

## DAY 9

1) There are  $3n$  piles of coins of varying size, you and your friends will take piles of coins as follows: In each step, you will choose any 3 piles of coins (not necessarily consecutive). Of your choice, Alice will pick the pile with the maximum number of coins. You will pick the next pile with the maximum number of coins. Your friend Bob will pick the last pile. Repeat until there are no more piles of coins. Given an array of integers `piles` where `piles[i]` is the number of coins in the  $i$ th pile. Return the maximum number of coins that you can have.

**Example 1:**

**Input:** `piles = [2,4,1,2,7,8]`

**Output:** 9

### CODE:

```
def maxCoins(piles):
    piles.sort(reverse=True)
    max_coins = 0
    for i in range(1, len(piles) * 2 // 3, 2):
        max_coins += piles[i]
    return max_coins

piles = [2, 4, 1, 2, 7, 8]
print(maxCoins(piles))
```

### **OUTPUT:**

9

2) You are given a 0-indexed integer array `coins`, representing the values of the coins available, and an integer `target`. An integer `x` is obtainable if there exists a subsequence of coins that sums to `x`. Return the minimum number of coins of any value that need to be added to the array so that every integer in the range `[1, target]` is obtainable. A subsequence of an array is a new non-empty array that is formed from the original array by deleting some (possibly none) of the elements without disturbing the relative positions of the remaining elements.

**Example 1:**

**Input:** `coins = [1,4,10]`, `target = 19`

**Output:** 2

### **CODE:**

```
def min_coins_to_reach_target(coins, target):
    coins.sort()
    current_max = 0
    count_added = 0
    i = 0
    while current_max < target:
        if i < len(coins) and coins[i] <= current_max + 1:
            current_max += coins[i]
            i += 1
        else:
            current_max += (current_max + 1)
            count_added += 1
    return count_added

coins = [1, 4, 10]
target = 19
result = min_coins_to_reach_target(coins, target)
print(f"Output: {result}")
```

### **OUTPUT:**

Minimum number of coins to be added: 2

**3) You are given an integer array jobs, where jobs[i] is the amount of time it takes to complete the ith job. There are k workers that you can assign jobs to. Each job should be assigned to exactly one worker. The working time of a worker is the sum of the time it takes to complete all jobs assigned to them. Your goal is to devise an optimal assignment such that the maximum working time of any worker is minimized. Return the minimum possible maximum working time of any assignment.**

**Example 1:**

**Input: jobs = [3,2,3], k = 3**

**Output: 3**

**CODE:**

```
def canAssign(jobs, k, limit):  
    # Array to store the workload of each worker  
    workloads = [0] * k  
  
    def backtrack(i):  
        if i == len(jobs):  
            return True  
  
        for j in range(k):  
            if workloads[j] + jobs[i] <= limit:  
                workloads[j] += jobs[i]  
                if backtrack(i + 1):  
                    return True  
                workloads[j] -= jobs[i]  
  
            if workloads[j] == 0:  
                break  
  
        return False  
  
    return backtrack(0)  
  
def minMaxWorkingTime(jobs, k):  
    jobs.sort(reverse=True)  
    left, right = max(jobs), sum(jobs)  
  
    while left < right:  
        mid = (left + right) // 2  
        if canAssign(jobs, k, mid):  
            right = mid # Try for a smaller possible time  
        else:
```

```
        left = mid + 1 # Increase the working time limit
    return left
jobs = [3, 2, 3]
k = 3
result = minMaxWorkingTime(jobs, k)
print(f"The minimum possible maximum working time is: {result}")
```

**OUTPUT:**

The minimum possible maximum working time is: 3

4) We have n jobs, where every job is scheduled to be done from `startTime[i]` to `endTime[i]`, obtaining a profit of `profit[i]`. You're given the `startTime`, `endTime` and `profit` arrays, return the maximum profit you can take such that there are no two jobs in the subset with overlapping time range. If you choose a job that ends at time X you will be able to start another job that starts at time X.

**Example 1:**

**Input:** `startTime = [1,2,3,3]`, `endTime = [3,4,5,6]`, `profit = [50,10,40,70]`

**Output:** 120

**Explanation:** The subset chosen is the first and fourth job.

### **CODE:**

```
from bisect import bisect_right

def jobScheduling(startTime, endTime, profit):
    # Combine start time, end time, and profit into a single list of jobs
    jobs = sorted(zip(startTime, endTime, profit), key=lambda x: x[1])
    dp = [0] * len(jobs)
    start = [job[0] for job in jobs]
    end = [job[1] for job in jobs]
    profit = [job[2] for job in jobs]
    dp[0] = profit[0]
    for i in range(1, len(jobs)):
        last_non_conflicting = bisect_right(end, start[i]) - 1
        include_profit = profit[i]
        if last_non_conflicting != -1:
            include_profit += dp[last_non_conflicting]
        dp[i] = max(dp[i-1], include_profit)
    return dp[-1]

startTime = [1, 2, 3, 3]
endTime = [3, 4, 5, 6]
profit = [50, 10, 40, 70]
result = jobScheduling(startTime, endTime, profit)
print(f"The maximum profit is: {result}")
```

### **OUTPUT:**

The maximum profit is: 120

**5) Given a graph represented by an adjacency matrix, implement Dijkstra's Algorithm to find the shortest path from a given source vertex to all other vertices in the graph. The graph is represented as an adjacency matrix where `graph[i][j]` denote the weight of the edge from vertex `i` to vertex `j`. If there is no edge between vertices `i` and `j`, the value is Infinity (or a very large number).**

**Test Case 1:**

**Input:**

**`n = 5`**

**`graph = [[0, 10, 3, Infinity, Infinity], [Infinity, 0, 1, 2, Infinity], [Infinity, 4, 0, 8, 2], [Infinity, Infinity, Infinity, 0, 7], [Infinity, Infinity, Infinity, 9, 0]]`**

**`source = 0`**

**Output: [0, 7, 3, 9, 5]**

### **CODE:**

```
import heapq
```

```
def dijkstra(graph, source):
```

```
    n = len(graph)
```

```
    dist = [float('inf')] * n
```

```
    dist[source] = 0
```

```
    pq = [(0, source)]
```

```
    while pq:
```

```
        current_dist, u = heapq.heappop(pq)
```

```
        if current_dist > dist[u]:
```

```
            continue
```

```
            for v in range(n):
```

```
                if graph[u][v] != float('inf'): # If there is an edge from u to v
```

```
                    distance = current_dist + graph[u][v]
```

```
                    if distance < dist[v]:
```

```
                        dist[v] = distance
```

```
                        heapq.heappush(pq, (distance, v)) # Push the new distance to the queue
```

```
    return dist
```

```
n = 5
```

```
graph = [
```

```
    [0, 10, 3, float('inf'), float('inf')],
```

```
    [float('inf'), 0, 1, 2, float('inf')],
```

```
    [float('inf'), 4, 0, 8, 2],
```

```
    [float('inf'), float('inf'), float('inf'), 0, 7],
```

```
[float('inf'), float('inf'), float('inf'), 9, 0]
]
source = 0
shortest_paths = dijkstra(graph, source)
print(f"The shortest distances from vertex {source} are: {shortest_paths}")
```

**OUTPUT:**

The shortest distances from vertex 0 are: [0, 7, 3, 9, 5]

6) Given a graph represented by an edge list, implement Dijkstra's Algorithm to find the shortest path from a given source vertex to a target vertex. The graph is represented as a list of edges where each edge is a tuple (u, v, w) representing an edge from vertex u to vertex v with weight w.

**Test Case 1:**

**Input:**

**n = 6**

**edges = [(0, 1, 7), (0, 2, 9), (0, 5, 14), (1, 2, 10), (1, 3, 15),  
(2, 3, 11), (2, 5, 2), (3, 4, 6), (4, 5, 9) ]**

**source = 0**

**target = 4**

**Output: 20**

**CODE:**

```
import heapq

def dijkstra(n, edges, source, target):

    graph = {i: [] for i in range(n)}

    for u, v, w in edges:

        graph[u].append((v, w))

        graph[v].append((u, w)) # Since the graph is undirected

    dist = [float('inf')] * n

    dist[source] = 0

    pq = [(0, source)] # Priority queue stores (distance, vertex)

    while pq:

        current_dist, u = heapq.heappop(pq) # Get the vertex with smallest distance

        if u == target:

            return current_dist # Return distance when the target is reached

        if current_dist > dist[u]:

            continue # If we have already found a shorter path, skip

        for v, weight in graph[u]:

            distance = current_dist + weight

            if distance < dist[v]:

                dist[v] = distance

                heapq.heappush(pq, (distance, v)) # Push the new distance to the queue

    return float('inf')

n = 6

edges = [
```



```
(0, 1, 7), (0, 2, 9), (0, 5, 14),  
(1, 2, 10), (1, 3, 15),  
(2, 3, 11), (2, 5, 2),  
(3, 4, 6),  
(4, 5, 9)  
]  
source = 0  
target = 4  
shortest_path = dijkstra(n, edges, source, target)  
print(f"The shortest path from vertex {source} to vertex {target} is: {shortest_path}")
```

**OUTPUT:**

The shortest path from vertex 0 to vertex 4 is: 20

**7) Given a set of characters and their corresponding frequencies, construct the Huffman Tree and generate the Huffman Codes for each character.**

**Test Case 1:**

**Input:**

**n = 4**

**characters = ['a', 'b', 'c', 'd']**

**frequencies = [5, 9, 12, 13]**

**Output: [('a', '110'), ('b', '10'), ('c', '0'), ('d', '11')]**

**CODE:**

```
import heapq
```

```
class HuffmanNode:
```

```
    def __init__(self, char=None, freq=0):
```

```
        self.char = char
```

```
        self.freq = freq
```

```
        self.left = None
```

```
        self.right = None
```

```
    def __lt__(self, other):
```

```
        return self.freq < other.freq
```

```
def build_huffman_tree(characters, frequencies):
```

```
    heap = []
```

```
    for i in range(len(characters)):
```

```
        node = HuffmanNode(characters[i], frequencies[i])
```

```
        heapq.heappush(heap, node)
```

```
    while len(heap) > 1:
```

```
        left = heapq.heappop(heap)
```

```
        right = heapq.heappop(heap)
```

```
        merged = HuffmanNode(None, left.freq + right.freq)
```

```
        merged.left = left
```

```
        merged.right = right
```

```
        heapq.heappush(heap, merged)
```

```
    return heap[0]
```

```
def generate_huffman_codes(root):
```

```
    codes = {}
```

```
    def _generate_codes(node, current_code):
```

```

    if not node:
        return
        if node.char is not None:
            codes[node.char] = current_code
            _generate_codes(node.left, current_code + '0')
            _generate_codes(node.right, current_code + '1')
        _generate_codes(root, "")
    return codes

def huffman_encoding(characters, frequencies):
    # Step 1: Build Huffman Tree
    huffman_tree_root = build_huffman_tree(characters, frequencies)
    huffman_codes = generate_huffman_codes(huffman_tree_root)
    return huffman_codes

n = 4
characters = ['a', 'b', 'c', 'd']
frequencies = [5, 9, 12, 13]
huffman_codes = huffman_encoding(characters, frequencies)
output = [(char, code) for char, code in huffman_codes.items()]
print(output)

```

## OUTPUT:

```
[('c', '0'), ('b', '10'), ('a', '110'), ('d', '111')]
```

**8) Given a Huffman Tree and a Huffman encoded string, decode the string to get the original message.**

**Test Case 1:**

**Input:**

**n = 4**

**characters = ['a', 'b', 'c', 'd']**

**frequencies = [5, 9, 12, 13]**

**encoded\_string = '1101100111110'**

**Output: "abacd"**

**CODE:**

```
import heapq
```

```
class HuffmanNode:
```

```
    def __init__(self, char=None, freq=0):
```

```
        self.char = char
```

```
        self.freq = freq
```

```
        self.left = None
```

```
        self.right = None
```

```
    def __lt__(self, other):
```

```
        return self.freq < other.freq
```

```
def build_huffman_tree(characters, frequencies):
```

```
    heap = []
```

```
    for i in range(len(characters)):
```

```
        node = HuffmanNode(characters[i], frequencies[i])
```

```
        heapq.heappush(heap, node)
```

```
    while len(heap) > 1:
```

```
        left = heapq.heappop(heap)
```

```
        right = heapq.heappop(heap)
```

```
        merged = HuffmanNode(None, left.freq + right.freq)
```

```
        merged.left = left
```

```
        merged.right = right
```

```
        heapq.heappush(heap, merged)
```

```
    return heap[0] # Return the root of the Huffman Tree
```

```
def decode_huffman(root, encoded_string):
```

```
    decoded_message = []
```

```

current_node = root
for bit in encoded_string:
    if bit == '0':
        current_node = current_node.left
    else:
        current_node = current_node.right
    if current_node.char is not None:
        decoded_message.append(current_node.char)
        current_node = root # Go back to the root for the next set of bits
return ''.join(decoded_message)

def huffman_decoding(characters, frequencies, encoded_string):
    huffman_tree_root = build_huffman_tree(characters, frequencies)
    decoded_message = decode_huffman(huffman_tree_root, encoded_string)
    return decoded_message

n = 4
characters = ['a', 'b', 'c', 'd']
frequencies = [5, 9, 12, 13]
encoded_string = '1101100111110'
decoded_message = huffman_decoding(characters, frequencies, encoded_string)
print(decoded_message)

```

## OUTPUT:

"abacd"

9) Given a list of item weights and the maximum capacity of a container, determine the maximum weight that can be loaded into the container using a greedy approach. The greedy approach should prioritize loading heavier items first until the container reaches its capacity.

**Test Case 1:**

**Input:**

**n = 5**

**weights = [10, 20, 30, 40, 50]**

**max\_capacity = 60**

**Output: 50**

**CODE:**

```
def max_weight(weights, max_capacity):  
    weights.sort(reverse=True)  
    total_weight = 0  
    for weight in weights:  
        if total_weight + weight <= max_capacity:  
            total_weight += weight  
        else:  
            break  
    return total_weight  
  
n = 5  
weights = [10, 20, 30, 40, 50]  
max_capacity = 60  
result = max_weight(weights, max_capacity)  
print(result)
```

**OUTPUT:**

50

**10) Given a list of item weights and a maximum capacity for each container, determine the minimum number of containers required to load all items using a greedy approach. The greedy approach should prioritize loading items into the current container until it is full before moving to the next container.**

**Test Case 1:**

**Input:**

**n = 7**

**weights = [5, 10, 15, 20, 25, 30, 35]**

**max\_capacity = 50**

**Output: 4**

**CODE:**

```
def min_containers(weights, max_capacity)

    weights.sort()

    container_count = 0

    current_capacity = 0

    for weight in weights:

        if current_capacity + weight > max_capacity:

            container_count += 1

            current_capacity = weight

        else:

            current_capacity += weight

    if current_capacity > 0:

        container_count += 1

    return container_count

n = 7

weights = [5, 10, 15, 20, 25, 30, 35]

max_capacity = 50

result = min_containers(weights, max_capacity)

print(result)
```

**OUTPUT:**

**4**

**11) Given a graph represented by an edge list, implement Kruskal's Algorithm to find the Minimum Spanning Tree (MST) and its total weight.**

**Test Case 1:**

**Input:**

**n = 4**

**m = 5**

**edges = [ (0, 1, 10), (0, 2, 6), (0, 3, 5), (1, 3, 15), (2, 3, 4) ]**

**Output:**

**Edges in MST: [(2, 3, 4), (0, 3, 5), (0, 1, 10)]**

**Total weight of MST: 19**

**CODE:**

```
class UnionFind:
```

```
    def __init__(self, n):
```

```
        self.parent = list(range(n))
```

```
        self.rank = [0] * n
```

```
    def find(self, u):
```

```
        if self.parent[u] != u:
```

```
            self.parent[u] = self.find(self.parent[u]) # Path compression
```

```
        return self.parent[u]
```

```
    def union(self, u, v):
```

```
        root_u = self.find(u)
```

```
        root_v = self.find(v)
```

```
        if root_u != root_v:
```

```
            # Union by rank
```

```
            if self.rank[root_u] > self.rank[root_v]:
```

```
                self.parent[root_v] = root_u
```

```
            elif self.rank[root_u] < self.rank[root_v]:
```

```
                self.parent[root_u] = root_v
```

```
            else:
```

```
                self.parent[root_v] = root_u
```

```
                self.rank[root_u] += 1
```

```
            return True
```

```
        return False
```

```
def kruskal(n, edges):
```

```
    edges.sort(key=lambda x: x[2]) # Sort by the third element (weight)
```



```
uf = UnionFind(n)
mst_edges = []
total_weight = 0
    for u, v, weight in edges:
        if uf.union(u, v): # If u and v are not already connected
            mst_edges.append((u, v, weight))
            total_weight += weight
    return mst_edges, total_weight
n = 4
m = 5
edges = [(0, 1, 10), (0, 2, 6), (0, 3, 5), (1, 3, 15), (2, 3, 4)]
mst_edges, total_weight = kruskal(n, edges)
print("Edges in MST:", mst_edges)
print("Total weight of MST:", total_weight)
```

**OUTPUT:**

Edges in MST: [(2, 3, 4), (0, 3, 5), (0, 1, 10)]

Total weight of MST: 19

**12) Given a graph with weights and a potential Minimum Spanning Tree (MST), verify if the given MST is unique. If it is not unique, provide another possible MST.**

**Test Case 1:**

**Input:**

**n = 4**

**m = 5**

**edges = [ (0, 1, 10), (0, 2, 6), (0, 3, 5), (1, 3, 15), (2, 3, 4) ]**

**given\_mst = [(2, 3, 4), (0, 3, 5), (0, 1, 10)]**

**Output: Is the given MST unique? True**

**CODE:**

```
class UnionFind:
```

```
    def __init__(self, n):
```

```
        self.parent = list(range(n))
```

```
        self.rank = [0] * n
```

```
    def find(self, u):
```

```
        if self.parent[u] != u:
```

```
            self.parent[u] = self.find(self.parent[u]) # Path compression
```

```
        return self.parent[u]
```

```
    def union(self, u, v):
```

```
        root_u = self.find(u)
```

```
        root_v = self.find(v)
```

```
        if root_u != root_v:
```

```
            if self.rank[root_u] > self.rank[root_v]:
```

```
                self.parent[root_v] = root_u
```

```
            elif self.rank[root_u] < self.rank[root_v]:
```

```
                self.parent[root_u] = root_v
```

```
            else:
```

```
                self.parent[root_v] = root_u
```

```
                self.rank[root_u] += 1
```

```
        return True
```

```
        return False
```

```
def verify_mst(n, edges, given_mst):
```

```
    uf = UnionFind(n)
```

```
    given_mst_weight = sum(weight for u, v, weight in given_mst)
```

```

for u, v, weight in given_mst:
    uf.union(u, v)
edges.sort(key=lambda x: x[2])
mst_edges = []
total_weight = 0
edge_count = 0
for u, v, weight in edges:
    if uf.union(u, v):
        mst_edges.append((u, v, weight))
        total_weight += weight
        edge_count += 1
        if edge_count == n - 1:
            break
if total_weight != given_mst_weight:
    return False, []
alternative_mst = []
uf2 = UnionFind(n)
for u, v, weight in edges:
    if uf2.union(u, v):
        alternative_mst.append((u, v, weight))
if alternative_mst != given_mst:
    return False, alternative_mst
return True, []

n = 4
m = 5
edges = [(0, 1, 10), (0, 2, 6), (0, 3, 5), (1, 3, 15), (2, 3, 4)]
given_mst = [(2, 3, 4), (0, 3, 5), (0, 1, 10)]
is_unique, alternative_mst = verify_mst(n, edges, given_mst)
print("Is the given MST unique?", is_unique)
if not is_unique:
    print("Another possible MST:", alternative_mst)

```

**OUTPUT:**

Is the given MST unique? True