

Using the HFS+ Journal For Deleted File Recovery

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Using the HFS+ Journal For Deleted File Recovery

Aaron Burghardt Adam Feldman

Introduction



- Client-sponsored assignment
- Tasked to replace existing deleted file recovery tools
- Increase automation and improve accuracy

Catalog File Records

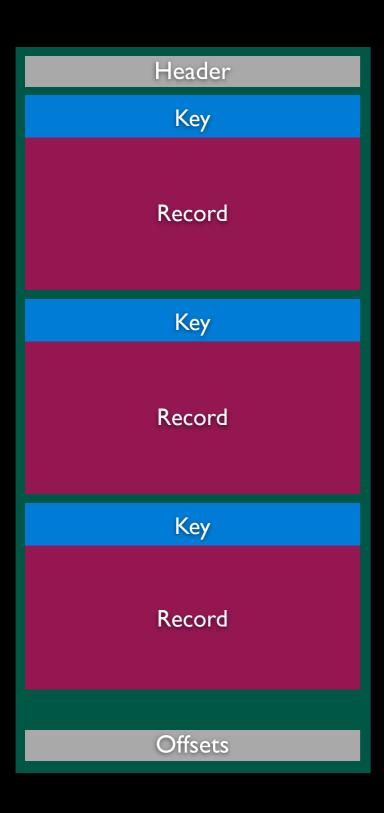
Key

Cat. File Rec.

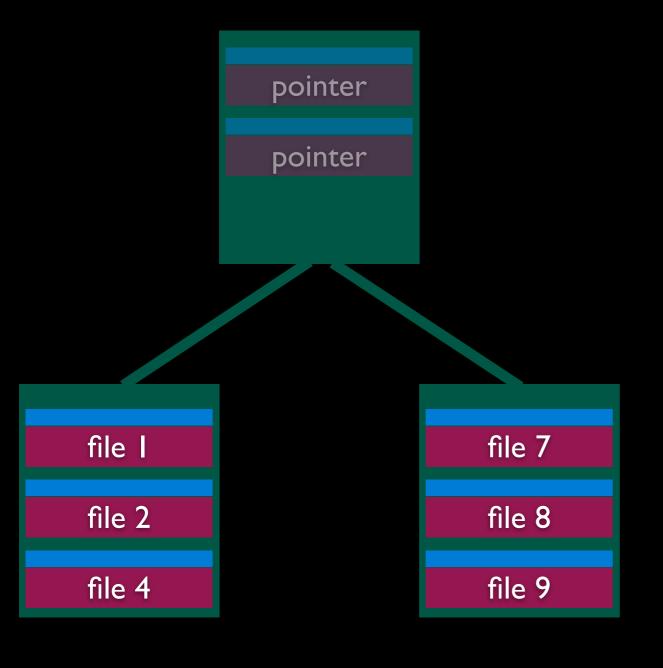
CNID
Create Date
Mod Date
Access Date
Owner ID
Group ID
Unix permissions
Extents

- Catalog Node ID (unique, like an inode)
- Create, mod, access times
- Owner and group IDs
- Unix permissions
- Extent Records (i.e., fragment descriptors)
- Stored adjacent to Key in B-tree node
- Key = Parent CNID + file name

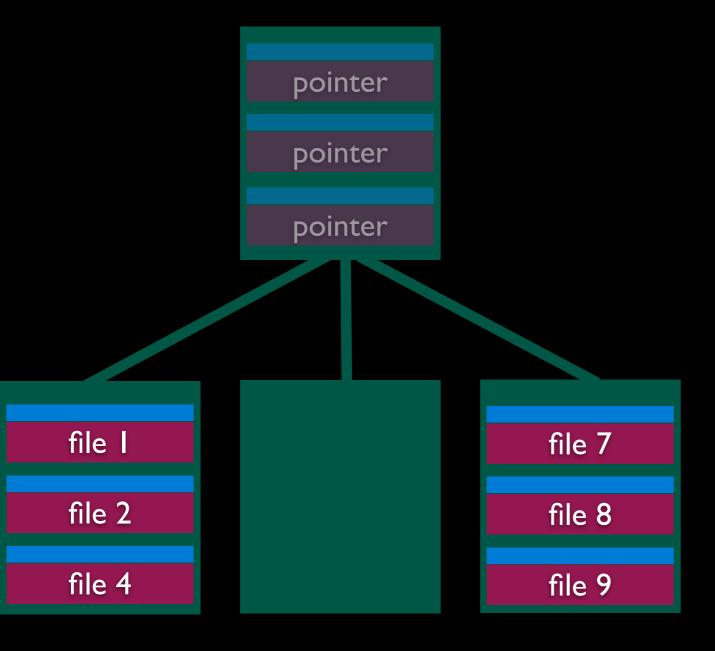
B-tree Node



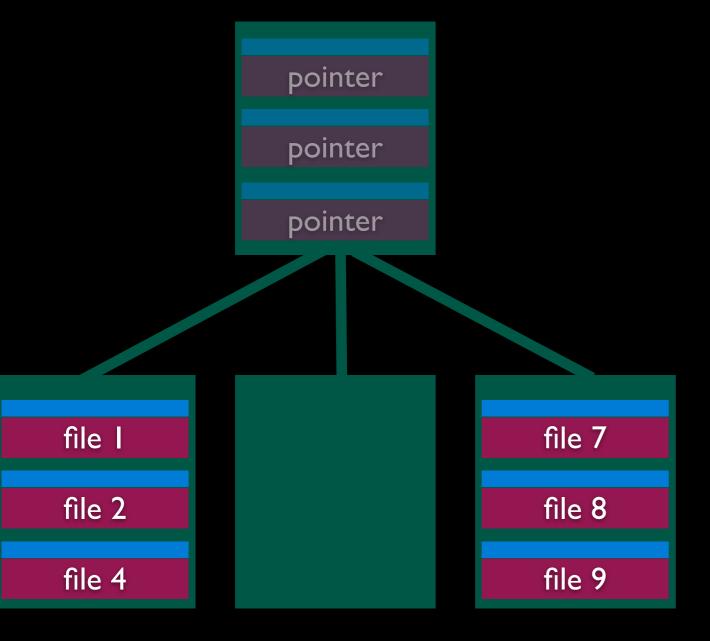
- Typically 8 KB
- Records and keys vary in size
- Records/Keys packed top-tobottom



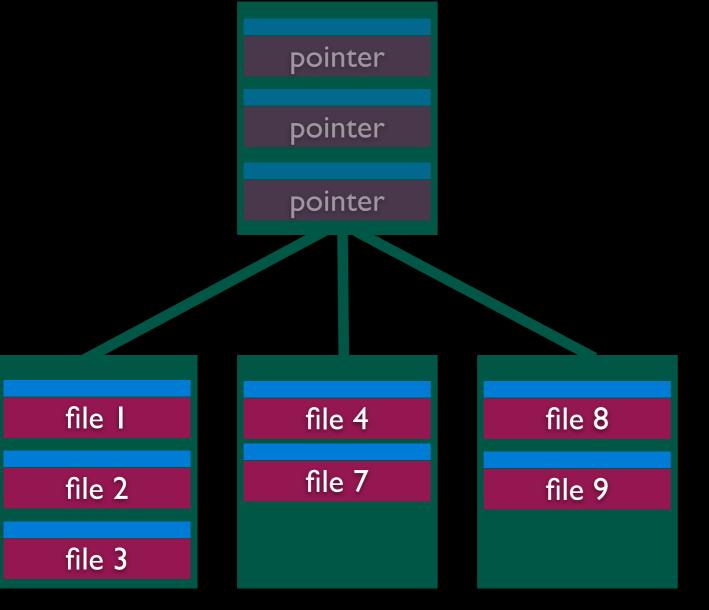
- Nodes are organized in a tree
- Records always maintained in sorted order
- Creation and deletion of files causes records to rearrange



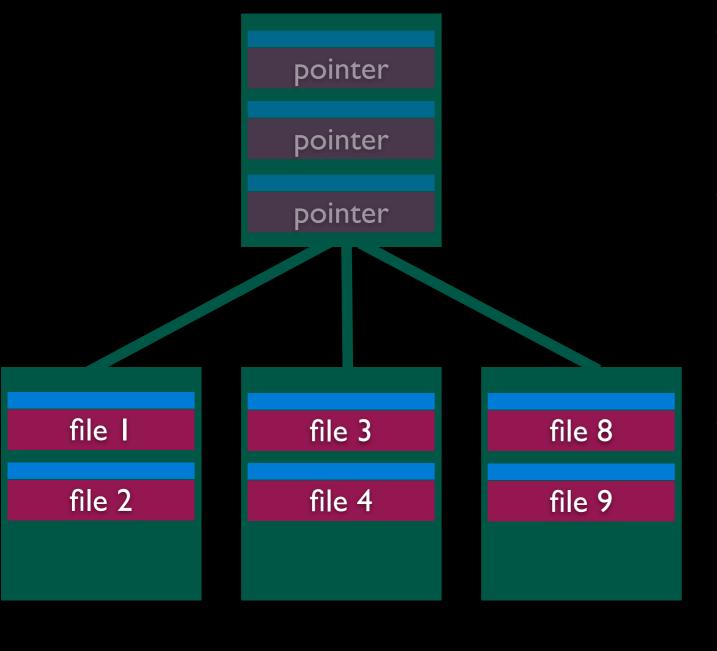
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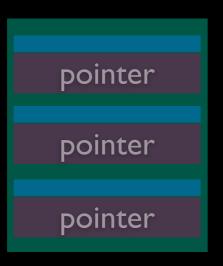


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File System Files



- Volume Bitmap
- Catalog File
- Extents Overflow
- Extended Attributes

file I

file 3

file 8 file 9

File System Files



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File System Files

Volume Bitmap

talog File

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Role of Journal

Volume Bitmap

Catalog File

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Role of Journal

Volume Bitmap

Catalog File

Extents Overflow

Extended Attributes

- Introduced in Mac OS X 10.2
- Records pending changes to metadata
- Collects related changes in transactions
- Sector/Block-oriented
- Allocation: 8 MB + 8
 MB per 100 GB vol. size

Role of Journal

Volume Bitmap

Catalog File

Extents Overflow

Extended Attributes

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Key Points

Volume Bitmap

Catalog File

Extents Overflow

Extended Attributes

Key Points

Volume Bitmap

Catalog File

Extents Overflow

Extended Attributes

- B-tree nodes are recorded as whole unit
- Catalog File, Extents
 Overflow, Extended
 Attributes are B-trees:
 must distinguish
- A Catalog File Record may appear in the journal due to unrelated changes

Recovery Approach

- I. Begin at logical start of journal file
- 2. Scan until a B-tree node is found
 - No header signature
 - Sanity checks used to validate
- 3. Iterate node records
 - a. Search the active Catalog File for each Catalog File Record
 - b. If not found, conclude it is a deleted file

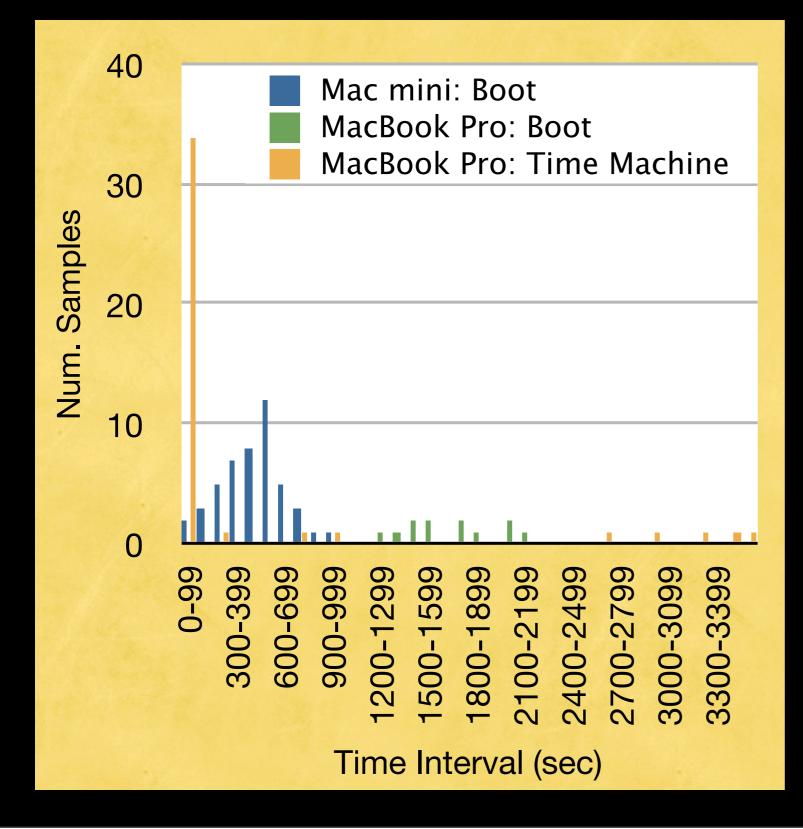
Recovery Approach (cont'd)

- 4. Cache the Catalog File Record in the deleted file cache:
 - Replace duplicate (by CNID) record
- 5. Score the recoverability:
 - Check current in-use status of blocks
 - Good: all blocks unused
 - Partial: first block(s) not in use
 - Poor: first block(s) in use

Test Configurations

- Goal: establish typical "window of opportunity"
- Two test configurations:
 - Mac mini
 - MacBook Pro
- Mixture of use cases:
 - Boot volumes
 - Secondary volumes
 - Time Machine

Lifetime of Data in Journal



- Boot volume: 5 min to 30 min
- Secondary: can be several hours or more
- Time Machine: idle between backups, approximately 30 sec during a backup

Empirical Results

Volume	Good	Partial	Poor	Total
MBP: Boot	59	0	8	67
MBP:Time Mach	3	0	0	3
Mini: Boot	10	0	4	4
Mini: FireWire	32	0	87	119
Mini: FireWire	4	0	22	36
Mini: Flash	4	0	21	162

Limitations

- Data in journal is short-lived
- Evidence of the file must be in the journal prior to it being deleted
 - Deleted status determined by deduction
 - Can't predict if a deleted file is detectable
- Path may not be recoverable
- Only has 8 extent records
- Time of deletion unknown

Summary

- Effective for recently deleted files
- Recovers files and metadata with high accuracy
- Limited by short window of opportunity and the need for record to exist in journal prior to deletion
- Complementary to file carving

Thank you

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Questions?