

Mathematical Induction Chapter 5.1-5.4

5.1

Proof Using Mathematical Induction - Summation Formulas

Mathematical Induction

To prove $P(x)$ is true for $x \in \mathbb{Z}^+$, where $P(x)$ is a propositional function, we complete two steps:

1) Basis Step - Verify $P(1)$ is true

2) Inductive Step - Verify if $P(k)$ is true, then $P(k+1)$ is true $\forall k \in \mathbb{Z}^+$

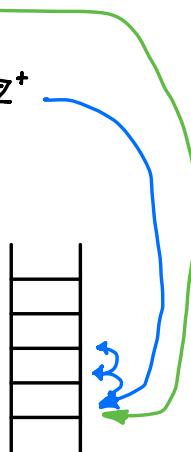
* Inductive Hypothesis: $P(k)$ is true

* Must Show: $P(k) \rightarrow P(k+1)$

Conclusion: $P(x)$ is true $\forall k \in \mathbb{Z}^+$

Other instructors might
not write this

like a ladder
with one foot rung



Proving a Summation Formula

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}$$

Show $1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$

Let $P(n) : 1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$

① Basis: Prove $P(1)$ is true

$$P(1) = 1 = \frac{1(1+1)}{2} \quad 1 = \frac{1(1)}{2} \quad 1 = 1 \checkmark$$

~ Prove Basis Step
~ easy its just math

② Inductive step: $P(k) \rightarrow P(k+1)$

Inductive Hypothesis: $1 + 2 + 3 + \dots + k = \frac{k(k+1)}{2}$

* Must Show: $1 + 2 + 3 + \dots + k + (k+1) = \frac{(k+1)((k+1)+1)}{2} = \frac{(k+1)(k+2)}{2}$

Exactly what we wanted

$$1 + 2 + 3 + \dots + k + (k+1) = \frac{k(k+1)}{2} + (k+1) \quad \begin{matrix} \downarrow \\ \text{write as fraction} \end{matrix} \quad \begin{matrix} 1^{\text{st}} \text{ Step: Added } (k+1) \text{ to both sides} \\ \end{matrix}$$

$$= \frac{k(k+1)}{2} + \frac{2(k+1)}{2} \quad \begin{matrix} \downarrow \\ \text{2nd Step: Simplify solution} \end{matrix}$$

$$= \frac{(k+1)(k+2)}{2}$$

∴ $P(n)$ is true
 $\forall n \in \mathbb{Z}$

Problem 2:

Use mathematical induction to show that for all non-negative integer n , $1+2+2^2+\dots+2^n = 2^{n+1}-1$
 $n > 0$

$$1+2+2^2+\dots+2^n = 2^{n+1}-1$$

$$\text{Let } P(n) : 1+2+2^2+\dots+2^n = 2^{n+1}-1$$

① Basis: Prove $P(0)$ is true

$$\begin{aligned} 2^0 &= 2^{0+1}-1 \\ 1 &= 1 \checkmark \end{aligned}$$

② Inductive step: $P(k) \rightarrow P(k+1)$

$$\text{Inductive Hypothesis: } 1+2+2^2+\dots+2^k = 2^{k+1}-1$$

$$\star \text{ Must Show: } 1+2+2^2+\dots+2^k + 2^{k+1} = 2^{(k+1)+1} = \boxed{2^{k+2}-1} \quad \leftarrow \begin{array}{l} \text{We want to get to} \\ \text{this} \end{array}$$

$$1+2+2^2+\dots+2^k + 2^{k+1} = 2^{k+1}-1 + 2^{k+1} \quad \sim \text{Add to both sides}$$

$$\begin{aligned} &= 2(2^{k+1}) - 1 \\ &\quad \nearrow \text{Can add exponent} \\ &= \boxed{2^{k+2}-1} \end{aligned}$$

$\therefore P(n)$ is True
 for all $n \geq 0$
 $n \in \mathbb{Z}$

Problem 3

Conjecture and prove a summation formula for the sum of the first n positive odd integers.

$n = 1$	$1 = 1$	1^2	Looking for a Pattern
$n = 2$	$1+3 = 4$	2^2	
$n = 3$	$1+3+5 = 9$	3^2	
$n = 4$	$1+3+5+7 = 16$	4^2	
\vdots	\vdots	\vdots	
$n = n$	$1+3+5+\dots+(2n-1) = n^2$		$\sim P(n)$

Prove $1+3+5+\dots+(2n-1) = n^2$

Let $P(n)$: the proposition that $1+3+5+\dots+(2n-1) = n^2$

① Basis: Prove $P(1)$ is true.

$$\begin{array}{l} 1 = (1)^2 \\ 1 = 1 \quad \checkmark \end{array}$$

② Inductive step: $P(k) \rightarrow P(k+1)$

Inductive Hypothesis: $1+3+5+\dots+(2k-1) = k^2$

$$\text{Show: } 1+3+5+\dots+(2k-1) + (2k+1) = (k+1)^2$$

if we added the next odd integer, which is +2.

$$1+3+5+\dots+(2k-1) + (2k+1) = k^2 + (2k+1)$$

$$1+3+5+\dots+(2k-1) + (2k+1) = (k+1)(k+1) \quad - \text{Factor}$$

$$1+3+5+\dots+(2k-1) + (2k+1) = (k+1)^2$$

$\therefore P(n)$ is true.

Proof Using Mathematical Induction - Inequalities

5.1

Proving an Inequality:

Prove $n < 2^n \quad \forall n \in \mathbb{Z}^+$ using mathematical induction.

Let $P(n): n < 2^n \quad \forall n \in \mathbb{Z}^+$

① Basis: $P(1)$

$$P(1): \begin{array}{l} 1 < 2^1 \\ 1 < 2 \end{array} \checkmark$$

② Inductive $P(k) \rightarrow P(k+1)$

IH: $k < 2^k$

Show: $k+1 < 2^{k+1}$

$$k+1 < 2^k + 1$$

- ~ Since it's an inequality that don't have same rules as an equals
- ~ We just have to make sure the inequality still holds

We know that: $1 < 2^k \quad \forall k \in \mathbb{Z}^+$

~ Proved this in Basis Step ①

$$k+1 < 2^k + 1 < 2^k + 2^k = 2(2^k)$$

- How can we do this?

B/c $1 < 2^k$

$$\boxed{k+1 < 2^{k+1}}$$

- it's ok to do this b/c were not breaking any mathematical law.

- Replacing, inequality still works!

Example 2:

Prove $2^n < n!$ $\forall n \in \mathbb{Z}^+$ and $n \geq 4$.

Let $P(n) : 2^n < n!$

① Basis $P(4)$

$$2^4 < 4! \rightarrow 16 < 24$$

$$16 < 24 \quad \checkmark$$

② Inductive $P(k) \rightarrow P(k+1)$

$$\text{IH: } 2^k < k!$$

$$\text{SHOW: } 2^{k+1} < \underline{(k+1)!}$$

Trying to get here!

$$2^k \cdot 2^1 < k! \cdot 2$$

$$k+1 > 2$$

$$2^{k+1} < \underline{2k!}$$

$$2^{k+1} < (k+1)k!$$

- Just have to make sure the inequality still holds

$$(k+1)(k)(k-1)\dots(2)(1) = (k+1)!$$

$$\boxed{2^{k+1} < (k+1)!}$$

Proof Using Mathematical Induction - Divisibility

5.1

Use mathematical induction to prove $7^{n+2} + 8^{2n+1}$ is divisible by 57 for all non-negative integers n.

Let $P(n)$: $7^{n+2} + 8^{2n+1}$ is divisible by 57

① Basis: $P(0)$

$$P(0): 7^{0+2} + 8^{2(0)+1} = 7^2 + 8 = 49 + 8 = \underline{\underline{57}} \quad \sim \text{is divisible by 57}$$

② Inductive:

I H: $7^{k+2} + 8^{2k+2}$ is divisible by 57

Show: $7^{k+3} + 8^{2k+3}$ is divisible by 57

$7 \cdot 7^{k+2} + 8^{2k+2} \cdot 8^1 = 7^{k+3} + 8^{2k+3}$

\sim Shown we can turn the I H into what were suppose to Show.

If you multiply out

$= 7 \cdot 7^{k+2} + (7 + 57) 8^{2k+1}$

- Lets rewrite the right-hand side.

$= 7(7^{k+2} + 8^{2k+1}) + 57 \cdot 8^{2k+1}$

- By Inductive Hypothesis this is divisible by 57

- Then this portion is also divisible by 57

Since $7^{k+2} + 8^{2k+1}$ is divisible by 57 and $57 \cdot 8^{2k+1}$ is divisible by 57,

then $7^{k+3} + 8^{2k+3}$ is divisible by 57.

$\therefore 7^{n+2} + 8^{2n+1}$ is divisible by 57 for all $n \geq 0$, $n \in \mathbb{Z}$

The Well-Ordering Principle and Strong Induction

5.2

The Well-Ordering Principle

Every non-empty set of non-negative integers has a least element

$$\begin{array}{ll} \mathbb{Z}^+ & n = 1 \\ n \in \mathbb{Z}^+ \quad \underline{n \geq 4} & n = q \\ \text{non-negative int} & n = 0 \end{array}$$

Prove that every amount of postage of \$ 0.12 or more can be formed
Using \$ 0.04 and \$ 0.05 stamps.

$P(n)$ is the statement that postage of n -cents can be formed
Using 4-cent and 5-cent stamps if $n \geq 12$.

Basis: Show postage of 12-cents can be made

Inductive: Show if $P(k)$ is true, then $P(k+1)$ is true for $k \geq 12$

Revisiting Recursive Definitions

Fibonacci Numbers

Recall the set of Fibonacci numbers:

$$0, 1, 1, 2, 3, 5, 8, 13, 21, \dots$$

How each number found?

add 2 previous terms

① Initial Conditions

$$f_0 = 0$$

$$f_1 = 1$$

② Function

$$f_n = f_{n-1} + f_{n-2}, \quad n \geq 2$$

Recursive Defined Functions

Basis Step: Specifies the value of the function for the first term(s)

Recursive Step: Gives a rule for finding subsequent values using a previous value(s) beginning at those defined in the basis step.

If f is defined recursively by $f(0) = 2$ and $f(n+1) = 3f(n) - 1$, find $f(1)$, $f(2)$, $f(3)$, and $f(4)$.

Give a recursive definition for a^n for $a \in \mathbb{R}$ and non-negative and $n \in \mathbb{Z}^+$

Give a recursive definition of the sequence $\{a_n\}, n=1, 2, 3, \dots$ if $a_n = 2n + 1$

Recursively Defined Sets and structures

Basis Step: Specifies an initial collection of elements

Recursive Step: Gives rules for forming new elements from those already in the set.

Recursively Defined Set

Let S be a subset of the integers defined recursively by:

Basis step:

Recursive Step:

List the elements of S produced by the first 5 application of the recursive definition.

Recursively Defined Structure

A set of rooted trees, where a rooted tree consists of a set of vertices

Example of Rooted Trees Defn:

Basis Step:

Recursive Step 1:

Step 2:

Structural Induction

Recursive Algorithms

An algorithm is called recursive if it solves a problem by reducing it to an instance of the same problem with a smaller input.