1 Inductive Proofs

Prove each of the following claims by induction

Claim 1. The sum of the first n odd numbers is n^2 . That is, $\sum_{i=1}^{n} (2i-1) = n^2$

$$P(n): \sum_{i=1}^{n} (2i-1) = n^2$$

Let us examine the claim for n=1 or P(1):

$$\begin{array}{l} \sum\limits_{i=1}^{1}(2i-1)=1^{2}\\ (2*1-1)=1^{2}\\ 1=1 \end{array}$$

LHS = RHS

The claim is true for n=1, P(1) is true

Assume true for the case of n=k, or P(k) $\sum_{i=1}^{k} (2i-1) = k^2$

Now let us examine if the proposition holds true for k+1

Prove P(k+1) is true:

$$\sum_{i=1}^{k+1} (2i-1) = (k+1)^2$$

$$LHS = \sum_{i=1}^{k} (2i-1) + 2(k+1) - 1$$

$$= k^2 + 2k + 2 - 1$$

$$= k^2 + 2k + 1$$

$$= (k+1)^2 = RHS$$

LHS = RHS

P(k+1) is true when P(k) is true

Therefore P(n) is true for all $n \ge 1$ by principle of mathematical induction

Claim 2.
$$\sum_{i=1}^{n} \frac{1}{2^i} = 1 - \frac{1}{2^n}$$

$$P(n): \sum_{i=1}^{n} \frac{1}{2^i} = 1 - \frac{1}{2^n}$$

Let us examine the claim for n = 1 or P(1):

$$\sum_{i=1}^{1} \frac{1}{2^i} = 1 - \frac{1}{2^1}$$

$$= \frac{1}{2} = 1 - \frac{1}{2}$$
The second s

The claim is true for n = 1, P(1) is true Assume true for the case of n = k, or P(k)

$$\sum_{i=1}^{k} \frac{1}{2^i} = 1 - \frac{1}{2^k}$$

Now let us examine if the proposition holds true for k+1

Prove P(k+1) is true:

LHS =
$$\sum_{i=1}^{k+1} \frac{1}{2^i}$$

= $\sum_{i=1}^{k} \frac{1}{2^i} + \frac{1}{2^{k+1}}$
= $1 - \frac{1}{2^k} + \frac{1}{2 * 2^k}$
= $1 - \frac{1}{2^k} + \frac{1}{2} \frac{1}{2^k}$
= $1 - \frac{1}{2^k} (1 - \frac{1}{2})$
= $1 - \frac{1}{2^k} * \frac{1}{2}$
= $1 - \frac{1}{2^{k+1}}$ = RHS
LHS = RHS
P(k+1) is true when

P(k+1) is true when P(k) is true

Therefore P(n) is true for all $n \ge 1$ by principle of mathematical induction

Claim 3.
$$\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$$

$$P(n): \sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$$

Let us examine the claim for n=1 or P(1):
$$\sum_{i=1}^{1} i^2 = \frac{1(1+1)(2*1+1)}{6}$$

$$LHS = 1$$

$$LHS = 1$$

$$RHS = 1 = LHS$$

The claim is true for n=1, P(1) is true

Assume true for the case of n=k, or P(k)

Assume true for the case of fi-k, of 1 (k)
$$\sum_{i=1}^{k} i^2 = \frac{k(k+1)(2k+1)}{6}$$
 Now let us examine if the proposition holds true for k+1

Now let us examine if the proposition holds true:

Prove
$$P(k+1)$$
 is true:
$$P(k+1): \sum_{i=1}^{k+1} i^2 = \frac{(k+1)(k+2)(2(k+1)+1)}{6}$$

$$LHS = \sum_{i=1}^{k+1} i^2$$

$$= (k+1)^2 + \sum_{i=1}^{k} i^2$$

$$= (k+1)^2 + \frac{k(k+1)(2k+1)}{6}$$
Expanding this, we get:

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\frac{2k^3+8k^2+13k+6}{6} RHS = \frac{(k+1)(k+2)(2(k+1)+1)}{6} = \frac{(k+1)(k+2)(2k+3)}{6} Expanding this, we get: \frac{2k^3+8k^2+13k+6}{6} = LHS LHS = RHS P(k+1) is true when P(k) is true Therefore P(n) is true for all n \geq 1 by principle of mathematical induction
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2 Recursive Invariants

The function minEven, given below in pseudocode, takes as input an array A of size n of numbers. It returns the smallest even number in the array. If no even numbers appear in the array, it returns positive infinity $(+\infty)$. Using induction, prove that the minEven function works correctly. Clearly state your recursive invariant at the beginning of your proof.

```
Function minEven(A,n)
  If n = 0 Then
    Return +
    Set best To minEven(A,n-1)
    If A[n-1] < best And A[n-1] is even Then
      Set best To A[n-1]
    EndIf
    Return best
  EndIf
EndFunction
   Let us examine the function minEven for an array A of size 0
   In this case minEven(A,0) will return +\infty
   Hence the function minEven works correctly for n=0
   This is the recursive invariant P(0)
   Let us assume that the minEven function works correctly for an array A of
size k and returns m
   This is P(k) and is assumed to be true.
   We have to prove that P(K+1) is true given P(K) is true.
   Assume that one more element E is added to the array A increasing its size
to k+1
```

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This element E will be added in the position A[n-1]
E's value can have three cases:
Case 1: E is odd
Case 2: E is Even and E ≥ m
Case 3: E is Even and E < m
Let us now run minEven(A,K+1):
k > 0, so we go to the Else branch
Set best To minEven(A,n-1): best = m

If A[n-1] < best And A[n-1] is even Then - Case 3
    Set best To A[n-1]: best = E
EndIf
Return best : return m for Case 1 and Case 2, return k for Case 3</pre>
```

So $\mathtt{minEven}$ works correctly for array of size k+1 if it works for correctly for size k

for all cases of the added element /textttE Hence P(K+1) is true given P(K) is true. Hence this function works correctly for all arrays of integers using the principle of mathematical induction.