

Hoare Logic: An Introduction

What is Hoare Logic?

- ▶ A formal system for reasoning about program correctness.
- ▶ Uses **Hoare triples** of the form

$$\{P\}C\{Q\}$$

where:

- ▶ P is the **precondition** (what must be true before execution).
- ▶ C is the **command** (the program to execute).
- ▶ Q is the **postcondition** (what is guaranteed after execution).

Hoare Triples: Meaning

$$\{x = 2\}(x := x + 1)\{x = 3\}$$

or

$$\begin{array}{c} \{x = 2\} \\ x := x + 1 \\ \{x = 3\} \end{array}$$

- ▶ Before executing $x := x + 1$, x is 2.
- ▶ After execution, x is guaranteed to be 3.

The Axioms and Rules of Hoare Logic

Assignment Rule

If P holds after replacing x with E , then P holds after assignment.

$$\{P[x := E]\} x := E \{P\}$$

Example

$$\{y + 1 = 5\} x := y + 1 \{x = 5\}$$

The Rules (Continued)

Consequence Rule

If P' implies P , and Q implies Q' , then the triple is valid.

$$\frac{\{P\}C\{Q\}, P' \rightarrow P, Q \rightarrow Q'}{\{P'\}C\{Q'\}}$$

Composition Rule

If $C1$ ensures Q and $C2$ ensures R , then their sequence ensures R .

$$\frac{\{P\}C1\{Q\}, \{Q\}C2\{R\}}{\{P\}C1; C2\{R\}}$$

Hoare Logic for Conditionals

If-Else Rule

The postcondition Q holds regardless of which branch executes.

$$\frac{\{P \wedge B\} C1 \{Q\}, \{P \wedge \neg B\} C2 \{Q\}}{\{P\} \text{if } B \text{ then } C1 \text{ else } C2 \{Q\}}$$

Example:

$$\begin{array}{c} \{x > 0\} \\ \text{if } x > 0 \text{ then } y := x \text{ else } y := 0 \\ \{y \geq 0\} \end{array}$$

Hoare Logic for Loops

While Loop Rule

$$\frac{\{P \wedge B\} C \{P\}}{\{P\} \text{while } B \text{ do } C \{P \wedge \neg B\}}$$

- ▶ **Loop invariant** P remains true before and after each iteration.
- ▶ Ensures correctness of loops by proving invariance.

Example: Proving a Simple Loop

Program:

```
{ x = n, y = 0 }  
while y < n do  
  y := y + 1;  
  x := x - 1  
{ y = n }
```

Loop Invariant:

```
{ x + y = n }
```

- ▶ Initially true ($y = 0, x = n \rightarrow x + y = n$).
- ▶ Maintained after each iteration.
- ▶ When loop exits ($x = 0, y = n$) holds.

Summary

- ▶ **Hoare Logic** allows formal reasoning about program correctness.
- ▶ Uses **Hoare triples** $\{P\}C\{Q\}$ to specify behavior.
- ▶ Key rules:
 - ▶ Assignment
 - ▶ Composition
 - ▶ Conditionals
 - ▶ Loops
- ▶ Loop invariants are crucial for proving correctness.