

BinComp: A Stratified Approach to Compiler Provenance Attribution

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Presented At

The Digital Forensic Research Conference **DFRWS 2015 USA** Philadelphia, PA (Aug 9th - 13th)

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Agenda

- Introduction
- □ Use cases
- Motivating Example
- □ BinComp Approach
- Evaluation
- Comparison
- □ Discussion and Limitations



Introduction

- □ Compiler provenance attribution encompasses
 - Compiler family
 - Compiler version
 - Optimization level
 - Compiler-related functions
- Why compiler provenance attribution?
 - It is a necessary component at pre-processing stage for binary analysis
 - □ It is important in digital forensics
 - It provides information about the process by which a malware binary is produced

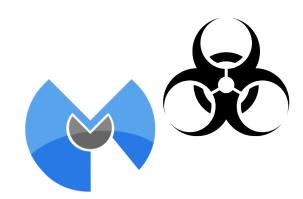


Use Cases

Compiler identification can be useful for

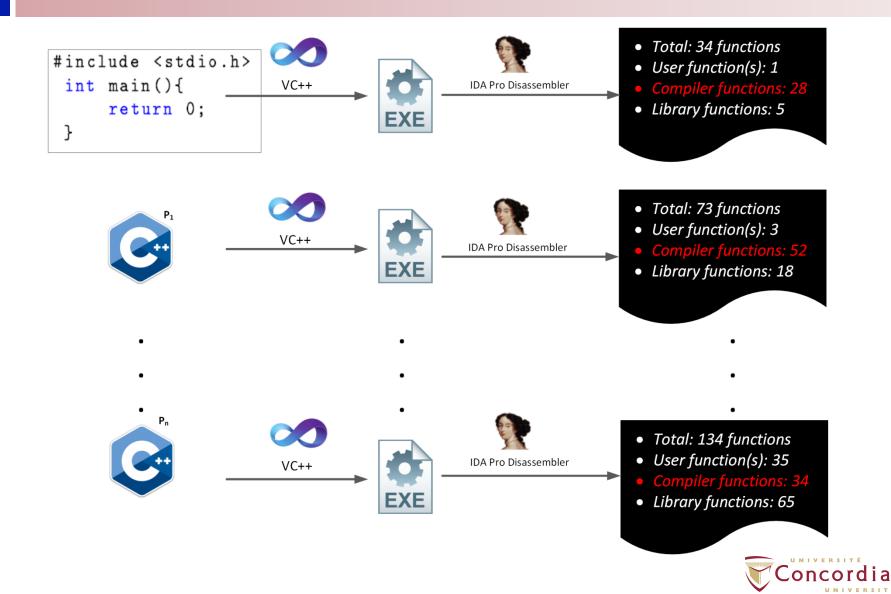
- Tool chain malware analysis in binaries
 - Clustering malware samples based on families (assuming each family is using a specific compiler/linker setting)
 - Reasoning about the evolution of malware binaries
- Library identification, authorship attribution, clone detection, etc.







Motivating Example



Extract Compiler-related Functions



```
['___p__environ', '_strlen', '_glob_match', '_rewinddir',
'___mingwthr_run_key_dtors.part.0', '__cexit', '_strncpy', '_signal',
'_memcpy', '_GetCommandLineA@0', '__pei386_runtime_relocator',
'__errno', '___mingw_globfree', '_register_frame_ctor', '_readdir',
'__setmode', '_puts', '_SetUnhandledExceptionFilter@4', ...
```



VC++

```
['___tmainCRTStartup', '___ValidateImageBase', '?
__CxxUnhandledExceptionFilter@@YGJPAU_EXCEPTION_POINTERS@@
@Z', '__dllonexit', '_controlfp_s', '_invoke_watson', '_amsg_exit', '?
terminate@@YAXXZ', '__setdefaultprecision', '__SEH_epilog4', '_lock',
'__onexit', 'start', '_pre_cpp_init', '_XcptFilter', '_init, ...
```



ICC

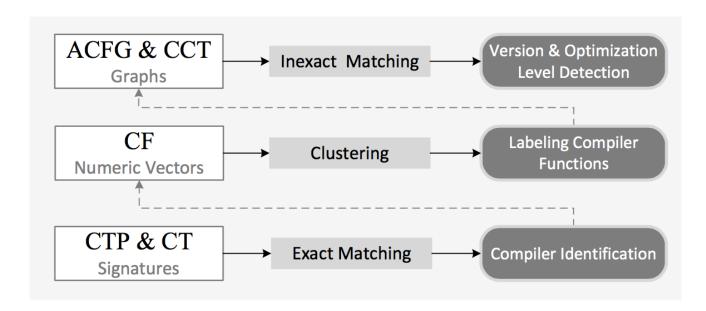
```
['__cxa_atexit', '.strncat', 'getenv@@GLIBC_2.2.5',
'_ZNSt8ios_base4InitD1Ev', '__intel_cpu_features_init', '.setenv',
'_ZNSt8ios_base4InitC1Ev', '.fprintf', '.catopen',
'catopen@@GLIBC_2.2.5', '.term_proc', 'vsprintf@@GLIBC_2.2.5',
'_start', 'strlen', 'strncpy', 'strchr', ...
```



BinComp Approach

□ A layered approach

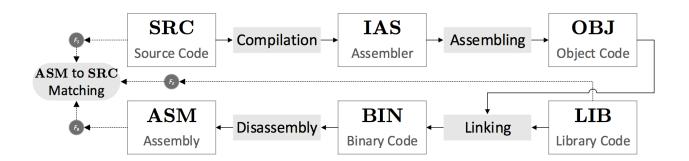
- Compiler identification (syntax features)
- Labeling compiler functions (semantic features)
- Version & optimization level detection (structural & semantic features)





1st layer: Compiler Identification

- □ Collecting a set of known source code
- Observing the compiled outputs
- Extracting the features
 - CTP (Compiler Transformation Profile)
 - CT (Compiler Tags)
- Exact matching





1st layer: Compiler Identification

- CTP (Compiler Transformation Profile)
 - How source-level data and control structures are reflected in the assembly output
 - For example, corresponding assembly code of if/else compiled with VS will be cmp or test, then jcc.
- CT (Compiler Tag)
 - Compilers may embed certain tags in the form of strings or constants
 - For example, GCC writes a .comment section that contains the GCC version string



2nd layer: Compiler Function Labeling

- □ Signature generation
 - Extracting CF (Compiler Functions)
 - Numerical vectors such as number of instructions, type of registers, etc.
 - Symbolic vectors: such as function names, function prototypes, etc.
- □ Signature detection
 - Computing the similarity

n Function Name, ne Demangled Name, i Imported Function, k Call List (From), ke Demangled Calls, ix Number of Imported Functions, kx Number of Calls c Constant, s String, cx Number of Constants, sx Number of Strings p Function Prototype, a Function Argument, r Return Type, g Number of a, b Size of Arguments, gx RES. Number of g m Number of Instructions, I Size of Local Variables, f Function Flags, o Code References (From), z Function Size, mx RES. Number of m, ox Number of o, cc Cyclomatic Complexity, bx Number of Basic Blocks d Dictionary (Malware Tag), t API Tag, dx Number of d, tx Number of t av Anti-VM, reg General Register, mem Memory Reference, bix [Base +Index], bid [Base+Index+Displacement], imm Immediate, ifa Immediate Far Address, ina Immediate Near Address, fpp FPP Register, ctr Control Register, dbr Debug Register, trr Trace Register DTR Data Transfer, DTO Data Transfer Address Object, FLG Flag Manipulation, DTC Data Transfer Conversion, ATH Binary Arithmetic, LGC Logical, CTL Control Transfer, INO Input Output, INT Interrupt and System, FLT Floating, MSC Misc.



2nd layer: Compiler Function Labeling

Example of CF features

COMP	OPL	Symbolic Function ID	DTR,	DTO,	FLG,	ATH,	LGC,	CTL,	INO,	INT,	FLT,	REG,	MEM,	IMM,	IFA,	INA
VS	OP2	@security_check_cookie@4	0,	0,	0,	0,	0,	4,	0,	0,	0,	001,	001,	000,	000,	002
VS	OP2	tmainCRTStartup	3,	0,	1,	2,	1,	6,	0,	0,	0,	077,	028,	017,	000,	025
GCC	OP0	mingw_CRTStartup	3,	1,	1,	3,	6,	7,	0,	0,	0,	216,	015,	068,	000,	071
GCC	OP0	gcc_register_frame	3,	0,	1,	1,	0,	3,	0,	0,	0,	029,	001,	014,	000,	009

Instruction Categories.

Feature	Description	Feature	Description
DTR	Data Transfer	INA	Indirect Near Address
INO FLT	Input/Output Float Point	DTO FLG	Data Transfer Object Flag Manipulation
REG	Registers	LGC	Logical Instructions
MEM	Memory	CTL	Control Instructions
IMM INT	Immediate Value Interrupt/System	IFA ATH	Indirect Far Address Arithmetic Instructions

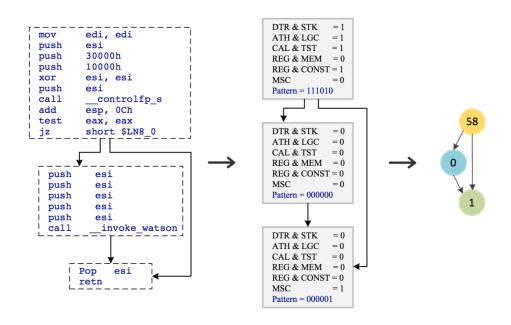
$$d_J(t_i, t_j) = \frac{S(t_i \wedge t_j)}{S(t_i \vee t_j)}$$

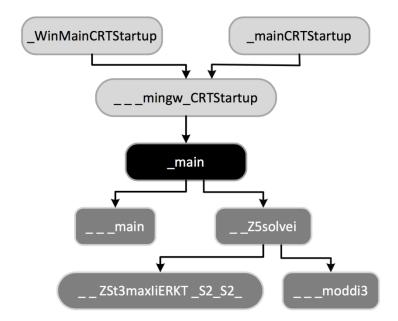
t_i, t_i: fingerprint vectors generated from the candidate function pairs



3rd layer: Version & Optimization

- Extracting the features
 - ACFG (Annotated Control Flow Graph)
- CCT (Compiler Constructor Terminator)







3rd layer: Version & Optimization

- According to our experiments ACFG and CCT are the best features to detect the version and optimization
 - The different versions can affect the ACFG
 - The different optimization levels affect the CCT
 - For instance, the CCT for full optimized is a subset of CCT for the no-optimization code

Data set compilation settings.

Compiler	Version	Optimization		
GCC	3.4	00		
	4.4	02		
ICC	10	00		
	11	02		
VS	2010	00		
	2012	02		
XCODE	5.1	00		
	6.1	O2s		



Evaluation

- □ Dataset
 - ■Four free open-source projects (SQLite, zlib, libpng, and openSSL)
 - ■Google Code Jam (232 files)
 - ■Students Code Projects (993 files)





Evaluation

□ Results

Accuracy for variations of compiler versions.

Compiler	Version	Accuracy		
GCC	3.4.x	86%		
	4.4.x	89%		
ICC	10.x	83%		
	11.x	90%		
VS	2010	70%		
	2012	71%		
XCode	5.x	78%		
	6.1	74%		

Accuracy for variations of compiler optimization levels.

Compiler	Optimization	Average Accuracy			
GCC	00, 02	91%			
ICC	O0, O2	89%			
VS	00, 02	95%			



Comparison

	IDA Pro Disassembler	Rosenblum et al	BinComp		
Features	Entry point signatures	Syntax (n-gram)	Syntax, Semantic, Structural		
Detection Method	Signature based	Classification	Exact/Inexact matching		
Compilers	VS, Delphi, Fortran	GCC, VS, ICC	VS, GCC, ICC, XCODE		
Required data set	Small	Large	Small		
Scalability					
Time efficient		×			
Identifying Ver.		8			
Identifying Opt.	× ×	×			
	× ×				



Discussion and Limitations

- BinComp requires few data set and it can be applied for any compiler
- □ BinComp is efficient in terms of time and scalability
- Limitations
 - The binary code is deobfuscated
 - \square Only Intel x86/x86-64 architecture is considered





Thank You!







Evaluation

- Accuracy
 - □ Precision (P)
 - Recall (R)

$$P = \frac{TP}{TP + FP}$$

$$R = \frac{TP}{TP + FN}$$

 Our application domain is much more sensitive to false positives than false negatives

$$F_{0.5} = 1.25 \cdot \frac{PR}{0.25P + R}$$



Neighbor hash graph kernel (NHGK)

- Neighbor hash graph kernel (NHGK)
 - Condense the information contained in a neighborhood into a single hash value
 - Label each node in the function call graph
 - each function is characterized by its numerical and symbolic feature vectors
- our method strives to model the composition of functions
 - The neighborhood of a function must be taken into account.
 - We compute a neighborhood hash over all of its direct neighbors in the function call graph

$$h(f_i) = shr_1(G(f_i)) \oplus \left(\oplus_{f_j \in N_{f_i}} G(f_j) \right)$$

- \square shr₁: a one-bit shift right operation
- : a bit-wise XOR on the binary labels

