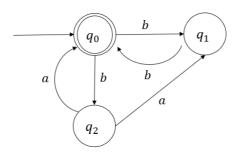
### **UNIT-II**

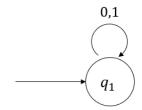
### PRACTICE QUESTIONS

- 1. Describe the following sets by regular expressions:
  - a. L is the set of all string of 0's and 1's ending with 00.
  - b. L is the set of all string of 0's and 1's beginning with 0 and ending with 1.
  - c.  $L = \{\epsilon, 11, 1111, 1111111, \dots \}$
- 2. Find a regular expression corresponding to each of the following subsets of  $\{a, b\}$ 
  - a. The set of all strings containing exactly 2  $\alpha$ 's
  - b. The set of all strings containing at least 2 a's
  - c. The set of all strings containing atmost 2 a's
  - d. The set of all strings containing the substring aa.
- 3. Find regular expressions representing the following sets:
  - a. The set of strings over {0,1} which have at most one pair of 0's or at most one pair of 1's.
  - b. The set of strings over  $\{a, b\}$  in which the number of occurrences of a is divisible by 3.
  - c. The set of strings over  $\{a, b\}$  in which there are at most two occurrences of b between any two occurrences of a.
  - d. The set of strings over  $\{a, b\}$  with three consecutive b's.
  - e. The set of strings over  $\{0,1\}$  beginning with  $\{0,0\}$ .
  - f. The set of strings over  $\{0,1\}$  ending with 00 and beginning with 1.
- 4. Find a regular expression consisting of all strings over  $\{a, b\}$  starting with any number of a's, followed by one or more b's, followed by one or more a's, followed by a single b, followed by any number of a's, followed by b and ending in any string of a's and b's.
- 5. Find the regular expression representing the set of all strings of the form
  - a.  $a^m b^n c^p$  where  $m, n, p \ge 1$
  - b.  $a^m b^{2n} c^{3p}$  where  $m, n, p \ge 1$
  - c.  $a^n b a^{2m} b^2$  where  $m \ge 0, n \ge 1$
- 6. Find all strings of length 5 or less in the regular set represented by the following regular expressions :
  - a.  $(ab + a)^*(aa + b)$
  - b.  $(a^*b + b^*a)^*a$
  - c.  $a^* + (ab + a)^*$
  - d. a(a+b)\*ab
  - e.  $a^*b + b * a$
  - f. (aa + b) \* (bb + a) \*

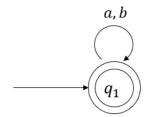
# 7. Covert the following NFA to DFA



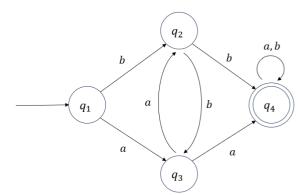
8. Find the set of strings over  $\Sigma = \{a, b\}$  recognized by the transition systems given below:



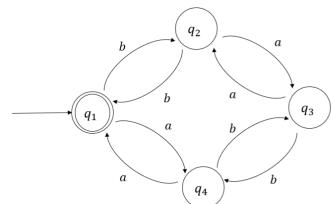
a.



b.

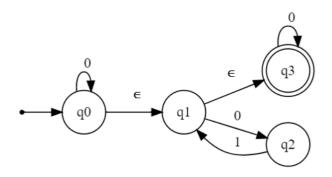


c.



d.

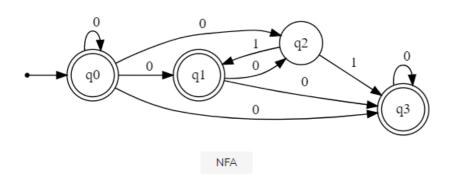
# 9. Convert the following $\epsilon - NFA$ to its equivalent NFA.



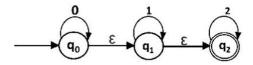
Answer:

**Transitions** 

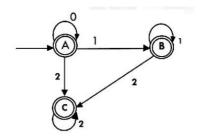
state	0	1
q0	{q0, q1, q2, q3}	-
q1	{q2, q3}	-
q2	-	{q1, q3}
q3	{q3}	-



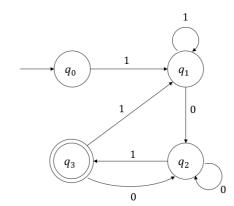
### 10. Convert the following $\epsilon - NFA$ to its equivalent NFA.



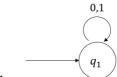
### Answer:



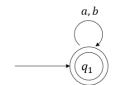
# 11. Find the regular expression corresponding to the automaton given below :



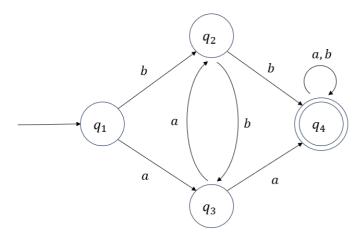
a.



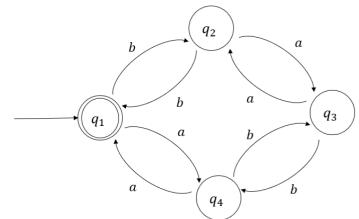
b.



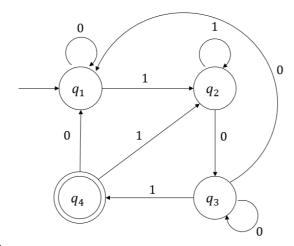
c.



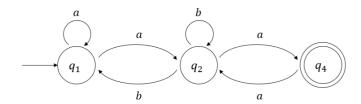
d.



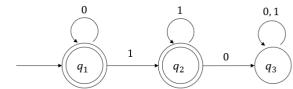
e.



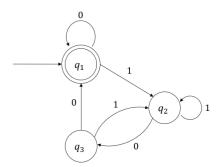
f.



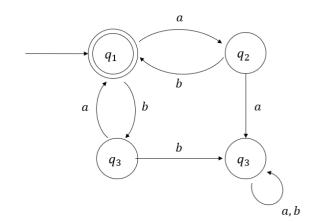
g.



h.



i.



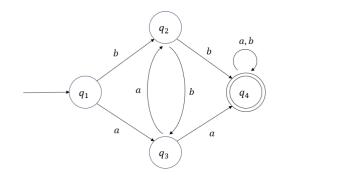
j.

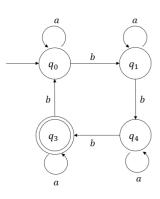
- 12. Construct a DFA with reduced states (minimized DFA) equivalent to the following regular expressions :
  - a.  $(0+1)^*(00+11)(0+1)^*$
  - b. 10 + (0 + 11)0 \* 1
- 13. Construct the transition systems equivalent to the following regular expressions :
  - a.  $(ab + a)^*(aa + b)$
  - b.  $(a^*b + b^*a)^*a$
  - c.  $a^* + (ab + a)^*$
  - d. a(a+b)\*ab
  - e.  $a^*b + b * a$
  - f. (aa + b) \* (bb + a) \*
  - g.  $(ab + c^*) * b$
  - h. a + bb + bab \* a
- 14. Construct a deterministic finite automaton corresponding to the following regular expressions
  - a.  $(ab + a)^*(aa + b)$
  - b.  $(a^*b + b^*a)^*a$
  - c.  $a^* + (ab + a)^*$
  - d.  $(a+b)^*abb$
- 15. Construct a finite automaton accepting all strings over {0,1} ending with 010 or 0010.
- 16. Construct a regular grammar G generating the regular set represented by  $a^*b(a+b)^*$
- 17. If a regular grammar G is given by  $S \to aS \mid a$ , find M accepting L(G).
- 18. Let  $G = (\{A_0, A_1, A_2, A_3\}, \{a, b\}, P, A_0)$ , where P consists of

$$A_0 \to aA_0|bA_1, A_1 \to aA_2|aA_3, A_3 \to a|bA_1|bA_3, A_3 \to b|bA_0.$$

Construct an NDFA accepting L(G).

- 19. Construct a finite automaton recognizing L(G), where G is the grammar  $S \to aS|bA|b$  and  $A \to aA|bS|a$ .
- 20. Find a regular grammar accepting the set recognized by the finite automaton given below:





21. Construct a deterministic finite automaton equivalent to the grammar

$$S \rightarrow aS|bS|aA$$
,  $A \rightarrow bB$ ,  $B \rightarrow aC$ ,  $C \rightarrow \epsilon$ 

22. Let  $G = (\{A_0, A_1, A_2, A_3\}, \{a, b\}, P, A_0\}$  where P consists of

$$A_0 \rightarrow aA_0|bA_1$$

$$A_1 \rightarrow aA_2 | aA_3$$

$$A_2 \rightarrow a|bA_1|bA_3$$

$$A_3 \rightarrow b|bA_0$$

Construct NFA accepting L(G).

- 23. Consider G whose productions are  $S \to aAS|a, A \to SbA|SS|ba$ . Show that  $S \stackrel{*}{\Rightarrow} aabbaa$  and construct a derivation tree whose yield is aabbaa.
- 24. Let G be the grammar  $S \to 0B|1A$ ,  $A \to 0|0S|1AA$ ,  $B \to 1|1S|0BB$ . For the string 00110101, find (a) the leftmost derivation, (b) the rightmost derivation, and (c) the derivation tree.
- 25. If G is a grammar  $S \to SbS|a$ . Show that G is ambiguous.
- 26. Consider the grammar G which has the following productions and answer the questions:

$$A \rightarrow a|Aa|bAA|AAb|AbA$$

- a. What is the start symbol of *G*?
- b. Is aaabb in L(G)?
- c. Is aaaabb in L(G)?
- d. Is abb in L(G)?
- 27. Consider the grammar *G* which has the following productions and state whether the given statements are true or false :

$$S \rightarrow aB|bA$$
,  $A \rightarrow aS|bAA|a$ ,  $B \rightarrow bS|aBB|b$ 

- a. L(G) is finite.
- b.  $abbbaa \in L(G)$
- c.  $aab \notin L(G)$
- d. L(G) has some strings of odd length.
- e. L(G) has some strings of even length.
- 28. Consider the grammar G whose productions are  $S \to aAS|a, A \to SbA|SS|ba$ .

Show that  $S \Rightarrow aabbaa$ . Also, construct a derivation tree.

- 29.  $S \rightarrow aAS|aSS|\epsilon$ ,  $A \rightarrow SbA|ba$ . For the string aabaa, find
  - a. the leftmost derivation

- b. the rightmost derivation
- c. the parse tree
- 30. Find a derivation tree of a \* b + a \* b given that a \* b + a \* b is in L(G), where G is given by  $S \rightarrow S + S \mid S * S, S \rightarrow a \mid b$ .
- 31. A context free grammar G has the following productions:

$$S \rightarrow 0S0 \mid 1S1 \mid A$$
,  $A \rightarrow 2B3$ ,  $B \rightarrow 2B3 \mid 3$ 

Describe the language generated by the parameters.

32. Consider the following productions:

$$S \rightarrow aB \mid bA$$

$$A \rightarrow aS \mid bAA \mid a$$

$$B \rightarrow bS \mid aBB \mid b$$

For the string aaabbabbba, find

- a. the leftmost derivation
- b. the rightmost derivation
- c. the parse tree
- 33. Consider the grammar-

$$S \rightarrow bB / aA$$

$$A \rightarrow b / bS / aAA$$

$$B \rightarrow a / aS / bBB$$

For the string w = bbaababa, find-

- a. Leftmost derivation
- b. Rightmost derivation
- c. Parse Tree
- 34. Consider the grammar-

$$S \rightarrow 0B / 1A$$

$$A \rightarrow 0 / 0S / 1AA$$

$$B \rightarrow 1/1S/0BB$$

For the string w = 00110101, find-

- a. Leftmost derivation
- b. Rightmost derivation

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#### c. Parse Tree

- 35. Show that the grammar  $S \to a|abSb|aAb$ ,  $A \to bS|aAAb$  is ambiguous.
- 36. Show that the grammar  $S \to aB|Ab$ ,  $A \to aAB|a$ ,  $B \to ABb|b$  is ambiguous.
- 37. Show that the grammar  $S \to aAS|aSS|\epsilon$ ,  $A \to SbA|ba$  is ambiguous.
- 38. Find a reduced grammar equivalent to the grammar G whose productions are

$$S \to AB \mid CA$$
,  $B \to BC \mid AB$ ,  $A \to \alpha$ ,  $C \to \alpha B \mid b$ 

39. Find a reduced grammar equivalent to the grammar G whose productions are

$$S \to aAa$$
,  $A \to Sb|bCC|DaA$ ,  $C \to abb|DD$ ,  $E \to aC$ ,  $D \to aDA$ 

40. Find a reduced grammar equivalent to the grammar G whose productions are

$$S \to AC|B$$
,  $A \to a$ ,  $C \to c|Bc$ ,  $E \to aA|\epsilon$ 

41. Find a reduced grammar equivalent to the grammar G whose productions are

$$S \rightarrow AB, A \rightarrow a, B \rightarrow b, B \rightarrow C, E \rightarrow c$$

42. Find a reduced grammar equivalent to the grammar G whose productions are

$$S \to AB \mid CA, B \to BC \mid AB, A \to a, C \to aB \mid b$$

- 43. Find a reduced grammar equivalent to the grammar  $S \to aAa$ ,  $A \to bBB$ ,  $B \to ab$ ,  $C \to aB$ .
- 44. Given a grammar  $S \to AB$ ,  $A \to a$ ,  $B \to C \mid b$ ,  $C \to D$ ,  $D \to E$ ,  $E \to a$ , find an equivalent grammar which is reduced and has no unit productions.
- 45. Remove Unit Productions from the grammar whose productions are given below:

$$S \rightarrow XY, X \rightarrow a, Y \rightarrow Z|b, Z \rightarrow M, M \rightarrow N, N \rightarrow a$$

46. Remove Null Productions from the whose productions are given below:

$$S \to ABAC, A \to aA|\epsilon, B \to bB|\epsilon, C \to c$$

47. Remove Null Productions from the whose productions are given below:

$$S \rightarrow aS|AB, A \rightarrow \epsilon, B \rightarrow \epsilon, D \rightarrow b$$

- 48. Reduce the following grammar G to CNF. G is
  - a.  $S \rightarrow aAD, A \rightarrow aB|bAB, B \rightarrow b, D \rightarrow d$ .
  - b.  $S \rightarrow aAbB, A \rightarrow aA|a, B \rightarrow bB|b$ .
  - c.  $S \rightarrow 1A|0B, A \rightarrow 1AA|0S|0, B \rightarrow 0BB|1S|1$
  - d.  $S \rightarrow a|b|cSS$
  - e.  $S \rightarrow abSb|a|aAb$ ,  $A \rightarrow bS|aAAb$
  - f.  $S \rightarrow ASA|bA, A \rightarrow B|S, B \rightarrow C$
- 49. Reduce the following grammar to GNF.
  - a.  $S \rightarrow SS, S \rightarrow 0S1|01$

b. 
$$S \rightarrow AB, A \rightarrow BSB, A \rightarrow BB, B \rightarrow aAb, B \rightarrow a, A \rightarrow b$$

c. 
$$S \rightarrow A0, A \rightarrow 0B, B \rightarrow A0, B \rightarrow 1$$

- 50. Construct a PDA accepting all palindromes over  $\{a, b\}$ .
- 51. Construct a PDA accepting  $L = \{a^i b^j c^k | i = j \text{ or } j = k\}$  by final state.
- 52. Construct a PDA accepting  $L = \{a^n b^m a^n | m, n \ge 1\}$  by final state.
- 53. Construct a PDA accepting  $L = \{a^n b^{2n} | n \ge 1\}$  by final state.
- 54. Construct a PDA accepting  $L = \{a^n b^m c^n | m, n \ge 1\}$  by final state.
- 55. Construct a PDA accepting  $L = \{a^m b^n | m > n, n \ge 1\}$  by final state.

\* \* \* \* \* \* \* \* \* \*