



CS5003
Quiz 0110

Foundations of C.S.

Spring, 2023

PRINT NAME: _____

SIGN: _____

1. (6 pts) Let G be the grammar

$$\begin{aligned} G : S &\rightarrow ABCABC \mid \lambda \\ A &\rightarrow abc \mid a \\ B &\rightarrow bca \mid b \\ C &\rightarrow bcbc \mid aC \mid \lambda \end{aligned}$$

Convert to an equivalent grammar, H , in Chomsky Normal Form.

♣ This is one way

$$\begin{aligned} G' : S &\rightarrow WW \mid \lambda \\ A &\rightarrow U_a Y \mid a \\ B &\rightarrow T U_a \mid b \\ C &\rightarrow Y Y \mid U_a C \mid \lambda \\ W &\rightarrow AX \quad (W \text{ produces } ABC) \\ X &\rightarrow BC \\ Y &\rightarrow U_b U_c \quad (\text{Occurs several times.}) \\ U_a &\rightarrow a \\ U_b &\rightarrow b \\ U_c &\rightarrow c \end{aligned}$$

This is still not in Chomsky Normal Form because of the λ -rule in C . $\text{Null}(G') = \{S, C\}$, giving

$$\begin{aligned} G'' : S &\rightarrow WW \mid \lambda \\ A &\rightarrow U_a Y \mid a \\ B &\rightarrow T U_a \mid b \\ C &\rightarrow Y Y \mid U_a C \mid a \\ W &\rightarrow AX \\ X &\rightarrow BC \mid C \\ Y &\rightarrow U_b U_c \\ U_a &\rightarrow a \\ U_b &\rightarrow b \\ U_c &\rightarrow c \end{aligned}$$

Now, removing the one chain rule gives

$$\begin{aligned} H : S &\rightarrow WW \mid \lambda \\ A &\rightarrow U_a Y \mid a \\ B &\rightarrow T U_a \mid b \\ C &\rightarrow Y Y \mid U_a C \mid a \\ W &\rightarrow AX \\ X &\rightarrow BC \mid Y Y \mid U_a C \mid a \\ Y &\rightarrow U_b U_c \\ U_a &\rightarrow a \\ U_b &\rightarrow b \\ U_c &\rightarrow c \end{aligned}$$



2. (2 pts) Professor Practicus defines Context-free grammar to be Practical if every variable is reachable and terminable.

Is a Practical Grammar a normal form?

Justify why or why not.

♣ Of course. Just use the procedure for removing useless symbols to get an equivalent Practical Grammar. ♣

3. (2 pts) Define a grammar G which is both in Chomsky Normal Form and in Greibach Normal Form and for which $|L(G)| > 3$, or prove on the back of the sheet why it cannot be done.

♣ No sentential form can both be in the form AB and $aBCD\dots$, so the only rules allowed are $V \rightarrow a$. And all variables except S are unreachable. So we have to take something like

$$G : S \rightarrow a \mid b \mid c \mid \lambda.$$

for which $|L(G)| = |\{a, b, c, \lambda\}| = 4 > 3$. ♣