Haskell: Assignment 1

Notes:

- For all of the problems below, you *must* write a type signature for any function/value you write in your code.
- You should turn in the **first two problems** on Canvas in a .hs file that has the functions you define for each part.
- You should *neatly write your answer* to **Problem 3** on a sheet of paper and turn it in on TODO:when??
- You can test your code yourself by loading your .hs file into ghci.

Problem 1 – Factorial

For each of the following sub-problems, you are going to write a factorial function that meets a specific requirement. The following is relevant to both parts:

- The input to the function is an Int, n
- The output of the function is the mathematical factorial of the input, n!, also an Int
- You do not have to handle negative input values assume the input is 0 or greater.

a. Guard Expression

Write the factorial function using a *guard expression*. You may want to refer to the fibonacci function fib in the lecture notes.

b. Pattern Matching

Write the factorial function using pattern matching (and no guards). You may want to refer to the second fibonacci function fib' from the lecture notes.

Problem 2 - Maybe

Haskell does not have an untyped empty value like NULL in C++, nil in Ruby, None in Python, null in Java, etc. Instead, Haskell uses data types and polymorphism to reach a similar end. The type used for this is Maybe, which is defined as:

```
data Maybe a = Nothing | Just a
```

Nothing is Haskell's equivalent of NULL. But this type definition looks kind of funny, primarily because of the a.

Recall Haskell's general list type, [a]. The a is a type variable, so one way to read [a] is "a list of elements of any type". The a in Maybe a is also a type variable. So, just as we can have a list of Ints with the type [Int], we can also have a Maybe Int.

Now turning to Maybe's constructors, we have Nothing and Just a.

- Nothing is simply a constructor with no arguments by default, its typing binding is Nothing :: Maybe a, but within a particular context it may be Nothing :: Maybe Int, Nothing :: Maybe String, etc.
- The Just a constructor tells us that we can use the Just constructor with an *argument of any type* (hence the type variable a).

Let's look at a few examples using Maybe's constructors.

If we bind Nothing to a symbol n and we don't assign a type, like so:

```
n = Nothing
```

Then the compiler will infer n's type to be Maybe a, since it has no way to narrow down the type variable a any further.

We could also force ${\tt n}$ to be of a more specific type:

```
n :: Maybe Int
n = Nothing

or
n = Nothing :: Maybe Int
```

The compiler can infer a bit more about a binding that uses Just:

```
j = Just "hi"
```

In this case, Haskell will infer that j's type binding must be j:: Maybe String, since we provided a String argument to the Just data constructor.

To see where Maybe might be useful, consider the function head :: [a] -> a, which returns the first element of a list. What if the input list is empty? head

wouldn't be able to return a value of the expected type. In fact, if you called head [] in ghci, you'd get an error.

Now imagine a function headMaybe with the type binding:

```
headMaybe :: [a] -> Maybe a
```

Example function calls:

```
headMaybe [1, 2, 3]
=> Just 1
headMaybe ['a', 'b', 'c']
=> Just 'a'
headMaybe []
=> Nothing
```

Since the return value is Maybe a, we can return a Just <value> whenever the list is non-empty, and when it is empty we simply return Nothing.

Define the function headMaybe so that it behaves as described above.

Hints:

- All you should use on the right-hand side (RHS) of your definition are the constructors for Maybe as well as the head function mentioned above.
- You'll have two cases to define this function for: the empty list, and all other lists. To pattern match against the empty list, you can use headMaybe [] = ...

Problem 3 – Reduction

Recall the Peano number data type example from lecture:

```
data Peano = Zero | Succ Peano
    deriving Show
```

In lecture, we wrote a particular definition for a function add that added two Peanos together. Our definition was longer than it needed to be, though. Here's a partial definition for a different form of the add function:

```
add :: Peano -> Peano -> Peano
add Zero p = p
```

- a. Add *one* more case to the add function's definition that completes the definition and will successful add all Peano numbers. Write the line below. *Hint: Think about the two fundamental cases in a recursive function.*
- b. Using your completed definition of add, write out the reduction steps for the expression add two one, where two = Succ (Succ Zero) and one = Succ Zero. Be sure to define your reduction rules before you do the reduction steps (there should be 4 rules in this case).

Rules:

Steps: