





we need to add two new types to the symbol table (VAL, REF) also add symbols and reserved words

how do we know whether a parameter is VAL/REF in the ST?

have an array of bools for each procedure/function parameter:

0/F->VAL; 1/T->REF

table[i].params[j] where i is the procedure/function ST index and j is the parameter index

we should not be able to pass constants or numbers to REF parameters

parameter loading

value: LOD (normal load)

reference: LDI (a new load indirect)

parameter storing

value: STO (normal store)

reference: STI (a new store indirect)

calling:

caller / callee	VAL	REF
VAL	LOD	LDA
REF	LDI	LOD
VAR	LOD	LDA
CONST	LIT	
NUM	LIT	

LDA (a new load address)

note how VAL and VAR are similar—we can simply treat them the same!

so we need 3 new instructions

remember that base(x) finds the base, x levels down existing instructions:

```
STI: s[s[base(table[i].level)+table[i].adr]] = s[t];
    t--;
LDI: t++;
    s[t] = s[s[base(table[i].level)+table[i].adr]];
LDA: t++;
    s[t] = base(table[i].level)+table[i].adr;
```

we must allow declaration of value and reference parameters they just behave like normally declared variables

the first parameter is at an offset of 3 in the segment, the second at 4, etc

```
additional declared variables in the block increment from there
       it's hard to parse for parameters in the normal block syntax diagram
              because the level of parameters is incorrect
              and also the variable address offset is only valid for the current block
       so let's just move that part to the top of block and allow "entrance" only if lev>0
              which means we're in the current procedure or function
              we also move the semicolon after the procedure/function ident to the top of block
              now we can just use dx as normal to set the correct address!
              one additional thing is to set the parameter boolean of the procedure/function properly
                      we do this when we enter the parameters in the ST
                             if (sym == valsym)
                             {
                                    enter(value);
                                    table[tabinx0].params[dx-4] = false;
                             }
                      why dx-4?
                             dx starts at 3
                             bool params starts at 0 (we need to shift it back)
                             but dx increments right after we add a val/ref parameter
                             so if we modify the param array right after adding a parameter
                                    we need dx-3-1 = dx-4
                      of course we could reverse the two instructions and it would be dx-3 (right?)
              for the designer (geek), we put it where it is most efficient and easiest
              for the user (ignoramus), we put it where s/he expects it to be
we then have to figure out how to handle storing to val/ref within procedures and functions
       this is done in statement
              we modify ident := <expression> to also allow assigning to VAL/REF
              if the ident is variable/value: STO lev-table[i].level, table[i].adr
              if the ident is a function: STO 0, -1 (so long as we're in the proper function!)
              if the ident is REF: STI lev-table[i].level, table[i].adr
we also need to figure out how to handle loading of val/ref within procedures and functions
       this is done in factor
              we modify ident to allow loading VAL/REF
              if the ident is variable/value: LOD lev-table[i].level, table[i].adr
              if the ident is const: LIT 0, table[i].val
              if the ident is ref: LDI lev-table[i].level, table[i].adr
we need to allow calling procedures and functions while passing parameters
       we do this in both statement (for procedures) and factor (functions)
              in either case, we preliminarily INT 0, 3 to temporarily make space for SL, DL, RA
                      this is like a "fake call" just to load parameters
              then we LOD, LIT, LDI, or LDA depending on type
              this really depends on two things:
                      what type is the variable we're trying to pass
                      what type is the procedure/function parameter (receiver)
              if the receiver is VAL then just call expression (we only need either LOD or LDI)
              if the receiver is REF then we use LDA for VAL, VAR and LOD for REF
```

once we're done loading parameters, we decrement the stack to undo the fake call

INT 0, -(3+# parameters)