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Home work 1

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1.10

For  $T(n)$  to be  $O(f(n))$ ,  $f(n)$  should be the upper bound of the growth rate of  $f(n)$  should be the upper bound of the growth rate of  $T(n)$

$T(n)$  is  $\Omega(g(n))$  mean there exists  $T(n) \geq c \cdot g(n)$

When  $i=1$

$f_1(n) = n^2$  is  $O(n^2)$

$f_2(n) = n^2 + 1000 \cdot n$  is  $O(n^2)$

$f_3(n)$  is  $\Omega(f_3(n))$

$f_4(n)$  is  $O(n^3)$

1.12 Using Big Oh notation, the worst case running times

a.  $n^3 \cdot t$  is  $O(n^3)$

b.  $t \cdot j \cdot n - (i+1) \cdot (n-1)$  is  $O(n^2)$

c.  $j \cdot (n-i) \cdot n$  is  $O(n^2)$

d.  $T(n) = 2 \cdot T(n-1)$

$T(1) = 2 \cdot T(0)$

$T(2) = 2 \cdot T(1) = 4 \cdot T(0)$

$T(3) = 2 \cdot T(2) = 8 \cdot T(0)$

Therefore  $T(n) = 2 \cdot T(n-1)$  is  $O(2^n)$