



## Heuristic Decisions in Static Analysis



**Astrée** **DOOP** **TAJS** **SAFE**

- Practical static analyzers involve many heuristics
  - Which procedures should be analyzed context-sensitively?
  - Which relationships between variables should be tracked?
  - Which program parts to analyze unsoundly or soundly?, etc
- Designing a good heuristic is an art
  - Usually done by trials and error: nontrivial and suboptimal

## Automatically Generating Heuristics from Data

- Automate the process: use data to make heuristic decisions in static analysis
- Automatic:** little reliance on analysis designers
- Powerful:** machine-tuning outperforms hand-tuning
- Stable:** can be generated for arbitrary programs

## Selective Context-Sensitivity

```

1 class D{} class E{}
2 class C{
3   void dummy(){}
4   Object id1(Object v){return id2(v);}//4
5   Object id2(Object v){return v;}
6 class B{
7   void m(){
8     C c = new C();
9     D d = (D)c.id1(new D());//query1 //9
10    E e = (E)c.id1(new E());//query2 //10
11    c(dummy());}//11
12 public class A{
13   public static void main(String[] args){
14     B b = new B();
15     b.m();}//15
16   b.m();}//16
  
```

- Context-insensitivity fails to prove the queries
  - 2-call-site-sensitivity succeeds but not scale
- solution**

Apply 2-call-sens:{C.id2}  
 Apply 1-call-sens:{C.id1}  
 Apply insens:{B.m, C.dummy}

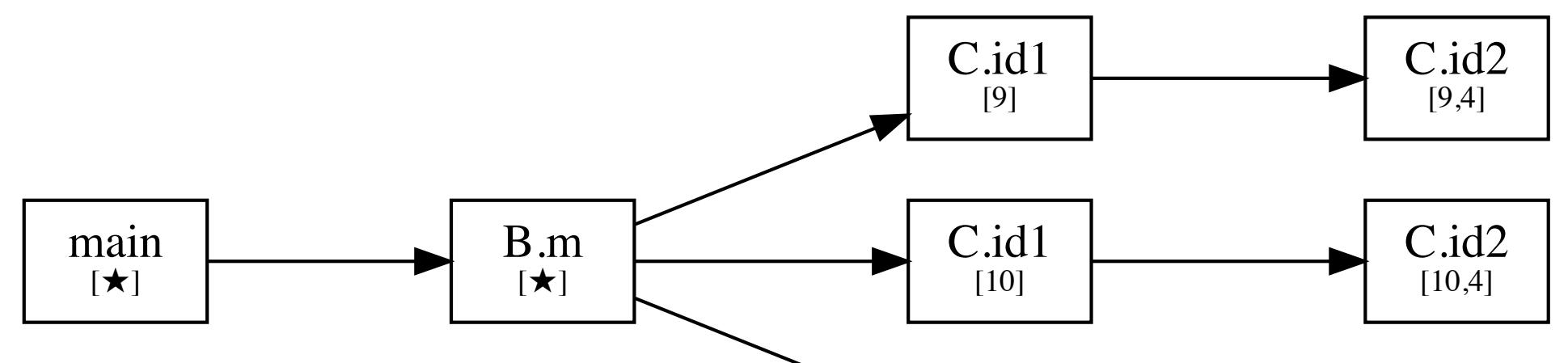


Figure 1: call graph of the solution

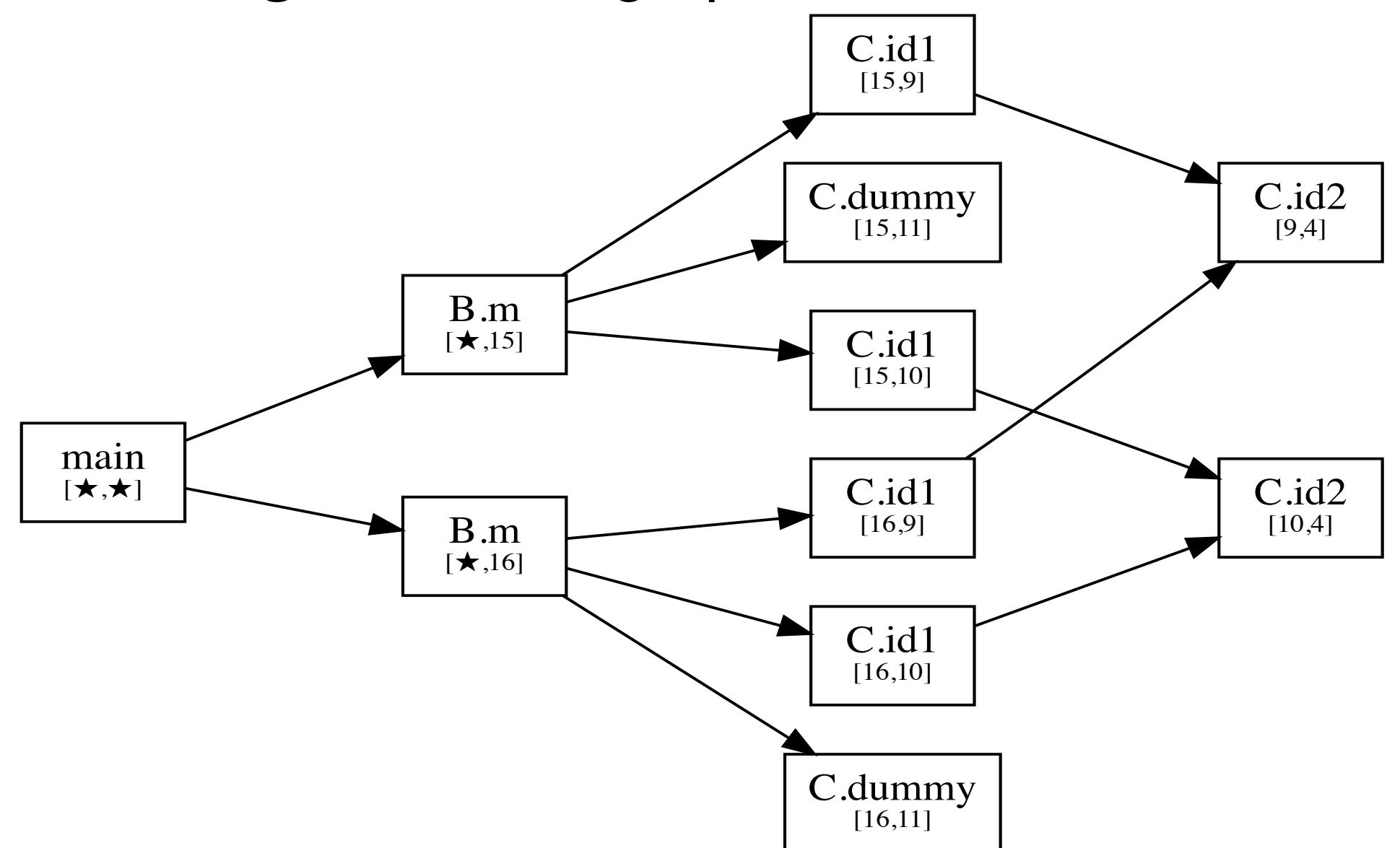
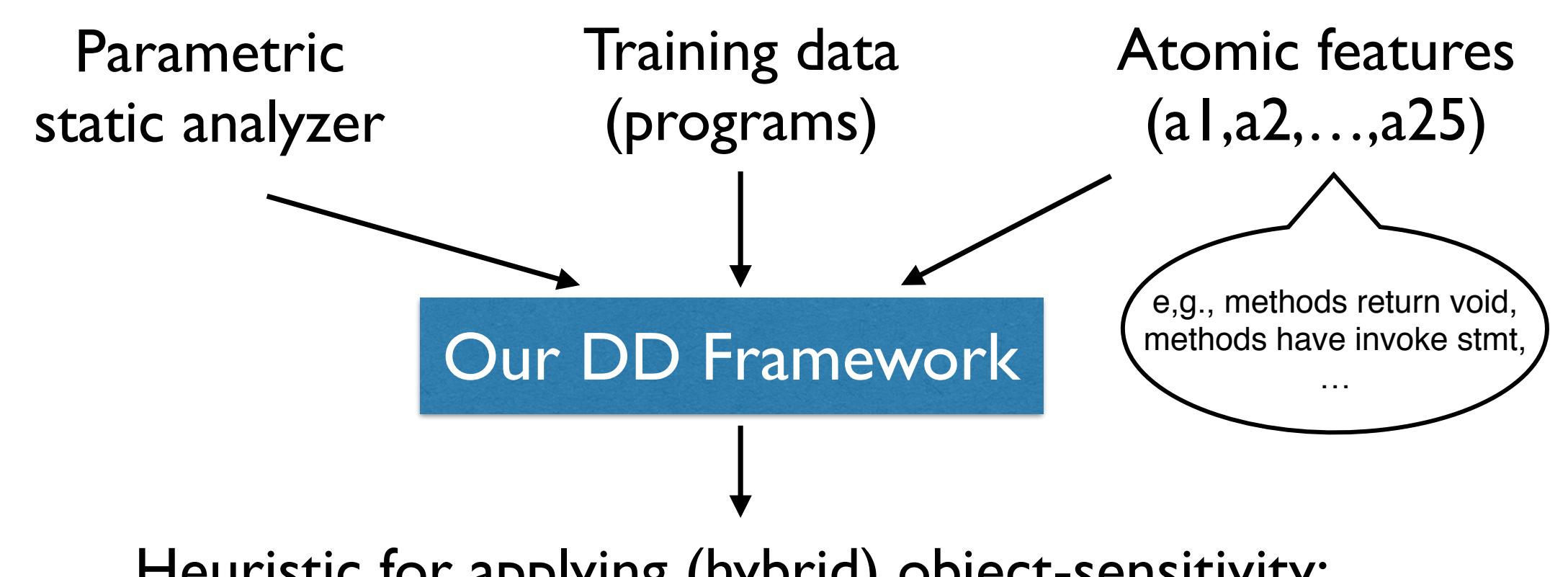


Figure 2: call graph of 2-call-site sensitive

### Challenge: How to decide?

⇒ Data-Driven approach

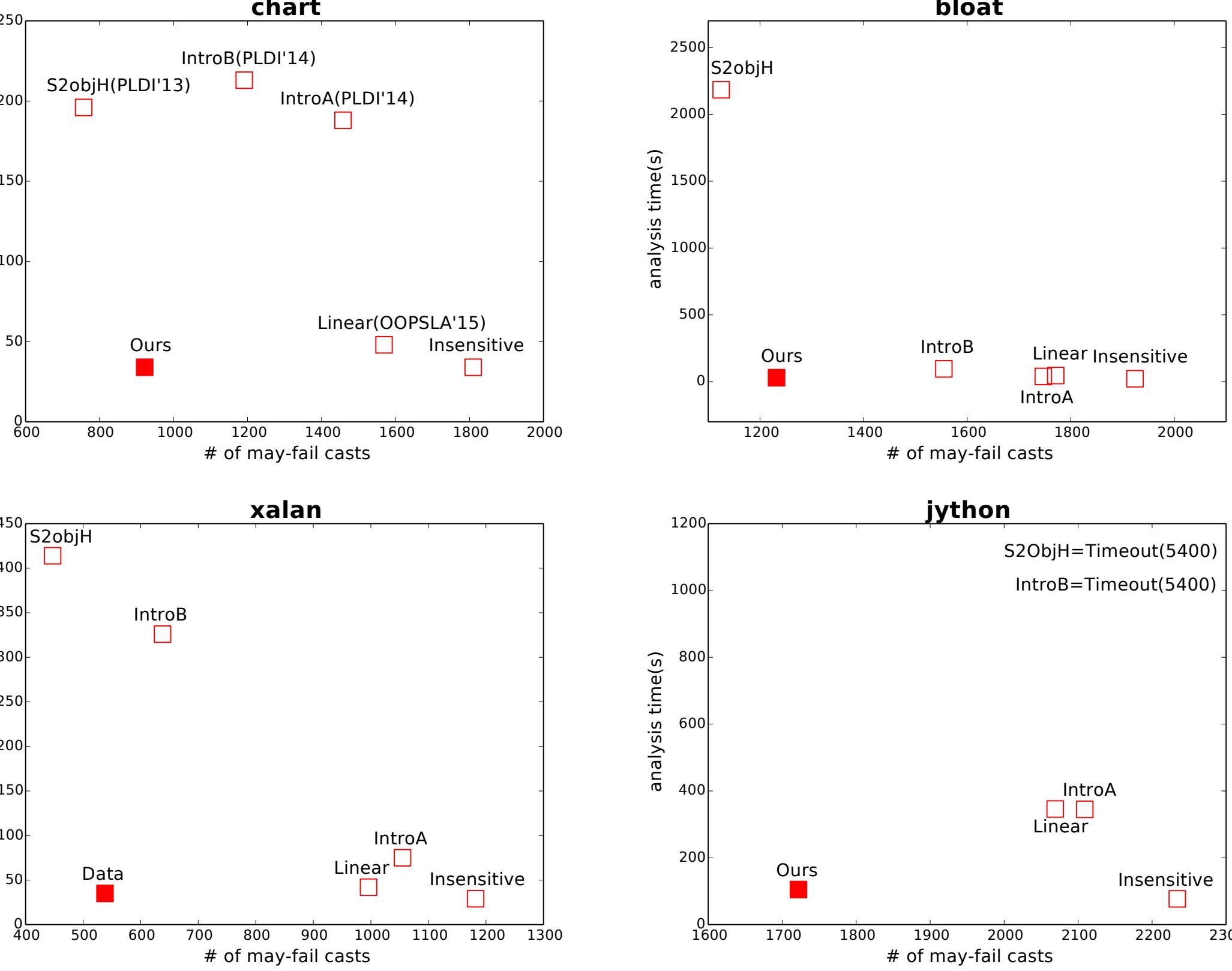
## Data-Driven Ctx-Sensitivity



I2: Methods that require 2-object-sensitivity  
 1  $\wedge$   $\neg$ 3  $\wedge$   $\neg$ 6  $\wedge$  8  $\wedge$   $\neg$ 9  $\wedge$   $\neg$ 16  $\wedge$   $\neg$ 17  $\wedge$   $\neg$ 18  $\wedge$   $\neg$ 19  $\wedge$   $\neg$ 20  $\wedge$   $\neg$ 21  $\wedge$   $\neg$ 22  $\wedge$   $\neg$ 23  $\wedge$   $\neg$ 24  $\wedge$   $\neg$ 25  
 I1: Methods that require 1-object-sensitivity  
 $(1 \wedge \neg 3 \wedge \neg 4 \wedge \neg 7 \wedge \neg 8 \wedge \neg 9 \wedge \neg 15 \wedge \neg 16 \wedge \neg 17 \wedge \neg 18 \wedge \neg 19 \wedge \neg 20 \wedge \neg 21 \wedge \neg 22 \wedge \neg 23 \wedge \neg 24 \wedge \neg 25) \vee$   
 $(\neg 3 \wedge \neg 4 \wedge \neg 7 \wedge \neg 8 \wedge \neg 9 \wedge \neg 10 \wedge \neg 11 \wedge \neg 12 \wedge \neg 13 \wedge \neg 16 \wedge \neg 17 \wedge \neg 18 \wedge \neg 19 \wedge \neg 20 \wedge \neg 21 \wedge \neg 22 \wedge \neg 23 \wedge \neg 24 \wedge \neg 25) \vee$   
 $(\neg 3 \wedge \neg 9 \wedge \neg 13 \wedge \neg 14 \wedge \neg 16 \wedge \neg 17 \wedge \neg 18 \wedge \neg 19 \wedge \neg 20 \wedge \neg 21 \wedge \neg 22 \wedge \neg 23 \wedge \neg 24 \wedge \neg 25) \vee$   
 $(1 \wedge \neg 2 \wedge \neg 3 \wedge \neg 4 \wedge \neg 5 \wedge \neg 6 \wedge \neg 7 \wedge \neg 8 \wedge \neg 9 \wedge \neg 10 \wedge \neg 13 \wedge \neg 15 \wedge \neg 16 \wedge \neg 17 \wedge \neg 18 \wedge \neg 19 \wedge \neg 20 \wedge \neg 21 \wedge \neg 22 \wedge \neg 23 \wedge \neg 24 \wedge \neg 25)$

## Performance

- Training with 4 small programs from DaCapo, and applied to 6 large programs
- Machine-tuning outperforms hand-tuning



## Key Contributions

We achieve the improvement with two key ideas.

- A new expressive model(Disjunctive Model)
- Learning algorithm for new model

## Disjunctive Model

Disjunctive model expresses set with DNF form.

- Method : Features      Goal = { $M_1, M_4$ }  
 $M_1 : a_1 a_2$       Disjunctive Model(Possible):  
 $M_2 : a_1$        $(a_1 \wedge a_2) \vee (\neg a_1 \wedge \neg a_2)$   
 $M_3 : a_2$       Linear Model(Impossible):  
 $M_4 :$        $C_1 * a_1 + C_2 * a_2$

Figure 3: Disjunctive vs Linear

With { $a_1, a_2$ }, Disjunctive model can express the target methods, but Linear model cannot.

## Learning Algorithm

Let  $\Pi = \{f_1, \dots, f_k\}$  be parameters. Each  $f_i$  expresses methods to be assigned with depth  $i$ . We assign deeper depth if a method is in both  $f_i$  and  $f_j$  ( $i \neq j$ ). We learn  $\Pi$  by solving the following problem.

### Optimization problem

Find parameter  $\Pi = \langle f_1, \dots, f_k \rangle$  that minimizes the cost of analysis while satisfying precision constraint over training set.

### Challenge

Assuming that  $|S|$  is the space of possible boolean formulas over which we learn, search space of original problem is  $|S|^k$ . We reduce the search space into  $k * |S|$  by decomposing the original problem into  $k$  subproblems( $\Psi_1 \sim \Psi_k$ ). Each  $f_i$  is obtained from  $\Psi_k$  to  $\Psi_1$ .

### Decomposed problem $\Psi_i$

Let  $\Pi = \langle \text{true}, \text{true}, \dots, \text{true}, f_i, f_{i+1}, \dots, f_k \rangle$ . Find formula  $f_i$  that makes  $\Pi$  minimize the cost while satisfying precision constraint over training set.

### Learning Algorithm for $\Psi_i$

To solve  $\Psi_i$ , we made a greedy algorithm. Let  $\{a_1, \dots, a_n\}$  be atomic features. Our algorithm proceeds in the following steps:

- $f_i$  starts from disjunctions of  $2n$  clauses :  
 $f_i = a_1 \vee \neg a_1 \vee \dots \vee a_n \vee \neg a_n$
- Choose the most expensive clause  $c_i$  to refine.
- Strengthen the clause  $c_i$  by conjoining an decent atom  $a_k$  with  $c_i$ :  $f'_i = c_1 \vee \dots \vee (c_i \wedge a_k) \vee \dots \vee c_j$ .
- Check if  $f'_i$  satisfies precision constraint. If it is,  $f_i = f'_i$ .
- Repeat 2~4 until  $f_i$  cannot be refined.