

65 points total

1.

$\frac{3}{5}$  for process, no tree

$\frac{1}{5}$  for process, incorrect tree

$\frac{2}{5}$  for only showing order, no tree

$\frac{1}{5}$  for tree, no work

$\frac{5}{5}$  if tree is right, work or not

a. (5 points)

(indegree, outdegree)

A: (1, 2)

B: (1, 3)

C: (2, 2)

D: (2, 2)

E: (3, 2)

F: (2, 0)

G: (1, 1)

b. BFS tree (5 points)



$Q = \{B\}$

$D = \{\}$

**Process B**

$Q = \{C, E, G\}$

$D = \{B\}$

**Process C**

$Q = \{E, G, D\}$

$D = \{B, C\}$

**Process E**

$Q = \{G, D, F\}$

$D = \{B, C, E\}$

**Process G**

$Q = \{D, F\}$

$D = \{B, C, E, G\}$

**Process D**

$Q = \{F, A\}$

$D = \{B, C, E, G, D\}$

**Process F**

$Q = \{A\}$

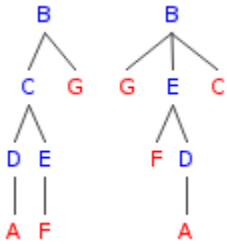
$D = \{B, C, E, G, D, F\}$

**Process A**

Q = {}

D = {B, C, E, G, D, F, A}

c. DFS Tree (5 points)



S = {B}

D = {}

**Process B**

S = {C, E, G}

D = {B}

**Process G**

S = {C, E}

D = {B, G}

**Process E**

S = {C, D, F}

D = {B, G, E}

**Process F**

S = {C, D}

D = {B, G, E, F}

**Process D**

S = {C, A}

D = {B, G, E, F, D}

**Process A**

S = {C}

D = {B, G, E, F, D, A}

**Process C**

S = {}

D = {B, G, E, F, D, A, C}

{B, C, D, A, F, E, G}

2.

a. (10 Points... -2 pts/incorrect path)

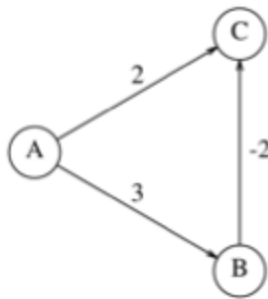
A to ...	Shortest Weighted Path	Weighted Length
B	A => B	5
C	A => C	3
D	A => B => G => E => D	9
E	A => B => G => E	7

F	$A \Rightarrow B \Rightarrow G \Rightarrow E \Rightarrow F$	8
G	$A \Rightarrow B \Rightarrow G$	6

b. (5 points... -1 pt/incorrect path)

B to ...	Shortest Unweighted	Unweighted Length
A	$B \Rightarrow C \Rightarrow D \Rightarrow A$	3
C	$B \Rightarrow C$	1
D	$B \Rightarrow C \Rightarrow D$	2
E	$B \Rightarrow E$	1
F	$E \Rightarrow E \Rightarrow F$	2
G	$B \Rightarrow G$	1

3. (5 points)



4. (10 points)

(b) Use an array  $numEdges$  such that for any vertex  $u$ ,  $numEdges[u]$  is the shortest number of edges on a path of distance  $d_u$  from  $s$  to  $u$  known so far. Thus  $numEdges$  is used as a tiebreaker when selecting the vertex to mark. As before,  $v$  is the vertex marked known, and  $w$  is adjacent to  $v$ .

If  $d_v + c_{v,w} = d_w$ , then change  $p_w$  to  $v$  and  $numEdges[w]$  to  $numEdges[v]+1$  if  $numEdges[v]+1 < numEdges[w]$ .

If  $d_v + c_{v,w} < d_w$ , then update  $p_w$  and  $d_w$ , and set  $numEdges[w]$  to  $numEdges[v]+1$ .

5. (5 points) Check diagonals of resulting tables to see if there are any negative values.

6.

a. (10 points)

D	1	2	3	4	5
1	0	-10	-9	-7	-3
2	inf	0	2	4	8
3	inf	-1	0	3	7
4	inf	-3	-2	0	5

<b>5</b>	inf	-7	-6	-4	0
----------	-----	----	----	----	---

<b>P</b>	<b>1</b>	<b>2</b>	<b>3</b>	<b>4</b>	<b>5</b>
<b>1</b>	-	3	4	5	1
<b>2</b>	-	-	4	5	2
<b>3</b>	-	3	-	5	2
<b>4</b>	-	3	4	-	2
<b>5</b>	-	3	4	5	-

b. (5 points)

<b>Path S/E</b>	<b>1</b>	<b>2</b>	<b>3</b>	<b>4</b>	<b>5</b>
<b>1</b>	-	15432 (-10)	1543 (-9)	154 (-7)	15 (-3)
<b>2</b>	-	-	2543 (2)	254 (4)	25 (8)
<b>3</b>	-	32 (-1)	-	3254 (3)	325 (7)
<b>4</b>	-	432 (-3)	43 (-2)	-	4325 (5)
<b>5</b>	-	5432 (-7)	543 (-6)	54 (-4)	-