



# Discrete Time and Race Conditions

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# Discreteness

Race conditions are where order is important but ambiguous

This is a modeling phenomenon

# Discreteness



Types

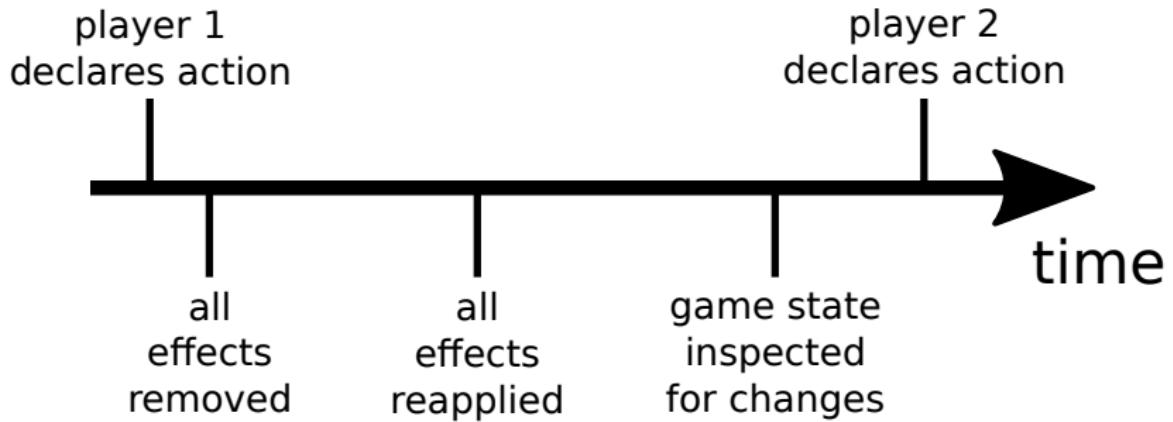
Abilities

Power/Toughness

# Discreteness



# Discreteness



# Discreteness

## 613. Interaction of Continuous Effects

613.1. The values of an object's characteristics are determined by starting with the actual object. For a card, that means the values of the characteristics printed on that card. For a token or a copy of a spell or card, that means the values of the characteristics defined by the effect that created it. Then all applicable continuous effects are applied in a series of layers in the following order:

613.1a *Layer 1*: Copy effects are applied. See rule 706, “Copying Objects.”

613.1b *Layer 2*: Control-changing effects are applied.

613.1c *Layer 3*: Text-changing effects are applied. See rule 612, “Text-Changing Effects.”

613.1d *Layer 4*: Type-changing effects are applied. These include effects that change an object's card type, subtype, and/or supertype.

613.1e *Layer 5*: Color-changing effects are applied.

613.1f *Layer 6*: Ability-adding effects, ability-removing effects, and effects that say an object can't have an ability are applied.

613.1g *Layer 7*: Power- and/or toughness-changing effects are applied.

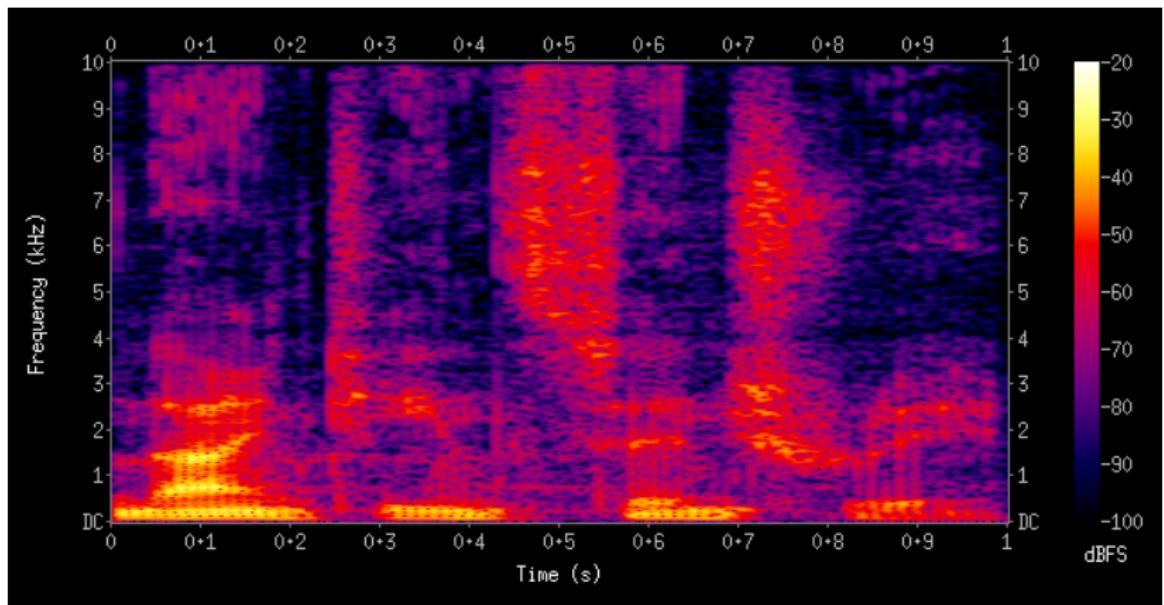


# Discreteness

Race conditions occur because of  
discretization of continuous systems

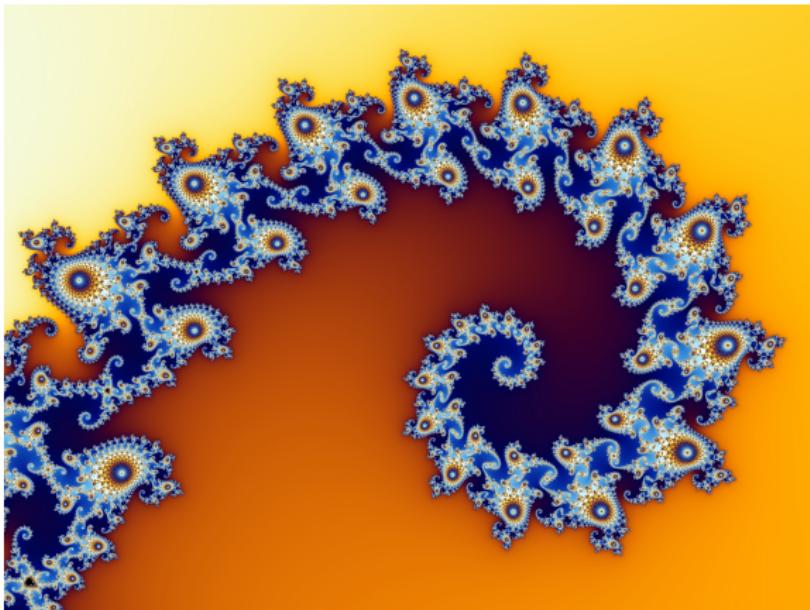
So why not deal in continuous systems?

# Continuousness



In general, continuous systems lack representation

# Continuousness



At best, we define processes then discretize them

# Continuousness

## Example

Select two random reals between 0 and 1

The probability of selecting the same real both times is 0

$$P = \lim_{x \rightarrow \infty} \frac{1}{x} = 0$$

# Continuousness

How does one select a random real?

We could define an algorithm, but it's nonterminating

So this example lacks implementation



# Referential Transparency

A pure expression can be replaced with its result without affecting the program

This lacks implementation as well



# Referential Transparency

Referential transparency is the single-write strategy

In other words, dependency elimination so no race conditions



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What should we do with mutable state?

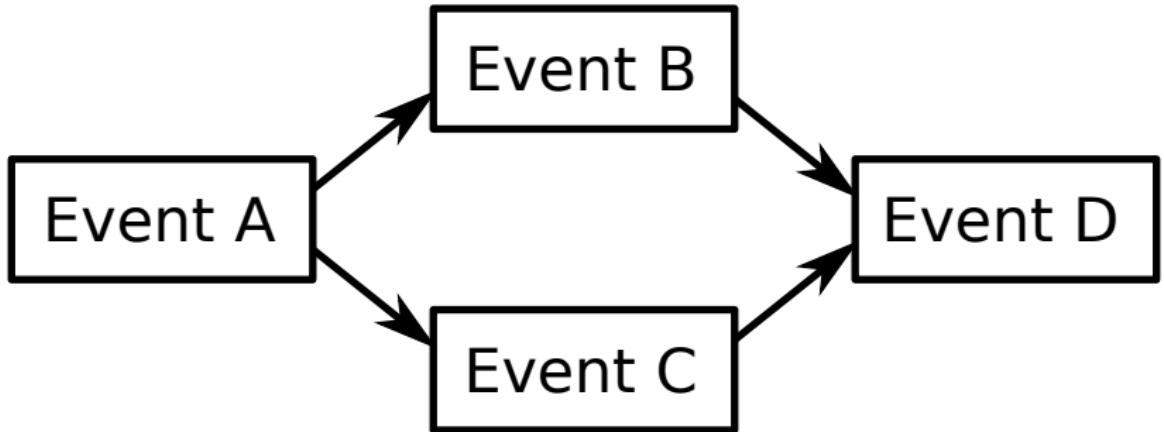


# Dependency Graphs

In discrete time, order isn't defined below temporal resolution

This is how we get race conditions

# Dependency Graphs



# Dependency Graphs

- ▶ Language syntax should be isomorphic to dependency edges
  - statements of the form
    - (event)
    - (stateful reads and pure transforms)
    - (write target)

# Dependency Graphs

- ▶ Language syntax should be isomorphic to dependency edges
  - statements of the form
    - (event)
    - (stateful reads and pure transforms)
    - (write target)
- ▶ Compiler should forbid diamond patterns
- ▶ Requires language-level implementation
- ▶ Style should adapt accordingly

# Dependency Graphs

Is this too strict?

Is this scalable?

What concessions can we make? What other nodes/subgraphs are safe?

We need a language first...



# Conclusion

Race conditions are the biggest problem

Referential transparency solves huge classes of them

Language-level dependency analysis may work



# Conclusion

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