

Data Types

Note: You can ignore the lines after each data declaration that say `deriving ...` – these are there to make it easier for you to play with the examples. If you're curious, though, those lines have to do with type classes.

Simple Data Declaration

Basic data type declaration:

```
data Bool' = True' | False'
    deriving (Show, Eq)
```

- `Bool'` is the name of our data type
- `True'` and `False'` are the *data constructors*

Note: `Bool` is already defined in Haskell, which is why the example codes adds a `'` to all of the types/constructors/functions in this section.

Then we can use the constructors to create instances of this type:

```
true :: Bool'
true = True'

false :: Bool'
false = False'
```

Let's write a few of the functions that are common for booleans:

1. `not'`
2. `and'`, `or'`
3. `xor'`

`not`

```
not' :: Bool' -> Bool'
not' True' = False'
not' False' = True'
```

`and`

First pass:

```
and' :: Bool' -> Bool' -> Bool'
and' False' False' = False'
and' False' True' = False'
```

```
and' True' False' = False'
and' True' True' = True'
```

Better:

```
and' :: Bool' -> Bool' -> Bool'
and' True' True' = True'
and' _ _ = False'
```

or

First pass:

```
or' :: Bool' -> Bool' -> Bool'
or' False' False' = False'
or' False' True' = True'
or' True' False' = True'
or' True' True' = True'
```

Better:

```
or' :: Bool' -> Bool' -> Bool'
or' False' False' = False'
or' _ _ = True'
```

xor

First pass:

```
xor' :: Bool' -> Bool' -> Bool'
xor' False' False' = True'
xor' False' True' = False'
xor' True' False' = False'
xor' True' True' = True'
```

```
xor' :: Bool' -> Bool' -> Bool'
xor' False' False' = True'
xor' True' True' = True'
xor' _ _ = False'
```

We can reduce this by one more line (of actual logic), but not with pattern matching (assumes we had a == equality comparison function for `Bool'`s):

```
xor' :: Bool' -> Bool' -> Bool'
xor' b1 b2
  | b1 == b2 = True'
  | otherwise = False'
```

Purpose of this example:

- Pattern matching can *only* compare *concrete values of a given input*
- Pattern matching can be very concise, but is limited in its application

Data Type with Parameters

To add parameters to a data constructor, simply list each of the *types* of the parameters after the constructor name:

```
-- Person has Name and Age
data Person = Person String Int
```

Note that the *data constructor* (RHS) may have the same name as the *data type* (LHS).

`String` and `Int` are vague, though. Let's clarify by declaring a couple of *type synonyms*.

```
data Person = Person Name Age
    deriving Show
type Name = String
type Age = Int
```

Then we can create a `Person`:

```
bigL = Person "Lamont Coleman" 24
```

Question

What if we wanted to pattern match on a `Person`?

Answer: We do the same thing we did with other data types: we pattern match against *concrete values of the correct type*.

```
bob = Person "Bob" 50

trueIfBob50 :: Person -> Bool
trueIfBob50 (Person "Bob" 50) = True
trueIfBob50 (Person _    _    ) = False

trueIfBob :: Person -> Bool
trueIfBob (Person "Bob" _) = True
trueIfBob (Person _    _) = False

trueIfOld :: Person -> Bool
trueIfOld (Person _ age) = age > 24
```

```

trueIfOldBob :: Person -> Bool
trueIfOldBob (Person "Bob" age) = age > 24
trueIfOldBob (Person _ _ ) = False

-- `head` returns first element of string/list
firstInitial :: Person -> Char
firstInitial (Person name _) = head name

```

Note: Whenever we don't care about a particular input value, we use an `_` to say "I don't care what the input value to this parameter is, and I'm not going to use it on the RHS of the function so don't bother giving it a name".

Tuples vs. Data Constructors

Note that we could also use a tuple to make a single data type out of multiple values. For example, we could define a `Point` as a type synonym of a tuple of `Float`s:

```

type Point = (Float, Float)

point :: Point
point = (1.1, 2.4)

```

The data constructor analog would look like:

```

data Point = Point Float Float
    deriving Show

point :: Point
point = Point 1.1 2.4

```

The difference comes down to judgement and specific cases, but here are the main thoughts to keep in your mind:

- A *type synonym* is simply a different name for a type, so the synonym is interchangeable with the original type.
- A good rule of thumb is to make anything in your code that is especially meaningful/important into its own data type (rather than a synonym).

Type Synonym Example

Let's look at a case where a type synonym fits very well. Imagine we had the `Point` data type we defined above, along with these functions:

```

getX :: Point -> Float
getX (Point x _) = x

```

```
getY :: Point -> Float
getY (Point _ y) = y
```

Now let's say we want to change the `Float` arguments in the `Point` constructor to be `Ints` instead:

```
data Point = Point Int Int
```

However, this breaks our `getX` and `getY` functions, because they each try to extract a `Float` from a `Point`, *not* an `Int`.

A type synonym can help us here. Let's define the type `Coordinate` to mean the same thing as `Float`:

```
data Point = Point Coordinate Coordinate

type Coordinate = Float
```

Now our `Point` constructor accepts two `Coordinates`. Our `getX` and `getY` now look like:

```
getX, getY :: Point -> Coordinate
getX (Point x _) = x
getY (Point _ y) = y
```

If we came back later wanted to change the `Coordinate` type to be `Int` instead of `Float`, we would only have to change a single line:

```
type Coordinate = Int
```

And since the types of the `Point` constructor and the `getX`, `getY` functions refer to `Coordinate`, we don't have to change anything else and everything should still work.

This example was somewhat contrived though, as there very well may be a part of the code somewhere that is dependent on `Float` and changing the meaning of `Coordinate` to `Int` might break that. A very simple case:

```
point = Point 3.4 5.6
```

If we tried to declare this `point` value but `Point` was expecting two `Int` coordinates, then our code wouldn't compile.

Data Constructor Types

Like most of what we've dealt with in Haskell, data constructors have types.

Question

- a. What is the type of `True'`?

Answer: `Bool'`

(We can verify in `ghci` with `:t True'`.)

- b. What is the type of the *constructor* `Person`?

Consider how we created a `Person`:

```
bob = Person "Bob" 50
```

Answer: `Person :: Name -> Age -> Person`

This looks a lot like a function!

Recursive Data Types: Peano Numbers

The *peano numbers* are a way to represent integers using only a zero value and a successor function.

```
data Peano = Zero | Succ Peano
```

We have two data constructors here

- `Zero`, which has no arguments
- `Succ`, which takes a single argument of type `Peano`
 - This makes `Peano` a *recursively-defined data type*

Question

Given:

```
zero = Zero
```

- a. How do we create a value `one` that is the successor of `Zero`? How about `two`?

Answer:

```
one = Succ zero  
two = Succ one
```

- b. How else could we constructor `two`?

```
two = Succ (Succ zero)  
-- or  
two = Succ (Succ Zero)
```

Question

How can we (exhaustively) convert these into the numbers we're used to? Start with the type signature: `peanoToInt :: Peano -> Int`

Answer:

```
peanoToInt :: Peano -> Int
peanoToInt Zero = 0
peanoToInt (Succ Zero) = 1
peanoToInt (Succ (Succ Zero)) = 2
-- this could go on awhile...
```

Increment/decrement functions

Goal: Function `increment` that returns the successor of the input `Peano`.

```
increment :: Peano -> Peano
increment p = Succ p
```

Goal: Function `decrement` that returns the predecessor of the input `Peano` number. The predecessor of `Zero` is still `Zero` (we have no way to represent negatives).

```
decrement :: Peano -> Peano
decrement (Succ p) = p
decrement Zero    = Zero
```

add function for Peano numbers

Let's try something a little more fun.

Goal: Function `add` that returns the sum of two `Peanos`.

Question

- What should the type signature be?

```
add :: Peano -> Peano -> Peano
```

- What cases do we have for this function?

First pass:

```
add Zero Zero = undefined
add Zero (Succ p) = undefined
add (Succ p) Zero = undefined
add (Succ p1) (Succ p2) = undefined
```

How might we implement the bodies of these?

```
add Zero Zero = Zero
add Zero (Succ p) = Succ p
add (Succ p) Zero = Succ p
add (Succ p1) (Succ p2) = Succ (Succ (add p1 p2))
```

We can simplify that form into:

```
add Zero p = Zero
add p Zero = Zero
add (Succ p1) (Succ p2) = Succ (Succ (add p1 p2))
```