

Lecture 9:

Efficiently Evaluating DNNs

**Parallel Computing
Stanford CS149, Fall 2025**

Extreme efficiency challenge

Diverse collection of AI models (topology and size)

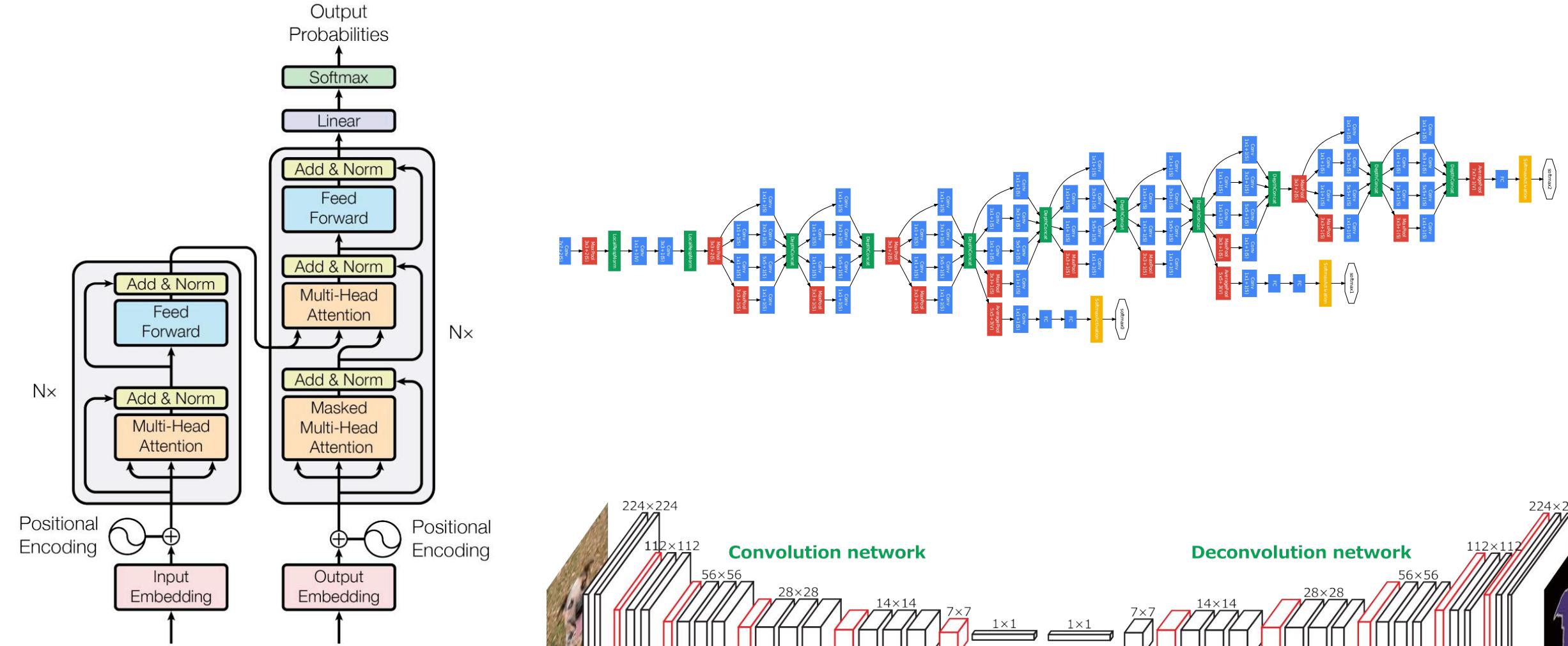


Table 1. MobileNet Body Architecture		
Type / Stride	Filter Shape	Input Size
Conv / s2	$3 \times 3 \times 3 \times 32$	$224 \times 224 \times 3$
Conv dw / s1	$3 \times 3 \times 32 \text{ dw}$	$112 \times 112 \times 32$
Conv / s1	$1 \times 1 \times 32 \times 64$	$112 \times 112 \times 32$
Conv dw / s2	$3 \times 3 \times 64 \text{ dw}$	$112 \times 112 \times 64$
Conv / s1	$1 \times 1 \times 64 \times 128$	$56 \times 56 \times 64$
Conv dw / s1	$3 \times 3 \times 128 \text{ dw}$	$56 \times 56 \times 128$
Conv / s1	$1 \times 1 \times 128 \times 128$	$56 \times 56 \times 128$
Conv dw / s2	$3 \times 3 \times 128 \text{ dw}$	$56 \times 56 \times 128$
Conv / s1	$1 \times 1 \times 128 \times 256$	$28 \times 28 \times 128$
Conv dw / s1	$3 \times 3 \times 256 \text{ dw}$	$28 \times 28 \times 256$
Conv / s1	$1 \times 1 \times 256 \times 256$	$28 \times 28 \times 256$
Conv dw / s2	$3 \times 3 \times 256 \text{ dw}$	$28 \times 28 \times 256$
Conv / s1	$1 \times 1 \times 256 \times 512$	$14 \times 14 \times 256$
5x Conv dw / s1	$3 \times 3 \times 512 \text{ dw}$ $1 \times 1 \times 512 \times 512$	$14 \times 14 \times 512$ $14 \times 14 \times 512$
Conv dw / s2	$3 \times 3 \times 512 \text{ dw}$	$14 \times 14 \times 512$
Conv / s1	$1 \times 1 \times 512 \times 1024$	$7 \times 7 \times 512$
Conv dw / s2	$3 \times 3 \times 1024 \text{ dw}$	$7 \times 7 \times 1024$
Conv / s1	$1 \times 1 \times 1024 \times 1024$	$7 \times 7 \times 1024$
Avg Pool / s1	Pool 7×7	$7 \times 7 \times 1024$
FC / s1	1024 \times 1000	$1 \times 1 \times 1024$
Softmax / s1	Classifier	$1 \times 1 \times 1000$

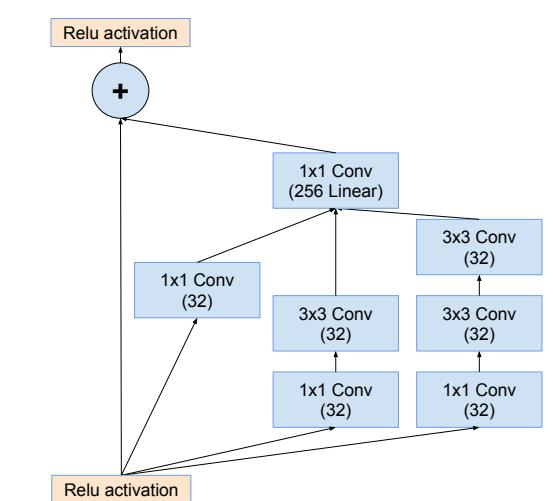
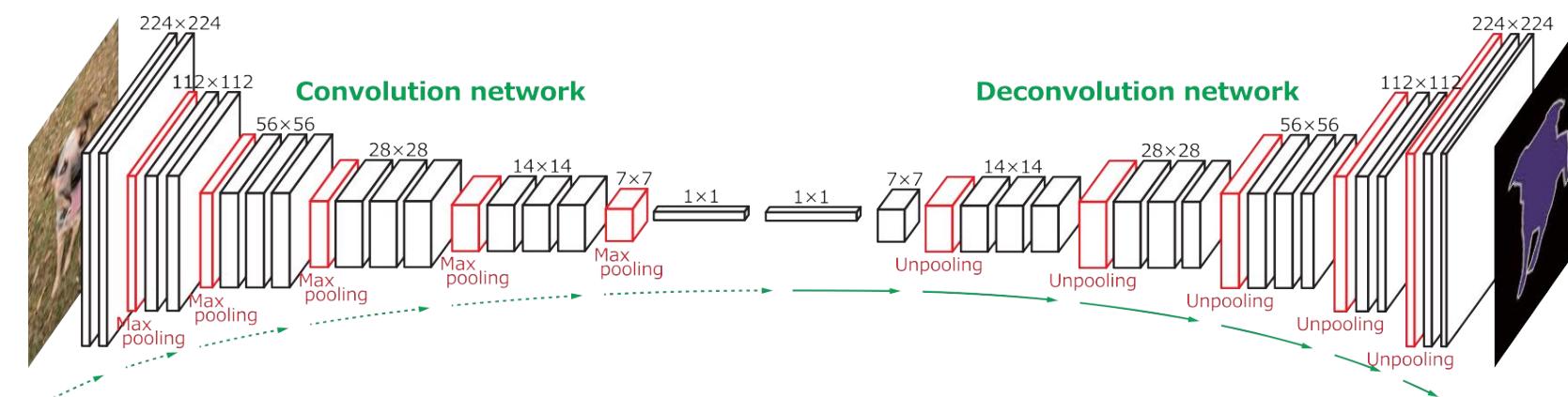
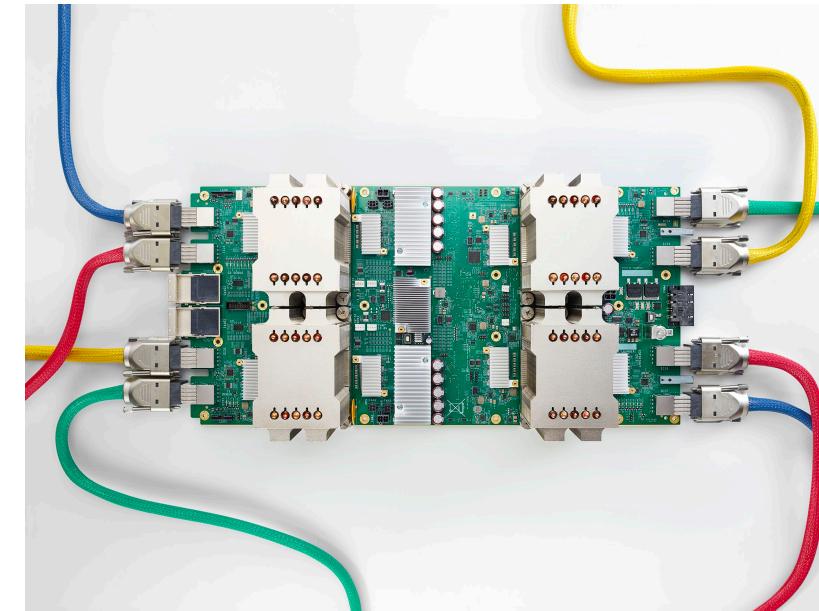
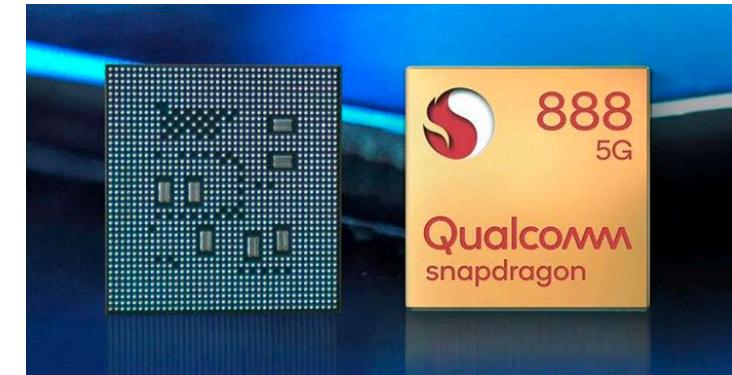
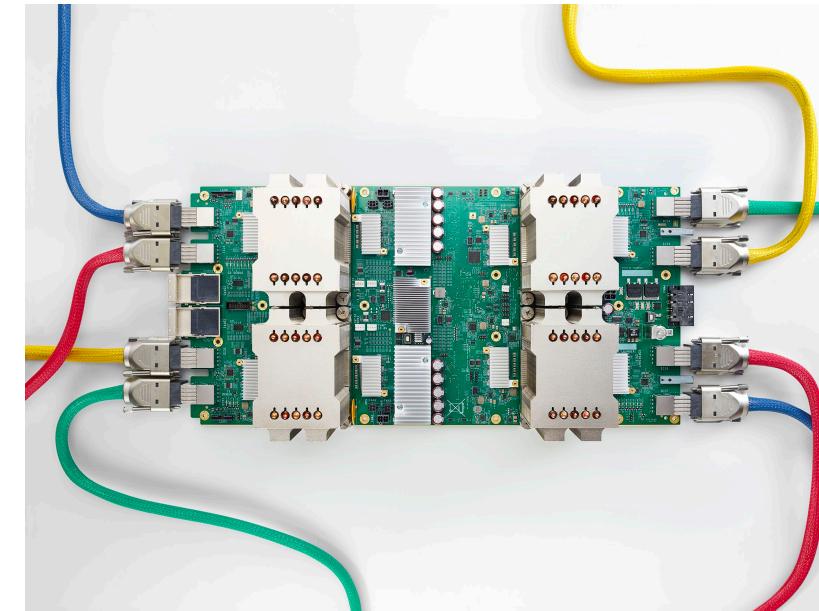


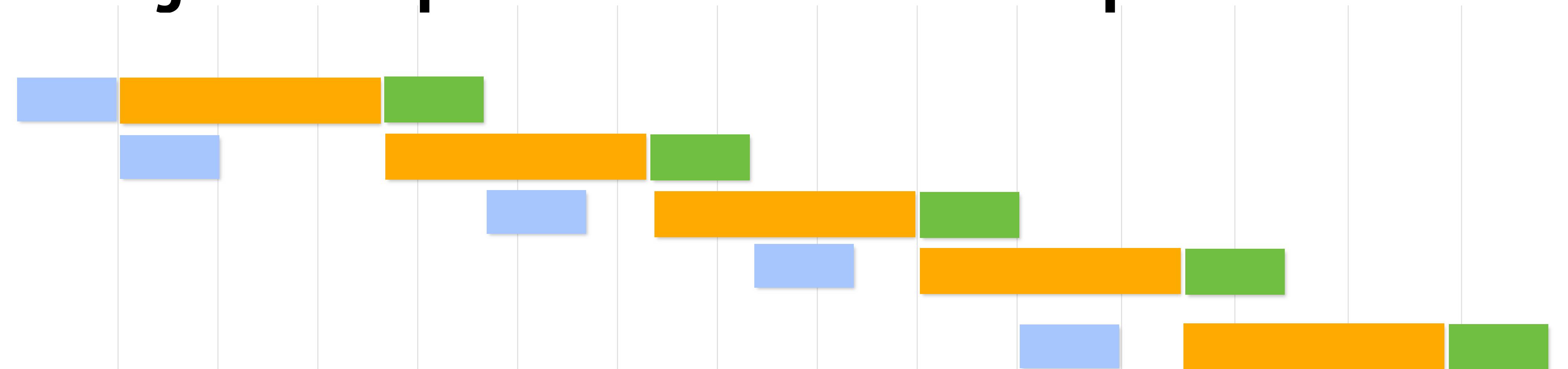
Figure 10. The schema for 35×35 grid (Inception-ResNet-A) module of Inception-ResNet-v1 network.

Many Target Devices



**Things you already know —
and should remember**

Pipelining to overlap data movement with computation

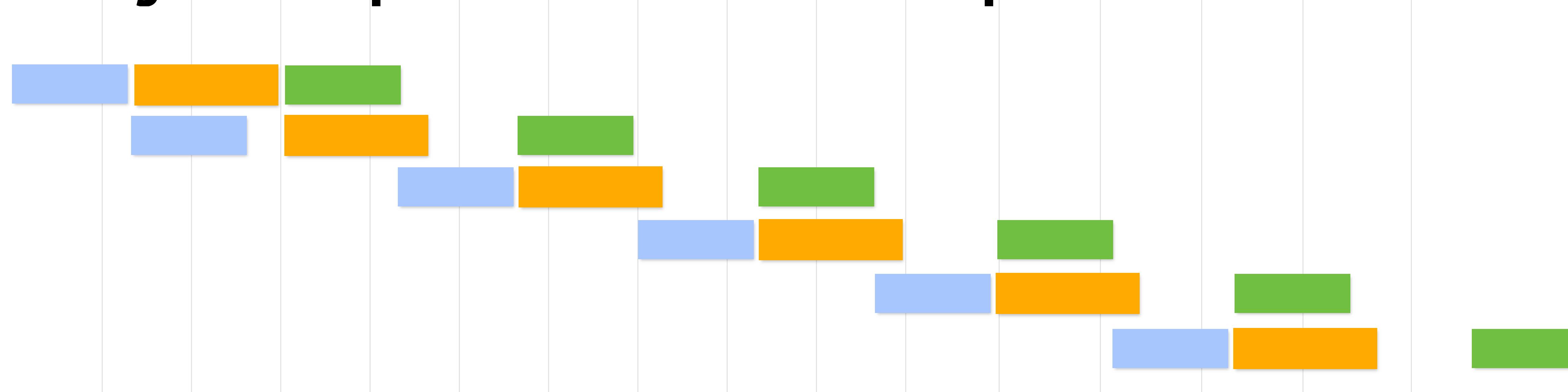


Question 1: is this program compute bound or BW bound?

- █ = Load data
- █ = Arithmetic operations
- █ = Store result

Question 2: what is the on chip storage cost of overlapping data movement with computation? (Hint: it's often called “double buffering”)

Pipelining to overlap data movement with computation

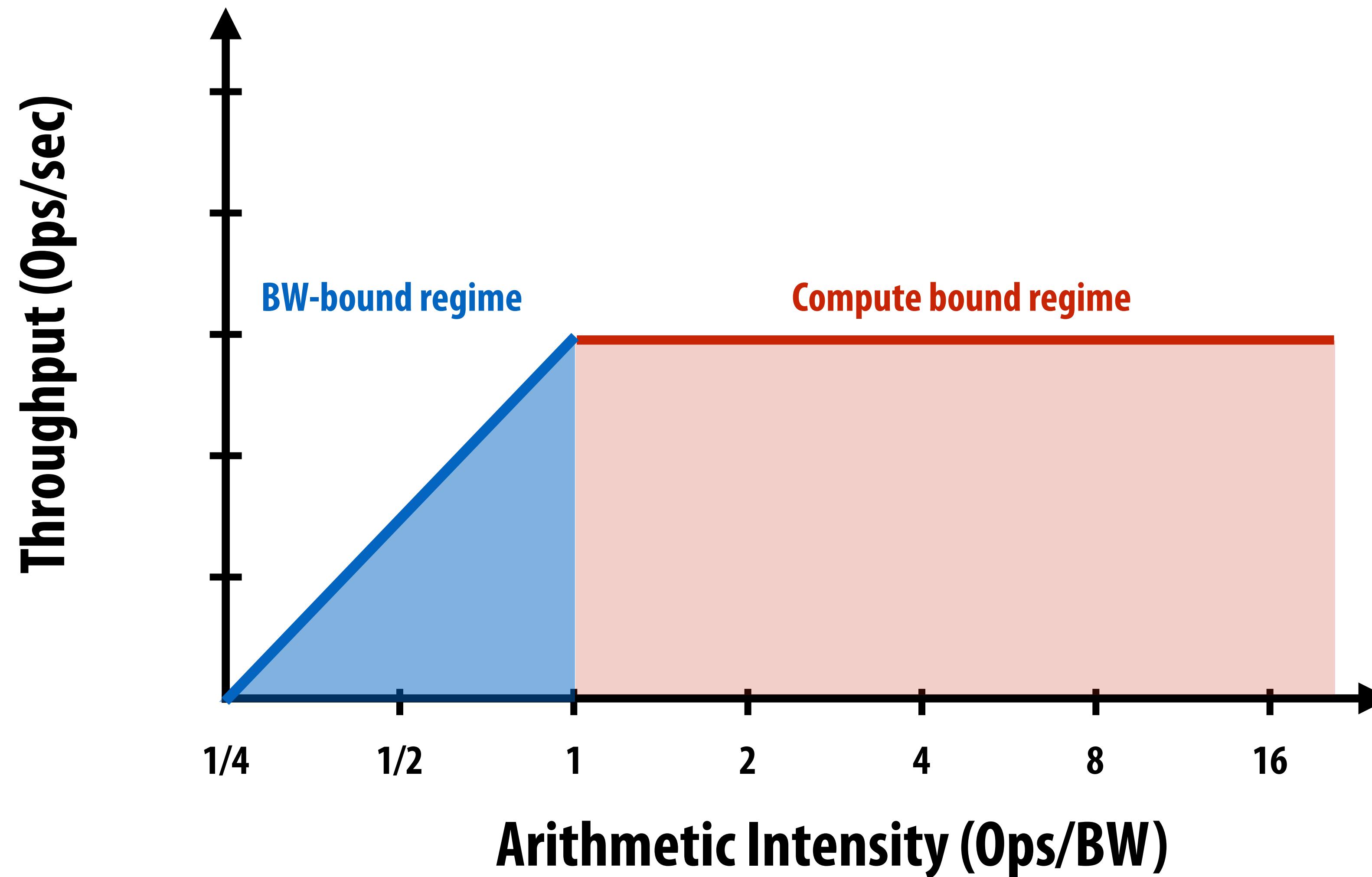


Question 1: is this program compute bound or BW bound?

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- █ = Arithmetic operations
- █ = Store result

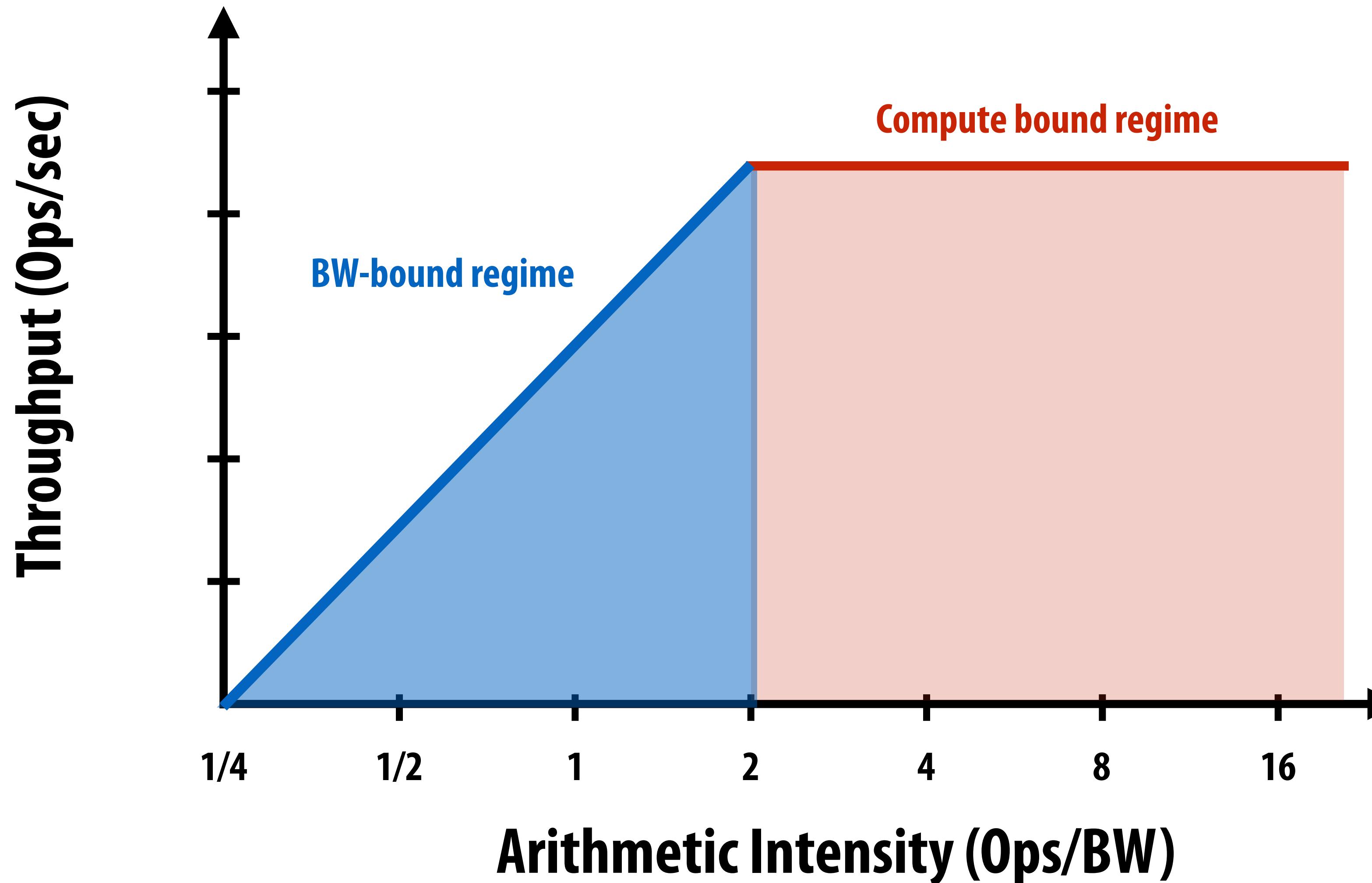
time

A roofline curve



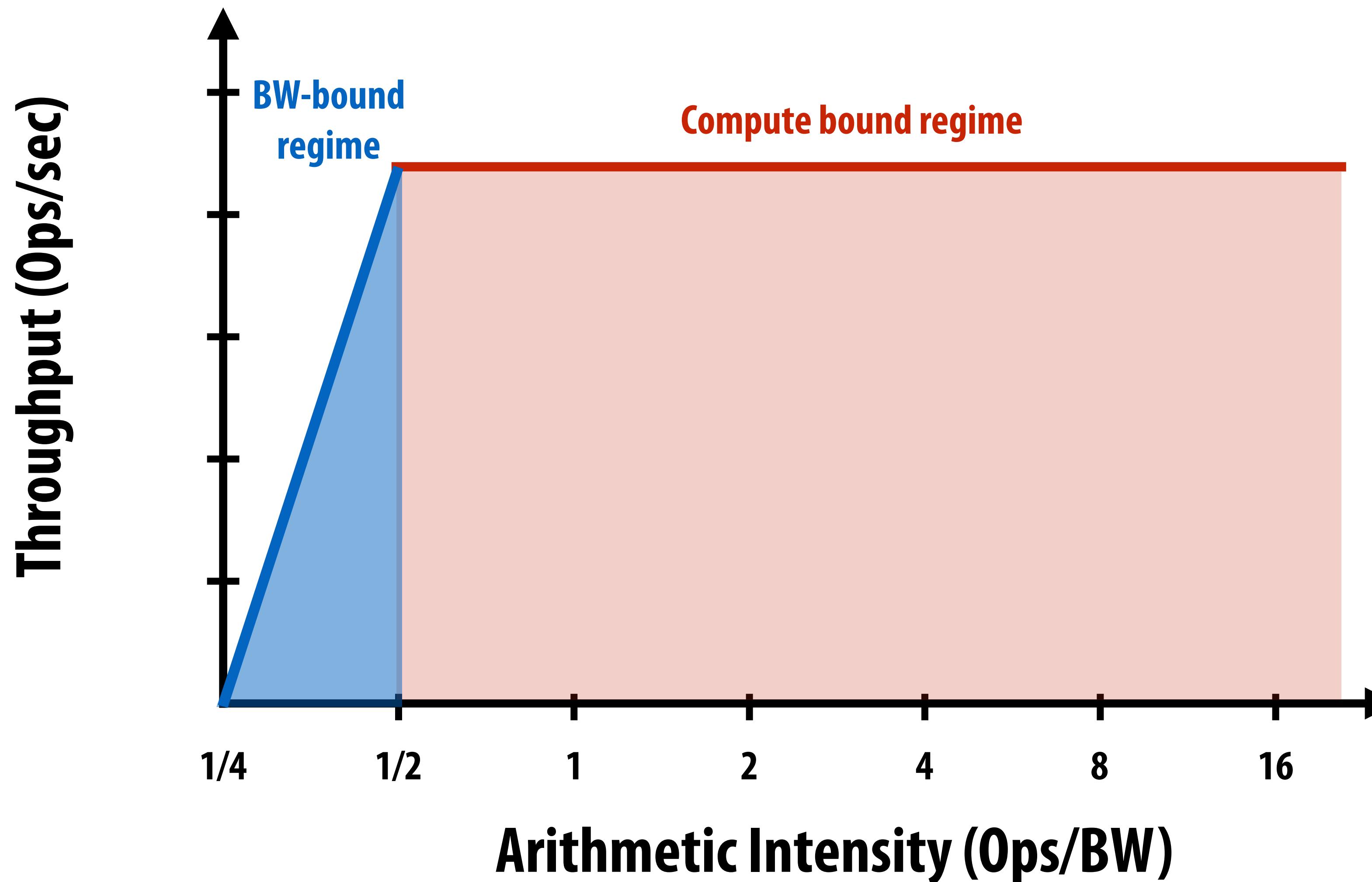
A roofline curve

Computer with the **same** memory system but **higher** peak compute capability



A roofline curve

Computer with the **higher-throughput** memory system and **higher** peak compute capability



Recall the loop fusion transformation: fuse multiple loops into one to increase a program's arithmetic intensity

The transformation of the code in program 1 to the code in program 2 is called “loop fusion”

Program 1

```
void add(int n, float* A, float* B, float* C) {  
    for (int i=0; i<n; i++)  
        C[i] = A[i] + B[i];  
}  
  
void mul(int n, float* A, float* B, float* C) {  
    for (int i=0; i<n; i++)  
        C[i] = A[i] * B[i];  
}  
  
float* A, *B, *C, *D, *E, *tmp1, *tmp2;  
  
// assume arrays are allocated here  
  
// compute E = D + ((A + B) * C)  
add(n, A, B, tmp1);  
mul(n, tmp1, C, tmp2);  
add(n, tmp2, D, E);
```

Two loads, one store per math op
(arithmetic intensity = 1/3)

Two loads, one store per math op
(arithmetic intensity = 1/3)

Overall arithmetic intensity = 1/3

Program 2

```
void fused(int n, float* A, float* B, float* C, float* D, float* E) {  
    for (int i=0; i<n; i++)  
        E[i] = D[i] + (A[i] + B[i]) * C[i];  
}  
  
// compute E = D + (A + B) * C  
fused(n, A, B, C, D, E);
```

Four loads, one store per 3 math ops
(arithmetic intensity = 3/5)

Review

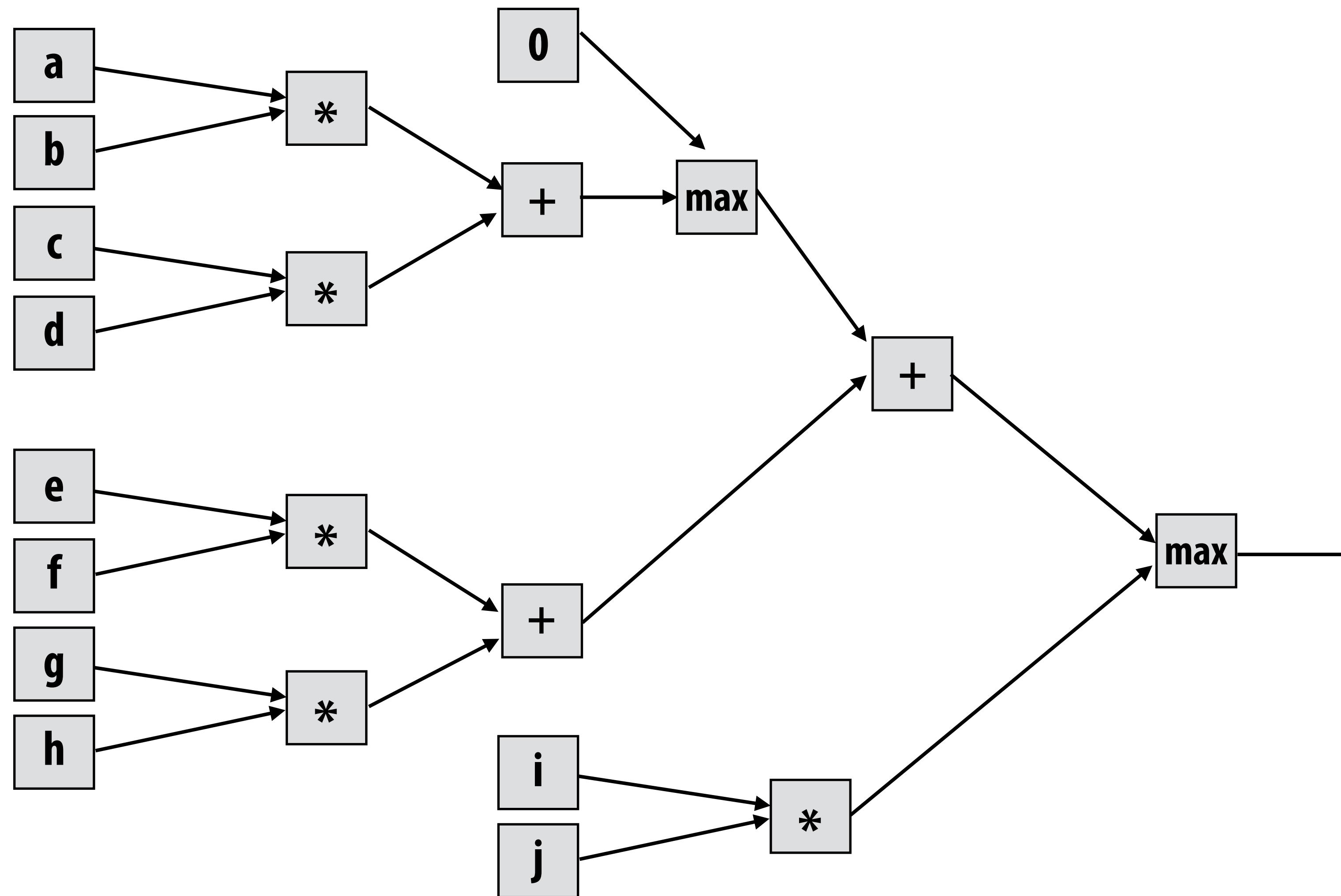
- When communication and computation are overlapped (hiding memory latency), the capabilities of the machine (ops throughput and communication bandwidth) AND the arithmetic intensity of the program determine if the program's overall instruction throughput is limited by available bandwidth (“bandwidth bound”) or by the machine's instruction processing capability (“compute bound”)
- Overlapping communication and computation costs footprint, since buffers for the data being processing AND the data being transferred need to be maintained on chip.
- Increasing arithmetic processing ability (“faster hardware”) makes a program more likely to be bandwidth bound
- Increasing a program's arithmetic intensity (“a program change”) makes a program more likely to be compute bound

If you know the previous slide, you know almost everything you need to know about the software side* of performance optimization of modern AI.

*** If you want to know the rest, wait for next class... and it basically amounts to (1) data movement costs energy, (2) chip resources used for on-chip storage are resources that cannot be used for compute, so minimize buffers as much as possible**

Mini-intro: Convolutional Neural Networks

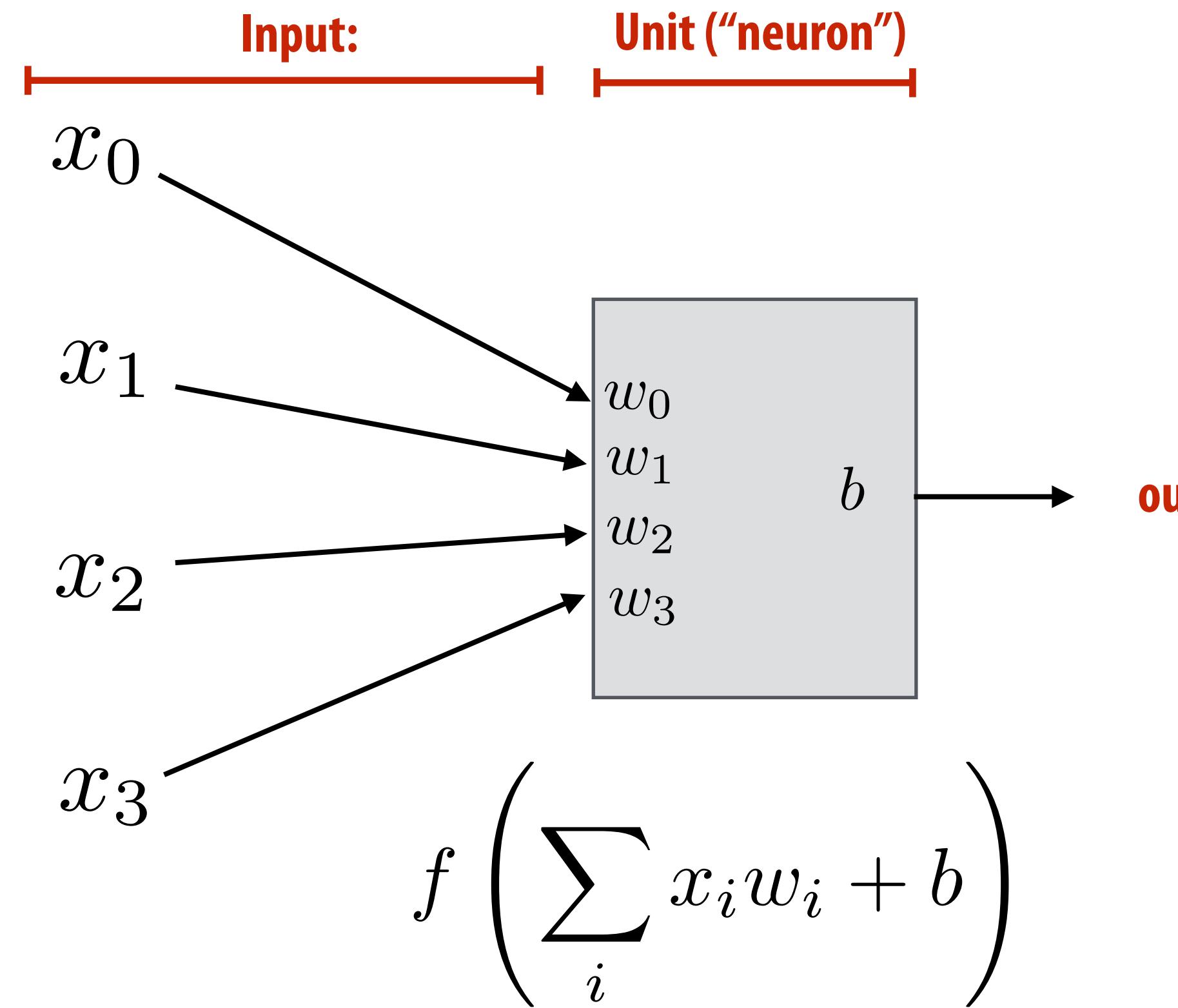
Consider the following expression


$$\max(\max(0, (a \cdot b) + (c \cdot d)) + (e \cdot f) + (g \cdot h), i \cdot j)$$

What is a deep neural network?

A basic unit:

Unit with n inputs described by $n+1$ parameters
(weights + bias)



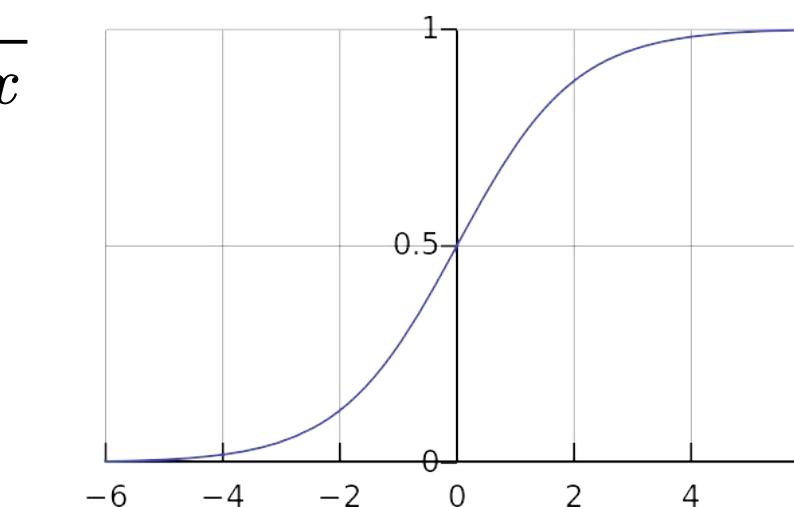
Example: rectified linear unit (ReLU)

$$f(x) = \max(0, x)$$

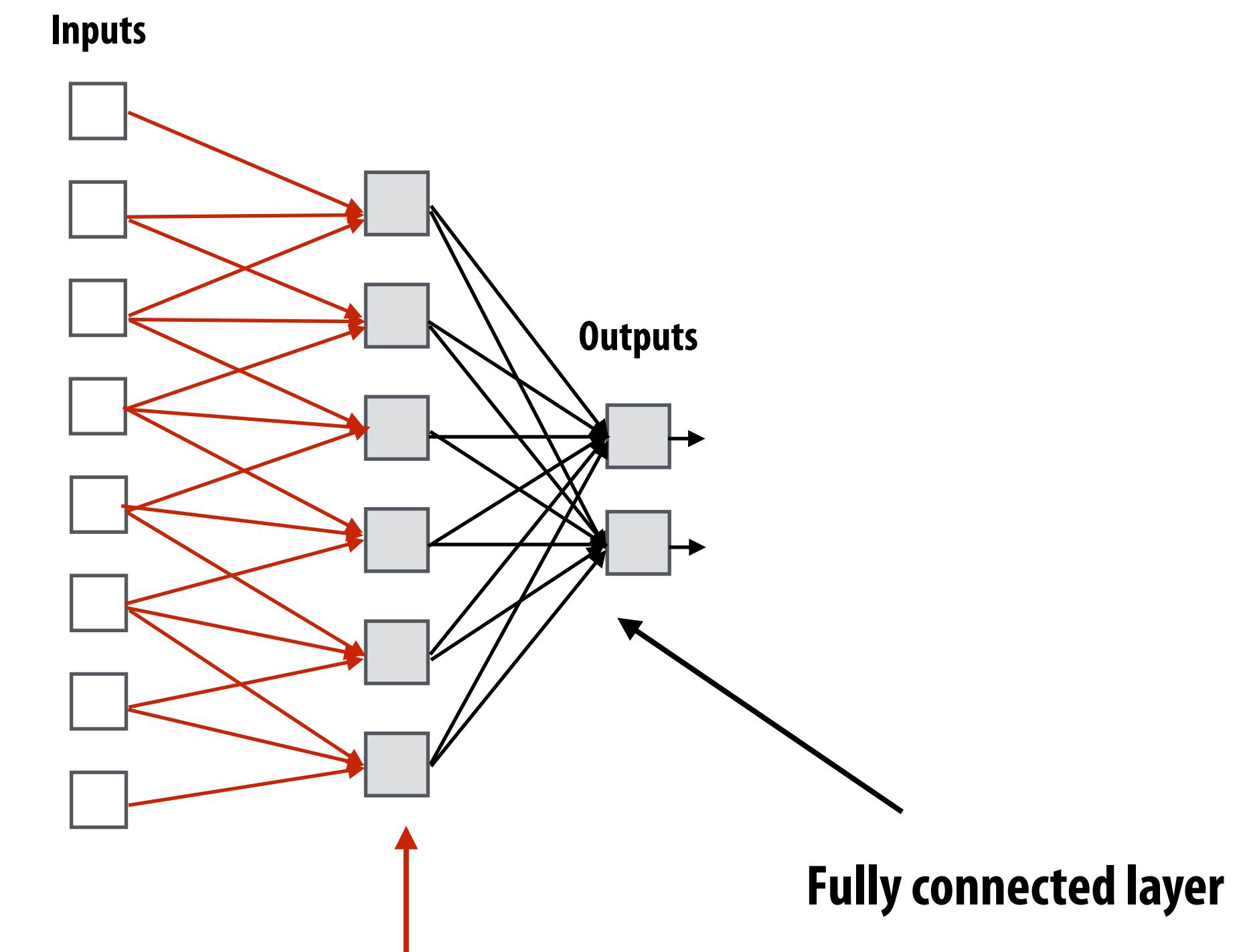
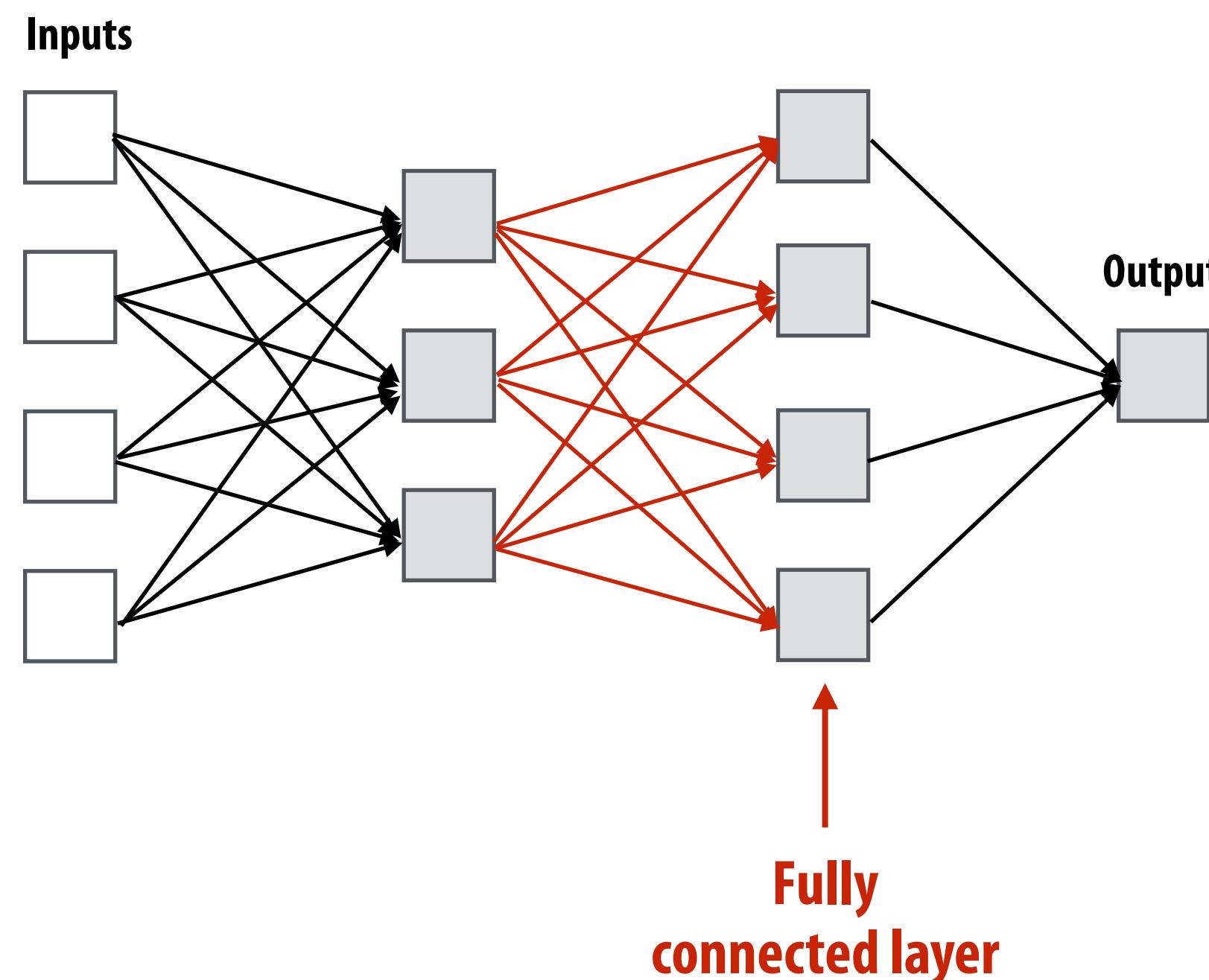
Basic computational interpretation:
It is just a circuit!

Machine learning interpretation:
Binary classifier: interpret output as the probability of one class

$$f(x) = \frac{1}{1 + e^{-x}}$$

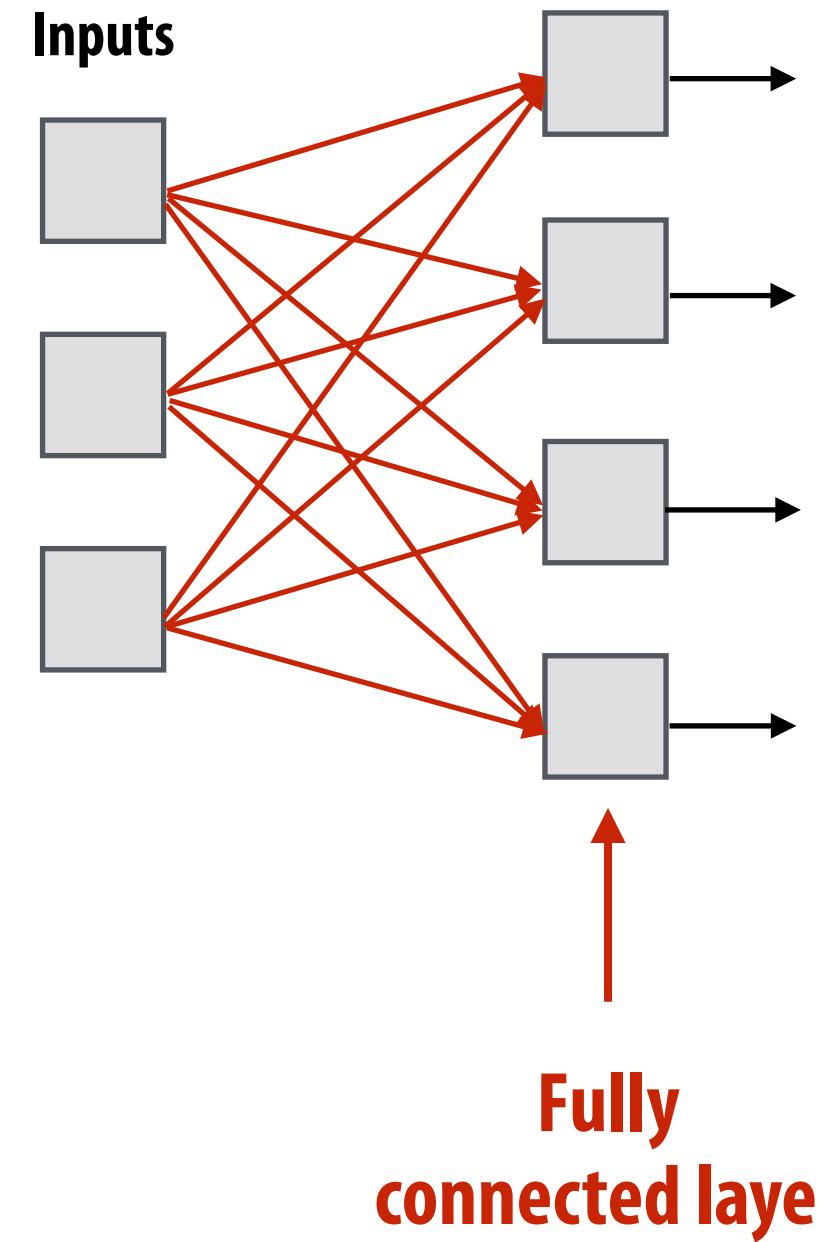


Deep neural network: topology



**Sparsely (locally) connected layer
(each unit only received inputs
from three input nodes)**

Fully connected layer as matrix-vector product



$$f \left(\begin{bmatrix} w_{00} & w_{01} & w_{02} \\ w_{10} & w_{11} & w_{12} \\ w_{20} & w_{21} & w_{22} \\ w_{30} & w_{31} & w_{32} \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \\ x_2 \end{bmatrix} + \begin{bmatrix} b_0 \\ b_1 \\ b_2 \\ b_3 \end{bmatrix} \right)$$

Assume $f()$ is the element-wise max function (ReLU)

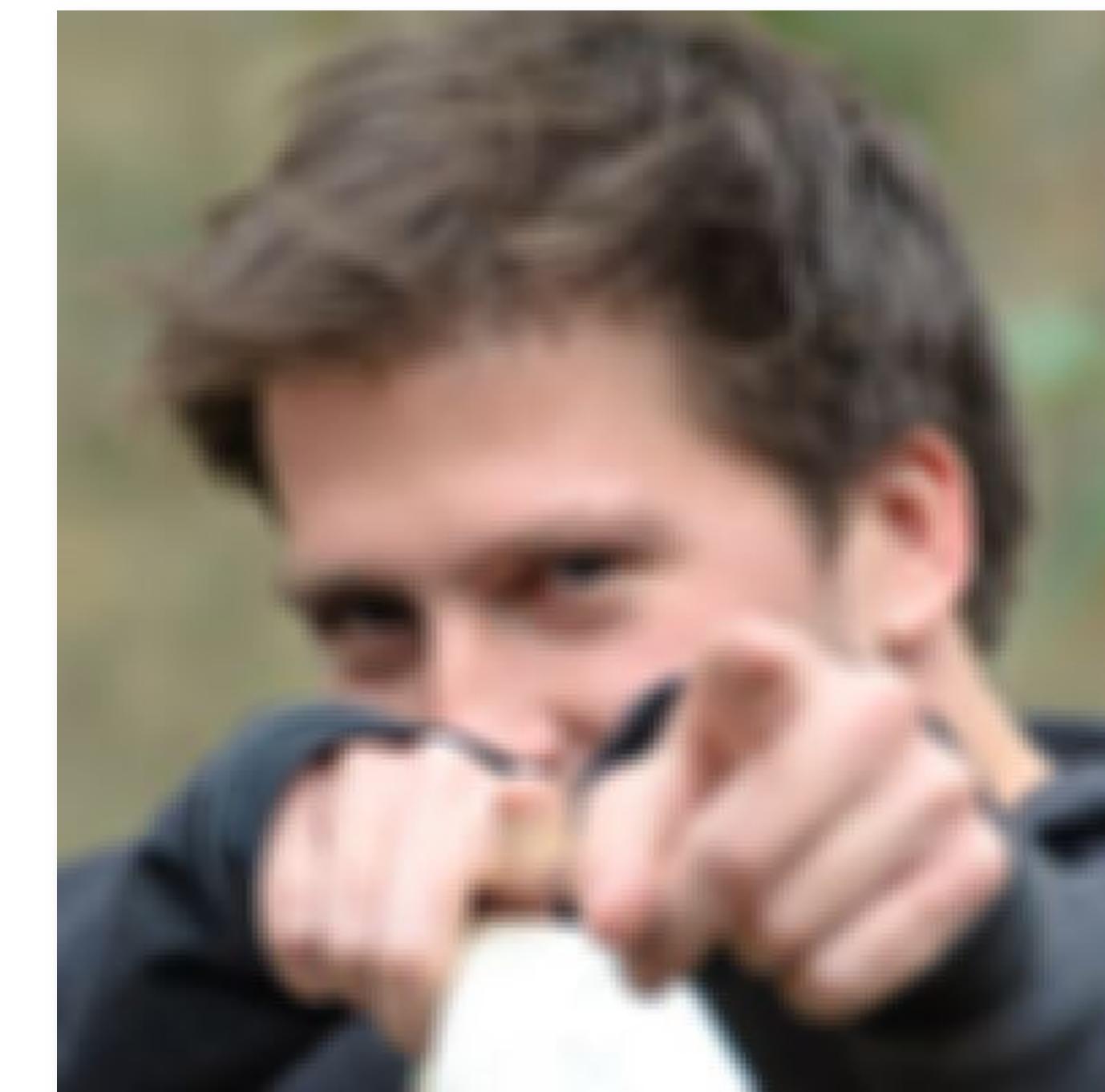
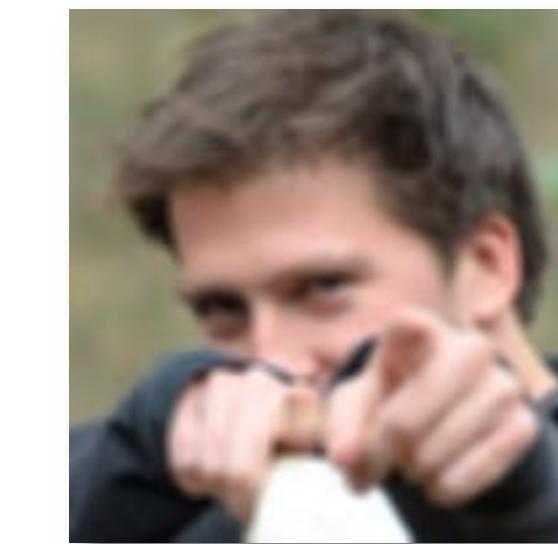
2D convolution: what does this C code do?

```
int WIDTH = 1024;
int HEIGHT = 1024;
float input[(WIDTH+2) * (HEIGHT+2)];
float output[WIDTH * HEIGHT];

float weights[] = {1.f/9, 1.f/9, 1.f/9,
                   1.f/9, 1.f/9, 1.f/9,
                   1.f/9, 1.f/9, 1.f/9};

for (int j=0; j<HEIGHT; j++) {
    for (int i=0; i<WIDTH; i++) {
        float tmp = 0.f;
        for (int jj=0; jj<3; jj++)
            for (int ii=0; ii<3; ii++)
                tmp += input[(j+jj)*(WIDTH+2) + (i+ii)] * weights[jj*3 + ii];
        output[j*WIDTH + i] = tmp;
    }
}
```

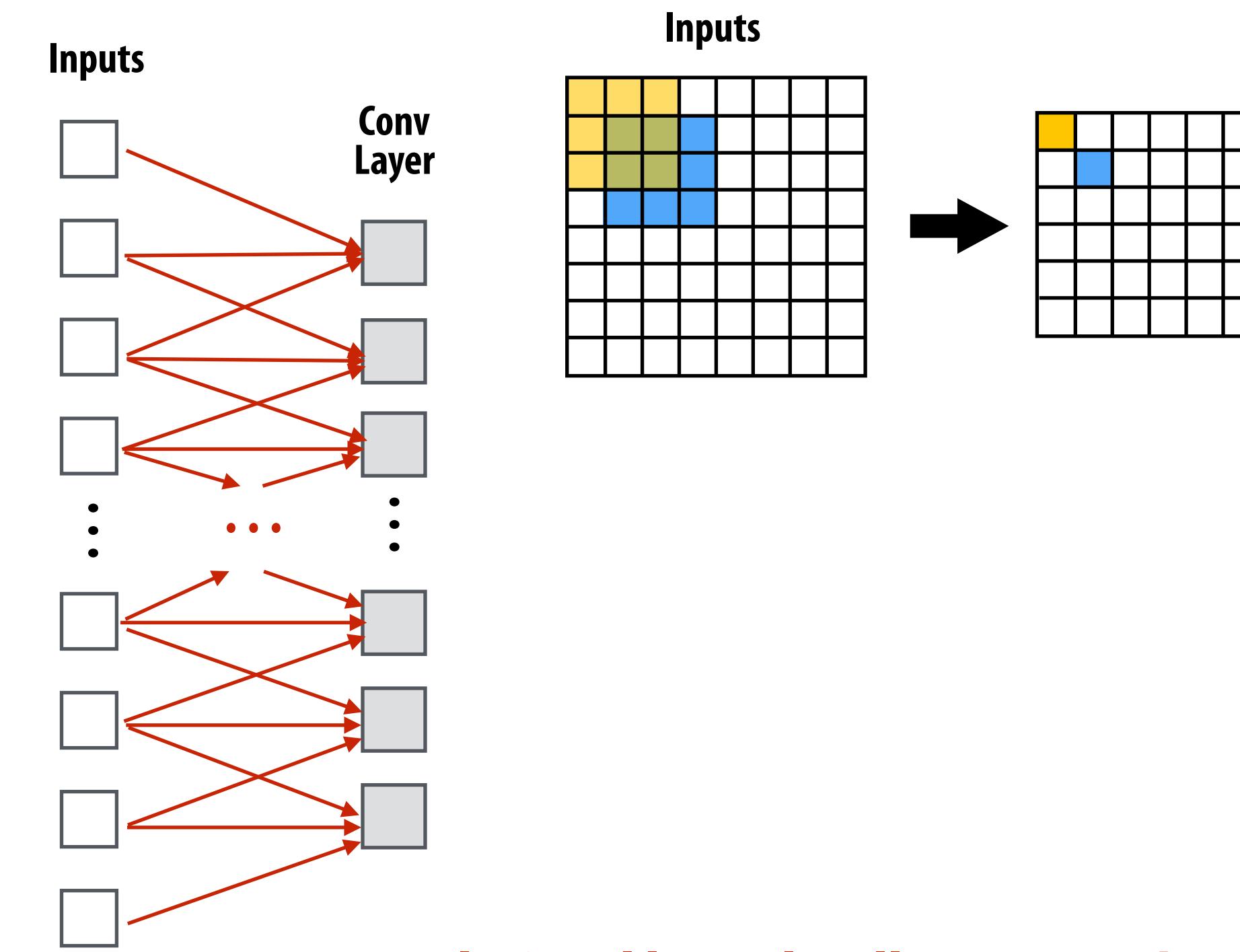
The code on the previous slide performed a 3x3 blur



(Zoomed view)

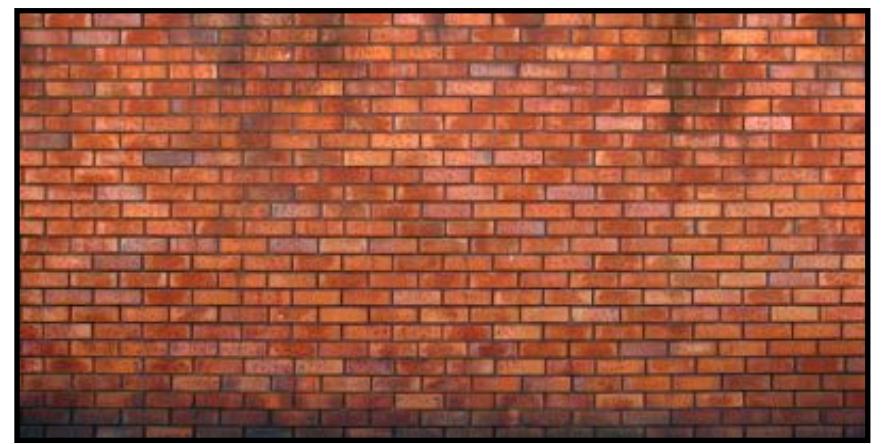
Image convolution (3x3 conv)

```
int WIDTH = 1024;  
int HEIGHT = 1024;  
float input[(WIDTH+2) * (HEIGHT+2)];  
float output[WIDTH * HEIGHT];  
  
float weights[] = {1.0/9, 1.0/9, 1.0/9,  
                   1.0/9, 1.0/9, 1.0/9,  
                   1.0/9, 1.0/9, 1.0/9};  
  
for (int j=0; j<HEIGHT; j++) {  
    for (int i=0; i<WIDTH; i++) {  
        float tmp = 0.f;  
        for (int jj=0; jj<3; jj++)  
            for (int ii=0; ii<3; ii++)  
                tmp += input[(j+jj)*(WIDTH+2) + (i+ii)] * weights[jj*3 + ii];  
        output[j*WIDTH + i] = tmp;  
    }  
}
```

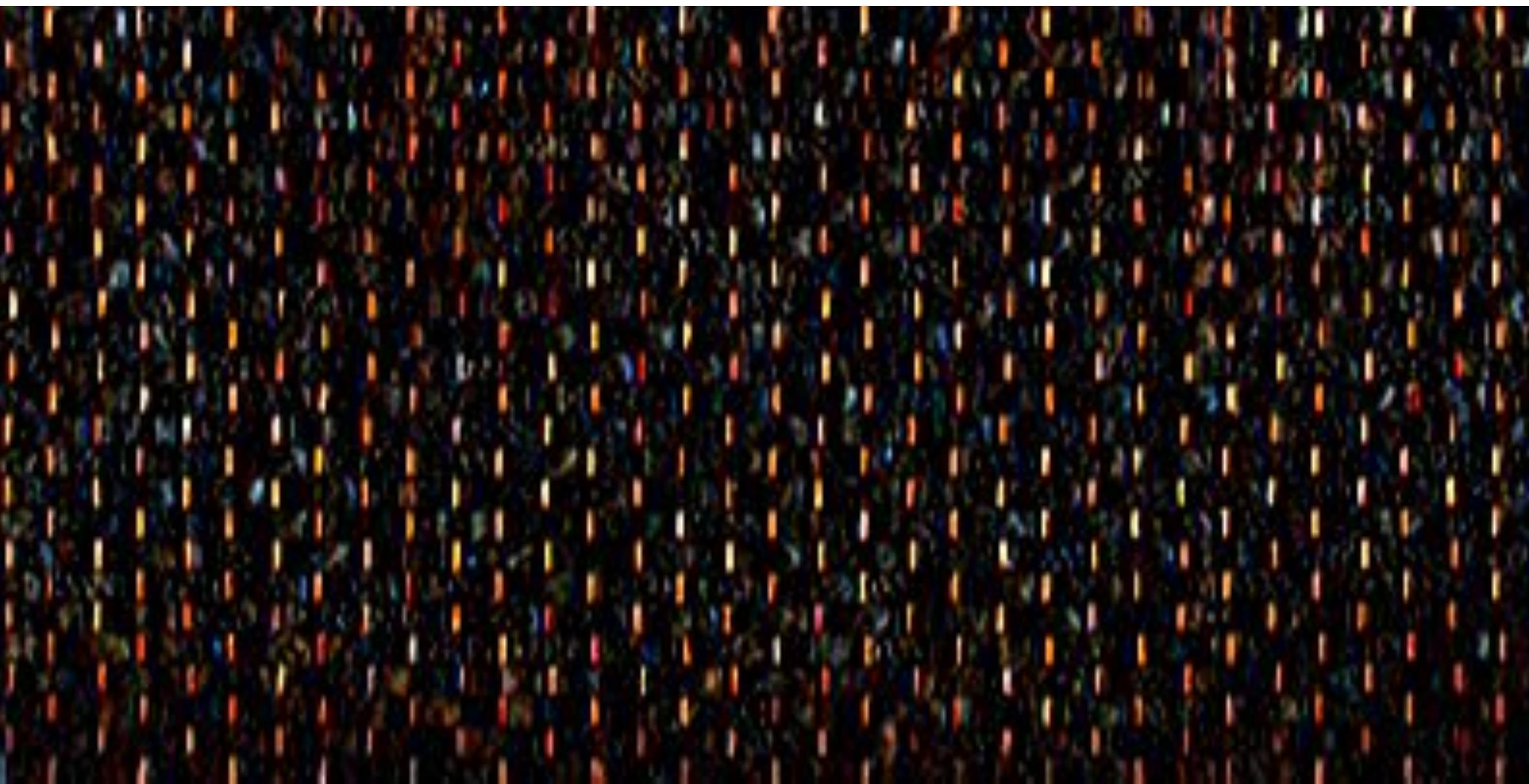


Convolutional layer: locally connected AND all units in layer share the same parameters (same weights + same bias):
(note: network illustration above only shows links for a 1D conv: a.k.a. one iteration of *ii* loop)

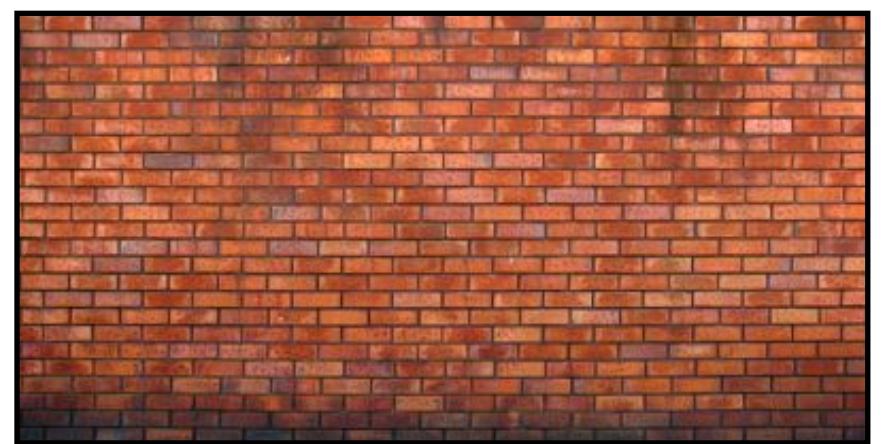
Gradient detection filters



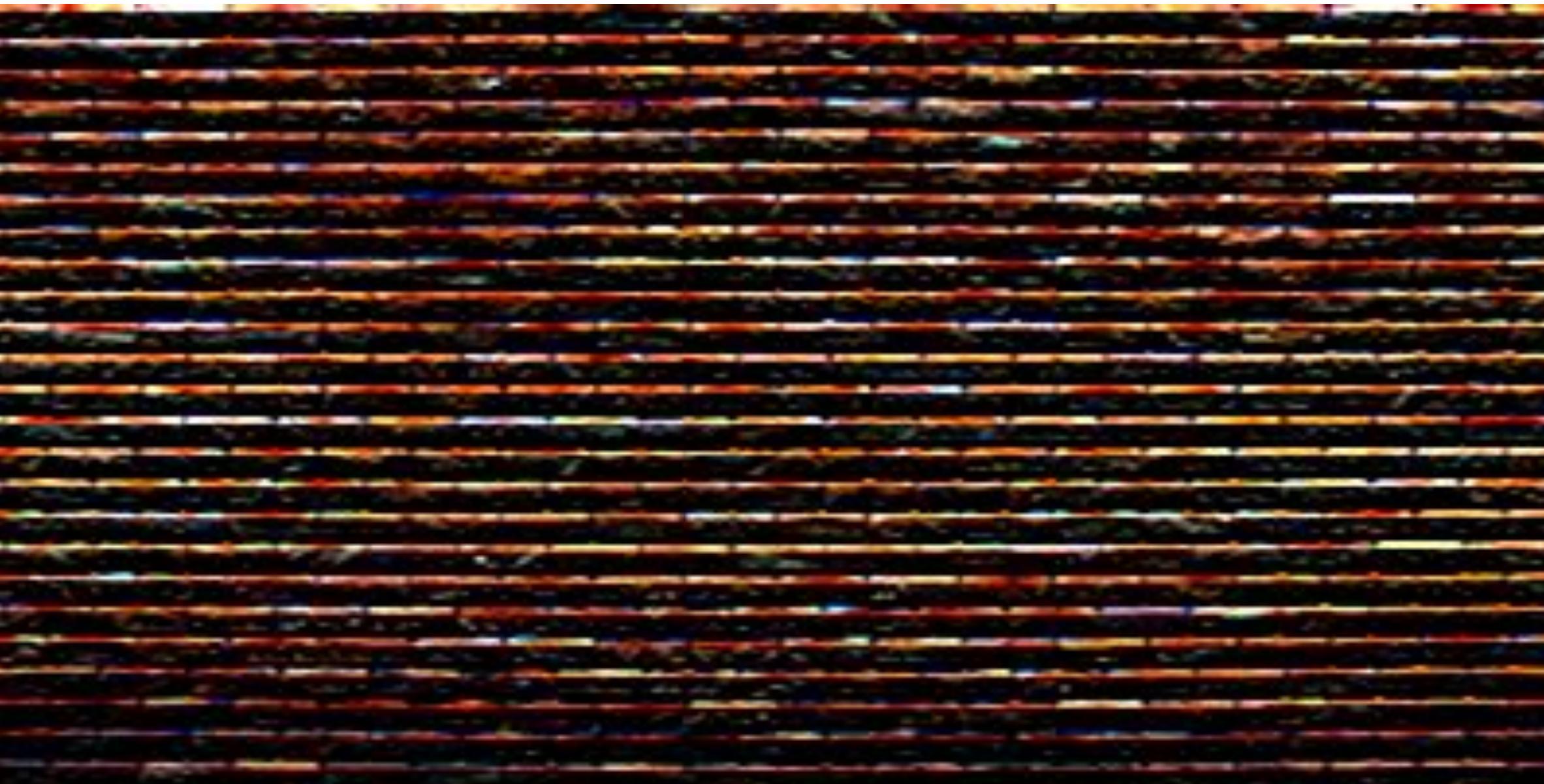
$$* \begin{bmatrix} -1 & 0 & 1 \\ -2 & 0 & 2 \\ -1 & 0 & 1 \end{bmatrix} =$$



Responds to horizontal gradients



$$* \begin{bmatrix} -1 & -2 & -1 \\ 0 & 0 & 0 \\ 1 & 2 & 1 \end{bmatrix} =$$



Responds to vertical gradients

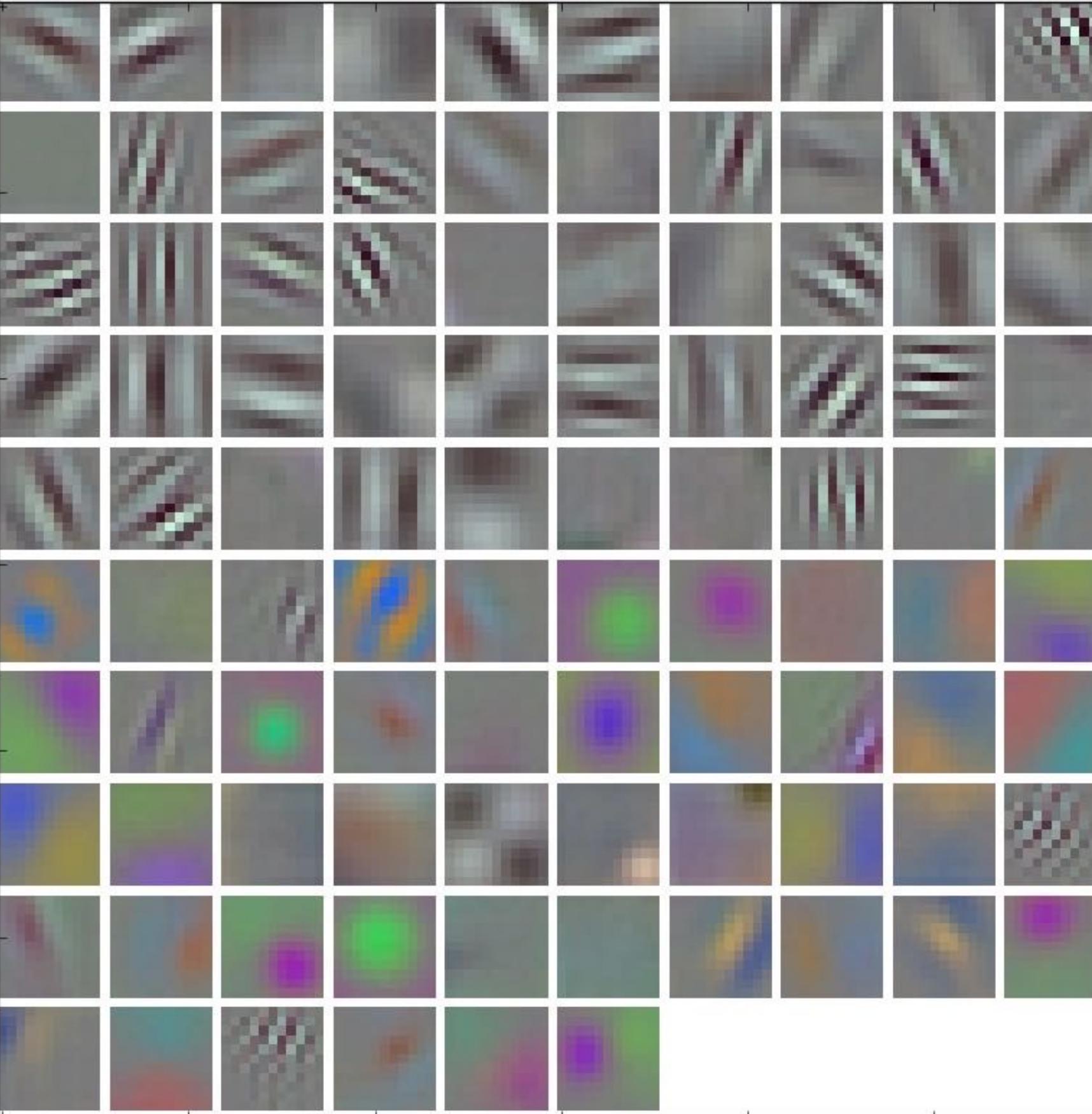
Note: you can think of a filter as a “detector” of a pattern, and the magnitude of a pixel in the output image as the “response” of the filter to the region surrounding each pixel in the input image

Applying many filters to an image at once

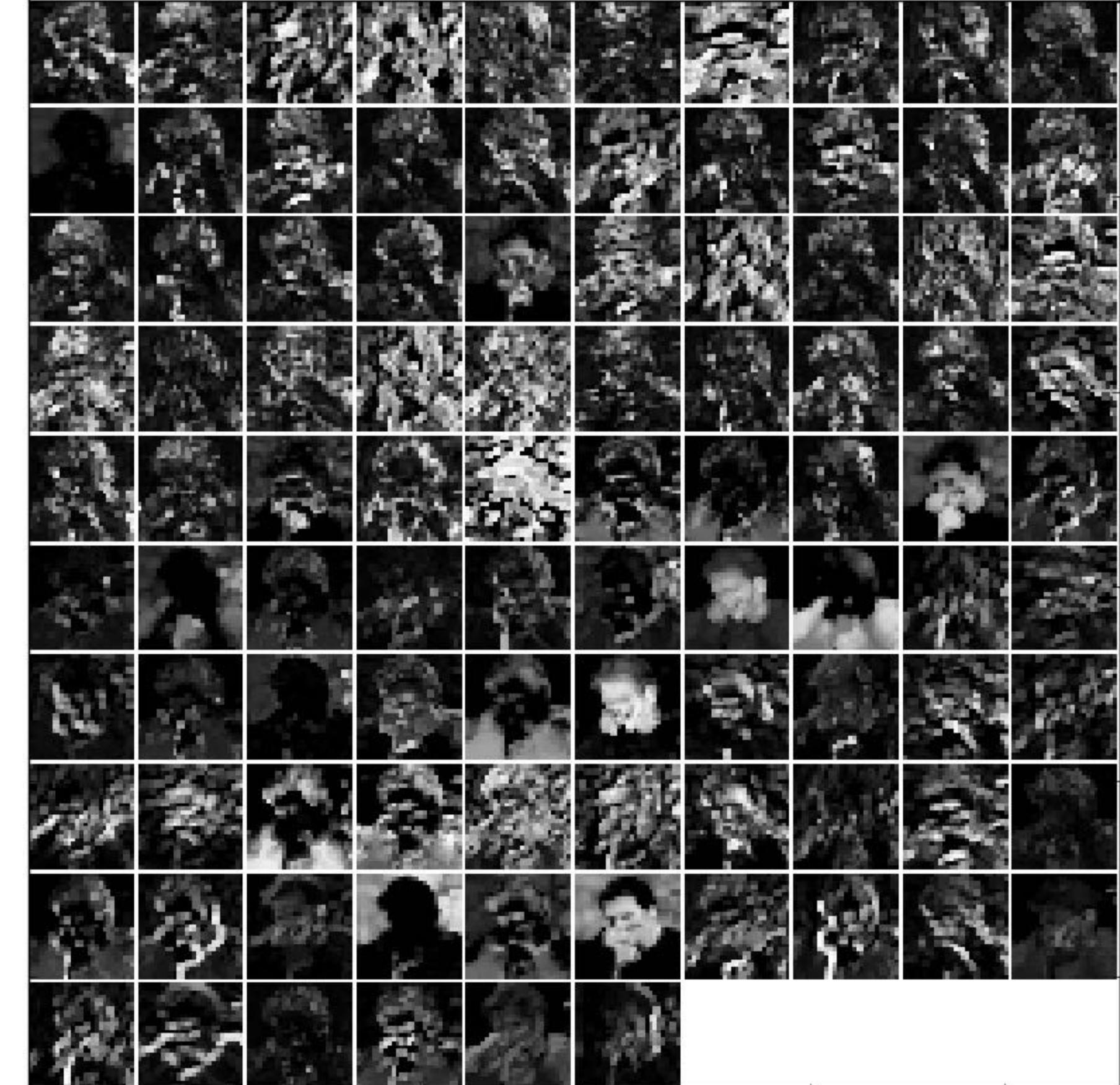
Input RGB image ($W \times H \times 3$)



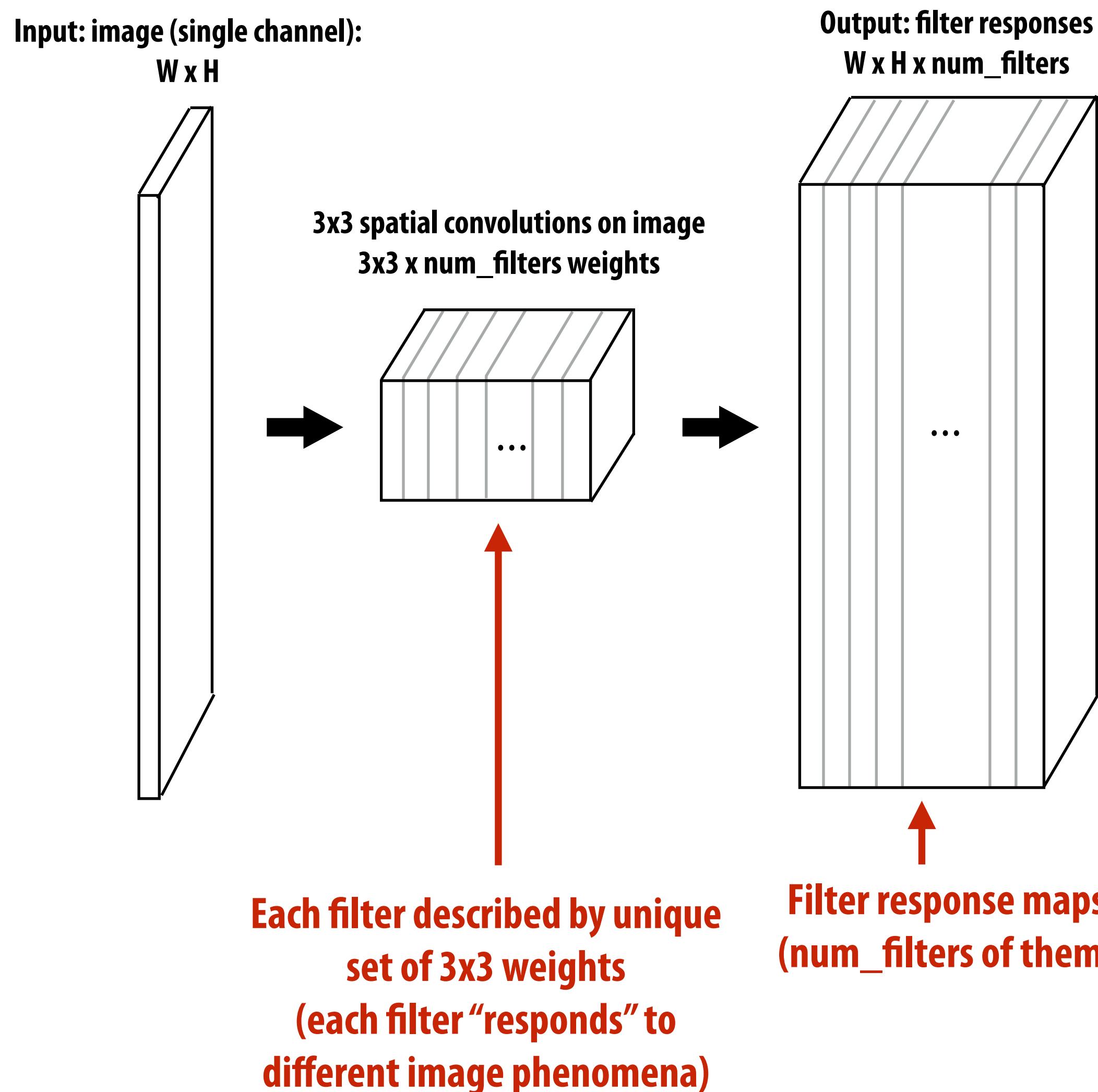
96 11x11x3 filters
(3D because they operate on RGB)



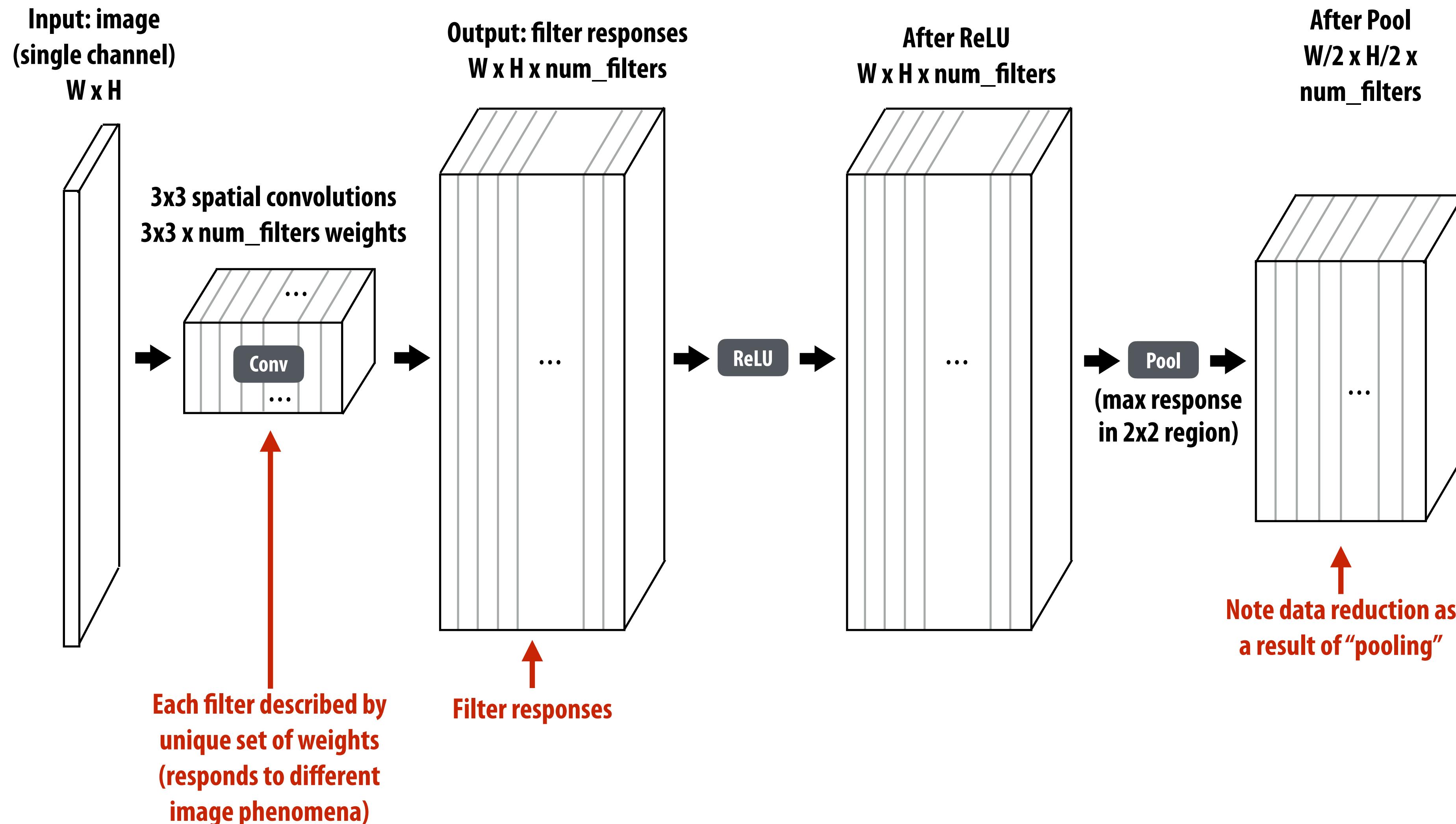
96 responses (normalized)



Applying many filters to an image at once



Adding additional layers



Efficiently implementing convolution layers

Direct implementation of conv layer (batched)

```
float input[IMAGE_BATCH_SIZE][INPUT_HEIGHT][INPUT_WIDTH][INPUT_DEPTH];           // input activations
float output[IMAGE_BATCH_SIZE][INPUT_HEIGHT][INPUT_WIDTH][LAYER_NUM_FILTERS];    // output activations
float layer_weights[LAYER_NUM_FILTERS][LAYER_CONVY][LAYER_CONVX][INPUT_DEPTH];
float layer_biases[LAYER_NUM_FILTERS];

// assumes convolution stride is 1
for (int img=0; img<IMAGE_BATCH_SIZE; img++)
    for (int j=0; j<INPUT_HEIGHT; j++)
        for (int i=0; i<INPUT_WIDTH; i++)
            for (int f=0; f<LAYER_NUM_FILTERS; f++) {
                float tmp = layer_biases[LAYER_NUM_FILTERS];
                for (int kk=0; kk<INPUT_DEPTH; kk++)          // sum over filter responses of input channels
                    for (int jj=0; jj<LAYER_FILTER_Y; jj++)   // spatial convolution (Y)
                        for (int ii=0; ii<LAYER_FILTER_X; ii++) // spatial convolution (X)
                            tmp += layer_weights[f][jj][ii][kk] * input[img][j+jj][i+ii][kk];
                output[img][j][i][f] = tmp;
            }
        }
```

Seven loops with significant input data reuse: reuse of filter weights (during convolution), and reuse of input values (across different filters)

3x3 convolution as matrix-vector product (“explicit gemm”)

Construct matrix from elements of input image

	X ₀₀	X ₀₁	X ₀₂	X ₀₃	...			
	X ₁₀	X ₁₁	X ₁₂	X ₁₃	...			
	X ₂₀	X ₂₁	X ₂₂	X ₂₃	...			
	X ₃₀	X ₃₁	X ₃₂	X ₃₃	...			
...				

Note: 0-pad matrix

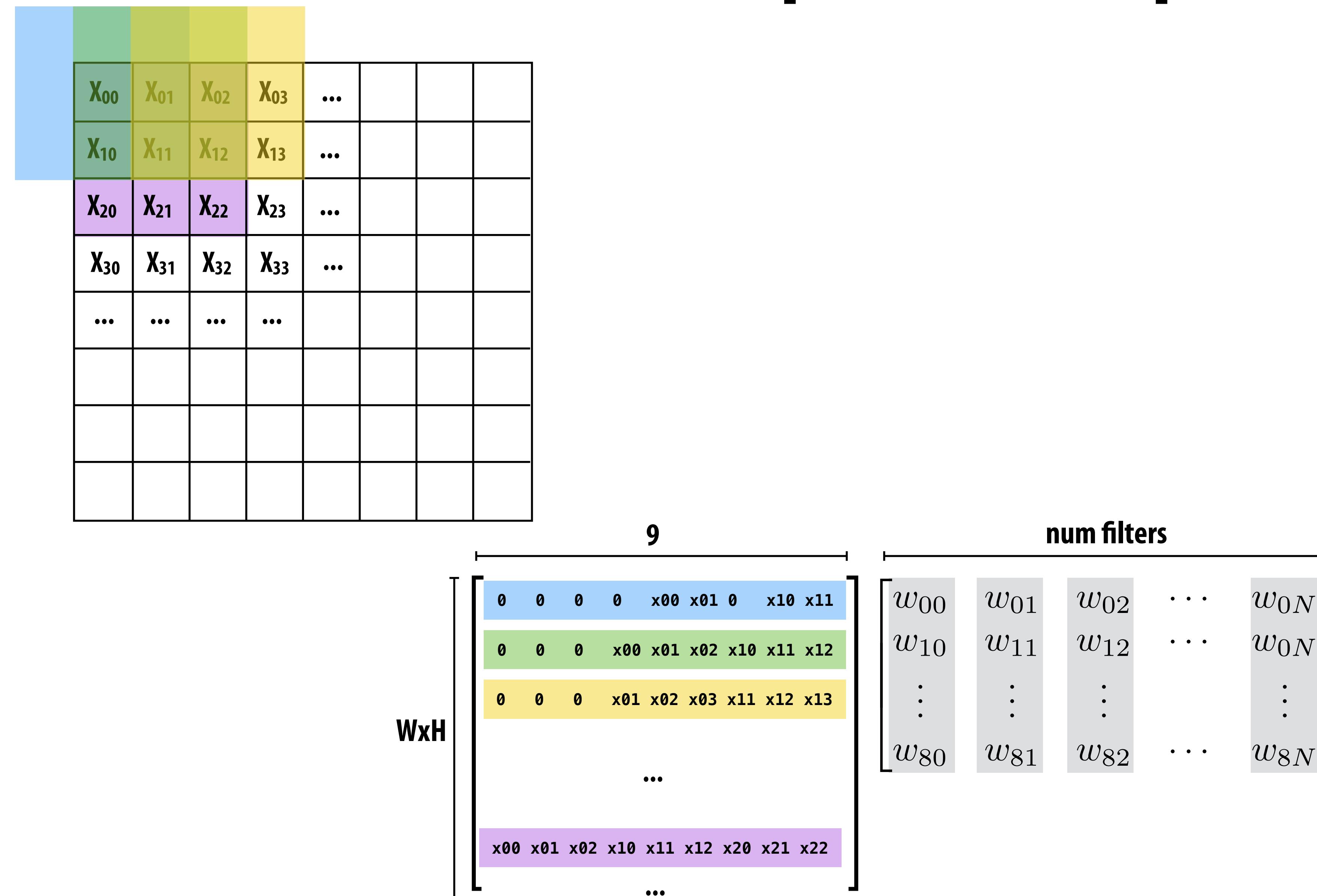
O(N) storage overhead for filter with N elements
Must construct input data matrix

$$\begin{matrix} & & & & & & & 9 \\ & & & & & & & \\ & \left[\begin{array}{ccccccccc} 0 & 0 & 0 & 0 & x_{00} & x_{01} & 0 & x_{10} & x_{11} \end{array} \right] & & & & & & \\ & \left[\begin{array}{ccccccccc} 0 & 0 & 0 & x_{00} & x_{01} & x_{02} & x_{10} & x_{11} & x_{12} \end{array} \right] & & & & & & \\ & \left[\begin{array}{ccccccccc} 0 & 0 & 0 & x_{01} & x_{02} & x_{03} & x_{11} & x_{12} & x_{13} \end{array} \right] & & & & & & \\ & & & & & & & \dots & \\ & & & & & & & & \\ & \left[\begin{array}{ccccccccc} x_{00} & x_{01} & x_{02} & x_{10} & x_{11} & x_{12} & x_{20} & x_{21} & x_{22} \end{array} \right] & & & & & & \\ & & & & & & & \dots & \\ & & & & & & & & \end{matrix}$$

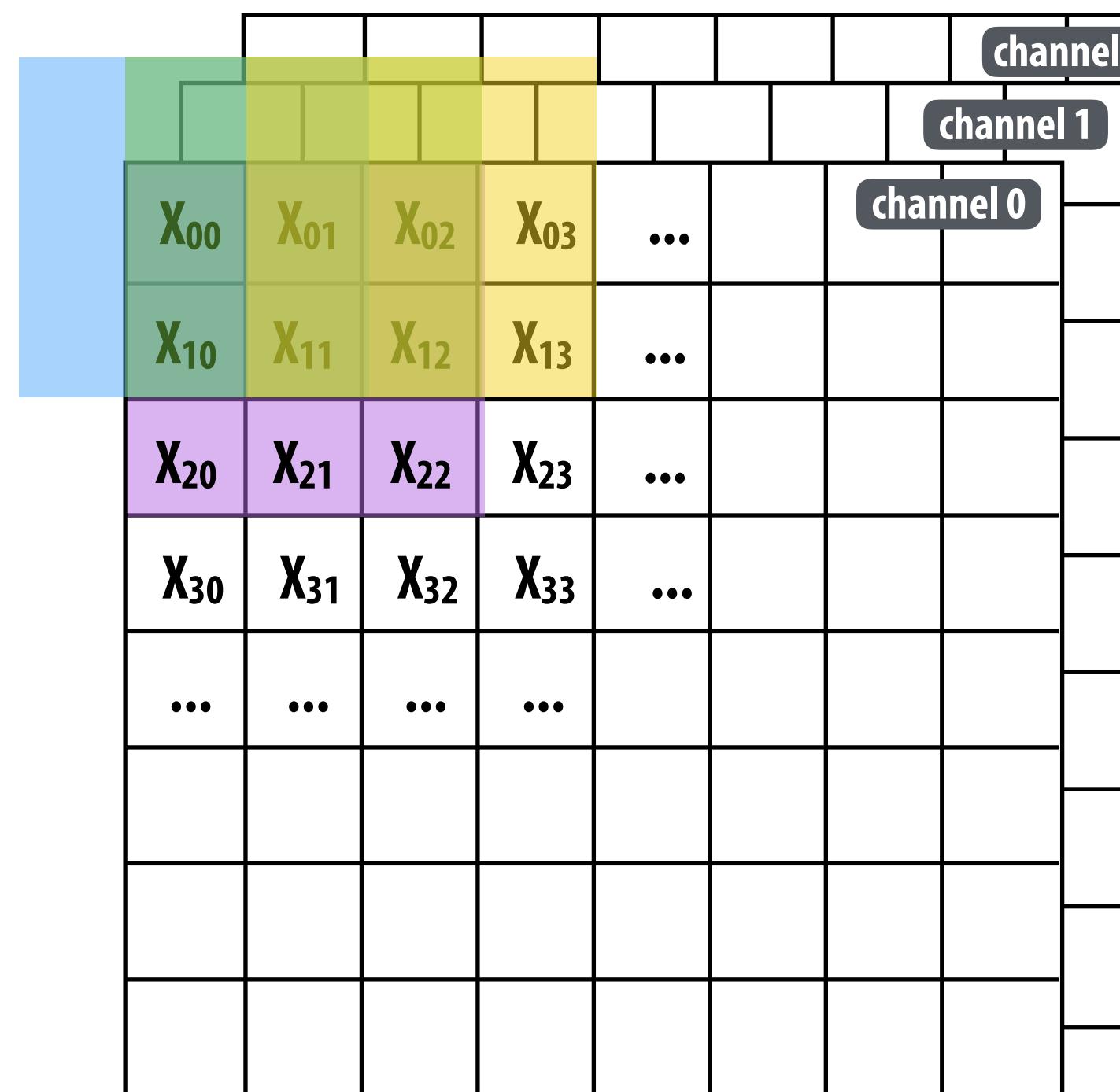
WxH

w_0
 w_1
 \vdots
 w_8

3x3 convolution as matrix-vector product (“explicit gemm”)

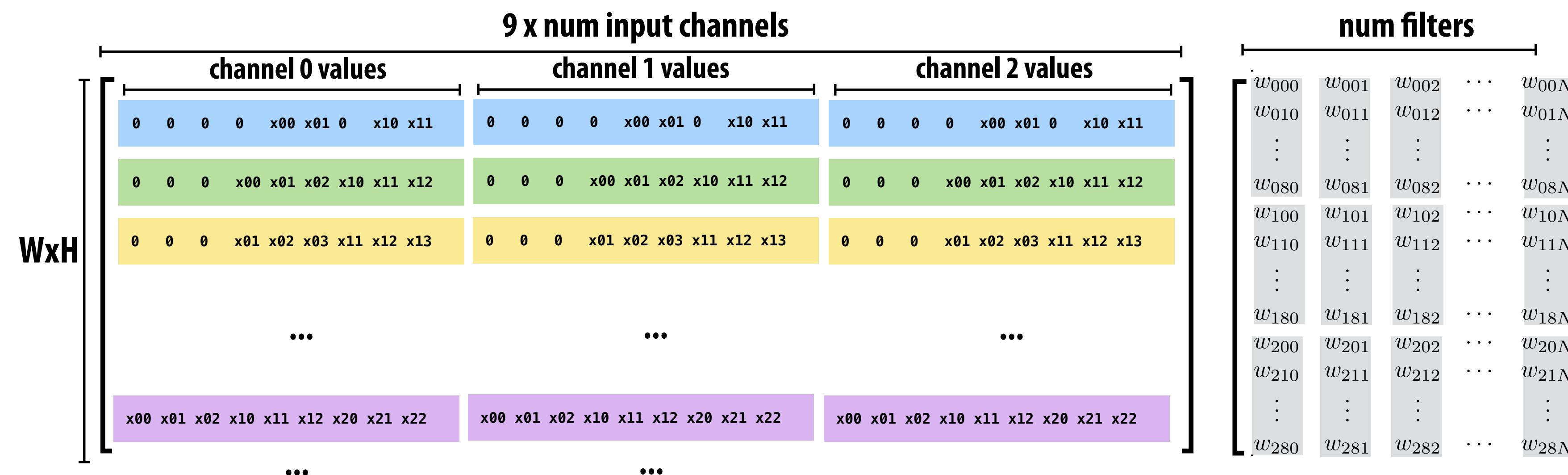


Multiple convolutions on multiple input channels



For each filter, sum responses over input channels

Equivalent to (3 x 3 x num_channels) convolution
on (W x H x num_channels) input data



Conv layer to explicit GEMM mapping

The convolution operation on 4D tensors can be mapped as matrix-multiply operation on 2D matrices

Convolution	GEMM
$y = \text{CONV}(x, w)$	$C = \text{GEMM}(A, B)$
$x[N, H, W, C]$: 4D activation tensor	→ $A[NPQ, RSC]$: 2D convolution matrix
$w[K, R, S, C]$: 4D filter tensor	→ $B[RSC, K]$: 2D filter matrix
$y[N, P, Q, K]$: 4D output tensor	→ $C[NPQ, K]$: 2D output matrix

Symbol reference:

Spatial support of filters: $R \times S$

Input channels: C

Number of filters: K

Batch size: N

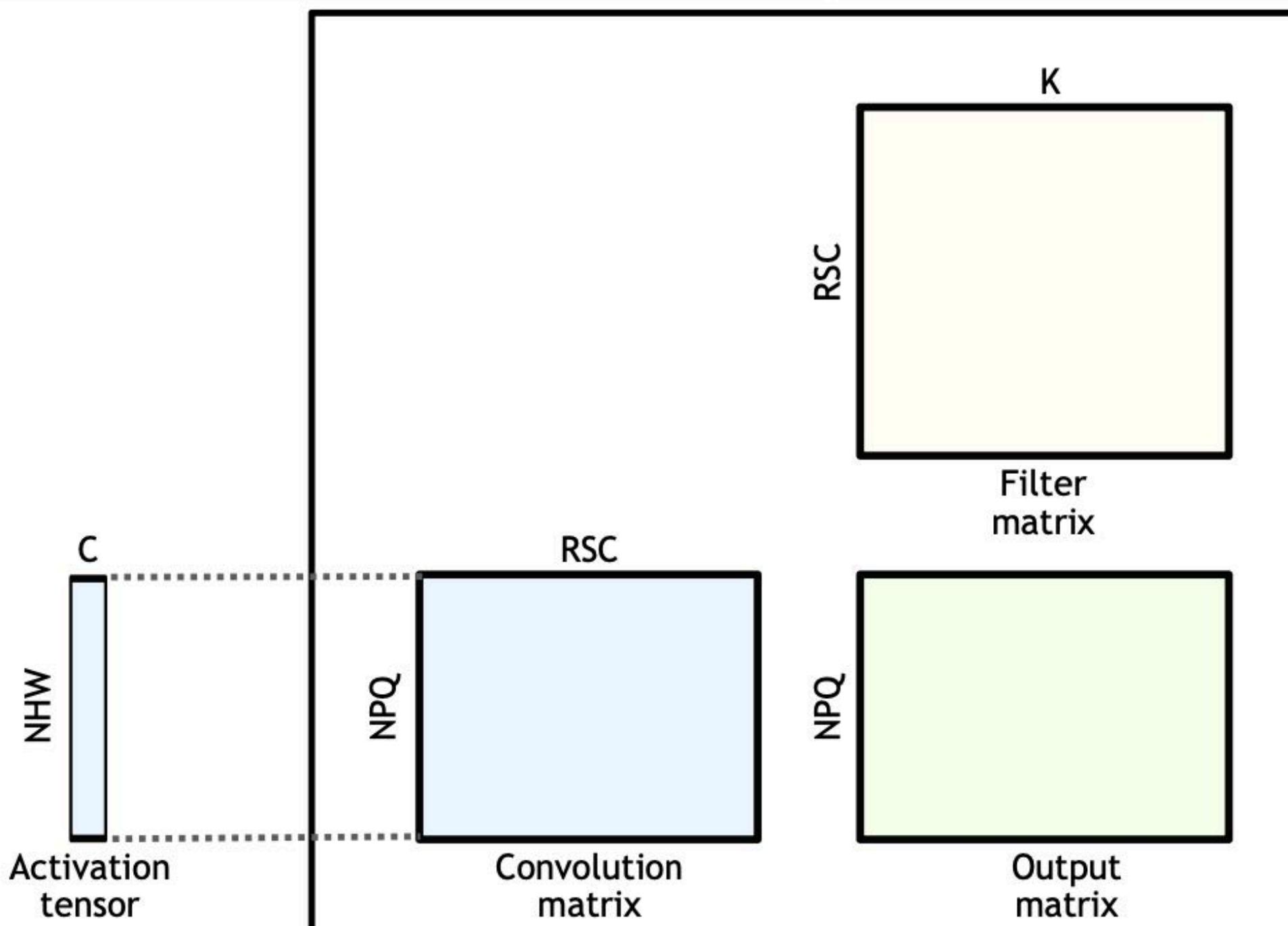
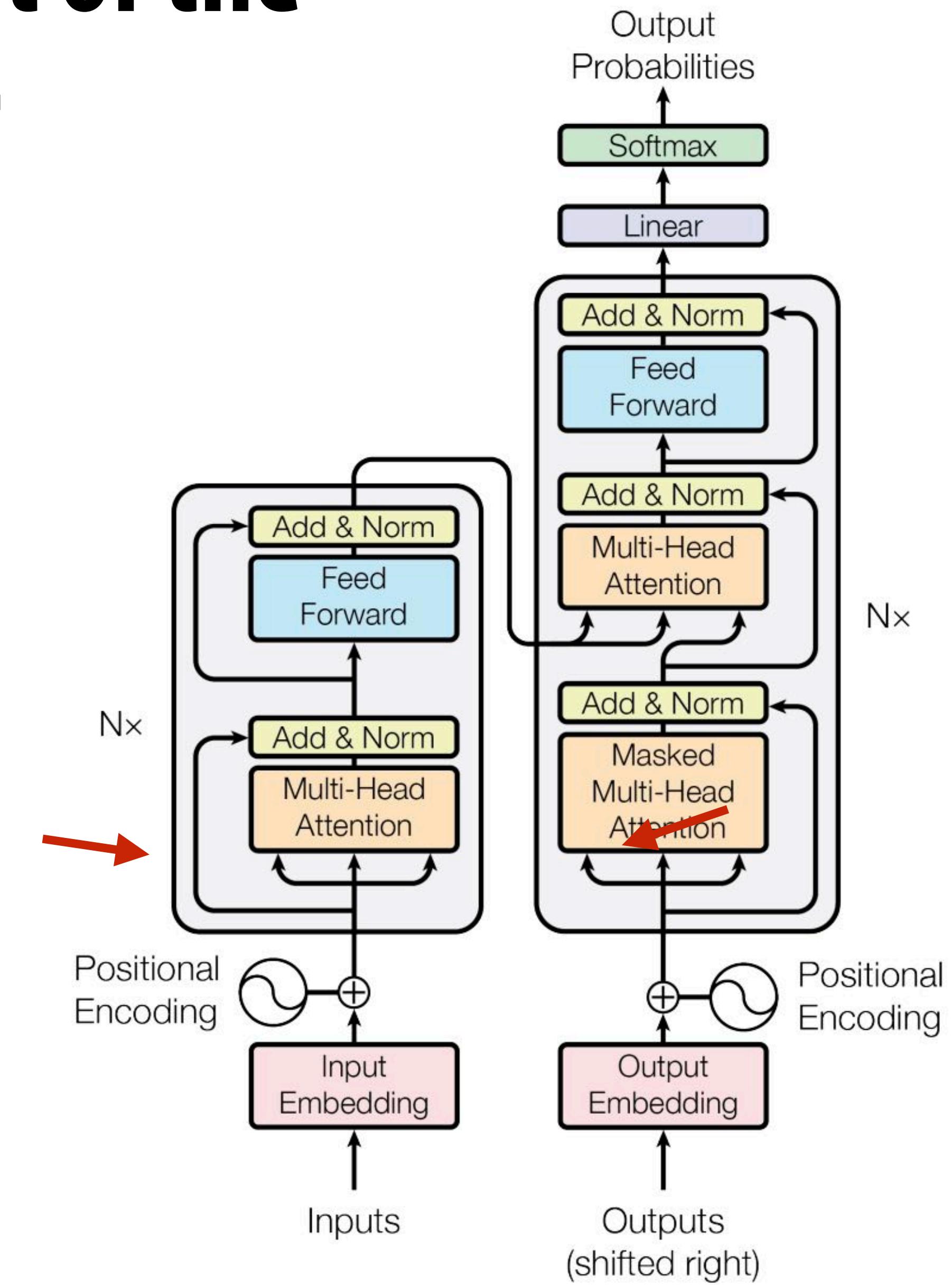


Image credit: NVIDIA

**Matrix multiplication is also at the heart of
the “attention” blocks of a transformer architecture**

Matrix multiplication is at the heart of the “attention” blocks of a transformer architecture

Sequence of tokens in, sequence of tokens out



The importance of dense matrix-matrix multiplication (GEMM) to modern AI

The kernel for...

- Fully-connected layers
- Convolutional layers
- The attention block of a transformer

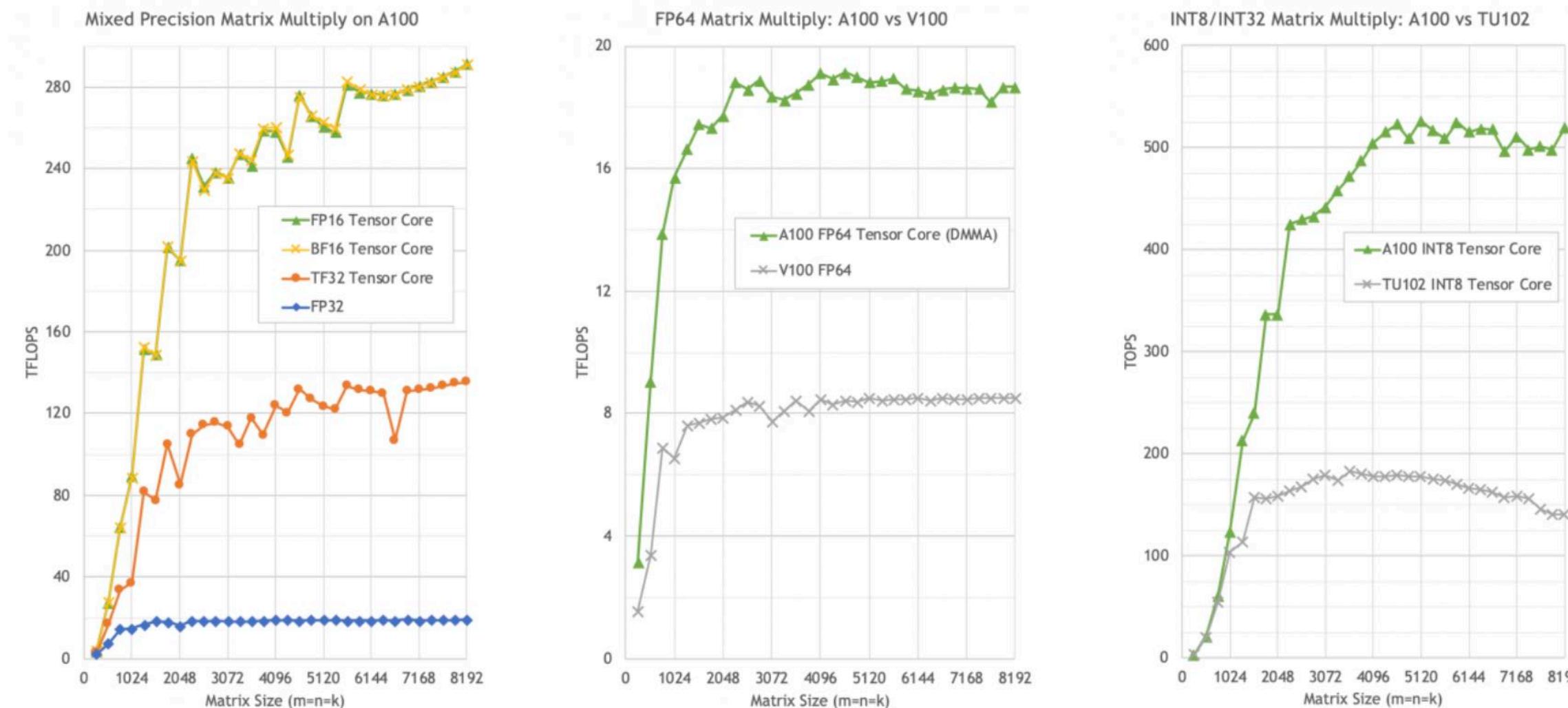
High performance implementations of GEMM exist

cuBLAS Performance

The cuBLAS library is highly optimized for performance on NVIDIA GPUs, and leverages tensor cores for acceleration of low and mixed precision matrix multiplication.

cuBLAS Key Features

- Complete support for all 152 standard BLAS routines
- Support for half-precision and integer matrix multiplication
- GEMM and GEMM extensions optimized for Volta and Turing Tensor Cores
- GEMM performance tuned for sizes used in various Deep Learning models
- Supports CUDA streams for concurrent operations



To use “off the shelf” libraries, must materialize input matrices. Intel® oneAPI Math Kernel Library

Intel®-Optimized Math Library for Numerical Computing

For convolutional layer implications, Increases DRAM traffic by a factor of $R \times S$

(To read input data from activation tensor and constitute “convolution matrix”)

Also requires large amount of additional storage

Optimized Library for Scientific Computing

- Enhanced math routines enable developers and data scientists to create performant science, engineering, or financial applications
- Core functions include BLAS, LAPACK, sparse solvers, fast Fourier transforms (FFT), random number generator functions (RNG), summary statistics, data fitting, and vector math
- Optimizes applications for current and future generations of Intel® CPUs, GPUs, and other accelerators
- Is a seamless upgrade for previous users of the Intel® Math Kernel Library (Intel® MKL)

Download as Part of the Toolkit

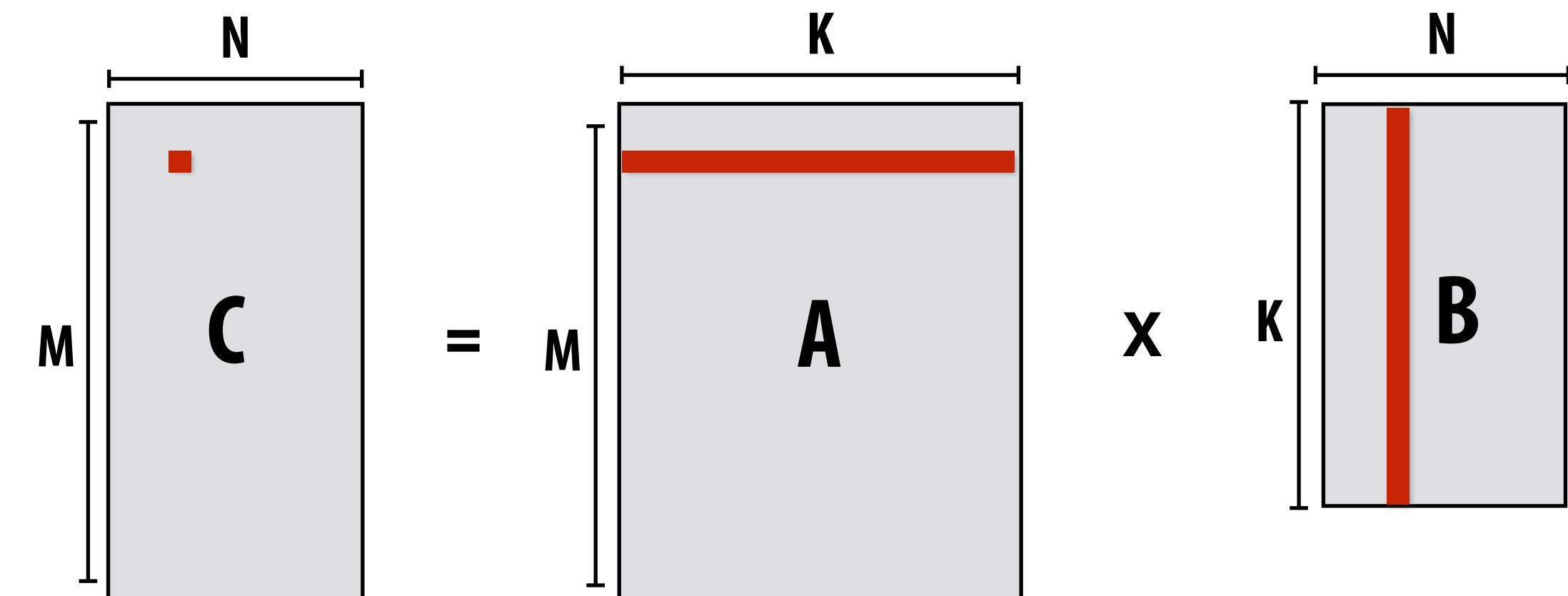
oneMKL is included in the Intel oneAPI Base Toolkit, which is a core set of tools and libraries for developing high-performance, data-centric applications across diverse architectures.

[Get It Now →](#)

Dense matrix multiplication

```
float A[M][K];
float B[K][N];
float C[M][N];

// compute C += A * B
#pragma omp parallel for
for (int j=0; j<M; j++)
    for (int i=0; i<N; i++)
        for (int k=0; k<K; k++)
            C[j][i] += A[j][k] * B[k][i];
```



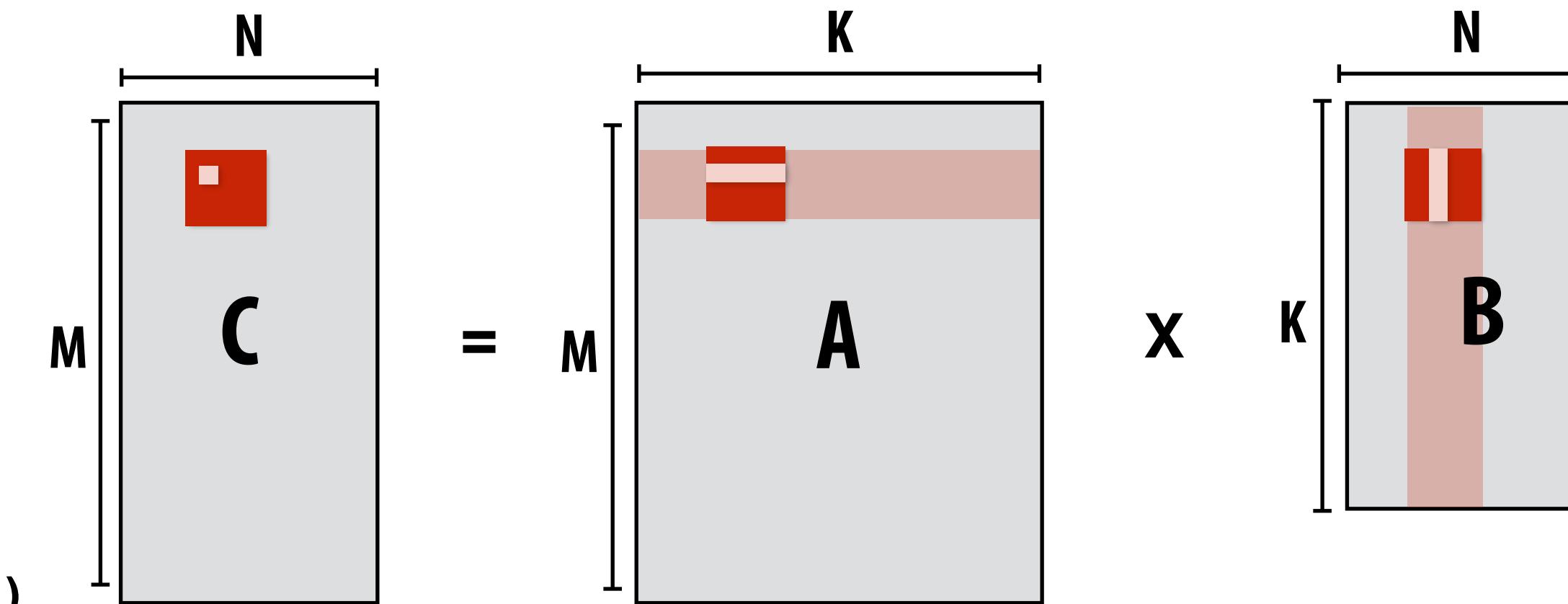
What is the problem with this implementation?

Low arithmetic intensity (does not exploit temporal locality in access to A and B)

Increasing arithmetic intensity by “blocking”

```
float A[M][K];
float B[K][N];
float C[M][N];

// compute C += A * B
#pragma omp parallel for
for (int jblock=0; jblock<M; jblock+=BLOCKSIZE_J)
    for (int iblock=0; iblock<N; iblock+=BLOCKSIZE_I)
        for (int kblock=0; kblock<K; kblock+=BLOCKSIZE_K)
            for (int j=0; j<BLOCKSIZE_J; j++)
                for (int i=0; i<BLOCKSIZE_I; i++)
                    for (int k=0; k<BLOCKSIZE_K; k++)
                        C[jblock+j][iblock+i] += A[jblock+j][kblock+k] * B[kblock+k][iblock+i];
```



Idea: compute partial result for block of C while required blocks of A and B remain in cache
(Assumes BLOCKSIZE chosen to allow block of A, B, and C to remain resident)

Self check: do you want as big a BLOCKSIZE as possible? Why?

Hierarchical blocked matrix mult

Exploit multiple levels of memory hierarchy (increase arithmetic intensity when considering multiple levels of memory hierarchy)

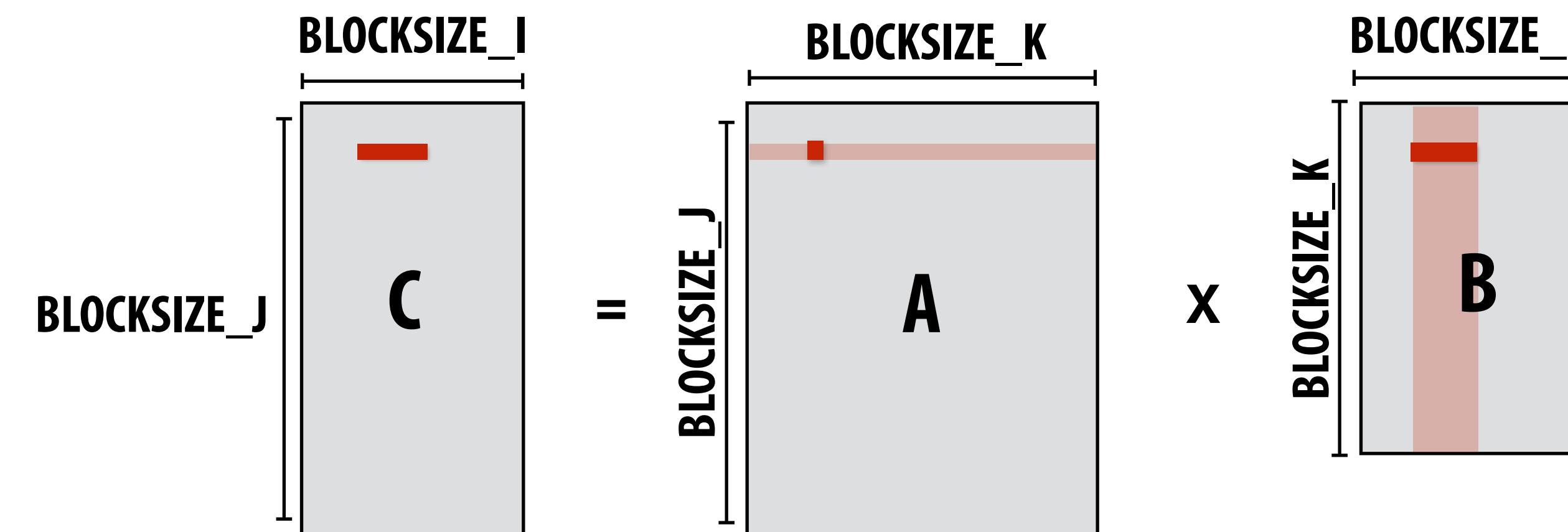
```
float A[M][K];
float B[K][N];
float C[M][N];

// compute C += A * B
#pragma omp parallel for
for (int jblock2=0; jblock2<M; jblock2+=L2_BLOCKSIZE_J)
    for (int iblock2=0; iblock2<N; iblock2+=L2_BLOCKSIZE_I)
        for (int kblock2=0; kblock2<K; kblock2+=L2_BLOCKSIZE_K)
            for (int jblock1=0; jblock1<L1_BLOCKSIZE_J; jblock1+=L1_BLOCKSIZE_J)
                for (int iblock1=0; iblock1<L1_BLOCKSIZE_I; iblock1+=L1_BLOCKSIZE_I)
                    for (int kblock1=0; kblock1<L1_BLOCKSIZE_K; kblock1+=L1_BLOCKSIZE_K)
                        for (int j=0; j<BLOCKSIZE_J; j++)
                            for (int i=0; i<BLOCKSIZE_I; i++)
                                for (int k=0; k<BLOCKSIZE_K; k++)
...
...
```

Not shown: final level of “blocking” for register locality...

Vectorized, blocked dense matrix multiplication (1)

Consider SIMD parallelism within a block



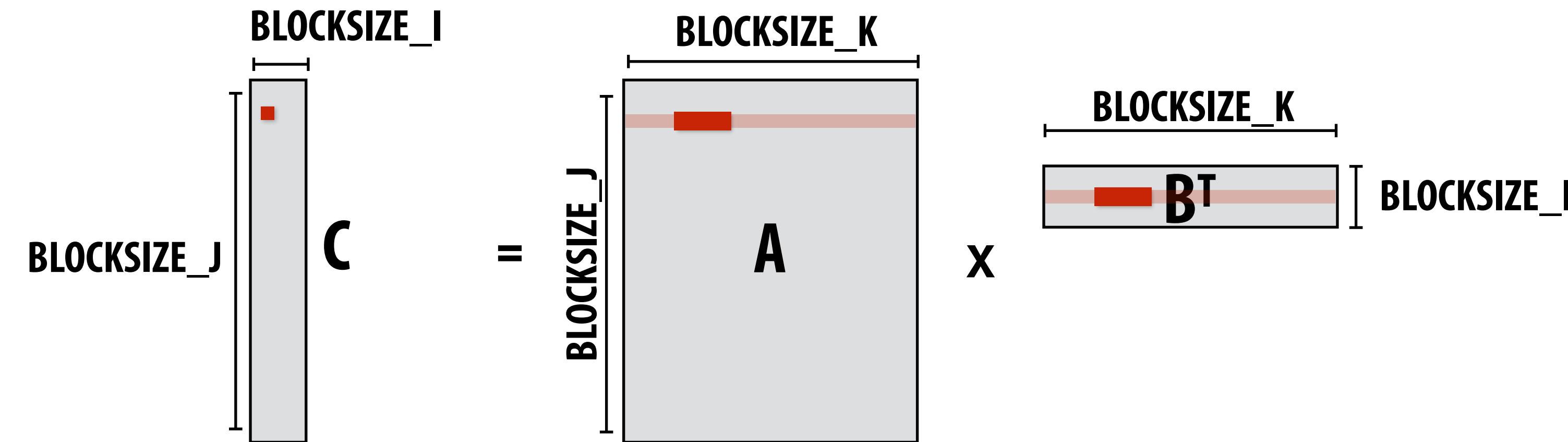
```
...
for (int j=0; j<BLOCKSIZE_J; j++) {
    for (int i=0; i<BLOCKSIZE_I; i+=SIMD_WIDTH) {
        simd_vec C_accum = vec_load(&C[jblock+j][iblock+i]);
        for (int k=0; k<BLOCKSIZE_K; k++) {
            // C = A*B + C
            simd_vec A_val = splat(&A[jblock+j][kblock+k]); // load a single element in vector register
            simd_muladd(A_val, vec_load(&B[kblock+k][iblock+i]), C_accum);
        }
        vec_store(&C[jblock+j][iblock+i], C_accum);
    }
}
```

Vectorize i loop

Good: also improves spatial locality in access to B

Bad: working set increased by SIMD_WIDTH, still walking over B in large steps

Vectorized, blocked dense matrix multiplication (2)



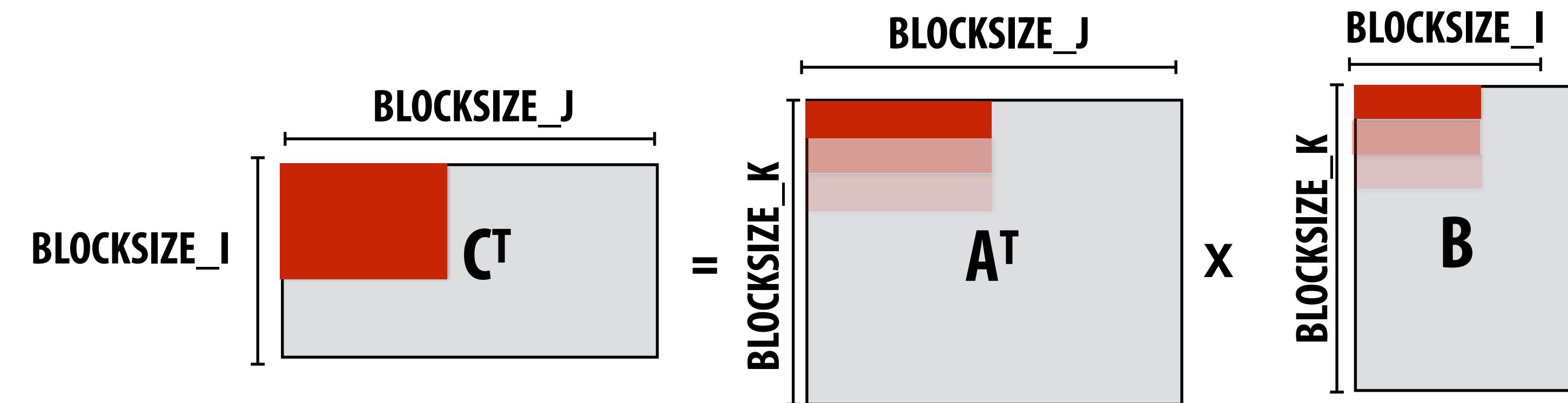
```
...
for (int j=0; j<BLOCKSIZE_J; j++)
    for (int i=0; i<BLOCKSIZE_I; i++) {
        float C_scalar = C[jblock+j][iblock+i];
        // C_scalar += dot(row of A, row of B)
        for (int k=0; k<BLOCKSIZE_K; k+=SIMD_WIDTH) {
            C_scalar += simd_dot(vec_load(&A[jblock+j][kblock+k]), vec_load(&Btrans[iblock+i][[kblock+k]));
        }
        C[jblock+j][iblock+i] = C_scalar;
    }
```

Assume i dimension is small. Previous vectorization scheme (1) would not work well.

Pre-transpose block of B (copy block of B to temp buffer in transposed form)

Vectorize innermost loop

Vectorized, blocked dense matrix multiplication (3)



```
// assume blocks of A and C are pre-transposed as Atrans and Ctrans
for (int j=0; j<BLOCKSIZE_J; j+=SIMD_WIDTH) {
    for (int i=0; i<BLOCKSIZE_I; i+=SIMD_WIDTH) {

        simd_vec C_accum[SIMD_WIDTH];
        for (int k=0; k<SIMD_WIDTH; k++) // load C_accum for a SIMD_WIDTH x SIMD_WIDTH chunk of C^T
            C_accum[k] = vec_load(&Ctrans[iblock+i+k][jblock+j]);

        for (int k=0; k<BLOCKSIZE_K; k++) {
            simd_vec bvec = vec_load(&B[kblock+k][iblock+i]);
            for (int kk=0; kk<SIMD_WIDTH; kk++) // innermost loop items not dependent
                simd_muladd(vec_load(&Atrans[kblock+k][jblock+j], splat(bvec[kk]), C_accum[kk]);
        }

        for (int k=0; k<SIMD_WIDTH; k++)
            vec_store(&Ctrans[iblock+i+k][jblock+j], C_accum[k]);
    }
}
```

Different layers of a single DNN may benefit from unique scheduling strategies (different matrix dimensions)

**Notice sizes of weights and activations in this network:
(and consider SIMD widths of modern machines).**

Ug for library implementers!

Table 1. MobileNet Body Architecture

Type / Stride	Filter Shape	Input Size
Conv / s2	$3 \times 3 \times 3 \times 32$	$224 \times 224 \times 3$
Conv dw / s1	$3 \times 3 \times 32$ dw	$112 \times 112 \times 32$
Conv / s1	$1 \times 1 \times 32 \times 64$	$112 \times 112 \times 32$
Conv dw / s2	$3 \times 3 \times 64$ dw	$112 \times 112 \times 64$
Conv / s1	$1 \times 1 \times 64 \times 128$	$56 \times 56 \times 64$
Conv dw / s1	$3 \times 3 \times 128$ dw	$56 \times 56 \times 128$
Conv / s1	$1 \times 1 \times 128 \times 128$	$56 \times 56 \times 128$
Conv dw / s2	$3 \times 3 \times 128$ dw	$56 \times 56 \times 128$
Conv / s1	$1 \times 1 \times 128 \times 256$	$28 \times 28 \times 128$
Conv dw / s1	$3 \times 3 \times 256$ dw	$28 \times 28 \times 256$
Conv / s1	$1 \times 1 \times 256 \times 256$	$28 \times 28 \times 256$
Conv dw / s2	$3 \times 3 \times 256$ dw	$28 \times 28 \times 256$
Conv / s1	$1 \times 1 \times 256 \times 512$	$14 \times 14 \times 256$
5× Conv dw / s1	$3 \times 3 \times 512$ dw	$14 \times 14 \times 512$
	$1 \times 1 \times 512 \times 512$	$14 \times 14 \times 512$
Conv dw / s2	$3 \times 3 \times 512$ dw	$14 \times 14 \times 512$
Conv / s1	$1 \times 1 \times 512 \times 1024$	$7 \times 7 \times 512$
Conv dw / s2	$3 \times 3 \times 1024$ dw	$7 \times 7 \times 1024$
Conv / s1	$1 \times 1 \times 1024 \times 1024$	$7 \times 7 \times 1024$
Avg Pool / s1	Pool 7×7	$7 \times 7 \times 1024$
FC / s1	1024×1000	$1 \times 1 \times 1024$
Softmax / s1	Classifier	$1 \times 1 \times 1000$

Matrix multiplication implementations

Optimization: do not materialize full matrix ("implicit gemm")

This is a naive implementation that does not perform blocking, but indexes into input weight and activation tensors.

Symbol reference:

Spatial support of filters: R x S

Input channels: C

Number of filters: K

Batch size: N

GEMM TRIPLE NEST LOOP

```
int GEMM_M = N * P * Q;
int GEMM_N = K;
int GEMM_K = R * S * C;

for (int gemm_m = 0; gemm_m < GEMM_M; ++gemm_m) {
    for (int gemm_n = 0; gemm_n < GEMM_N; ++gemm_n) {

        int n = gemm_m / (PQ);
        int npq_residual = gemm_m % (PQ);
        int p = npq_residual / Q;
        int q = npq_residual % Q;

        Accumulator accum = 0;
        for (int gemm_k = 0; gemm_k < GEMM_K; ++gemm_k) {

            int k = gemm_n;
            int crs_residual = gemm_k / C;
            int r = crs_residual / S;
            int s = crs_residual % S;
            int c = gemm_k % C;

            int h = h_bar(p, r);
            int w = w_bar(q, s);

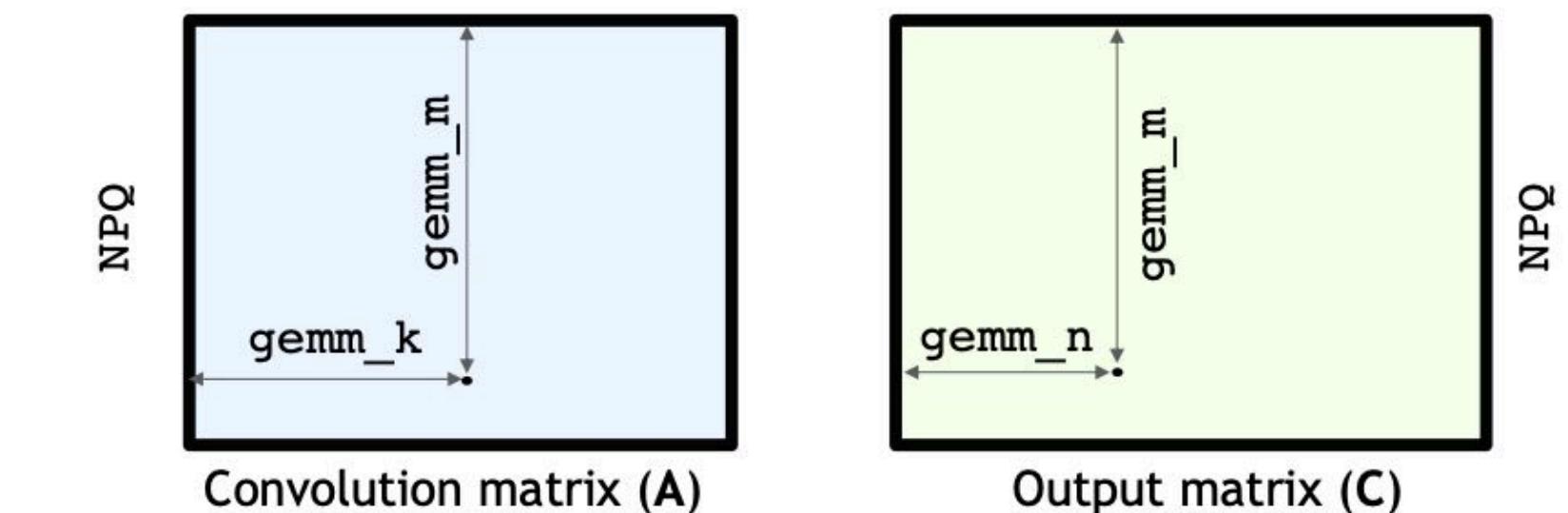
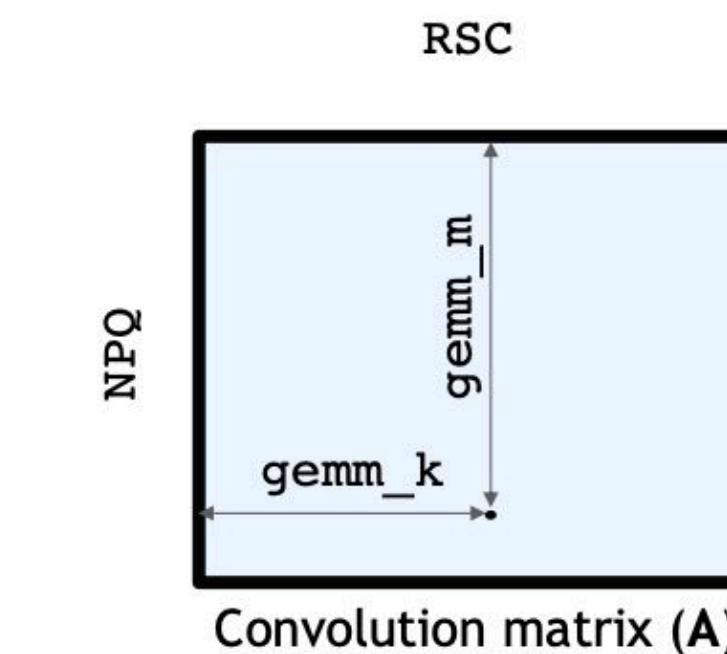
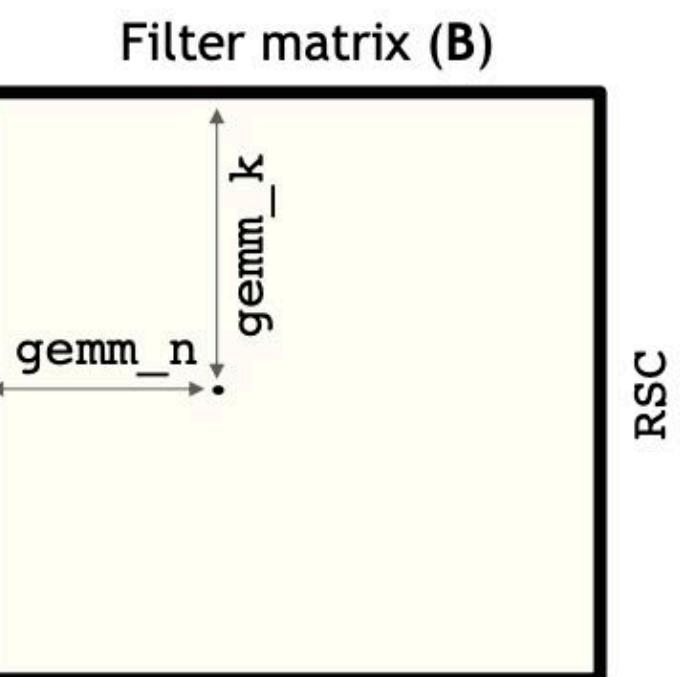
            ElementA a = activation_tensor.at({n, h, w, c});
            ElementB b = filter_tensor.at({k, r, s, c});
            accum += a * b;
        }
    }
}
```

```
C[gemm_m * K + gemm_n] = accum;
```

CONVOLUTION MAPPED TO GEMM

GEMM M-by-N-by-K dimensions

```
GEMM-M = NPQ
GEMM-N = K
GEMM-K = RSC
```



$$C[gemm_m, gemm_n] = \sum_{gemm_k=0}^{RSC-1} (A[gemm_m, gemm_k] * B[gemm_k, gemm_n])$$

Optimization: do not materialize full matrix (“implicit gemm”)

Better implementation:
materialize only a sub-block of the
convolution matrix at a time in
GPU on-chip “shared memory”

Forward Propagation (Fprop)	
$y = \text{CONV}(x, w)$	
$x[N, H, W, C]$: 4D activation tensor
$w[K, R, S, C]$: 4D filter tensor
$y[N, P, Q, K]$: 4D output tensor

Does not require additional off-chip storage and
does not increase required DRAM traffic.

Use well-tuned shared-memory based GEMM
routines to perform sub-block GEMM (see CUTLASS)

Symbol reference:

Output size: $P \times Q$

Spatial support of filters: $R \times S$

Input channels: C

Number of filters (output channels): K

Batch size: N

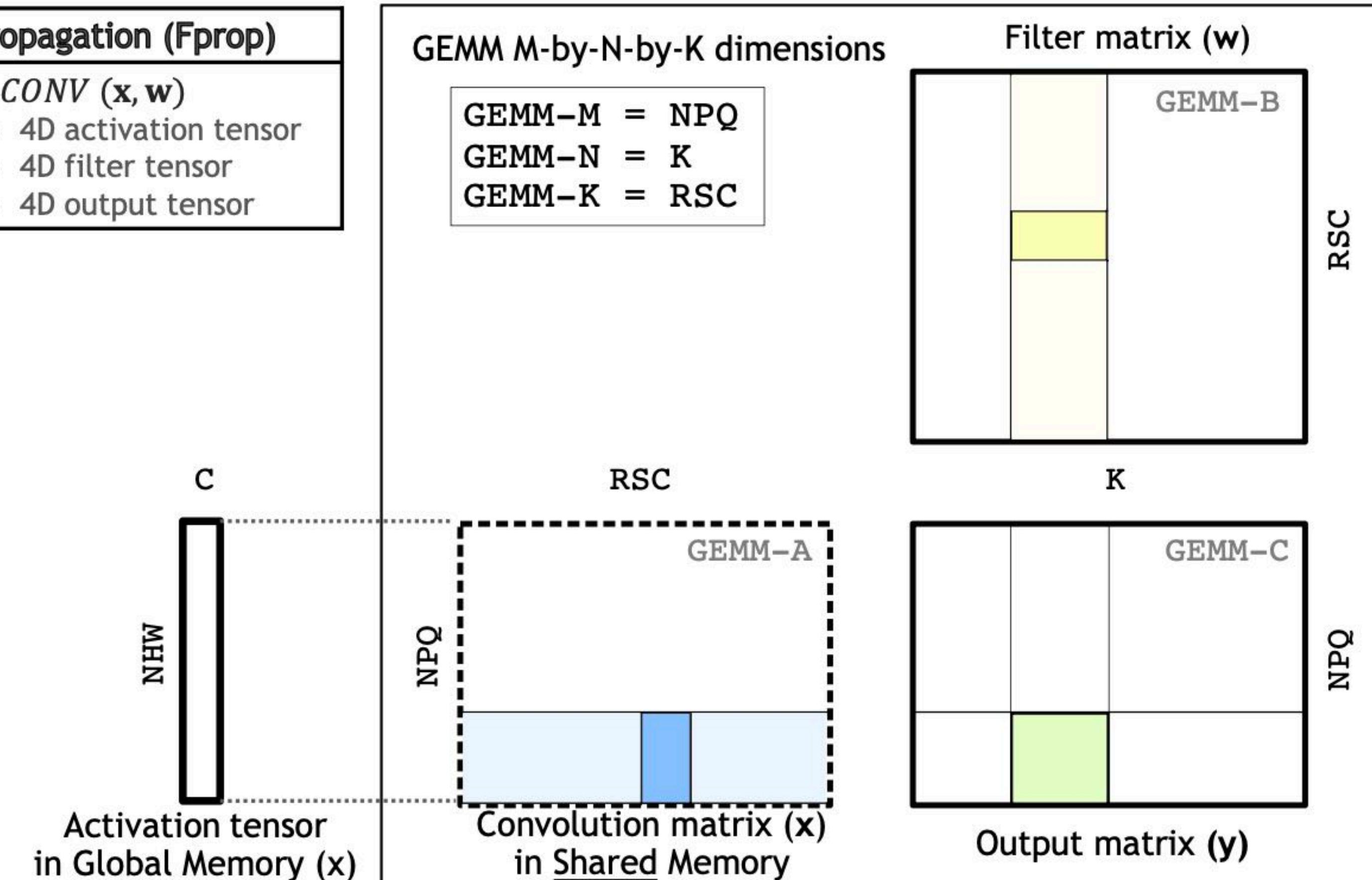


Image credit: NVIDIA

NVIDIA CUTLASS

Basic primitives/building block for implementing your custom high performance DNN layers. (e.g, unusual sizes that haven't been heavily tuned by cuDNN)

The screenshot shows the GitHub repository page for `NVIDIA/cutlass`. The top navigation bar includes links for Pull requests, Issues, Marketplace, and Explore. Below the header, there's a search bar and a link to the repository's README. The main content area displays the repository's structure with branches and tags. A list of recent commits is shown, with the most recent one being a merge from `sfjeng1999` (commit `#449`) that fixes code indentation. Below the commit list is the `README.md` file, which contains a diagram illustrating the tensor operations supported by CUTLASS. The diagram shows various tensors (F, M, N) being processed through different stages of computation, likely representing GEMM operations.

CUTLASS 2.8

CUTLASS 2.8 - November 2021

CUTLASS is a collection of CUDA C++ template abstractions for implementing high-performance matrix-multiplication (GEMM) and related computations at all levels and scales within CUDA. It incorporates strategies for hierarchical decomposition and data movement similar to those used to implement cuBLAS and cuDNN. CUTLASS decomposes these "moving parts" into reusable, modular software components abstracted by C++ template classes. These thread-wide, warp-wide, block-wide, and device-wide primitives can be specialized and tuned via custom tiling sizes, data types, and other algorithmic policy. The resulting flexibility simplifies their use as building blocks within custom kernels and applications.

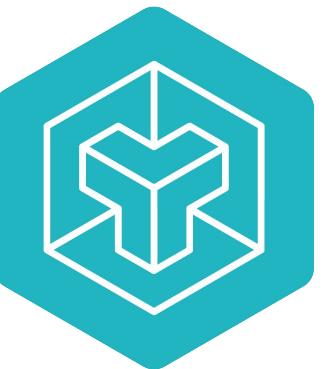
Fast (in-shared memory) GEMM

Fast WARP level GEMMs

Iterators for fast block loading/tensor indexing

Tensor reductions

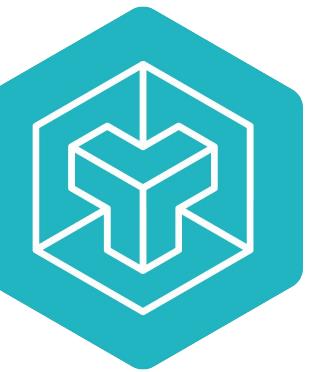
Etc.



- Language support for operations that load/store tensors
- Load “blocks” of data into GPU shared memory
- Perform data-parallel operations on those blocks

A simple blocked matrix multiplication

```
# Do in parallel
for m in range(0, M, BLOCK_SIZE_M):
    # Do in parallel
    for n in range(0, N, BLOCK_SIZE_N):
        acc = zeros((BLOCK_SIZE_M, BLOCK_SIZE_N), dtype=float32)
        for k in range(0, K, BLOCK_SIZE_K):
            a = A[m : m+BLOCK_SIZE_M, k : k+BLOCK_SIZE_K]
            b = B[k : k+BLOCK_SIZE_K, n : n+BLOCK_SIZE_N]
            acc += dot(a, b)
        C[m : m+BLOCK_SIZE_M, n : n+BLOCK_SIZE_N] = acc
```



Full Triton reference implementation: two levels of blocking

```

@triton.jit
def matmul_kernel(
    # Pointers to matrices
    a_ptr, b_ptr, c_ptr,
    # Matrix dimensions
    M, N, K,
    # The stride variables represent how much to increase the ptr by when moving by 1
    # element in a particular dimension. E.g. `stride_am` is how much to increase `a_ptr`
    # by to get the element one row down (A has M rows).
    stride_am, stride_ak, #
    stride_bk, stride_bn, #
    stride_cm, stride_cn,
    # Meta-parameters
    BLOCK_SIZE_M: tl.constexpr, BLOCK_SIZE_N: tl.constexpr, BLOCK_SIZE_K: tl.constexpr, #
    GROUP_SIZE_M: tl.constexpr, #
    ACTIVATION: tl.constexpr #
):
    """Kernel for computing the matmul C = A x B.
    A has shape (M, K), B has shape (K, N) and C has shape (M, N)
    """
    # -----
    # Map program ids `pid` to the block of C it should compute.
    # This is done in a grouped ordering to promote L2 data reuse.
    # See above `L2 Cache Optimizations` section for details.
    pid = tl.program_id(axis=0)
    num_pid_m = tl.cdiv(M, BLOCK_SIZE_M)
    num_pid_n = tl.cdiv(N, BLOCK_SIZE_N)
    num_pid_in_group = GROUP_SIZE_M * num_pid_n
    group_id = pid // num_pid_in_group
    first_pid_m = group_id * GROUP_SIZE_M
    group_size_m = min(num_pid_m - first_pid_m, GROUP_SIZE_M)
    pid_m = first_pid_m + ((pid % num_pid_in_group) % group_size_m)
    pid_n = (pid % num_pid_in_group) // group_size_m

```

```

# -----
# Add some integer bound assumptions.
# This helps to guide integer analysis in the backend to optimize
# load/store offset address calculation
tl.assume(pid_m >= 0)
tl.assume(pid_n >= 0)
tl.assume(stride_am > 0)
tl.assume(stride_ak > 0)
tl.assume(stride_bn > 0)
tl.assume(stride_bk > 0)
tl.assume(stride_cm > 0)
tl.assume(stride_cn > 0)

# -----
# Create pointers for the first blocks of A and B.
# We will advance this pointer as we move in the K direction
# and accumulate
# `a_ptrs` is a block of [BLOCK_SIZE_M, BLOCK_SIZE_K] pointers
# `b_ptrs` is a block of [BLOCK_SIZE_K, BLOCK_SIZE_N] pointers
# See above `Pointer Arithmetic` section for details
offs_am = (pid_m * BLOCK_SIZE_M + tl.arange(0, BLOCK_SIZE_M)) % M
offs_bn = (pid_n * BLOCK_SIZE_N + tl.arange(0, BLOCK_SIZE_N)) % N
offs_k = tl.arange(0, BLOCK_SIZE_K)
a_ptrs = a_ptr + (offs_am[:, None] * stride_am + offs_k[None, :] * stride_ak)
b_ptrs = b_ptr + (offs_k[:, None] * stride_bk + offs_bn[None, :] * stride_bn)

# -----
# Iterate to compute a block of the C matrix.
# We accumulate into a `[BLOCK_SIZE_M, BLOCK_SIZE_N]` block
# of fp32 values for higher accuracy.
# `accumulator` will be converted back to fp16 after the loop.
accumulator = tl.zeros((BLOCK_SIZE_M, BLOCK_SIZE_N), dtype=tl.float32)
for k in range(0, tl.cdiv(K, BLOCK_SIZE_K)):
    # Load the next block of A and B, generate a mask by checking the K dimension.
    # If it is out of bounds, set it to 0.
    a = tl.load(a_ptrs, mask=offs_k[None, :] < K - k * BLOCK_SIZE_K, other=0.0)
    b = tl.load(b_ptrs, mask=offs_k[:, None] < K - k * BLOCK_SIZE_K, other=0.0)
    # We accumulate along the K dimension.
    accumulator = tl.dot(a, b, accumulator)
    # Advance the ptrs to the next K block.
    a_ptrs += BLOCK_SIZE_K * stride_ak
    b_ptrs += BLOCK_SIZE_K * stride_bk
    # You can fuse arbitrary activation functions here
    # while the accumulator is still in FP32!
    if ACTIVATION == "leaky_relu":
        accumulator = leaky_relu(accumulator)
    c = accumulator.to(tl.float16)

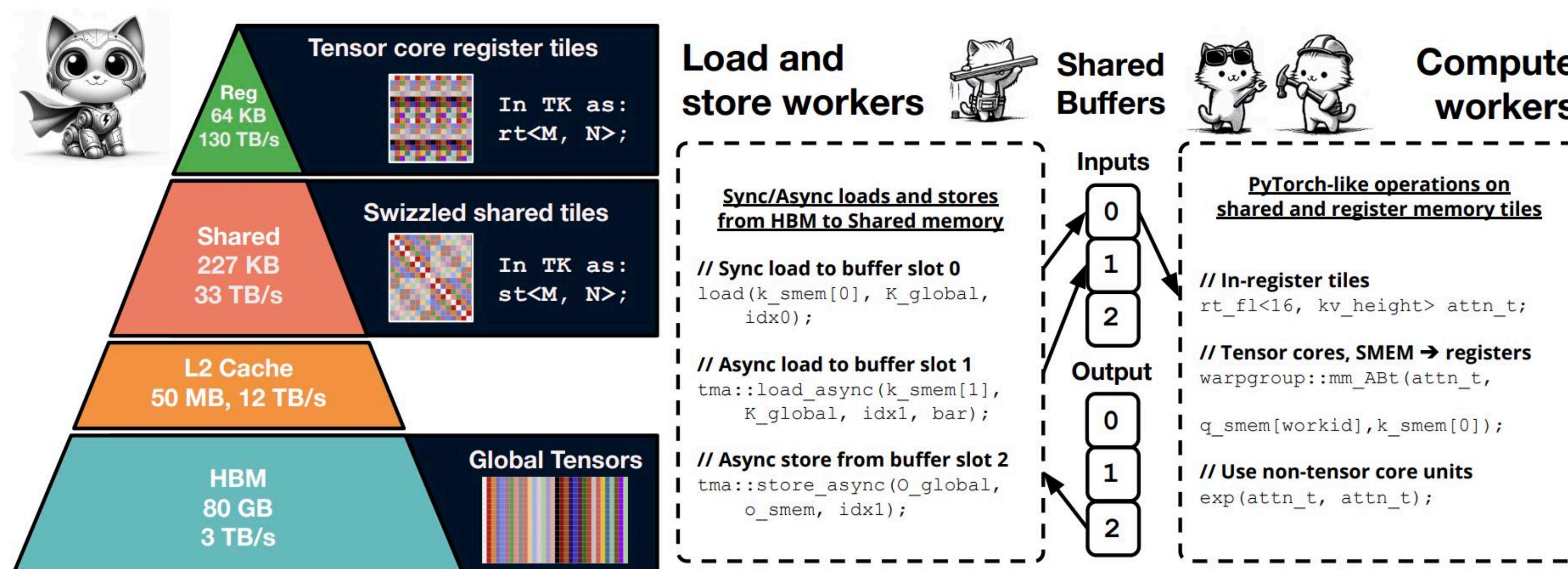
# -----
# Write back the block of the output matrix C with masks.
offs_cm = pid_m * BLOCK_SIZE_M + tl.arange(0, BLOCK_SIZE_M)
offs_cn = pid_n * BLOCK_SIZE_N + tl.arange(0, BLOCK_SIZE_N)
c_ptrs = c_ptr + stride_cm * offs_cm[:, None] + stride_cn * offs_cn[None, :]
c_mask = (offs_cm[:, None] < M) & (offs_cn[None, :] < N)
tl.store(c_ptrs, c, mask=c_mask)

```

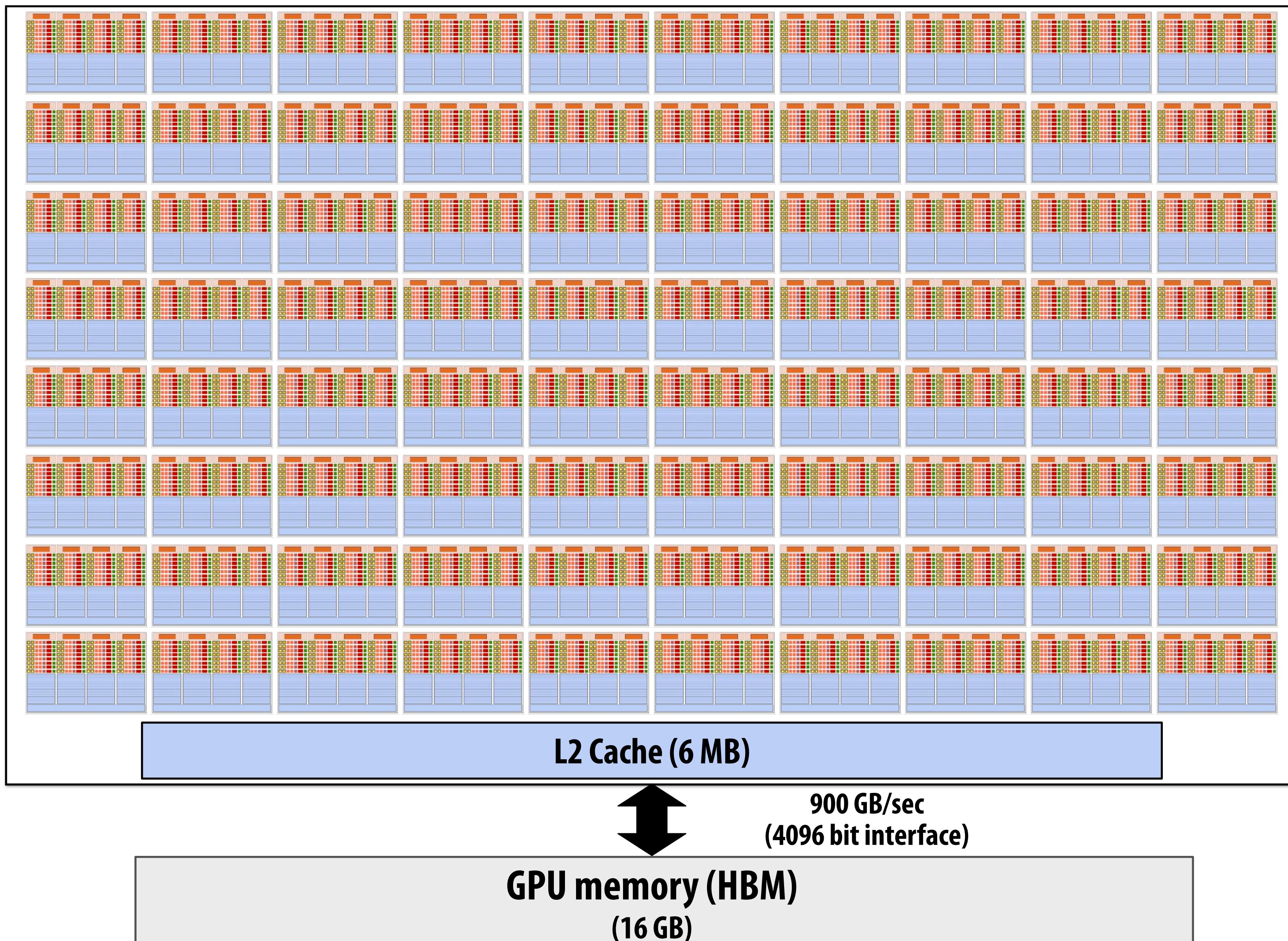
Thunderkittens



- CUDA library of useful tile-based programming primitives
- Intended to make advanced developers (CS149-level folks) more productive writing blocked code
 - Async load/store of tiles
 - Support for advanced memory layouts (blocked tiles, interleaved elements, etc.)



Recall: NVIDIA V100 GPU (80 SMs)



Many processing units and many tensor cores.

Need “a lot of parallel work” to fill the machine.

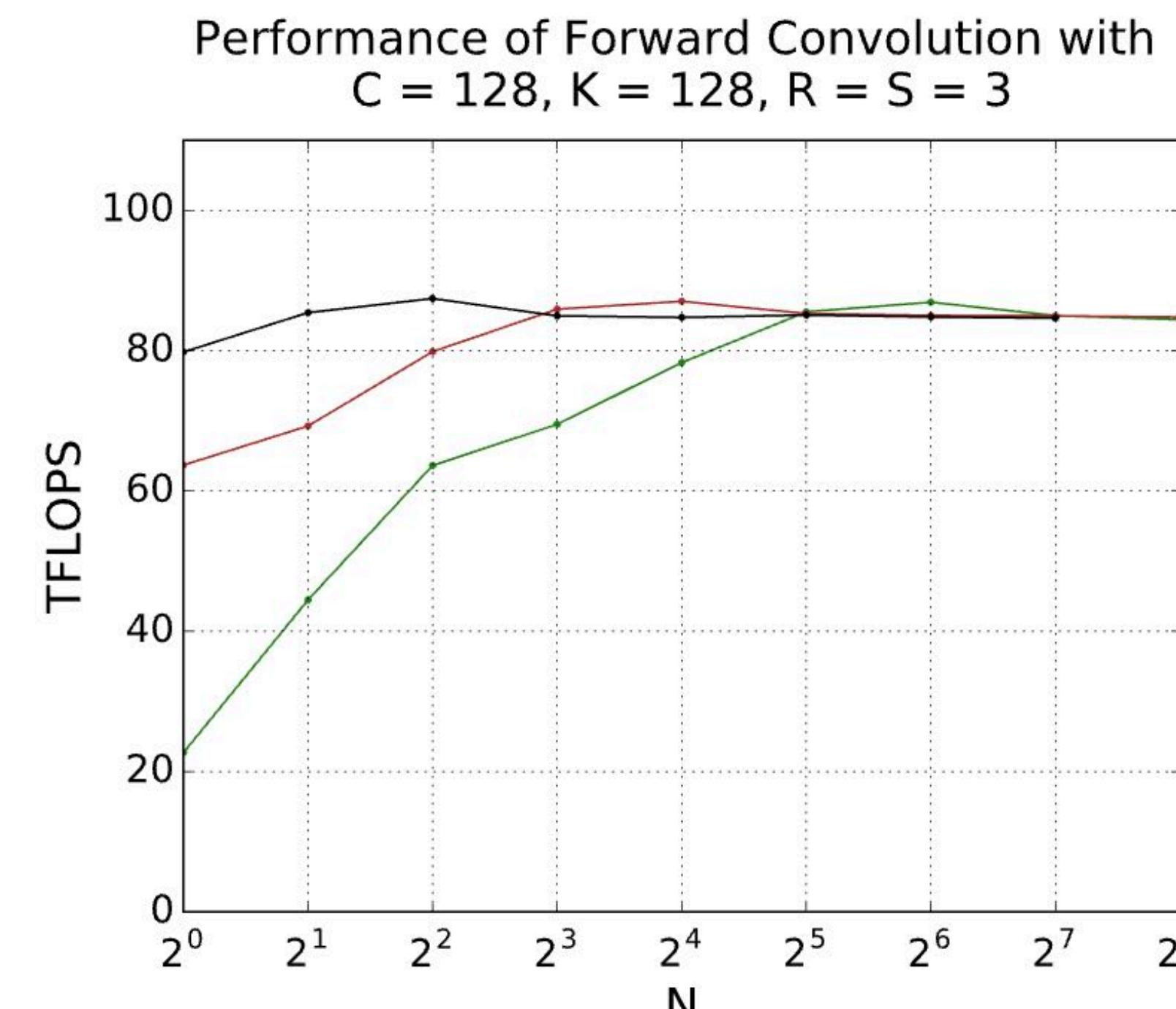
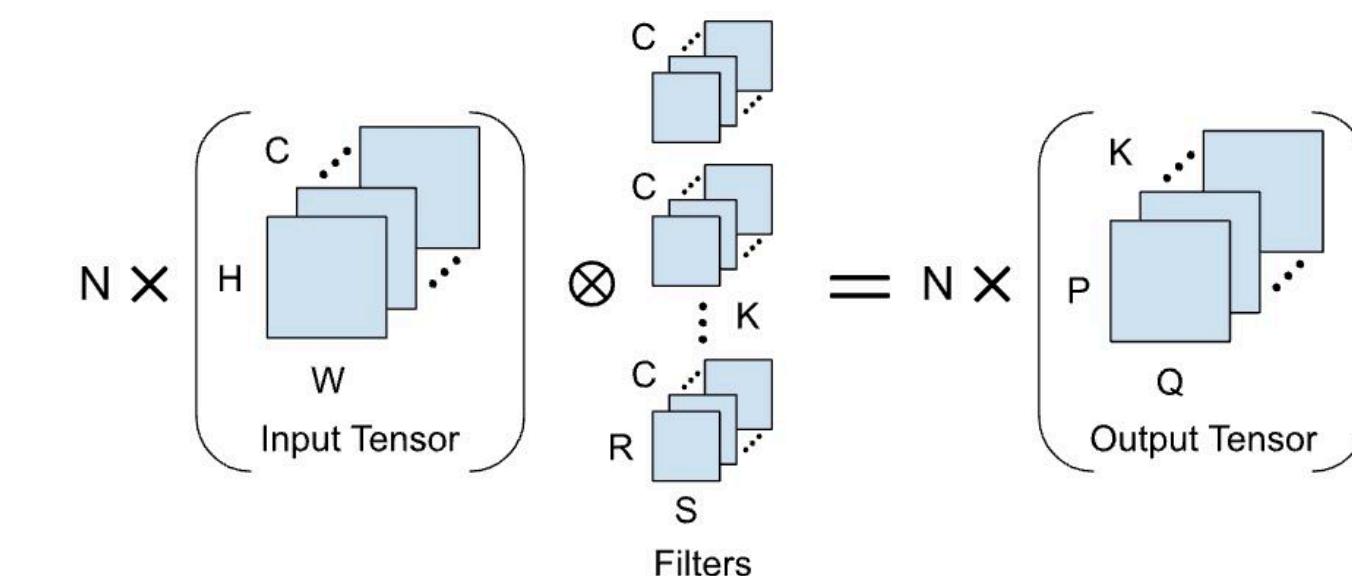
Higher performance with “more work”

$N=1, P=Q=64$ case:

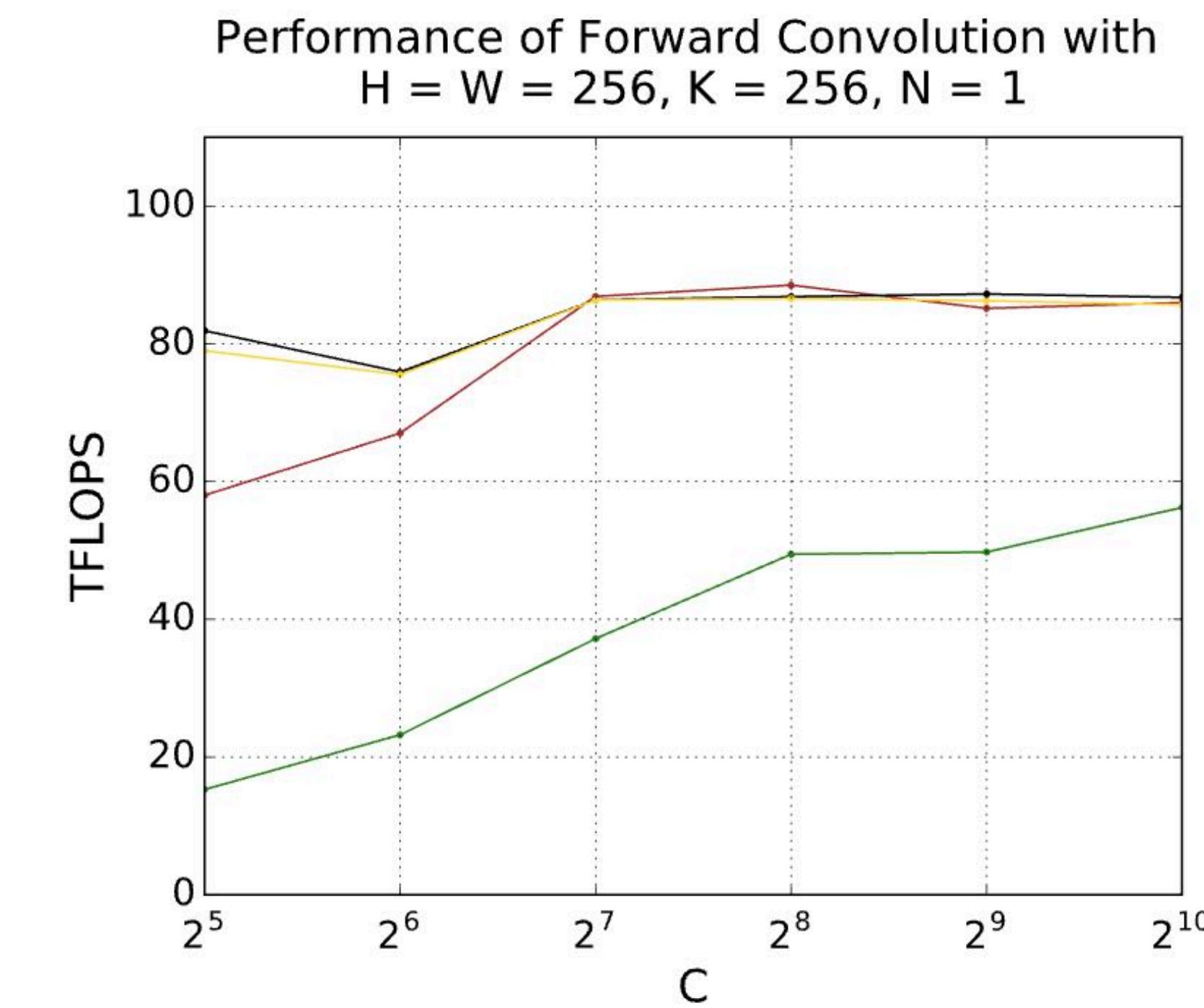
$64 \times 64 \times 128 \times 1 = 524K$ outputs = 2 MB of output data (float32)

$N=32, P=Q=256$ case:

$256 \times 256 \times 128 \times 32 = 256M$ outputs = 1 GB of output data (float32)



Legend:
● $P = Q = 64$
■ $P = Q = 128$
▲ $P = Q = 256$



Legend:
● $R = S = 1$
■ $R = S = 3$
▲ $R = S = 5$
◆ $R = S = 7$

Direct implementation

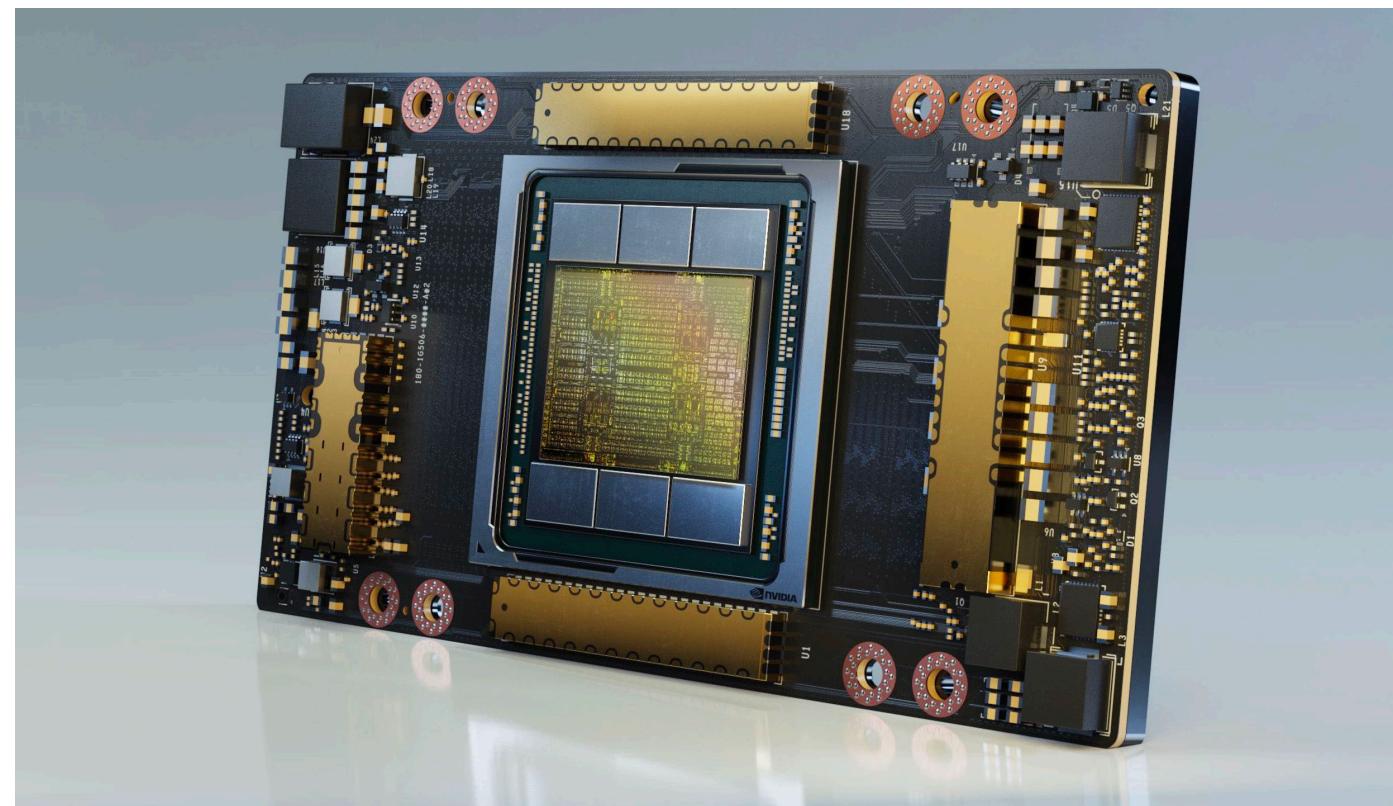
```
float input[IMAGE_BATCH_SIZE][INPUT_HEIGHT][INPUT_WIDTH][INPUT_DEPTH];           // input activations
float output[IMAGE_BATCH_SIZE][INPUT_HEIGHT][INPUT_WIDTH][LAYER_NUM_FILTERS];    // output activations
float layer_weights[LAYER_NUM_FILTERS][LAYER_CONVY][LAYER_CONVX][INPUT_DEPTH];
float layer_biases[LAYER_NUM_FILTERS];

// assumes convolution stride is 1
for (int img=0; img<IMAGE_BATCH_SIZE; img++)                                // for all images in batch
    for (int j=0; j<INPUT_HEIGHT; j++)
        for (int i=0; i<INPUT_WIDTH; i++)
            for (int f=0; f<LAYER_NUM_FILTERS; f++) {                      // for all output channels
                float tmp = layer_biases[LAYER_NUM_FILTERS];
                for (int kk=0; kk<INPUT_DEPTH; kk++)                         // combine filter responses from all input channels
                    for (int jj=0; jj<LAYER_FILTER_Y; jj++)                  // spatial convolution (Y)
                        for (int ii=0; ii<LAYER_FILTER_X; ii++)          // spatial convolution (X)
                            tmp += layer_weights[f][jj][ii][kk] * input[img][j+jj][i+ii][kk];
                output[img][j][i][f] = tmp;
            }
        }
```

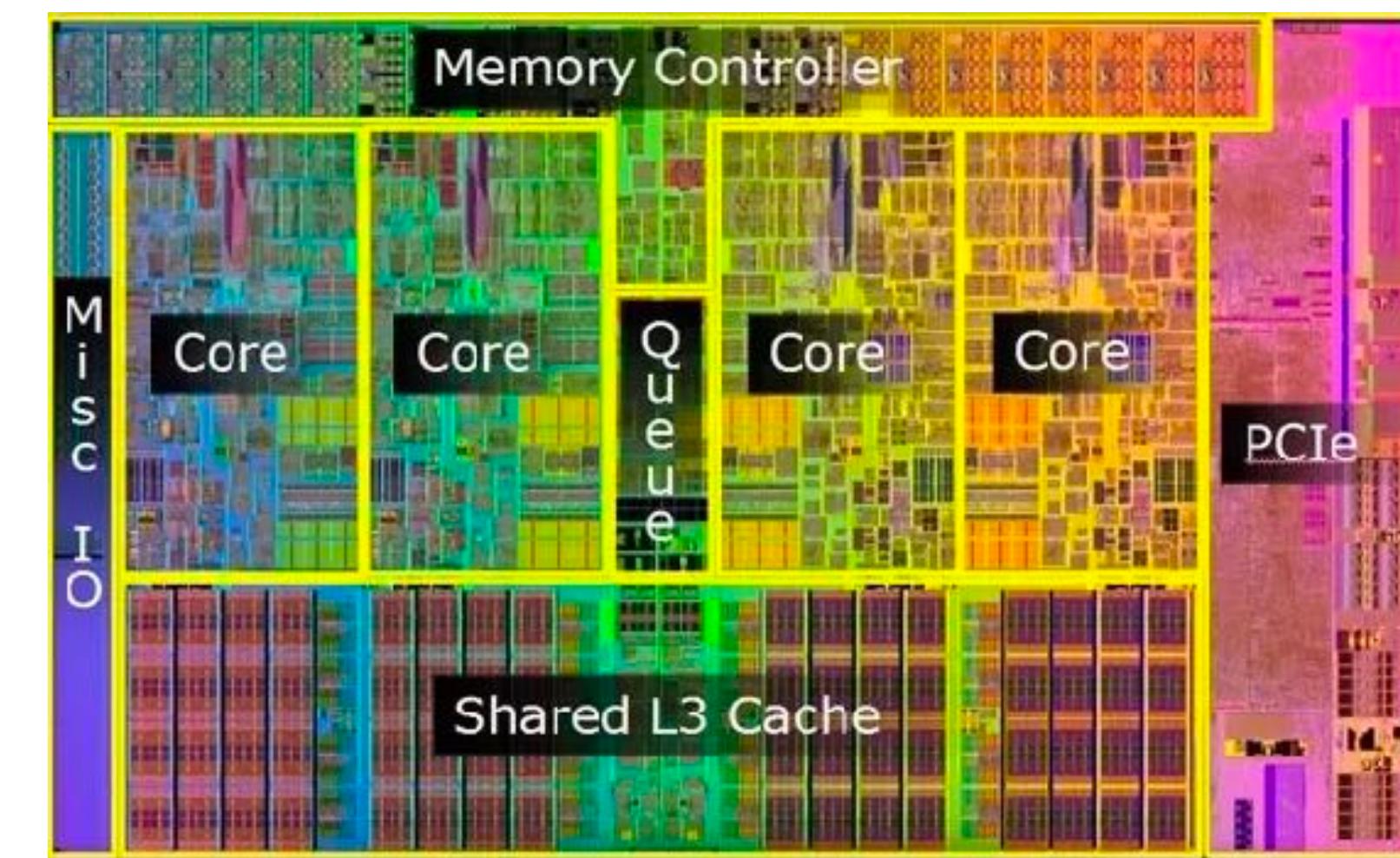
Or you can just directly implement this loop nest directly yourself.

Low-level chip libraries offer high-performance implementations of key DNN layers

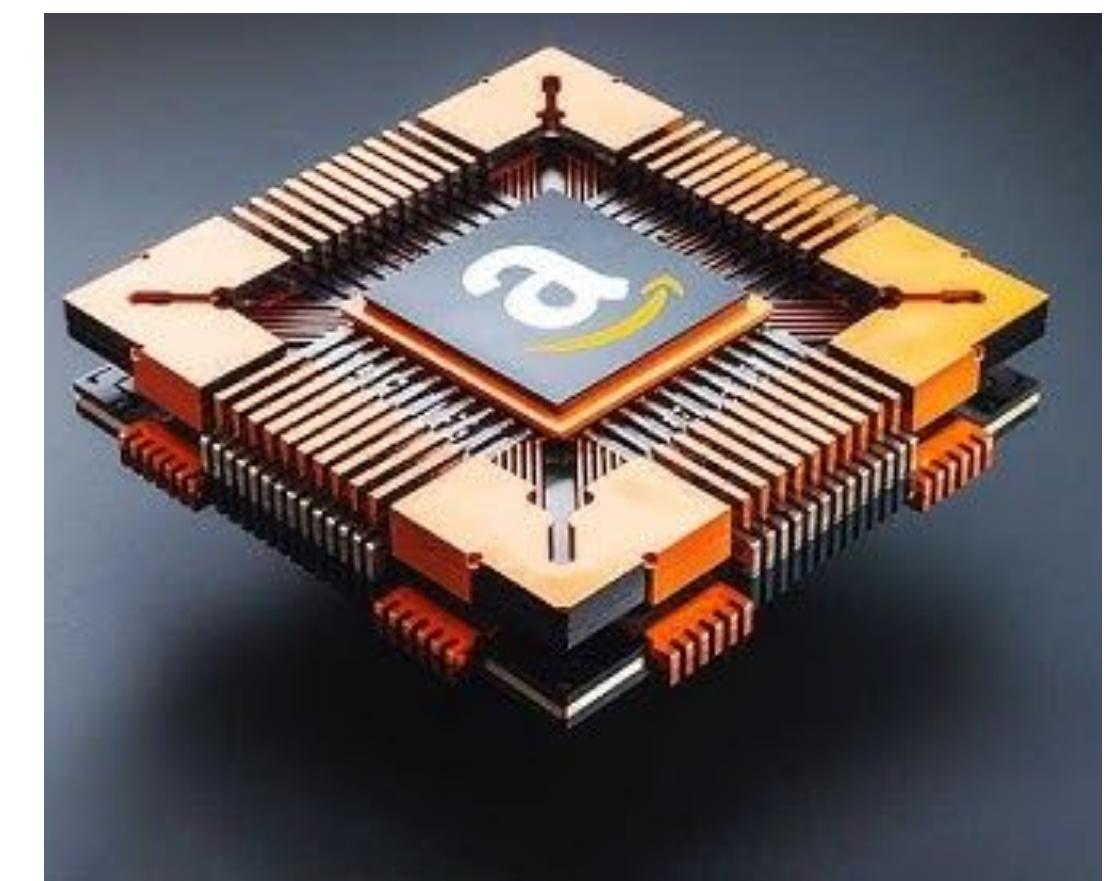
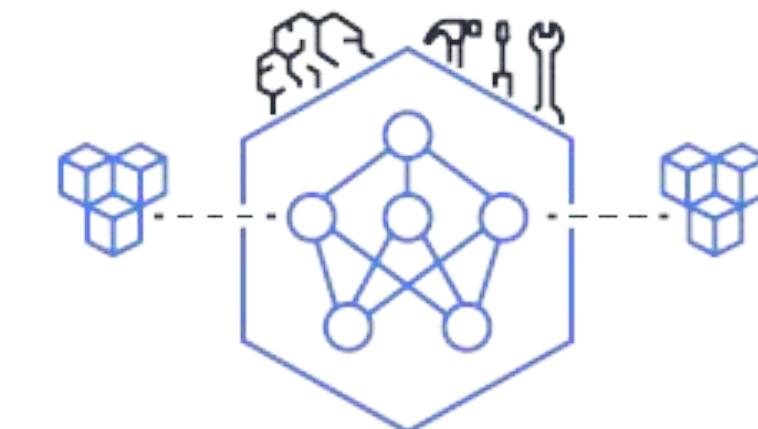
NVIDIA cuDNN



Intel® oneAPI Deep Neural Network Library



AWS NKI



Libraries offering high-performance implementations of key DNN layers



Convolution Layers

`nn.Conv1d`

Applies a 1D convolution over an input signal composed of several input planes.

`nn.Conv2d`

Applies a 2D convolution over an input signal composed of several input planes.

`nn.Conv3d`

Applies a 3D convolution over an input signal composed of several input planes.

`nn.ConvTranspose1d`

Applies a 1D transposed convolution operator over an input image composed of several input planes.

`nn.ConvTranspose2d`

Applies a 2D transposed convolution operator over an input image composed of several input planes.

`nn.ConvTranspose3d`

Applies a 3D transposed convolution operator over an input image composed of several input planes.

`nn.LazyConv1d`

A `torch.nn.Conv1d` module with lazy initialization of the `in_channels` argument.

`nn.LazyConv2d`

A `torch.nn.Conv2d` module with lazy initialization of the `in_channels` argument.

`nn.LazyConv3d`

A `torch.nn.Conv3d` module with lazy initialization of the `in_channels` argument.

`nn.LazyConvTranspose1d`

A `torch.nn.ConvTranspose1d` module with lazy initialization of the `in_channels` argument.

`nn.LazyConvTranspose2d`

A `torch.nn.ConvTranspose2d` module with lazy initialization of the `in_channels` argument.

Normalization Layers

`nn.BatchNorm1d`

Applies Batch Normalization over a 2D or 3D input.

`nn.BatchNorm2d`

Applies Batch Normalization over a 4D input.

`nn.BatchNorm3d`

Applies Batch Normalization over a 5D input.

`nn.LazyBatchNorm1d`

A `torch.nn.BatchNorm1d` module with lazy initialization.

`nn.LazyBatchNorm2d`

A `torch.nn.BatchNorm2d` module with lazy initialization.

`nn.LazyBatchNorm3d`

A `torch.nn.BatchNorm3d` module with lazy initialization.

`nn.GroupNorm`

Applies Group Normalization over a mini-batch of inputs.

`nn.SyncBatchNorm`

Applies Batch Normalization over a N-Dimensional input.

`nn.InstanceNorm1d`

Applies Instance Normalization.

`nn.InstanceNorm2d`

Applies Instance Normalization.

`nn.InstanceNorm3d`

Applies Instance Normalization.

`nn.LazyInstanceNorm1d`

A `torch.nn.InstanceNorm1d` module with lazy initialization of the `num_features` argument.

`nn.LazyInstanceNorm2d`

A `torch.nn.InstanceNorm2d` module with lazy initialization of the `num_features` argument.

Transformer Layers

`nn.Transformer`

A

`nn.TransformerEncoder`

T

`nn.TransformerDecoder`

T

`nn.TransformerEncoderLayer`

T

`nn.TransformerDecoderLayer`

T

Libraries offering high-performance implementations of key DNN layers

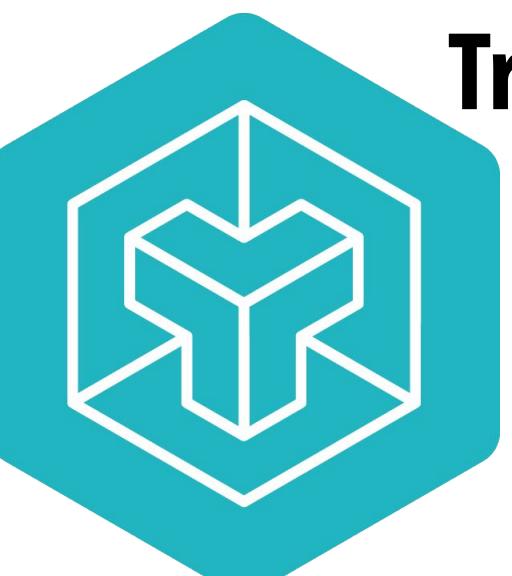
PyTorch

Convolution Layers

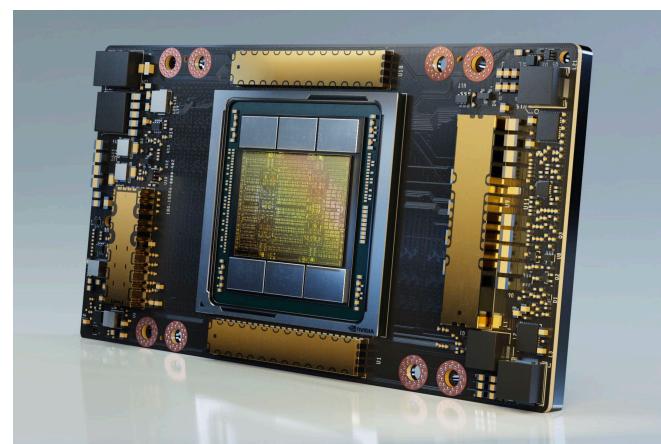
<code>nn.Conv1d</code>	Applies a 1D convolution over an input signal composed of several input planes.
<code>nn.Conv2d</code>	Applies a 2D convolution over an input signal composed of several input planes.
<code>nn.Conv3d</code>	Applies a 3D convolution over an input signal composed of several input planes.
<code>nn.ConvTranspose1d</code>	Applies a 1D transposed convolution operator over an input image composed of several input planes.
<code>nn.ConvTranspose2d</code>	Applies a 2D transposed convolution operator over an input image composed of several input planes.
<code>nn.ConvTranspose3d</code>	Applies a 3D transposed convolution operator over an input image composed of several input planes.
<code>nn.LazyConv1d</code>	A <code>torch.nn.Conv1d</code> module with lazy initialization of the <code>in_channels</code> argument.
<code>nn.LazyConv2d</code>	A <code>torch.nn.Conv2d</code> module with lazy initialization of the <code>in_channels</code> argument.
<code>nn.LazyConv3d</code>	A <code>torch.nn.Conv3d</code> module with lazy initialization of the <code>in_channels</code> argument.
<code>nn.LazyConvTranspose1d</code>	A <code>torch.nn.ConvTranspose1d</code> module with lazy initialization of the <code>in_channels</code> argument.
<code>nn.LazyConvTranspose2d</code>	A <code>torch.nn.ConvTranspose2d</code> module with lazy initialization of the <code>in_channels</code> argument.
<code>nn.LazyConvTranspose3d</code>	A <code>torch.nn.ConvTranspose3d</code> module with lazy initialization of the <code>in_channels</code> argument.



NVIDIA cuDNN



Triton



Example: CUDNN convolution

```
cudnnStatus_t cudnnConvolutionForward(
    cudnnHandle_t                      handle,
    const void*                         *alpha,
    const cudnnTensorDescriptor_t       xDesc,
    const void*                         *x,
    const cudnnFilterDescriptor_t       wDesc,
    const void*                         *w,
    const cudnnConvolutionDescriptor_t convDesc,
    cudnnConvolutionFwdAlgo_t          algo,
    void*                               *workSpace,
    size_t                             workSpaceSizeInBytes,
    const void*                         *beta,
    const cudnnTensorDescriptor_t       yDesc,
    void*                               *y)
```

Possible algorithms:

CUDNN_CONVOLUTION_FWD_ALGO_IMPLICIT_GEMM

This algorithm expresses the convolution as a matrix product without actually explicitly forming the matrix that holds the input tensor data.

CUDNN_CONVOLUTION_FWD_ALGO_IMPLICIT_PRECOMP_GEMM

This algorithm expresses convolution as a matrix product without actually explicitly forming the matrix that holds the input tensor data, but still needs some memory workspace to precompute some indices in order to facilitate the implicit construction of the matrix that holds the input tensor data.

CUDNN_CONVOLUTION_FWD_ALGO_GEMM

This algorithm expresses the convolution as an explicit matrix product. A significant memory workspace is needed to store the matrix that holds the input tensor data.

CUDNN_CONVOLUTION_FWD_ALGO_DIRECT

This algorithm expresses the convolution as a direct convolution (for example, without implicitly or explicitly doing a matrix multiplication).

CUDNN_CONVOLUTION_FWD_ALGO_FFT

This algorithm uses the Fast-Fourier Transform approach to compute the convolution. A significant memory workspace is needed to store intermediate results.

CUDNN_CONVOLUTION_FWD_ALGO_FFT_TILING

This algorithm uses the Fast-Fourier Transform approach but splits the inputs into tiles. A significant memory workspace is needed to store intermediate results but less than CUDNN_CONVOLUTION_FWD_ALGO_FFT for large size images.

CUDNN_CONVOLUTION_FWD_ALGO_WINOGRAD

This algorithm uses the Winograd Transform approach to compute the convolution. A reasonably sized workspace is needed to store intermediate results.

CUDNN_CONVOLUTION_FWD_ALGO_WINOGRAD_NONFUSED

This algorithm uses the Winograd Transform approach to compute the convolution. A significant workspace may be needed to store intermediate results.

Recall the loop fusion transformation: fuse multiple loops into one to increase a program's arithmetic intensity

Program 1

```
void add(int n, float* A, float* B, float* C) {
    for (int i=0; i<n; i++)
        C[i] = A[i] + B[i];
}

void mul(int n, float* A, float* B, float* C) {
    for (int i=0; i<n; i++)
        C[i] = A[i] * B[i];
}

float* A, *B, *C, *D, *E, *tmp1, *tmp2;

// assume arrays are allocated here

// compute E = D + ((A + B) * C)
add(n, A, B, tmp1);
mul(n, tmp1, C, tmp2);
add(n, tmp2, D, E);
```

Two loads, one store per math op
(arithmetic intensity = 1/3)

Two loads, one store per math op
(arithmetic intensity = 1/3)

Overall arithmetic intensity = 1/3

Program 2

```
void fused(int n, float* A, float* B, float* C, float* D, float* E) {
    for (int i=0; i<n; i++)
        E[i] = D[i] + (A[i] + B[i]) * C[i];
}

// compute E = D + (A + B) * C
fused(n, A, B, C, D, E);
```

Four loads, one store per 3 math ops
(arithmetic intensity = 3/5)

Memory traffic between operations

- Consider this sequence:



- Imagine the bandwidth cost of dumping 1 GB of conv outputs to memory, and then reading it back to just scale all the values, and then rereading to perform the pool! 😢

- Better solution:

- Per-element [scale+bias] operation can easily be performed per-element right after each element is computed by conv!
- And max pool's output can be computed once every 2x2 region of output is computed.



Fusing scale/bias with conv layer

```
float input[IMAGE_BATCH_SIZE][INPUT_HEIGHT][INPUT_WIDTH][INPUT_DEPTH];
float output[IMAGE_BATCH_SIZE][INPUT_HEIGHT][INPUT_WIDTH][LAYER_NUM_FILTERS];
float layer_weights[LAYER_NUM_FILTERS][LAYER_CONVY][LAYER_CONVX][INPUT_DEPTH];

// assumes convolution stride is 1
for (int img=0; img<IMAGE_BATCH_SIZE; img++) // for all images in batch
    for (int j=0; j<INPUT_HEIGHT; j++)
        for (int i=0; i<INPUT_WIDTH; i++) // for all output pixels
            for (int f=0; f<LAYER_NUM_FILTERS; f++) { // for all output channels
                float tmp = 0.0f;
                for (int kk=0; kk<INPUT_DEPTH; kk++) // filter combines responses from all input channels
                    for (int jj=0; jj<LAYER_FILTER_Y; jj++) // spatial convolution (Y)
                        for (int ii=0; ii<LAYER_FILTER_X; ii++) // spatial convolution (X)
                            tmp += layer_weights[f][jj][ii][kk] * input[img][j+jj][i+ii][kk];
                output[img][j][i][f] = tmp*scale + bias;
            }
        }
```

Exercise to class:

How would you “fuse” a max pool operation following this layer (max of 2x2 blocks of output matrix)?

Hint: how would you “block” the yellow loops?

Another example: softmax on rows of a matrix

softmax(S) is computing softmax over the rows of S

For a row x :

$$\text{softmax}(x) = \frac{f(x)}{l(x)}$$

Where:

$$f(x) = [e^{x_1 - m(x)} \quad e^{x_1 - m(x)} \quad \dots \quad e^{x_B - m(x)}]$$

$$m(x) = \max_i(x_i)$$

$$l(x) = \sum_i f(x)_i = \sum_i e^{x_i - m(x)}$$

Naive code:

```
def naive_softmax(x):
    """Compute row-wise softmax of X using native pytorch

    We subtract the maximum element in order to avoid overflows.
    this shift.
    """
    # read MN elements ; write M elements
    x_max = x.max(dim=1)[0]
    # read MN + M elements ; write MN elements
    z = x - x_max[:, None]
    # read MN elements ; write MN elements
    numerator = torch.exp(z)
    # read MN elements ; write M elements
    denominator = numerator.sum(dim=1)
    # read MN + M elements ; write MN elements
    ret = numerator / denominator[:, None]
    # in total: read 5MN + 2M elements ; wrote 3MN + 2M elements
    return ret
```

The problem is that an entire $M \times N$ “matrix” is read/written from memory each step. So the problem has low arithmetic intensity.

Another example: softmax on rows of a matrix

Naive code:

```
def naive_softmax(x):
    """Compute row-wise softmax of X using native pytorch

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    numerator = torch.exp(z)
    # read MN elements ; write M elements
    denominator = numerator.sum(dim=1)
    # read MN + M elements ; write MN elements
    ret = numerator / denominator[:, None]
    # in total: read 5MN + 2M elements ; wrote 3MN + 2M elements
    return ret
```

The problem is that an entire $M \times N$ “matrix” is read/written from memory each step. So the problem has low arithmetic intensity.

Reads $5MN + 2M$ elements, writes $3MN + 2M$ elements

“Fused” implementation:

```
@triton.jit
def softmax_kernel(output_ptr, input_ptr, input_row_stride, output_row_stride, n_rows, n_cols,
                   num_stages: tl.constexpr):
    # starting row of the program
    row_start = tl.program_id(0)
    row_step = tl.num_programs(0)
    for row_idx in tl.range(row_start, n_rows, row_step, num_stages=num_stages):
        # The stride represents how much we need to increase the pointer to advance 1 row
        row_start_ptr = input_ptr + row_idx * input_row_stride
        # The block size is the next power of two greater than n_cols, so we can fit each
        # row in a single block
        col_offsets = tl.arange(0, BLOCK_SIZE)
        input_ptrs = row_start_ptr + col_offsets
        # Load the row into SRAM, using a mask since BLOCK_SIZE may be > than n_cols
        mask = col_offsets < n_cols
        row = tl.load(input_ptrs, mask=mask, other=-float('inf'))
        # Subtract maximum for numerical stability
        row_minus_max = row - tl.max(row, axis=0)
        # Note that exponentiation in Triton is fast but approximate (i.e., think __expf in CUDA)
        numerator = tl.exp(row_minus_max)
        denominator = tl.sum(numerator, axis=0)
        softmax_output = numerator / denominator
        # Write back output to DRAM
        output_row_start_ptr = output_ptr + row_idx * output_row_stride
        output_ptrs = output_row_start_ptr + col_offsets
        tl.store(output_ptrs, softmax_output, mask=mask)
```

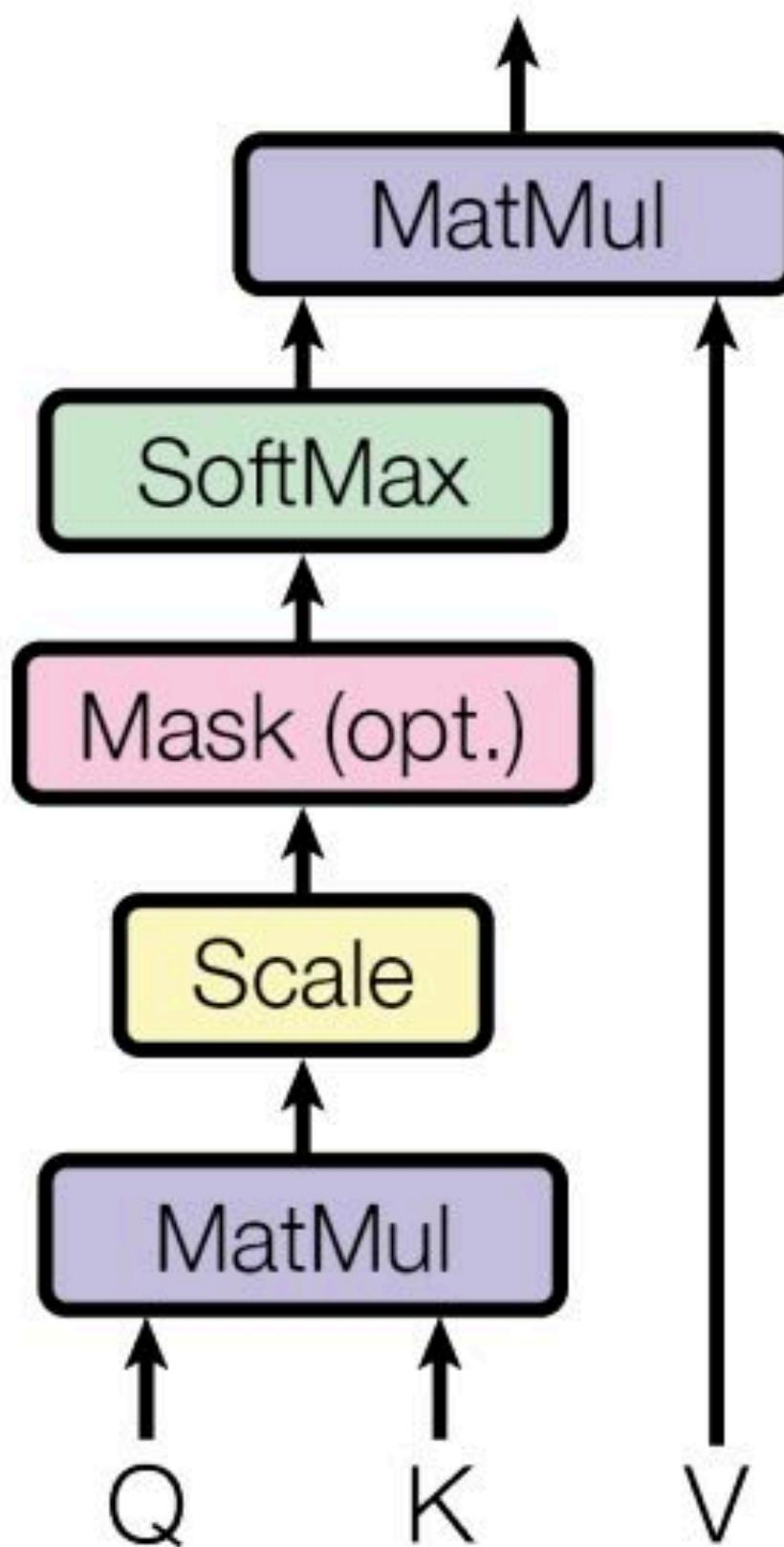
For each row:

Load row → compute entire softmax for a row → store row

Reads MN elements, writes MN elements,
assuming that working set for a single row fits in on-chip storage

**A good idea:
fusion trick for computing “attention”
in a modern transformer**

Attention module in a modern transformer



Let N be the length of the input sequence

Let d be the size of a feature embedding

Where

Let Q be a $N \times d$ matrix

Let K be a $N \times d$ matrix

Let V be a $N \times d$ matrix

Let $S = QK^\top \in \mathbb{R}^{N \times N}$

Let $P = \text{softmax}(S) \in \mathbb{R}^{N \times N}$

$\text{softmax}(S)$ is computing softmax over the rows of S

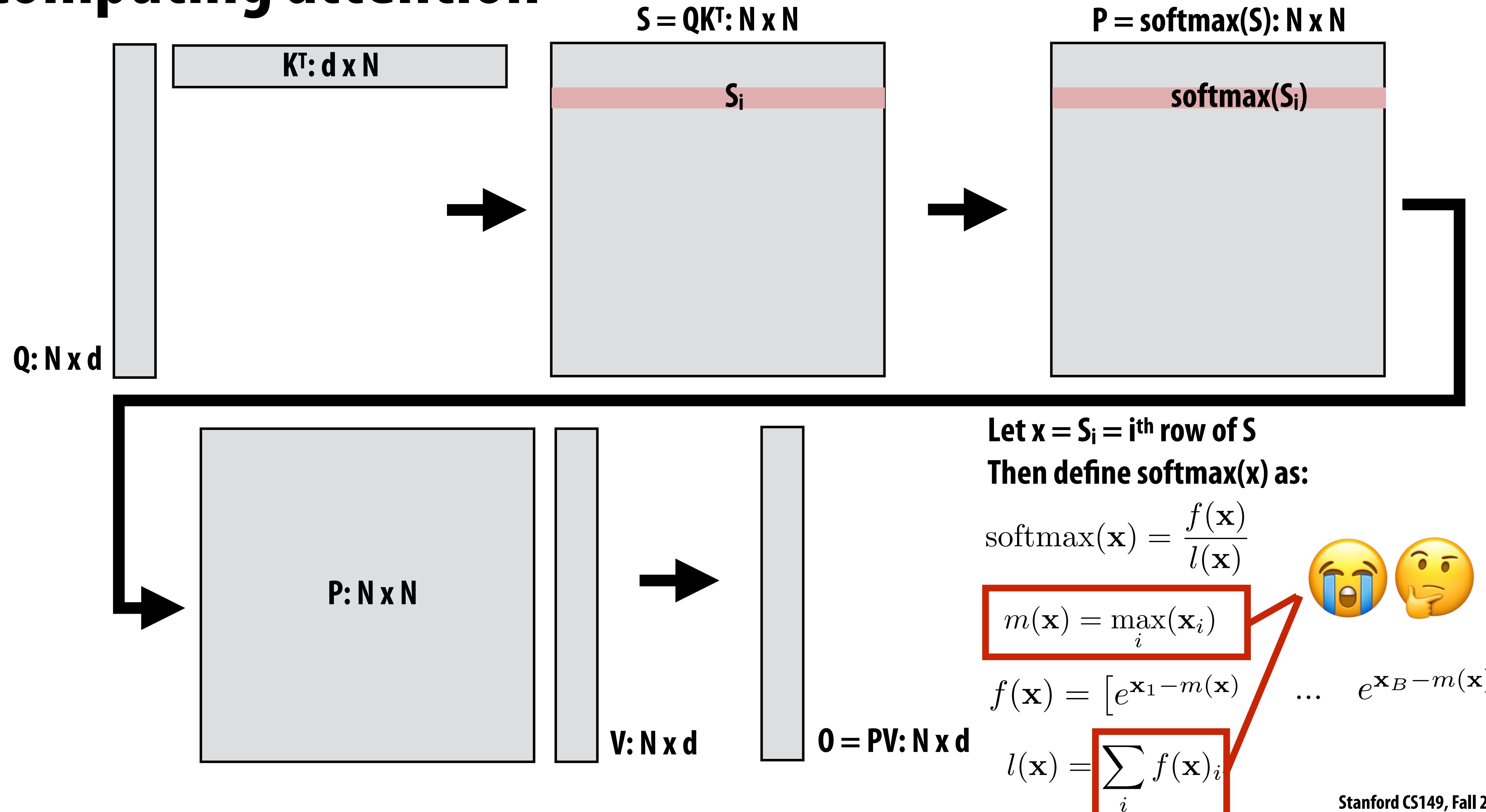
Let $O = PV \in \mathbb{R}^{N \times d}$

Notes:

N can be long for long sequences (e.g., thousands)

Naive implementation uses N^2 space! Trouble!!!

Computing attention



Let's look into softmax more closely...

$$\text{softmax}(\mathbf{x}) = \frac{f(\mathbf{x})}{l(\mathbf{x})} \quad \text{Where: } f(\mathbf{x}) = [e^{\mathbf{x}_1 - m(\mathbf{x})} \quad e^{\mathbf{x}_2 - m(\mathbf{x})} \quad \dots \quad e^{\mathbf{x}_B - m(\mathbf{x})}]$$

$$m(\mathbf{x}) = \max_i(\mathbf{x}_i)$$

$$l(\mathbf{x}) = \sum_i f(\mathbf{x})_i = \sum_i e^{\mathbf{x}_i - m(\mathbf{x})}$$

Let's break vector \mathbf{x} into chunks:

$$\mathbf{x} = [\mathbf{x}^{(1)} \quad \mathbf{x}^{(2)}]$$

$$\mathbf{x} = \boxed{\mathbf{x}^{(1)}} \quad \boxed{\mathbf{x}^{(2)}}$$

Now:

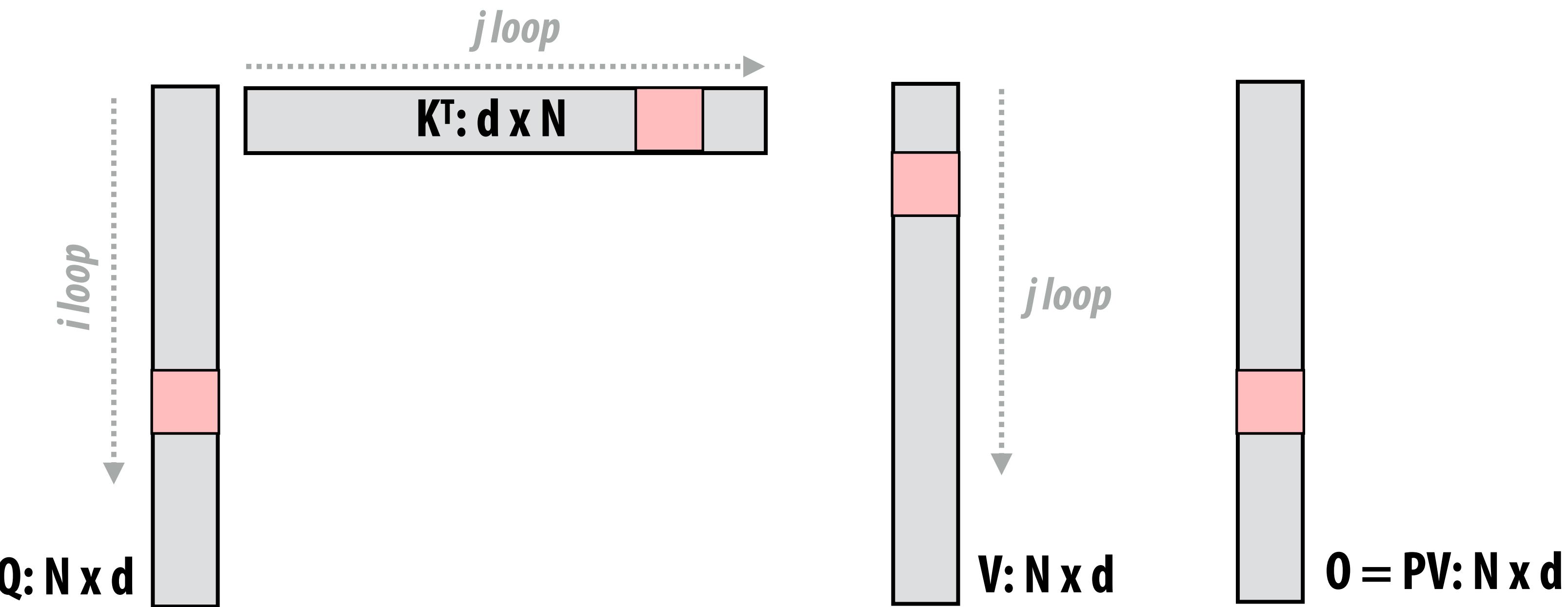
$$m(\mathbf{x}) = \max(m(\mathbf{x}^{(1)}), m(\mathbf{x}^{(2)}))$$

$$f(\mathbf{x}) = [e^{m(\mathbf{x}^{(1)}) - m(\mathbf{x})} f(\mathbf{x}^{(1)}) \quad e^{m(\mathbf{x}^{(2)}) - m(\mathbf{x})} f(\mathbf{x}^{(2)})]$$

$$l(\mathbf{x}) = e^{m(\mathbf{x}^{(1)}) - m(\mathbf{x})} l(\mathbf{x}^{(1)}) + e^{m(\mathbf{x}^{(2)}) - m(\mathbf{x})} l(\mathbf{x}^{(2)})$$

So softmax can be computed in chunks!

Fused attention



for each j :

for each i :

Load block $Q_i, K^T_j, V_j, 0_i$

Compute $S_{ij} = Q_i K^T_j$

Compute $M_{ij} = m(S_{ij})$, $P_{ij} = f(S_{ij})$, and $l_{ij} = l(S_{ij})$ (all functions operate row-wise on row-vectors)

Multiply $P_{ij}V_j$ and accumulate into 0_i with appropriate scalings (see previous slide for math)

Save memory footprint:
Never materialize N^2 matrix

Save memory bandwidth:
(high arithmetic intensity)

- Read 3 blocks (from Q, K, V)
- Do two matrix multiplies + a few row summations
- Accumulate into 0 block (which is resident in cache)

Note there is additional computation vs. the original version (must re-scale prior values of 0 each step of i-loop)

“Flash-Attention” in Thunderkittens

```
#include "kittens.cuh"

using namespace kittens;

constexpr int NUM_WORKERS = 4; // This kernel uses 4 worker warps per block, and 2 blocks per
template<int D> constexpr size_t ROWS = 16*(128/D); // height of each worker tile (rows)
template<int D, typename T=bf16, typename L=row_l> using qkvo_tile = rt<T, ROWS<D>, D, L>;
template<int D, typename T=float> using attn_tile = rt<T, ROWS<D>, ROWS<D>>;
template<int D> using shared_tile = st_bf<ROWS<D>, D>;
template<int D> using global_layout = gl<bf16, -1, -1, -1, D>; // B, H, g.Qg.rows specified at
template<int D> struct globals { global_layout<D> Qg, Kg, Vg, Og; };

template<int D> __launch_bounds__(NUM_WORKERS*WARP_THREADS, 1)
__global__ void attend_ker(const __grid_constant__ globals<D> g) {
    using load_group = kittens::group<2>; // pairs of workers collaboratively load k, v tiles
    int loadid = load_group::groupid(), workerid = kittens::warpid(); // which worker am I?
    constexpr int LOAD_BLOCKS = NUM_WORKERS / load_group::GROUP_WARPS;
    const int batch = blockIdx.z, head = blockIdx.y, q_seq = blockIdx.x * NUM_WORKERS + workerid;

    extern __shared__ alignment_dummy __shm[]; // this is the CUDA shared memory
    shared_allocator al((int*)&__shm[0]);
    // K and V live in shared memory. Here, we instantiate three tiles for a 3-stage pipeline.
    shared_tile<D> (&k_smem)[LOAD_BLOCKS][3] = al.allocate<shared_tile<D>, LOAD_BLOCKS, 3>();
    shared_tile<D> (&v_smem)[LOAD_BLOCKS][3] = al.allocate<shared_tile<D>, LOAD_BLOCKS, 3>();
    // We also reuse this memory to improve coalescing of DRAM reads and writes.
    shared_tile<D> (&qo_smem)[NUM_WORKERS] = reinterpret_cast<shared_tile<D>(&)[NUM_WORKERS]>(
        // Initialize all of the register tiles.
        qkvo_tile<D, bf16> q_reg, k_reg; // Q and K are both row layout, as we use mma_ABt.
        qkvo_tile<D, bf16, col_l> v_reg; // V is column layout, as we use mma_AB.
        qkvo_tile<D, float> o_reg; // Output tile.
        attn_tile<D, float> att_block; // attention tile, in float. (We want to use float wherever
        attn_tile<D, bf16> att_block_mma; // bf16 attention tile for the second mma_AB. We cast it
        typename attn_tile<D, float>::col_vec max_vec_last, max_vec, norm_vec; // these are column
        // each warp loads its own Q tile of 16x64
        if (q_seq*ROWS<D> < g.Qg.rows) {
            load(qo_smem[workerid], g.Qg, {batch, head, q_seq, 0}); // going through shared memory
            __syncwarp();
            load(q_reg, qo_smem[workerid]);
        }
    __syncthreads();
}
```

```
// temperature adjustment. Pre-multiplying by lg2(e), too, so we can use exp2 later.
if constexpr(D == 64) mul(q_reg, q_reg, __float2bfloat16(0.125f * 1.44269504089));
else if constexpr(D == 128) mul(q_reg, q_reg, __float2bfloat16(0.08838834764f * 1.44269504
// initialize flash attention L, M, and O registers.
neg_infty(max_vec); // zero registers for the Q chunk
zero(norm_vec);
zero(o_reg);
// launch the load of the first k, v tiles
int kv_blocks = g.Qg.rows / (LOAD_BLOCKS*ROWS<D>), tic = 0;
load_group::load_async(k_smem[loadid][0], g.Kg, {batch, head, loadid, 0});
load_group::load_async(v_smem[loadid][0], g.Vg, {batch, head, loadid, 0});
// iterate over k, v for these q's that have been loaded
for(auto kv_idx = 0; kv_idx < kv_blocks; kv_idx++, tic=(tic+1)%3) {
    int next_load_idx = (kv_idx+1)*LOAD_BLOCKS + loadid;
    if(next_load_idx*ROWS<D> < g.Kg.rows) {
        int next_tic = (tic+1)%3;
        load_group::load_async(k_smem[loadid][next_tic], g.Kg, {batch, head, next_load_idx});
        load_group::load_async(v_smem[loadid][next_tic], g.Vg, {batch, head, next_load_idx});
        load_async_wait<2>(); // next k, v can stay in flight.
    }
    else load_async_wait(); // all must arrive
    __syncthreads(); // Everyone's memory must be ready for the next stage.
    // now each warp goes through all of the subtiles, loads them, and then does the flash
    #pragma unroll LOAD_BLOCKS
    for(int subtile = 0; subtile < LOAD_BLOCKS && (kv_idx*LOAD_BLOCKS + subtile) < g.Qg.ro
        load(k_reg, k_smem[subtile][tic]); // load k from shared into registers
        zero(att_block); // zero 16x16 attention tile
        mma_ABt(att_block, q_reg, k_reg, att_block); // Q@K.T
        copy(max_vec_last, max_vec);
        row_max(max_vec, att_block, max_vec); // accumulate onto the max_vec
        sub_row(att_block, att_block, max_vec); // subtract max from attention -- now all
        exp2(att_block, att_block); // exponentiate the block in-place.
        sub(max_vec_last, max_vec_last, max_vec); // subtract new max from old max to find
        exp2(max_vec_last, max_vec_last); // exponentiate this vector -- this is what we n
        mul(norm_vec, norm_vec, max_vec_last); // and the norm vec is now normalized.
        row_sum(norm_vec, att_block, norm_vec); // accumulate the new attention block onto
        copy(att_block_mma, att_block); // convert to bf16 for mma_AB
        load(v_reg, v_smem[subtile][tic]); // load v from shared into registers.
        mul_row(o_reg, o_reg, max_vec_last); // normalize o_reg in advance of mma_AB'ing o
        mma_AB(o_reg, att_block_mma, v_reg, o_reg); // mfma onto o_reg with the local atte
    }
}
div_row(o_reg, o_reg, norm_vec);
__syncthreads();
if (q_seq*ROWS<D> < g.Qg.rows) { // write out o.
    store(qo_smem[workerid], o_reg); // going through shared memory improves coalescing of
    __syncwarp();
    store(g.Og, qo_smem[workerid], {batch, head, q_seq, 0});
}
```

Fusion in modern DNN frameworks

Old style: library writers hardcoded a few “fused” ops

```
cudnnStatus_t cudnnConvolutionBiasActivationForward(
    cudnnHandle_t                      handle,
    const void*                         *alpha1,
    const cudnnTensorDescriptor_t       xDesc,
    const void*                         *x,
    const cudnnFilterDescriptor_t      wDesc,
    const void*                         *w,
    const cudnnConvolutionDescriptor_t convDesc,
    cudnnConvolutionFwdAlgo_t          algo,
    void*                               *workSpace,
    size_t                             workSpaceSizeInBytes,
    const void*                         *alpha2,
    const cudnnTensorDescriptor_t       zDesc,
    const void*                         *z,
    const cudnnTensorDescriptor_t       biasDesc,
    const void*                         *bias,
    const cudnnActivationDescriptor_t   activationDesc,
    const cudnnTensorDescriptor_t       yDesc,
    void*                               *y)
```

This function applies a bias and then an activation to the convolutions or cross-correlations of [cudnnConvolutionForward\(\)](#), returning results in `y`. The full computation follows the equation $y = \text{act}(\alpha_1 \cdot \text{conv}(x) + \alpha_2 \cdot z + \text{bias})$.

Tensorflow:

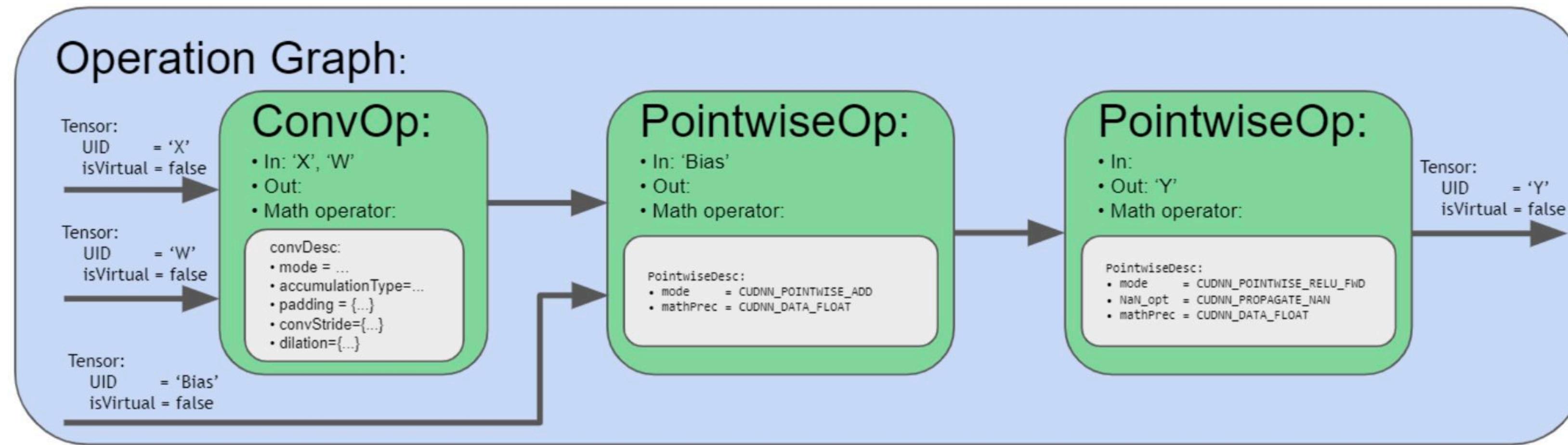
[tensorflow::ops::FusedBatchNorm](#)

Batch normalization.

[tensorflow::ops::FusedResizeAndPadConv2D](#)

Performs a resize and padding as a preprocess during a convolution.

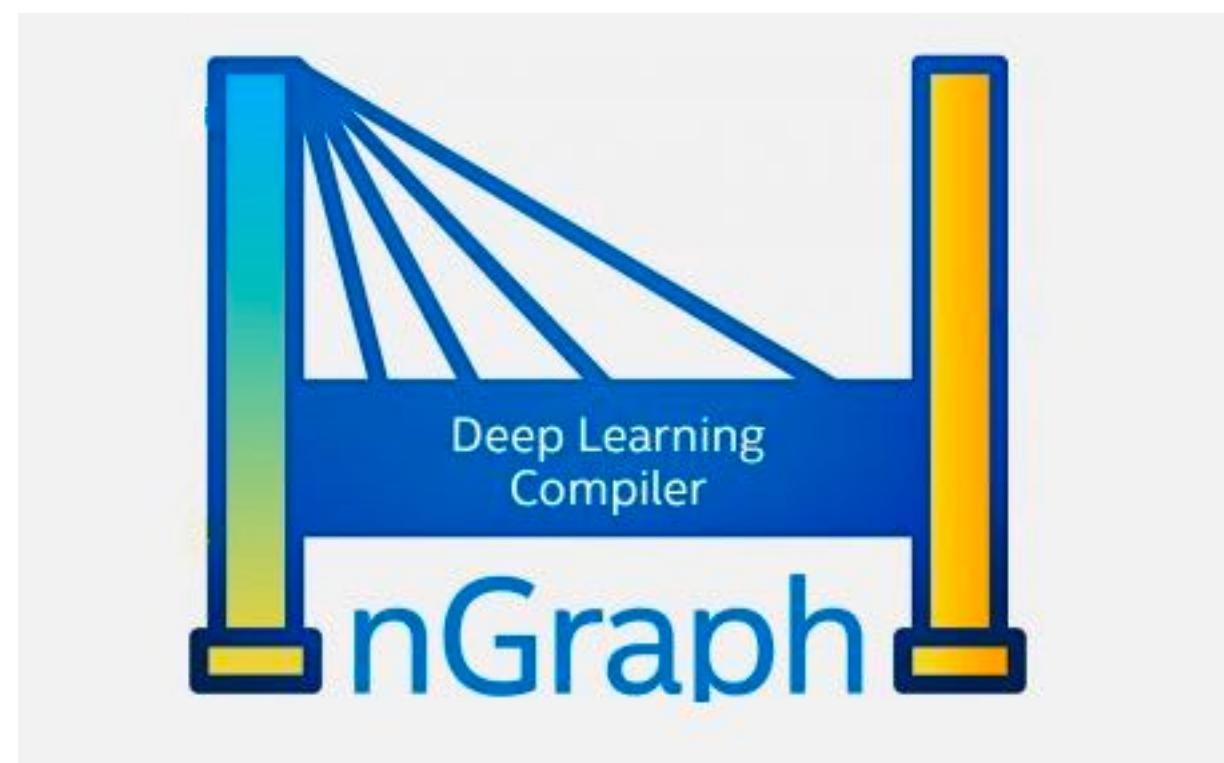
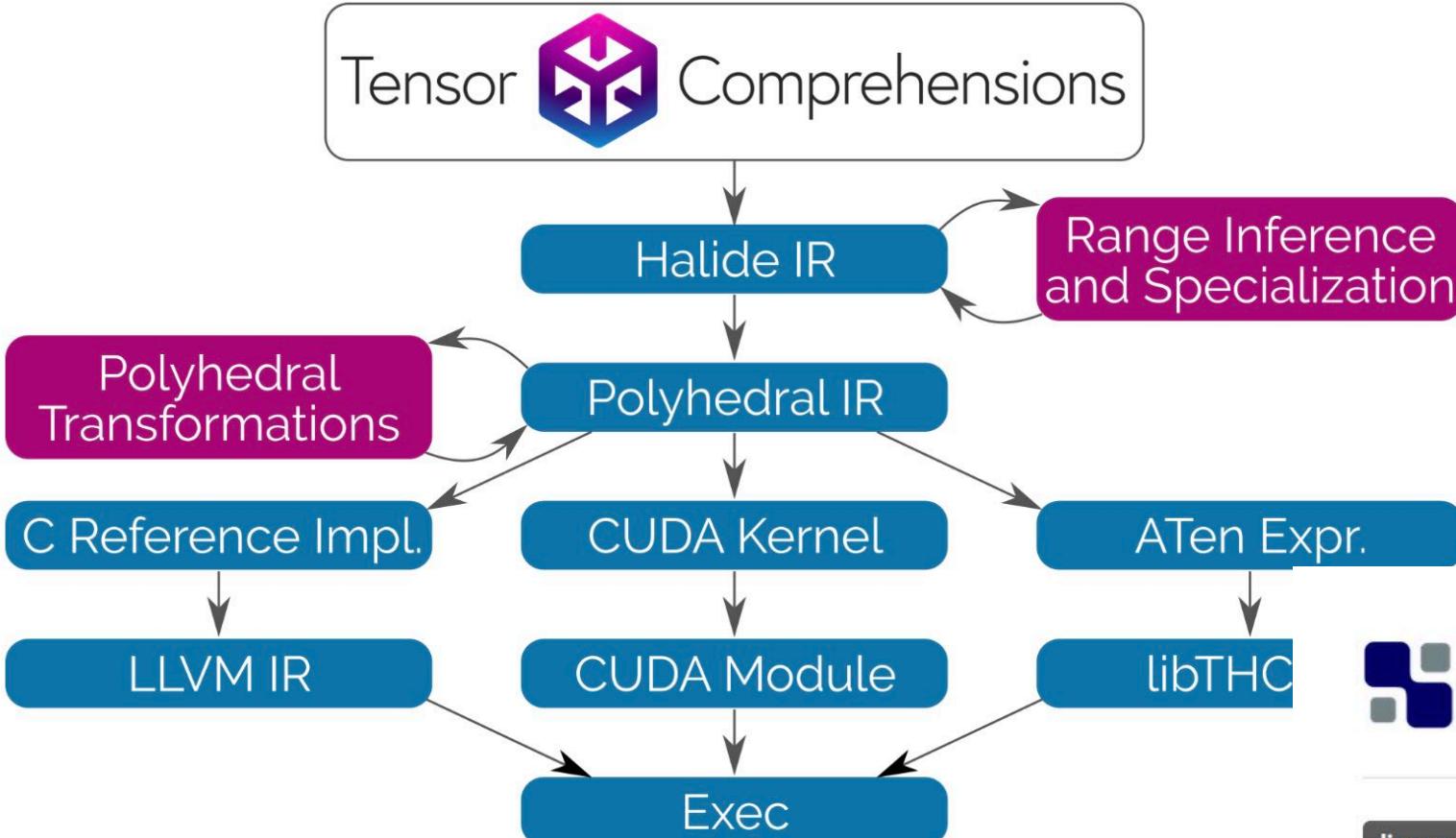
More flexible fusion example: CUDNN “backend”



Note for operation fusion use cases, there are two different mechanisms in cuDNN to support them. First, there are engines containing offline compiled kernels that can support certain fusion patterns. These engines try to match the user provided operation graph with their supported fusion pattern. If there is a match, then that particular engine is deemed suitable for this use case. In addition, there are also runtime fusion engines to be made available in the upcoming releases. Instead of passively matching the user graph, such engines actively walk the graph and assemble code blocks to form a CUDA kernel and compile on the fly. Such runtime fusion engines are much more flexible in its range of support. However, because the construction of the execution plans requires runtime compilation, the one-time CPU overhead is higher than the other engines.

Compiler generates new implementations that “fuse” multiple operations into a single node that executes efficiently (without runtime overhead or communicating intermediate results through memory)

Many compiler-based efforts to automatically schedule key DNN operations



NVIDIA TensorRT
Programmable Inference Accelerator

JAX: Autograd and XLA 

 tvm Open Deep Learning Compiler Stack

license Apache 2.0 build passing

[Documentation](#) | [Contributors](#) | [Community](#) | [Release Notes](#)

TVM is a compiler stack for deep learning systems. It is designed to close the gap between learning frameworks, and the performance- and efficiency-focused hardware backends. T' frameworks to provide end to end compilation to different backends. Checkout the [tvm sta](#) information.

Python

```
import torch

# Compile a function
@torch.compile
def my_function(x):
    return x * 2 + 1

# Compile a model
model = MyModel()
compiled_model = torch.compile(model)
```

 torch.compile



Introducing Triton: Open-source GPU programming for neural networks

Another trick: use of low precision values

- Many efforts to use low precision values for DNN weights and intermediate activations
- 16 bit and 8-bit values are common
- Now moving into 4 bit values
- In the extreme case: 1-bit ;-)

XNOR-Net: ImageNet Classification Using Binary Convolutional Neural Networks

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Abstract. We propose two efficient approximations to standard convolutional neural networks: Binary-Weight-Networks and XNOR-Networks. In Binary-Weight-Networks, the filters are approximated with binary values resulting in $32 \times$ memory saving. In XNOR-Networks, both the filters and the input to convolutional layers are binary. XNOR-Networks approximate convolutions using primarily binary operations. This results in $58 \times$ faster convolutional operations (in terms of number of the high precision operations) and $32 \times$ memory savings. XNOR-Nets offer the possibility of running state-of-the-art networks on CPUs (rather than GPUs) in real-time. Our binary networks are simple, accurate, efficient, and work on challenging visual tasks. We evaluate our approach on the ImageNet classification task. The classification accuracy with a Binary-Weight-Network version of AlexNet is the same as the full-precision AlexNet. We compare our method with recent network binarization methods, BinaryConnect and BinaryNets, and outperform these methods by large margins on ImageNet, more than 16% in top-1 accuracy. Our code is available at: <http://allenai.org/plato/xnornet>.

 NVIDIA. 2025-9-30

Pretraining Large Language Models with NVFP4

NVIDIA

Abstract. Large Language Models (LLMs) today are powerful problem solvers across many domains, and they continue to get stronger as they scale in model size, training set size, and training set quality, as shown by extensive research and experimentation across the industry. Training a frontier model today requires on the order of tens to hundreds of yottaflops, which is a massive investment of time, compute, and energy. Improving pretraining efficiency is therefore essential to enable the next generation of even more capable LLMs. While 8-bit floating point (FP8) training is now widely adopted, transitioning to even narrower precision, such as 4-bit floating point (FP4), could unlock additional improvements in computational speed and resource utilization. However, quantization at this level poses challenges to training stability, convergence, and implementation, notably for large-scale models trained on long token horizons.

In this study, we introduce a novel approach for stable and accurate training of large language models (LLMs) using the NVFP4 format. Our method integrates Random Hadamard transforms (RHT) to bound block-level outliers, employs a two-dimensional quantization scheme for consistent representations across both the forward and backward passes, utilizes stochastic rounding for unbiased gradient estimation, and incorporates selective high-precision layers. We validate our approach by training a 12-billion-parameter model on 10 trillion tokens – the longest publicly documented training run in 4-bit precision to date. Our results show that the model trained with our NVFP4-based pretraining technique achieves training loss and downstream task accuracies comparable to an FP8 baseline. For instance, the model attains an MMLU-pro accuracy of 62.58%, nearly matching the 62.62% accuracy achieved through FP8 pretraining. These findings highlight that NVFP4, when combined with our training approach, represents a major step forward in narrow-precision LLM training algorithms.

Code: Transformer Engine support for NVFP4 training.

Optimization techniques

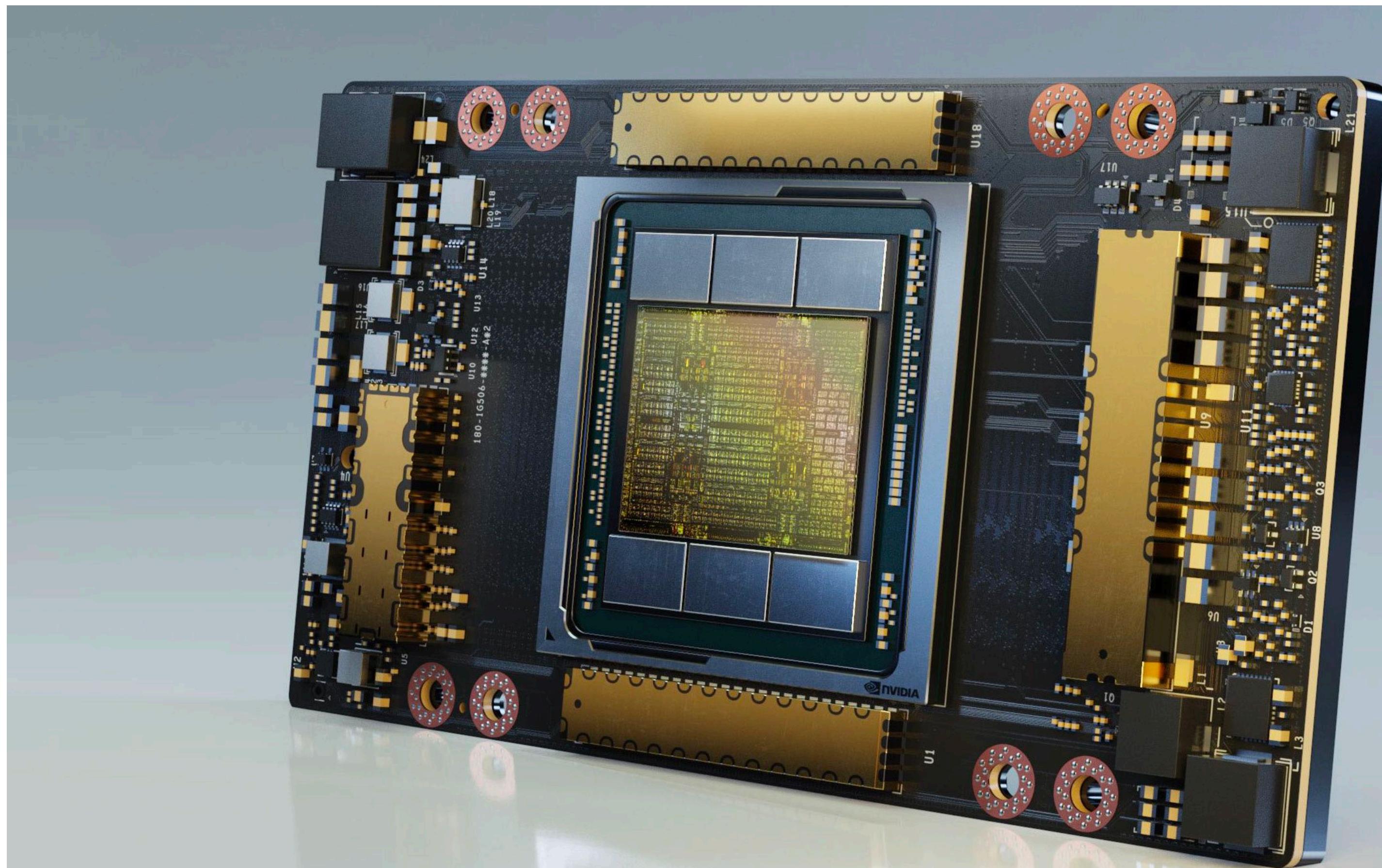
- **Better algorithms: manually designing better ML models**
 - **Common parameters: depth of network, width of filters, number of filters per layer, convolutional stride, etc.**
 - **Common to perform automatic search for efficient topologies**
- **Software optimization: Good scheduling of performance-critical operations**
 - **Loop blocking/tiling, fusion**
 - **Typically optimized manually by humans (but significant research efforts to automate scheduling)**
- **Forms of approximation: compressing models**
 - **Lower bit precision**

Why might a GPU be a good platform for DNN evaluation?

**consider: arithmetic intensity, SIMD, data-
parallelism, memory bandwidth requirements**

Deep neural networks on GPUs

- Many high-performance DNN implementations target GPUs
 - High arithmetic intensity matrix-matrix computations benefit from flop-rich GPU architectures
 - Highly-optimized library of kernels exist for GPUs (cuDNN)

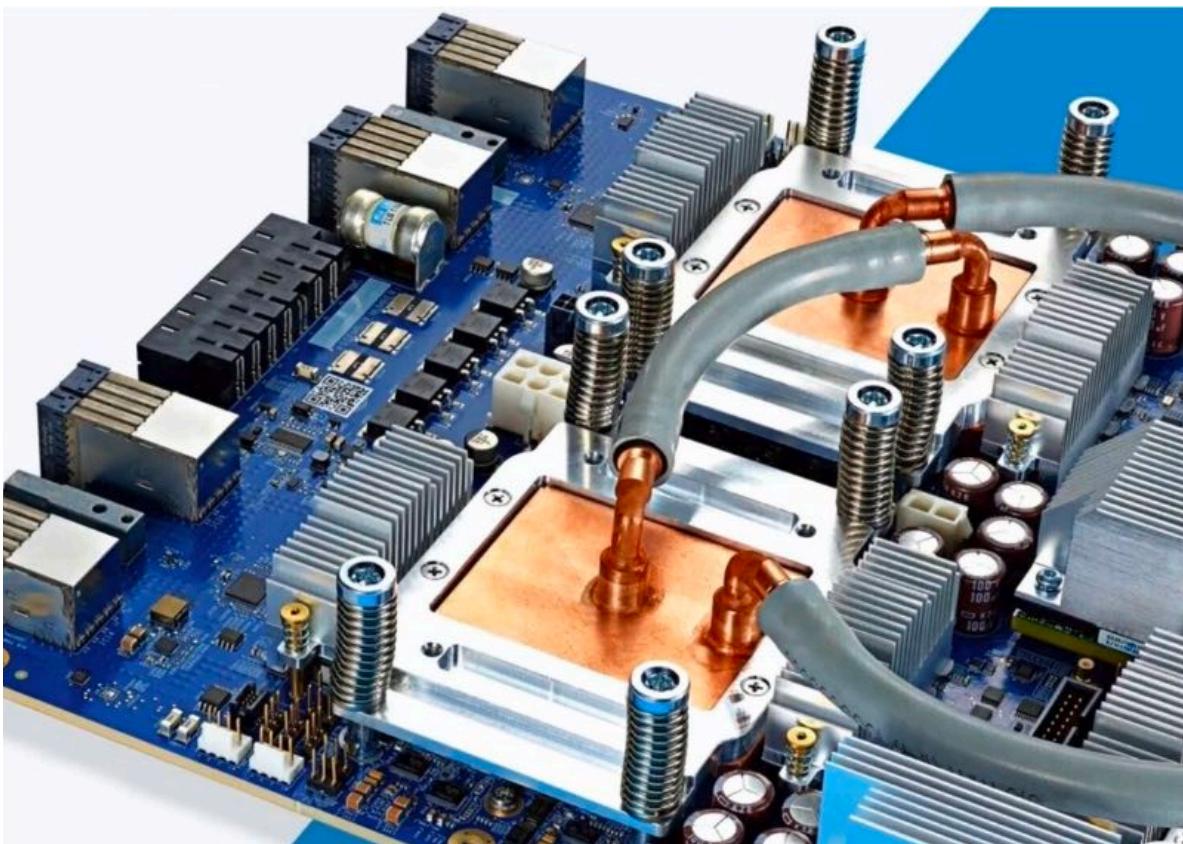


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Why might a GPU be a sub-optimal platform for DNN evaluation?

(Hint: is a general purpose processor really needed?)

Next time: maximizing efficiency via specialized hardware acceleration for DNN inference/training



Google TPU3



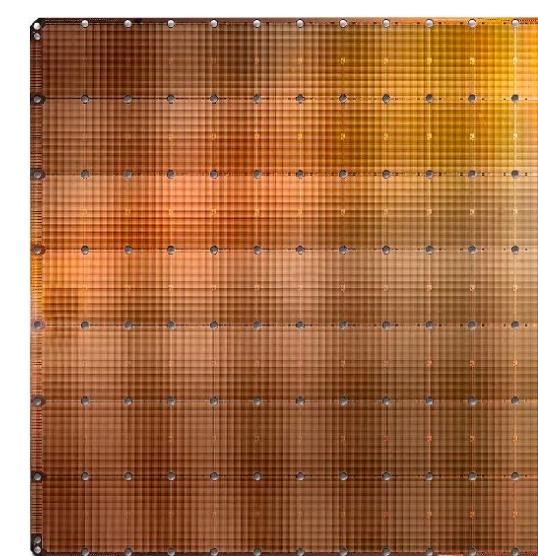
GraphCore IPU



Apple Neural Engine



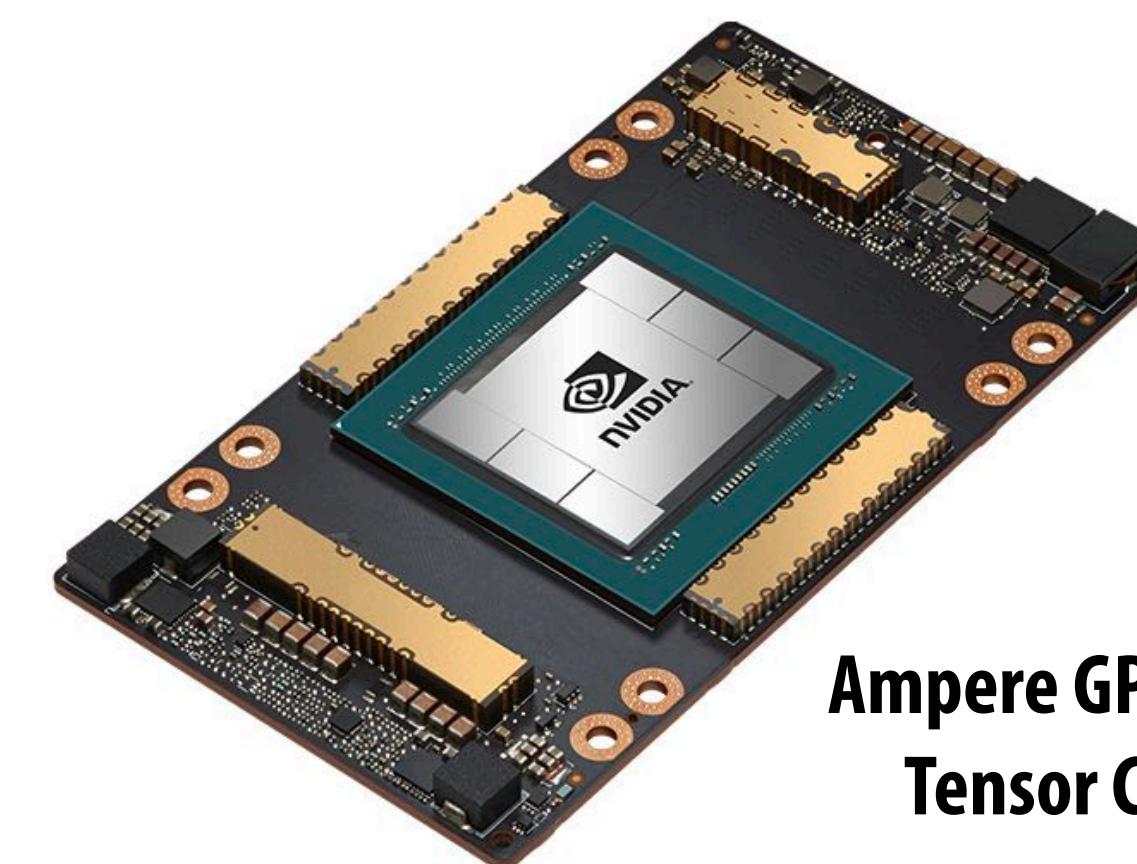
Intel Deep Learning
Inference Accelerator



Cerebras Wafer Scale Engine



SambaNova
Cardinal SN10



Ampere GPU with
Tensor Cores