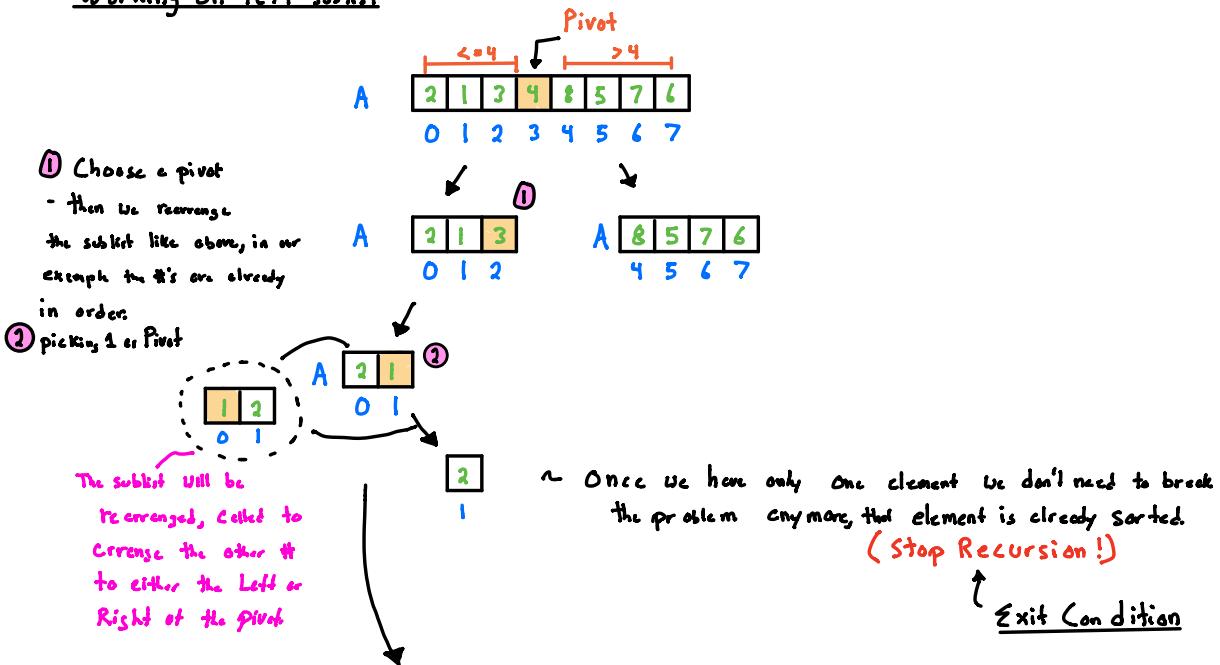
- We can break this problem into two sub-problems, these being sorting the array to the left of the pivot and sorting the array to the right of the pivot.



## Different than Merge Sort

- Unlike merge sort we do not need to create auxiliary arrays entirely, we can work on the same array. We just have to keep track of the start and end index of the segments.

Working on left sublist:



Then we work back through the problem.

## QuickSort Function Suedo Code

```
def QuickSort( A, start, end): Index  
    if start < end: ] exit condition  
        return  
        pIndex = Partition( A, start, end)  
        Recursion calls ↘ QuickSort( A, start, pIndex-1)  
        QuickSort( A, pIndex+1, end)
```

```
def Partition( A, Start, End):  
    pivot = A[End] ~ Selecting the Pivot  
    Partition-index = start  
  
    for i from start to end-1:  
        if (A[i] <= pivot):  
            A[i] = A[partition-index] ~ Swap  
            Partition-index = partition-index + 1  
  
    A[Partition Index] = A[End] ~ Swap  
    Return Partition-Index
```

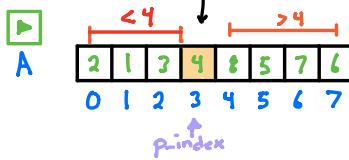
## Start

① QS(A, 0, 7) A

7	2	1	6	8	5	3	4
0	1	2	3	4	5	6	7

← Unsorted Array

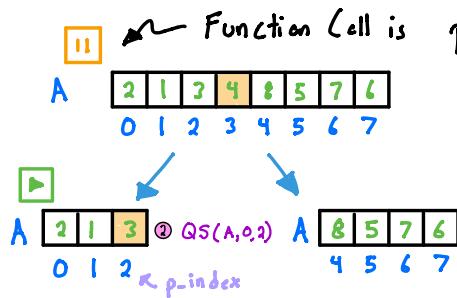
- ① Since this satisfies our condition  $\text{Start} < \text{End}$  ( $0 < 7$ ), we will call the Partition function, which will find our Pivot and swap values Left or Right depending on if they're higher or lower than the Pivot point. It then will return a p-index, which will be index 3, for this step.



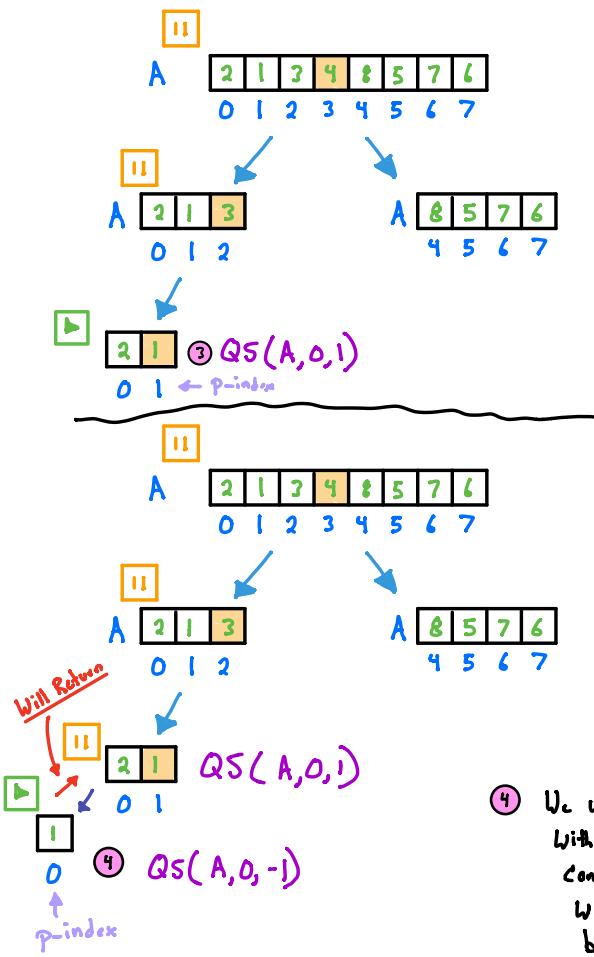
← We have 4 as our Pivot

- The Partition function will rearrange the array to this; it will return the p-index

- ② With our p-index returned we will call the `QuickSort(A, Start, pIndex-1)`, which will move our "End" index ( $3-1 = 2$ ), with our starting being 0, it will still satisfy our condition  $\text{Start} < \text{End}$  ( $0 < 2$ ), so we call the Partition function it takes in our array, end index value 0 to 2, finds a Pivot point of 3 and reverse things, and return  $p\text{-index} = 2$ .

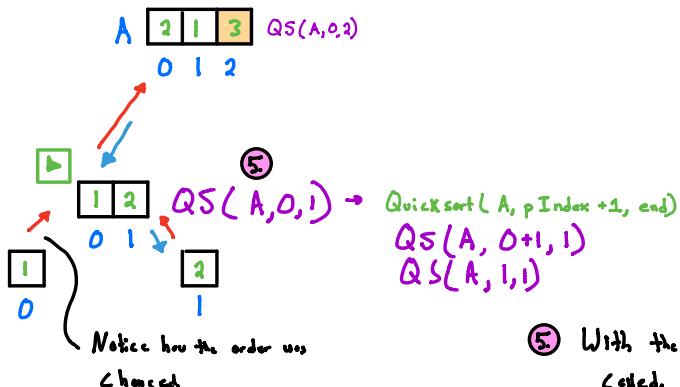


- Machine says it will come back once it's finished.

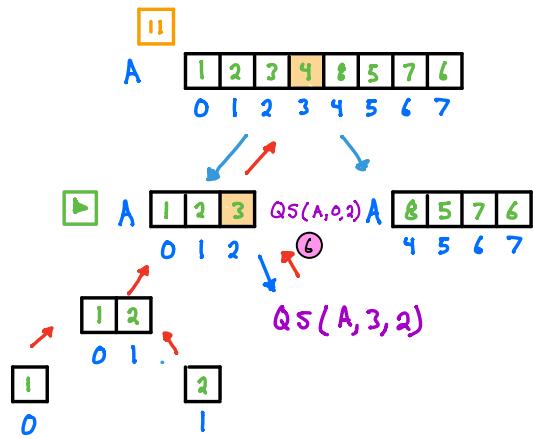
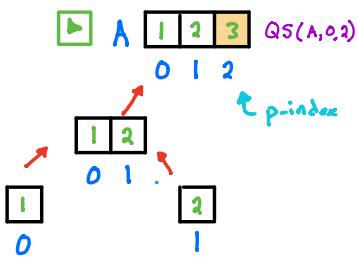


③ After the p-index is returned  $QuickSort(A, Start, pIndex-1)$ , is called. Our p-index goes  $(2-1)=1$ . ( $QS(A, 0, 1)$ ). This still satisfies the condition  $Start > End$  ( $1 > 0$ ). So we call the Partition function again. It will make the pivot become 1. and the p-index will become 0.

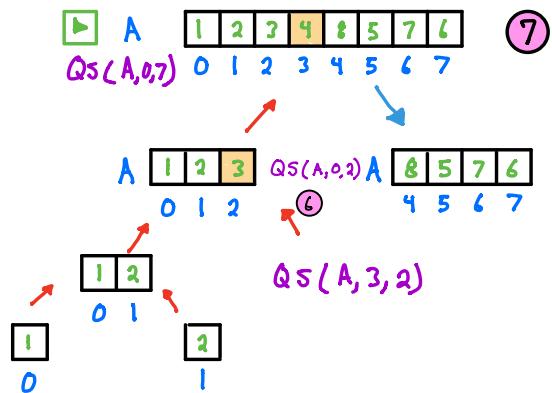
④ We we will call the  $QuickSort(A, Start, pIndex-1)$ , with  $QS(A, 0, 0-1) = QS(A, 0, -1)$ , our condition  $Start > End$  is false, so we will return, and control will return back to  $QS(A, 0, 1)$ .



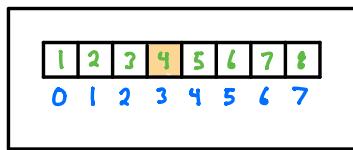
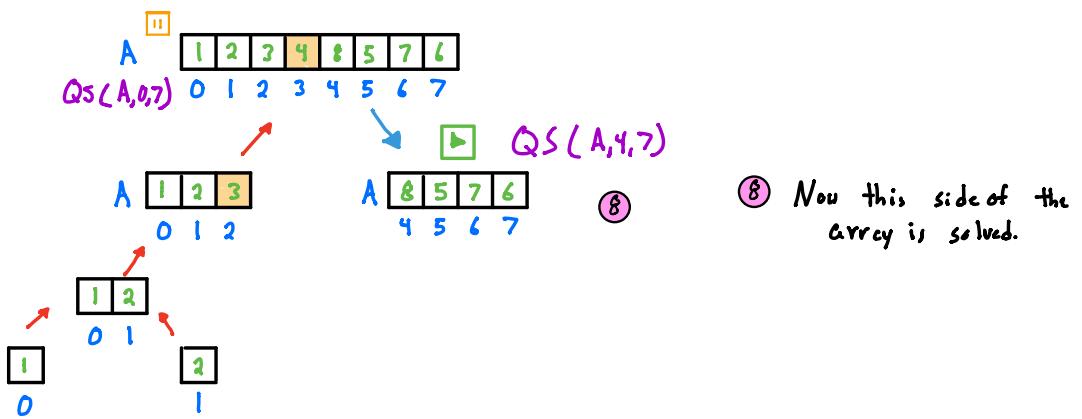
⑤ With the control back,  $QuickSort(A, pIndex + 1, end)$ , is called, but since  $Start < end$  is false, it will return and return control to  $QS(A, 0, 1)$ .



⑥ With Control back here,  
Quicksort(A, pIndex + 1, end) is  
called and QS(A,3,2) is  
invalid, so QS will return  
(→)

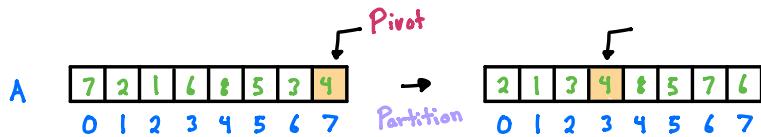


⑦ Our Original array at the top  
will now resume QS(A,0,7). Now  
it will make a call right of the  
pivot, Quicksort(A, pIndex + 1, end),  
which will be QS(A,4,7).

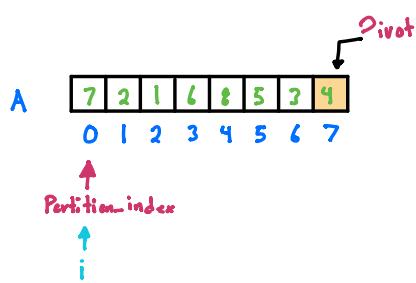


Sorted

## Partition Logic



- This is what we want, to select a pivot and to rearrange a list such that all the elements lesser than the pivot are to the left and all the elements greater than the pivot are to the right of it.



```
def Partition(A, Start, End):  
    pivot = A[End]      ~ Selecting the Pivot  
    Partition_index = Start  
  
    for i from start to end-1:  
        if (A[i] <= pivot):  
            A[i] = A[partition_index] ~ Swap  
            partition_index = partition_index + 1  
  
    A[Partition Index] = A[End] ~ Swap  
  
    Return Partition_Index
```

**Running Time of Quicktime**

Actual Code