

# The SMite Virtual Machine

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## Typographical conventions

Instructions and registers are shown in Typewriter font; interface calls are shown in **Bold** type.

Addresses are given in bytes and refer to the VM address space except where stated. Addresses are written in hexadecimal; hex numbers are prefixed with “0x”.

## 1 Introduction

SMite is a simple virtual machine for study and experiment. It is a stack machine, based on the more complex register machine [3]. This paper gives a full description of SMite.

SMite is conceptually (and usually in fact) a library, embedded in other programs; it supports a simple object module format.

## 2 Architecture

SMite’s address unit is the byte, which is eight bits. Words are `WORD_SIZE` bytes. The word is the size of the numbers and addresses on which SMite operates, and of the items placed on the stack. The size of the byte and range of word sizes allowed have been chosen with a view to making efficient implementation possible on the vast majority of current machine architectures.

Words have the bytes stored in big-endian or little-endian order, according to `ENDISM`.

### 2.1 Registers

The registers, each with its function, are set out in table 1.

The registers are word quantities.

To ease efficient implementation, the registers may only be accessed by instructions (see section 3.9); not all registers are accessible, and only a few are writable.

### 2.2 Memory

SMite’s memory is a contiguous sequence of bytes with addresses starting at 0. The address of a word is that of the byte in it with the lowest address.

Register	Function
PC	The Program Counter. Points to the next word from which I may be loaded.
I	The Instruction register. Contains instructions to be executed.
BAD	The invalid stack position, or invalid or unaligned memory address, that last caused an error.
STACK_DEPTH	The number of items on the stack.
ENDISM	The endianness of SMite: 0 = Little-endian, 1 = Big-endian.
WORD_SIZE	The number of bytes in a word. Must be in the range 2 to 32 inclusive, and a power of 2.

Table 1: Registers

## 2.3 Stack

The stack is a LIFO stack of words used for passing values to instructions and routines and for holding subroutine return addresses. To **push** an item on to the stack means to add a new item to the top of the stack, increasing the stack depth by 1; to **pop** an item means to reduce the stack depth by 1. Instructions that change the number of items on the stack implicitly pop their arguments and push their results.

## 2.4 Operation

Before SMite is started, ENDISM should be set to 0 or 1 according to the implementation, and WORD\_SIZE to the appropriate value. The other registers should be initialised to 0.

ENDISM and WORD\_SIZE must not change.

SMite is started by a call to the interface calls **run()** or **single\_step()** (see section 4.2). In the former case, the execution cycle is entered:

```
begin
    execute the instruction in the 6 least-significant bits of I
    shift I logically 6 bits to the right
repeat
```

In the latter case, the contents of the execution cycle is executed once, and control returns to the calling program.

Note that the calls **run()** and **single\_step()** do not perform the initialisation specified above.

## 2.5 Errors and termination

When SMite encounters certain abnormal situations, such as an attempt to access an invalid address, or divide by zero, an **error** is **raised**, and execution terminates; an **error code** is returned to the caller. The instruction has no effect, and PC is not advanced to the next instruction. If the error is a stack or memory access error, BAD is set to the stack position or address that caused the error.

Execution can also be terminated explicitly by performing a HALT instruction (see section 3.10).

Error codes are unsigned numbers. 0 to 128 are reserved for SMite's own error codes; the meanings of those that may be raised by SMite are shown in table 2.

Code	Meaning
0	<b>single_step()</b> has terminated without error.
1	Invalid opcode (see section 3.12).
2	Attempt to set STACK_DEPTH greater than STACK_SIZE. BAD is set to the extra number of words of stack space required.
3	Invalid stack read. BAD is set to the invalid stack position.
4	Invalid stack write. BAD is set to the invalid stack position.
5	Invalid memory read. BAD is set to the invalid address.
6	Invalid memory write. BAD is set to the invalid address.
7	Address alignment error: raised when an instruction expecting a word-aligned address is given a valid but non-aligned address.
8	Division by zero attempted (see section 3.6).
128	A HALT instruction was executed.

Table 2: Errors raised by SMite

### 3 Instruction set

The instruction set is listed below, with the instructions grouped according to function. The instructions are given in the following format:

NAME ( *before* - *after* )  
Description.

The first line consists of the name of the instruction. On the right is the stack effect, which shows the effect of the instruction on the stack. Underneath is the description.

**Stack effects** are written

( *before* - *after* )

where *before* and *after* are stack pictures showing the items on top of a stack before and after the instruction is executed. An instruction only affects the items shown in its stack effects. The brackets and dashes serve merely to delimit the stack effect and to separate *before* from *after*. **Stack pictures** are a representation of the top-most items on the stack, and are written

$$i_1 \ i_2 \dots i_{n-1} \ i_n$$

where the  $i_k$  are stack items, each of which occupies a whole number of words, with  $i_n$  being on top of the stack. The symbols denoting different types of stack item are shown in table 3.

Types are only used to indicate how instructions treat their arguments and results; SMite does not distinguish between stack items of different types. In stack pictures the most general argument types with which each instruction can

Symbol	Data type
<i>flag</i>	a Boolean flag, 1 for true and 0 for false
<i>byte</i>	byte
<i>n</i>	signed number
<i>u</i>	unsigned number
<i>n u</i>	number (signed or unsigned)
<i>x</i>	unspecified word
<i>addr</i>	address
<i>a-addr</i>	word-aligned address

Table 3: Types used in stack effects

be supplied are given; subtypes may be substituted. Using the phrase “ $i \Rightarrow j$ ” to denote “ $i$  is a subtype of  $j$ ”, table 4 shows the subtype relationships. The subtype relation is transitive.

$u \Rightarrow x$
$n \Rightarrow x$
$flag \Rightarrow u$
$byte \Rightarrow u$
$a-addr \Rightarrow addr \Rightarrow u$

Table 4: The subtype relation

Numbers are represented in twos complement form. *addr* consists of all valid virtual machine addresses. Numeric constants can be included in stack pictures, and are of type *n|u*.

Each type may be suffixed by a number in stack pictures; if the same combination of type and suffix appears more than once in a stack effect, it refers to identical stack items. Alternative *after* pictures are separated by “|”, and the circumstances under which each occurs are detailed in the instruction description.

The symbols  $i*x$  and  $j*x$  are used to denote different collections of zero or more words of any data type. Ellipsis is used for indeterminate numbers of specified types of word.

### 3.1 Instruction fetch

NEXT ( - )

Load the word pointed to by PC into I then add WORD\_SIZE to PC. If PC is not word-aligned, raise error 7, setting BAD to the contents of PC.

### 3.2 Stack manipulation

These instructions manage the stack:

POP ( x - )

Remove  $x$  from the stack.

DUP  $(x_u \dots x_0 \ u - x_u \dots x_0 \ x_u)$

Remove  $u$ . Copy  $x_u$  to the top of the stack.

SWAP  $(x_u \dots x_0 \ u - x_0 \ x_{u-1} \dots x_1 \ x_u)$

Exchange the top stack item with the  $u$ th. If  $u$  is zero, do nothing.

### 3.3 Numbers

LIT  $(- \ n)$

The word pointed to by PC is pushed on to the stack, and PC is incremented to point to the following word.

LIT\_PC\_REL  $(- \ n)$

Like LIT, except that the initial value of PC is added to the value pushed on to the stack.

### 3.4 Logic and shifts

Logic functions:

NOT  $(x_1 \ - \ x_2)$

Invert all bits of  $x_1$ , giving its logical inverse  $x_2$ .

AND  $(x_1 \ x_2 \ - \ x_3)$

$x_3$  is the bit-by-bit logical “and” of  $x_1$  with  $x_2$ .

OR  $(x_1 \ x_2 \ - \ x_3)$

$x_3$  is the bit-by-bit inclusive-or of  $x_1$  with  $x_2$ .

XOR  $(x_1 \ x_2 \ - \ x_3)$

$x_3$  is the bit-by-bit exclusive-or of  $x_1$  with  $x_2$ .

Shifts:

LSHIFT  $(x_1 \ u \ - \ x_2)$

Perform a logical left shift of  $u$  bit-places on  $x_1$ , giving  $x_2$ . Put zero into the least significant bits vacated by the shift. If  $u$  is greater than or equal to the number of bits in a word,  $x_2$  is zero.

RSHIFT  $(x_1 \ u \ - \ x_2)$

Perform a logical right shift of  $u$  bit-places on  $x_1$ , giving  $x_2$ . Put zero into the most significant bits vacated by the shift. If  $u$  is greater than or equal to the number of bits in a word,  $x_2$  is zero.

ARSHIFT  $(x_1 \ u \ - \ x_2)$

Perform an arithmetic right shift of  $u$  bit-places on  $x_1$ , giving  $x_2$ . Copy the original most-significant bits into the most significant bits vacated by the shift. If  $u$  is greater than or equal to the number of bits in a word, all the bits of  $x_2$  are the same as the original most-significant bit.

### 3.5 Comparison

These words compare two numbers (or, for equality tests, any two words) on the stack, returning a flag:

EQ (  $x_1 \ x_2$  - *flag* )

*flag* is true if and only if  $x_1$  is bit-for-bit the same as  $x_2$ .

LT (  $n_1 \ n_2$  - *flag* )

*flag* is true if and only if  $n_1$  is less than  $n_2$ .

ULT (  $u_1 \ u_2$  - *flag* )

*flag* is true if and only if  $u_1$  is less than  $u_2$ .

### 3.6 Arithmetic

These instructions consist of monadic and dyadic operators. All calculations are made without bounds or overflow checking, except as detailed for certain instructions.

Negation and addition:

NEGATE (  $n_1$  -  $n_2$  )

Negate  $n_1$ , giving its arithmetic inverse  $n_2$ .

ADD (  $n_1 | u_1 \ n_2 | u_2$  -  $n_3 | u_3$  )

Add  $n_2 | u_2$  to  $n_1 | u_1$ , giving the sum  $n_3 | u_3$ .

Multiplication and division (note that all division instructions raise error -4 if division by zero is attempted):

MUL (  $n_1 | u_1 \ n_2 | u_2$  -  $n_3 | u_3$  )

Multiply  $n_1 | u_1$  by  $n_2 | u_2$  giving the product  $n_3 | u_3$ .

DIVMOD (  $n_1 \ n_2$  -  $n_3 \ n_4$  )

Divide  $n_1$  by  $n_2$  using symmetric division, giving the single-word quotient  $n_3$  and the single-word remainder  $n_4$ . The quotient is rounded towards zero.

UDIVMOD (  $u_1 \ u_2$  -  $u_3 \ u_4$  )

Divide  $u_1$  by  $u_2$ , giving the single-word quotient  $u_3$  and the single-word remainder  $u_4$ .

### 3.7 Memory

These instructions fetch and store words and bytes to and from memory:

LOAD ( *a-addr* -  $x$  )

$x$  is the value stored at *a-addr*. If *a-addr* is not word-aligned, raise error 7, setting BAD to *a-addr*.

STORE (  $x$  *a-addr* - )

Store  $x$  at *a-addr*. If *a-addr* is not word-aligned, raise error 7, setting BAD to *a-addr*.

LOADB ( *addr* - *byte* )  
Fetch the byte stored at *addr*. The unused high-order bits are all zeroes.

STOREB ( *byte* *addr* - )  
Store *byte* at *addr*. Only one byte is transferred.

### 3.8 Control

These instructions implement unconditional and conditional branches, and subroutine call and return (subroutine return is BRANCH):

BRANCH ( *a-addr* - )  
Set PC to *a-addr*. Perform the action of NEXT.

BRANCHZ ( *flag* *a-addr* - )  
If *flag* is false then set PC to *a-addr*. Perform the action of NEXT.

CALL ( *a-addr*<sub>1</sub> - *a-addr*<sub>2</sub> )  
Exchange PC with the top stack value. Perform the action of NEXT.

### 3.9 Registers

GET\_WORD\_SIZE ( - *u* )  
*u* is the value of WORD\_SIZE.

GET\_STACK\_DEPTH ( - *u* )  
*u* is the value of STACK\_DEPTH, the number of items on the stack.

SET\_STACK\_DEPTH ( *u* - )  
Set STACK\_DEPTH to *u*.

### 3.10 Termination

This instruction terminates execution (see section 2.5):

HALT ( - )  
Terminate execution with error code 128.

### 3.11 Instruction encoding

Instructions are encoded as 6-bit opcodes; opcodes are packed into words, which are executed starting at the least-significant bits.

### 3.12 Instruction opcodes

Table 5 lists the instruction opcodes in numerical order. Other instruction opcodes are undefined.

Opcode	Instruction	Opcode	Instruction
0x00	NEXT	0x10	NEGATE
0x01	POP	0x11	ADD
0x02	DUP	0x12	MUL
0x03	SWAP	0x13	DIVMOD
0x04	LIT	0x14	UDIVMOD
0x05	LIT_PC_REL	0x15	LOAD
0x06	NOT	0x16	STORE
0x07	AND	0x17	LOADB
0x08	OR	0x18	STOREB
0x09	XOR	0x19	BRANCH
0x0a	LSHIFT	0x1a	BRANCHZ
0x0b	RSHIFT	0x1b	CALL
0x0c	ARSHIFT	0x1c	GET_WORD_SIZE
0x0d	EQ	0x1d	GET_STACK_DEPTH
0x0e	LT	0x1e	SET_STACK_DEPTH
0x0f	ULT	0x1f	HALT

Table 5: Instruction opcodes

## 4 External interface

SMite’s external interface comes in three parts. The calling interface allows SMite to be controlled by other programs. Implementations can provide extra instructions to access system facilities, previously written code, native libraries and so on. The object module format allows compiled code to be saved, reloaded and shared between systems.

### 4.1 Object module format

The object module starts with the ASCII codes of the letters “smite” followed by ASCII NUL (0x00), then the one-byte values of the `ENDISM` and `WORD_SIZE` registers of the system which saved the module, then a word (of the given endianness and size) containing the number of bytes the code occupies. Then follows the code.

Object modules have a simple structure, as they are only intended for loading an initial memory image into SMite.

### 4.2 Calling interface

The calling interface is difficult to specify with the same precision as the rest of SMite, as it may be implemented in any language. However, since only basic types are used, and the semantics are simple, it is expected that implementations in different language producing the same result will be easy to program. A Modula-like syntax is used to give the definitions here. Implementation-defined error codes must be documented, but are optional. All addresses passed as parameters must be word-aligned. An implementation of SMite must provide the following calls:



**run ()** : integer

Start SMite by entering the execution cycle as described in section 2.4. If an error is raised, the code is returned as the result.

**single\_step ()** : integer

Execute a single pass of the execution cycle. If an error is raised, it is returned; otherwise, return 128.

**load\_object (file, address)** : integer

Load the object module given by the file descriptor *file*, which may be a filename or some other specifier, to the VM address *address*. If the module is invalid, then return  $-2$ . If the module cannot be loaded by the current implementation (for example, the word size or `ENDISM` is incompatible), return  $-3$ . If the code will not fit into memory at the address given, or the address is out of range or unaligned, return  $-4$ . Otherwise load the code into memory, converting it if the endianness value is different from the current value of `ENDISM`. The result is the length of the code in bytes if successful, and  $-1$  if there is a filing system or other error.

**save\_object (file, address, length)** : integer

Save the area of memory specified by *address* and *length* to the file descriptor given by *file*. On success, 0 is returned. If the address or length are invalid, the return code is  $-2$ . If a file system error occurs, the return code is  $-1$ .

SMite must also provide access to its registers and address space through appropriate data objects.

## Acknowledgements

Martin Richards's demonstration of his BCPL-oriented Cintcode virtual machine [1] convinced me it was going to be fun working on virtual machines. He also supervised my BA dissertation project, Beetle [2], on which SMite is based.

## References

- [1] Martin Richards. Cintcode distribution, 2000. <https://www.cl.cam.ac.uk/~mr/BCPL.html>.
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- [3] Reuben Thomas. *Mite: a basis for ubiquitous virtual machines*. PhD thesis, University of Cambridge Computer Laboratory, November 2000. <https://rrt.sc3d.org/>.