Foundations of C.S.

Spring, 2023

PRINT NAME: \mathcal{SIGN} :

1. (5 pts) Let L(G) be the language of the grammar

$$\begin{array}{ccc} G:S & \rightarrow & AB \mid \lambda \\ & A & \rightarrow & aAa \mid bAb \mid cAc \mid \lambda \\ & B & \rightarrow & acacB \mid bcbcB \mid \lambda \end{array}$$

Set $M = \{w \in \{a, b, c\}^* \mid n_a(w), n_b(w), n_c(w) \text{ are all even}\}$. Prove that $L(G) \subseteq M$. Hint: You will need induction.

 \clubsuit We will prove by induction on the number of rules applied that every sentential w form has $n_a(w)$, $n_b(w)$, $n_c(w)$ all even.

Base Case: With no rules applied the sentential form is S, and $n_a(S) = n_b(S) = n_c(S) = 0$, which is even.

For the inductive step, let $n \ge 0$ be given and assume that all sentential forms generated after n moves have an even number of a's, b's and c's. Since the n+1 rule application adds 0 or 2 to the number of a's, b's and c's, every sentential form w' obtained after n+1 rules also has $n_a(w')$, $n_b(w')$, and $n_c(w')$ even.

So the result is true for all sentential forms by induction, and for all elements of L(G).

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2. (2 pts) Let $k \in \mathbb{N}$. Prove $a^{2k}(ac)^{2k} \in L(G)$. Hint: You will not need induction.

3. (3 pts) Let L(G) be the language of the grammar

$$\begin{array}{ccc} G:S & \rightarrow & AB \mid \lambda \\ & A & \rightarrow & aAa \mid bAb \mid cAc \mid \lambda \\ & B & \rightarrow & acacB \mid bcbcB \mid \lambda \end{array}$$

 $N = \{w \in \{a, b, c\}^* \mid w = w^R, \text{length}(w) \text{ is even.}\}$. On the back of this page, prove that $N \subseteq L(G)$.

Hint: You might need an induction.

 \clubsuit Every even length palindrome is of the form $w = uu^R$. We will use induction on the length of u to show that uAu^R is derivable from A.

Base Case: Length 0, $\lambda A \lambda^R = A$ is derivable from A.

Inductive Step: Let $n \geq 0$ be given and suppose that $A \stackrel{n}{\Rightarrow} uAu^R$ for any string u of length n. Every string of length n+1 is of the form ux, with $x \in \{a,b,c\}$ and u of length n. Then $A \stackrel{n}{\Rightarrow} uAu^R \Rightarrow u(xAx)u^R = uxAx^Ru^R = uxA(ux)^R$, as required.

So $A \stackrel{*}{\Rightarrow} uu^R$ for all strings u by induction.

So for any $w = uu^R$, we have $S \Rightarrow AB \stackrel{*}{\Rightarrow} uAu^RB \Rightarrow uu^RB \Rightarrow uu^R = w$ and $w \in L(G)$.