

→ Dynamic Programming

- Identify and solve subproblems.
- Use Optimal substructure to construct/build the solution to the larger problem.
  - Local optimal solution of the smaller subproblems.

Longest Common Subsequence

- If letters match  $m(i-1, j-1) + 1$     chk for greater. ✓

- If letters doesn't match  $m(i-1, j) \vee m(i, j-1)$     ✓

		i	n	t	r	a	n	e	t
	0	1	2	3	4	5	6	7	8
i	0	1	1	1	1	1	1	1	1
n	0	1	2	2	2	2	2	2	2
t	0	1	2	3	3	3	3	3	3
e	0	1	2	3	3	3	3	4	4
r	0	1	2	3	4	4	4	4	4
n	0	1	2	3	4	4	5	5	5
e	0	1	2	3	4	4	5	6	6
t	0	1	2	3	4	4	5	6	7

Back Tracking

Length = 7 of Longest Common Subsequence

→ internet  
→ intranet > intranet

Formula  $\rightarrow$  if ( $A[i] = B[i]$ )

$$MAT[i, j] = MAT[i-1, j-1] + 1$$

else

$$MAT[i, j] = \max \begin{cases} MAT[i-1, j] \\ MAT[i, j-1] \end{cases}$$

Time Complexity :-

$$O(n * m)$$

Example #1

$\begin{matrix} \text{net} & \text{er} & \text{int} \\ \text{in} & \text{er} & \text{net} \\ \text{nt} & & \end{matrix}$

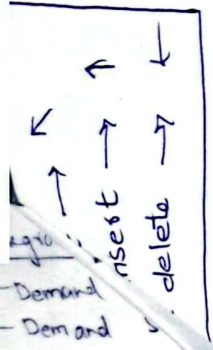
$\rightarrow$  Verify at home

Example #2

# through backtracking

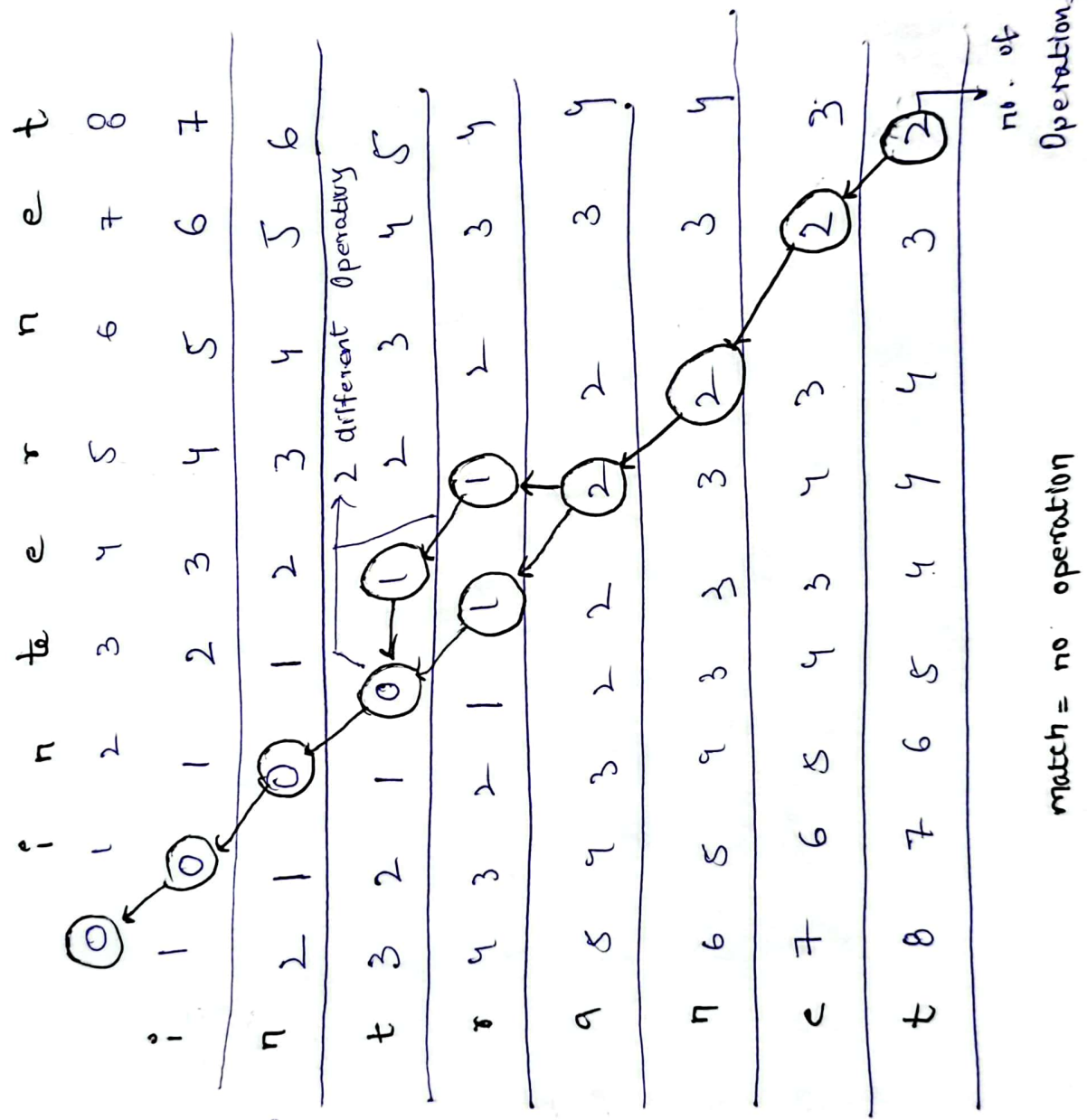
$\begin{matrix} \text{in} & \text{er} & \text{net} \\ \text{te} & \text{er} & \text{net} \end{matrix} \rightarrow \text{length} = 5$

		i	n	t	e	r	n	e	t
t	0	0	0	0	0	0	0	0	0
e	0	0	0	0	1	2	2	2	2
n	0	0	1	1	1	2	3	3	3
r	0	0	1	1	1	2	3	3	3
e	0	0	1	1	2	3	3	4	4
t	0	0	1	2	2	3	3	4	5
n	0	0	1	2	2	3	4	4	5
i	0	1	1	2	2	3	4	4	5



Efficiency of Demand Paging  
 - Less I/O and memory needed.

\* Minimum Edit Distance



in ter net  
 in t r a n s m i t