

Coffee forensics — Reconstructing data in IoT devices running Contiki OS

By:

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COFFEE FORENSICS

Reconstructing data in IoT devices running Contiki OS

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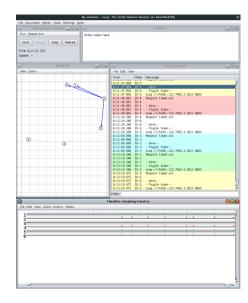
Introduction - Background

- IoT systems includes a variety of subsystems
 - Focus on resource constrained devices
- Many different operating systems and file systems
 - Linux, Windows, FreeRTOS, Contiki, QNX, ARM MBed OS, TinyOS, RiotOS, VxWorks, etc.
 - ► FAT, ext4, YAFFS, Reliance Edge/Nitro, Coffee, HRFS, LittleFS, etc.
- Contiki-NG, based on Contiki
 - Minimum requirements: 10 kB RAM, 100 kB ROM
 - 6LoWPAN, CoAP, MQTT, RPL
 - ...but not Matter/Thread
 - Used in the Thingsquare platform
 - Have been used by 5% of IoT developers



Introduction - Cooja

- A network simulator
- A device emulator
- Simple to set up using docker image
- Log status messages from nodes
- Can change network during simulation





Coffee File System - Coffee File System

- Designed for flash memory
- Low resource usage
- No \$MFT, FAT table or Superblock
- Have to scan all pages to build up internal view of FS
- File headers contain all information
- FS structures held in RAM, not written to disk
- No folders
- File name is unique identifier



Coffee File System - Coffee File System nomenclature

- ► File system consist of **sectors** which consist of **pages**
- ▶ A whole page must be written, can only **program** bits from "1" to "0"
- ▶ A whole sector must be **erased**, reset all bits to "1"
 - Done during garbage collection
 - **Active:** Pages that have been allocated and belong to existing files.
- **Obsolete:** Pages that have been active but currently not in use. They are available for garbage collection.
 - **Isolated:** Pages that have been active and belonging to a file starting in the previous sector. They exist as file fragments without a header at the start of a sector.
 - **Deleted:** Pages that are obsolete or isolated.
 - **Unused:** Pages that have not been written to after being erased by the garbage collector.



Coffee File System - Coffee File System structures

Size (bytes)	Field
2	LOG_PAGE
2	LOG_RECORDS
2	LOG_RECORD_SIZE
2	MAX_PAGES
1	DEPRECATED_EOF_HINT
1	FLAGS
16/40	NAME

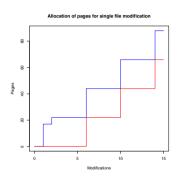
Value	Flag
0x01	HDR_FLAG_VALID
0x02	HDR_FLAG_ALLOCATED
0x04	HDR_FLAG_OBSOLETE
0x08	HDR_FLAG_MODIFIED
0x10	HDR_FLAG_LOG
0x20	HDR_FLAG_ISOLATED

00032600	37	02	00	00	00	00	11	00	00	Of	66	69	6c	65	30	30	7file00
00032610	33	2e	74	78	74	00	00	00	00	00	46	69	6с	65	33	20	3.txtFile3
00033700	00	00	00	00	00	00	05	00	00	17	66	69	6c	65	30	30	file00
00033710	33	2e	74	78	74	00	00	00	00	00	05	00	05	00	03	00	3.txt
00033720	03	00	46	69	6c	65	33	20	20	42	23	30	30	30	30	30	File3 B#00000



Coffee File System - File creation and modification

- ▶ Upon creation 17 pages are allocated for file
- Appends just added to end, as no already written bits are changed
- File size doubled when in need of more pages
- Any other modification adds a 5-page logfile
 - Contain a list indicating which pages are modified
 - A list of the modified pages
- New file/logfile created when log is full
 - Last version copied to new file, changes to log



Coffee File System - Coffee file system dynamics

- Files are written end-to-end
- Extended wear-leveling
- When allocation fails, garbage collection starts
 - No consecutive pages to allocate
- No reshuffling of pages to free up more sectors
- Deleted pages in sectors with live files survive
- Pages belonging to deleted files at start of sector are "isolated pages"
 - No information about header



Coffee File System - File system tests

- File operations tested in four tests: two for file writing and two for GC
- ► File writing testing file creation, deletion, appends and changes
 - Appends don't cause any log file to be written
 - Exceeding allocated length doubles allocation size
 - One writing operation to one page often caused two writes to log file
 - Last version of a base+log file compound is copied to the next base file
- Garbage collection
 - Garbage collection starts when file reaches end of file system
 - No sorting and rearranging of pages
 - ▶ When all sectors contain active pages, FS can't free up any sectors
 - ▶ Huge files spanning more than one sector no different from small files



Recreating file version history - File version history

- Files in same sector is written consecutively
- Files in different sectors can be from different GC runs
- ► Two approaches:
 - ► File content differences
 - Sector offset analysis
- Both methods have their limitations



Recreating file version history - File offset analysis

- ► Files are written consecutively, sector size is in general not divisible by the file size
- Assumption that all files have the same length
- ▶ A sequence of offsets for the first file in each consecutive sectors

Offset of first file

$$A_0 = 0 \tag{1}$$

$$A_n = \left[\left(A_{n-1} + \left\lceil \frac{S_S}{S_F} \right\rceil \right) \pmod{S_S} \right] \pmod{S_F}$$
 (2)

▶ Setting S_S to 256 and F_S to 22:

$$0 \rightarrow 8 \rightarrow 16 \rightarrow 2 \rightarrow 10 \rightarrow 18 \rightarrow 4 \rightarrow 12 \rightarrow 20 \rightarrow 6 \rightarrow 14 \rightarrow 0$$

Recreating file version history - File offset analysis results

- Created a mote that wrote two files: 1 s/w, 100 s/w
- Worked for most part...
- After GC, offset start at 0
- Assumption on same file size don't hold
- If file size is even, why an odd offset?

Starting file offsets

starting me onse								
Sector	First file							
0x00	0x00							
0x01	80x0							
0x06	0x04							
0x0d	0x01							



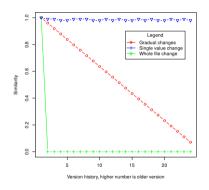
Recreating file version history - File content analysis

- For gradually changing files
 - Ring buffers
- Not good for files where the same part(s) of the file changes every time
- Several existing diff algorithms
 - Levenstein distance/ Edit distance
 - Myers' diff algorithm
- Knowledge about the writing pattern for selecting comparison algorithm
- In this study we used Myers diff algorithm



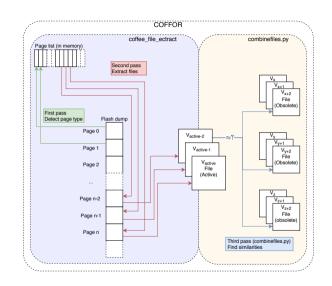
Recreating file version history - File content analysis results

- Testing 3 types of file changes between versions:
 - 1. New location changed
 - 2. A single value changed
 - 3. Whole file changed
- Only the gradually changed file can be reconstructed
- Ordering not dependent on distance between versions
 - ...Until whole file has changed



COFFOR - COFFOR

- COFfee FORensics
- 2 passes:
 - 1. Page type detection
 - Active/del base/log
 - ► File fragment
 - Zero-page
 - Isolated page
 - 2. File extraction
 - name_A/D+page_ver
- Comparing files:
 - Comparing last versions w/ given version
 - Radare2's radiff



COFFOR - Results

- All pages are correctly classified by COFFOR
- ► The file versions are correctly saved
- Should be tested on a wider variety of flash images
- radiff works well
- But only for a subset of files



Final thoughts - Conclusion and way forward

- Wrap-up
 - The artifacts of the Coffee File System
 - File version history ordering
 - COFFOR tool
- Do we really need to collect data from the devices?
- Can the version history be used for other types of data?
- https://github.com/jenspets/coffor



Thank you for your attention

