



Advancing the AFF4 to the challenges of Volatile Memory & Single Hashes

Dr. Bradley Schatz Director, Schatz Forensic

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Agenda

- AFF4 background
- What can be improved with hashing in AFF4?
- Our refined single hashing proposal: the Linear Chunk Hash





Traditional forensic imaging introduces unnecessary delays

- Physical acquisition delays analysis by hours
 - Doesn't scale to multi-core/RAID/SSD/NVMe
 - Doesn't allow non-linear writes (concurrent imaging & analysis)

- Possible solution: logical imaging/triage
 - loses context and unallocated





Traditional forensic imaging doesn't scale

- ZIP/Deflate compression is really CPU expensive
 - (~ 40MB/s/core for high entropy)

- Hashing is CPU expensive
 - (~ 600MB/s/core for 2016 i7)

- Linear bitstream hashing isn't parallelizable
 - Ceiling of around 600MB/s on current generation i7 CPU's





The AFF4 removes the bottleneck from forensic imaging

- Snappy compression is CPU inexpensive
 - (> 1000 MB/s/core)

- Hashing is still CPU expensive
 - (~ 600MB/s/core for 2016 i7)

- AFF4 block hashing is parallelizable
 - Seeing > 1000 MB/s on current generation CPU's



The AFF4 removes enables acquisition to occur at the same time as analysis

Non-linear, potentially partial compressed block stream

Hashing mechanism that enables incremental hash calculation





AFF4 Standard v1.0 Hashing

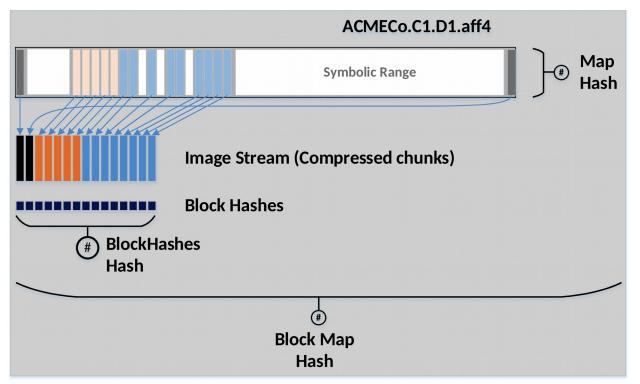
- Defines a hash construction called the Block Map Hash
 - Single hash for image

- Hybrid of:
 - Block (segment) hashing for concrete data chunks
 - Non concrete data chunks (eg. Sparse) protected by hash on map
 - Merkel tree like hash hierarchy





AFF4 Standard v1.0 Hashing





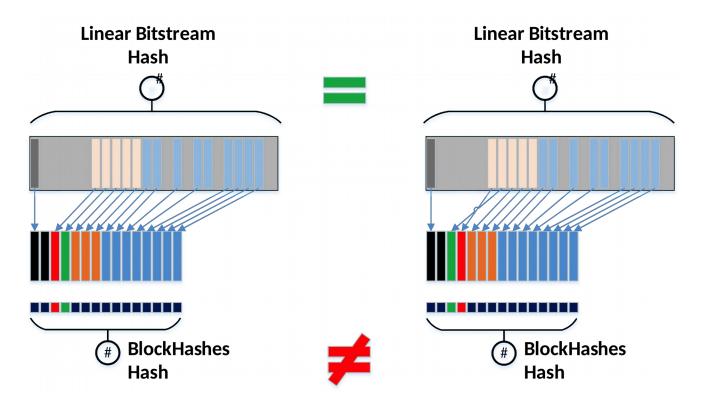
Desirable properties of hashing

	Linear Bitstream Hash	AFF4 Block Map Hash
Hash is concise	Yes	Yes
Hash of image matches hash calculated when imaging	Yes	Yes
Changed image results in failed hash match	Yes	Yes
Successive hashes of same source result in same hash	Yes	Maybe





Parallelism may lead to differing hashes





Can we generate a block based hash that is not subject to reordering?

- Hypothesis 1: Reorder the lower Block Hashes based on their upper address
 - Works well for images of disks with no read errors.
 - Doesn't work so well for:
 - Disk images with read errors
 - Memory images with chunk sizes > memory page size





Sub chunk size data introduces further complication

- High performance requires (chunk size > disk sector size)
 - IO Throughput 1MB for disk IO
 - Memory consumption- Maps
- ddrescue bad sector imaging algorithm requires chunks of 512B or 4KB when the chunk size might be many multiples of this
- RAM physical address space holes lead to sub chunk sized data runs



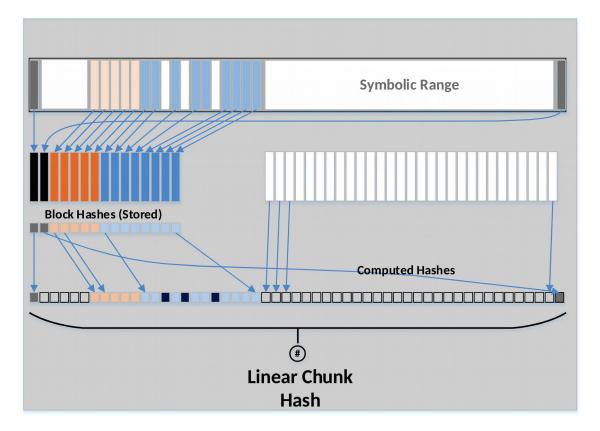


Hypothesis 1a: introduce sub chunk encoding rules

- 4k sector gets written to first 4k of chunk, rest zero filled.
- -or-
- 4k sectors coalesced

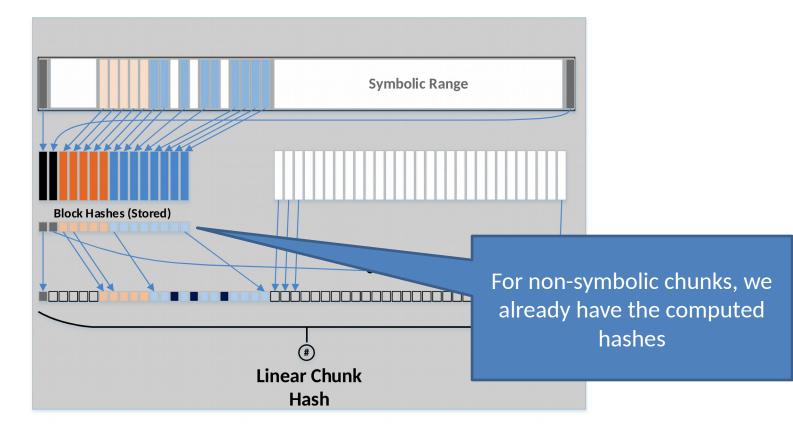






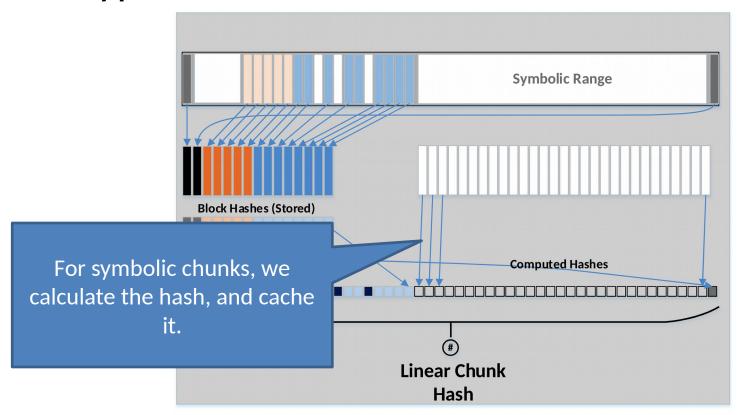






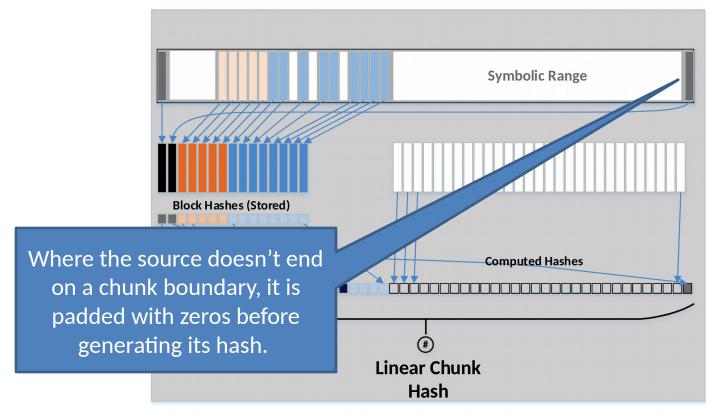








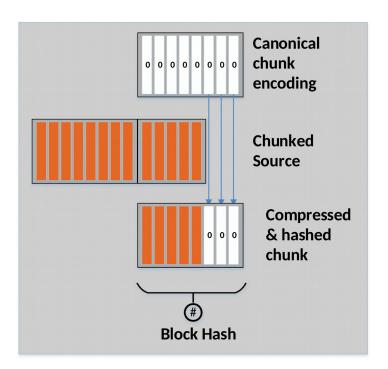






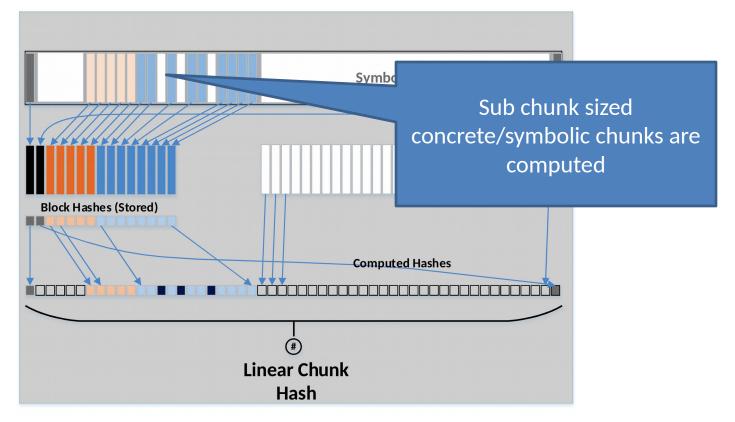


Sub chunk sized end of stream is zero filled



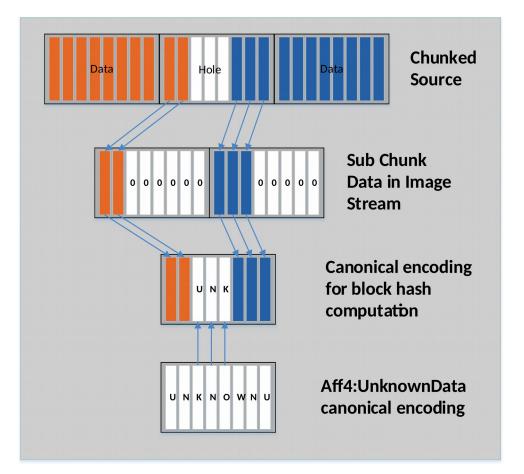












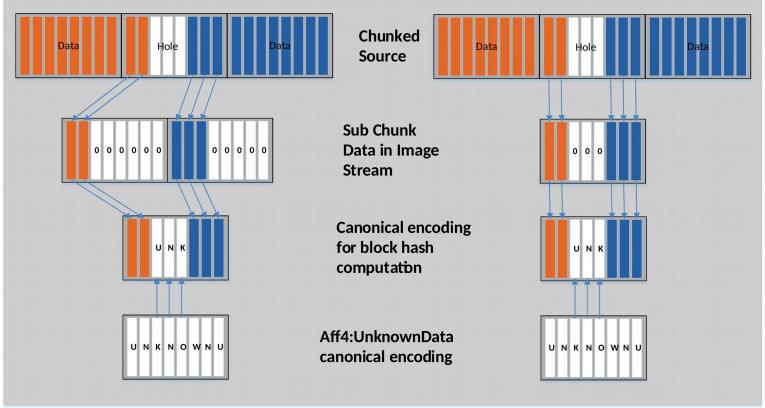
Sub chunk symbolic encoding rules

- Partial chunks stored padded
- Block hash computed from canonical block encoding



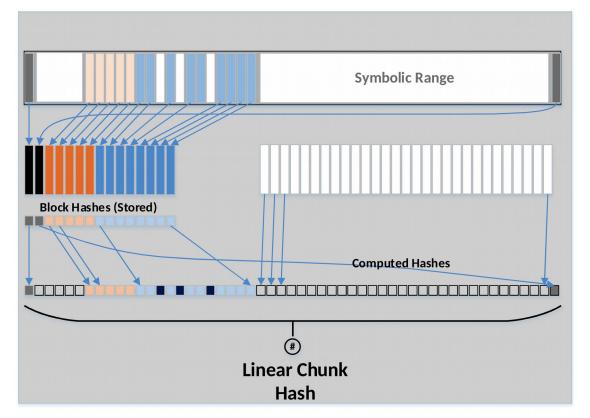


Alternate part chunk storage schemes result in the same hash











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Contact

Dr Bradley Schatz

https://evimetry.com/

bradley@evimetry.com

@blschatz

@wirespeed4n6