

# DAA ASSIGNMENT ON DYNAMIC PROGRAMMING

GROUP -8

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# PROBLEM STATEMENT

Given a matrix of random numbers, find maximum length Snake sequence and print it.

If multiple snake sequences exist with the maximum length, print any one of them.

A snake sequence is made up of adjacent numbers in the grid such that for each number, the number on the right or the number below it is +1 or -1 its value.

# ALGORITHM DESIGN

$dp[i][j]$  stores maximum snake sequence length starting at  $(i,j)$ .

From the question, it is clear that a snake starting at a position (say  $(i,j)$ ) can only move right or down.

We will traverse dp array in such a way that while we are at position  $(i,j)$ , we would have already calculated

$dp(i,j+1)$  and  $dp(i+1,j)$ . It's because that  $dp(i,j)$  requires  $dp(i+1,j)$  and  $dp(i,j+1)$ .

One possibility is to traverse from bottom-right to top-left.

we can move from

$$(i, j) \rightarrow (i, j + 1) \Leftrightarrow a[i][j] - a[i][j + 1] = (+|-)1$$

$$(i, j) \rightarrow (i + 1, j) \Leftrightarrow a[i][j] - a[i + 1][j] = (+|-)1$$

## POSSIBILITIES –

- 1) (!right possible AND !down possible),  $dp[i][j] = 1$  ( snake starts and ends at (i,j))
- 2) (right possible AND down possible),  $dp[i][j] = 1 + \max(dp[i][j+1], dp[i+1][j])$
- 3) (right possible AND !down possible),  $dp[i][j] = 1 + dp[i][j+1]$
- 4) (!right possible AND down possible),  $dp[i][j] = 1 + dp[i+1][j]$

## BASE CASES -

$dp[n][m]=1$  (no more cells right or down of  $(n,m)$ ).

right most column - snake can only go down.

bottom most row - snake can only go right.

Final Answer –

$$\max\{dp[i][j]\} \forall 1 \leq i \leq n, 1 \leq j \leq m$$

## Trace path –

To trace the path of maximum snake sequence, we maintain a matrix  $path[][]$  where  $path[i][j]$  points to either RIGHT or DOWN cell (whichever of  $dp(i+1, j)$ ,  $dp(i, j+1)$  yields maximum) or NONE, if it has to end there itself.



## TIME COMPLEXITY ANALYSIS

Number of sub problems= $n*m$ .

Time per subproblem= $O(1)$  (only finding maximum of 2 values).

Hence  $T(n,m)=(n*m)*O(1)$ .

$=O(n*m)$

## AUXILLARY SPACE

We have created `dp[][]` and `path[][]` of  $n*m$  size.

Hence auxillary space= $O(n*m)$



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