## Sparse Subspace Clustering

Based on Sparse Subspace Clustering: Algorithm, Theory, and Applications by Elhamifar and Vidal (2013)

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#### Outline

- 1 Motivation and Background
- 2 Set-up and Algorithm
- **3** Practical Extensions
- Theoretical Guarantees
- 6 Real Experiments

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Slides Based on  $Sparse\ Subspace\ Clustering:\ Algorithm,\ Theory,\ and\ Applications$ 

by Elhamifar and Vidal (2013, algorithm first appeared in 2009)

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## Data in low-dimensional subspaces

High dimensional data is often well-approximated by low-dimensional subspaces. For example:

- Feature trajectories of a rigidly moving object in a video
- 2 face images of a subject under varying illumination
- 3 a hand-written digit with different rotations, translations, and thicknesses

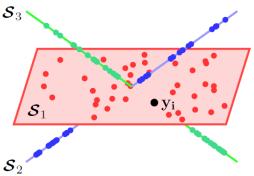
are all well-modeled by a low-dimensional subspace of the corresponding ambient space.

Thus when collecting data from multiple classes we would expect the data to lie in the union of multiple subspaces.

## Subspace clustering

#### Subspace clustering (at a high level)

Given many (noisy) data points drawn from a union of subspaces, find the subspaces and the assignments of each point to a subspace.



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## Ingredients

- **1** Ambient dimension:  $\mathbb{R}^d$
- **2 Subspaces:**  $\{S_{\ell}\}_{\ell=1}^n$ , linear subspaces of  $\mathbb{R}^d$
- **3 Subspace dimensions:**  $\{d_{\ell}\}_{\ell=1}^n$  the dimension of each subspace
- **4** Data:  $\{y_i\}_{i=1}^N$ , N data points lying in  $\bigcup_{\ell=1}^n S_\ell$ :

$$Y := [y_1, \dots, y_N] = [Y_1, \dots, Y_n]\Gamma,$$

where  $Y_{\ell} \in \mathbb{R}^{D \times N_{\ell}}$  is the matrix containing all the points lying in  $S_{\ell}$  and  $\Gamma$  is an (unknown) permutation matrix.

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where  $Y_{\ell} \in \mathbb{R}^{D \times N_{\ell}}$  is the matrix containing all the points lying in  $S_{\ell}$  and  $\Gamma$  is an (unknown) permutation matrix.

**Assume:**  $Y_{\ell}$  is rank  $d_{\ell}$  and  $N_{\ell} > d_{\ell}$ .

## Parsimonious Representation

#### Idea:

we should<sup>1</sup> be able to represent every point in a subspace as the linear combination of **just a few** other points in that same subspace.

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## Parsimonious Representation

#### Idea:

we should<sup>1</sup> be able to represent every point in a subspace as the linear combination of **just a few** other points in that same subspace.

Formally, given  $y_i \in S_{\ell}$  we can write

$$y_i = Yc_i$$
, where  $c_{i,i} = 0$ 

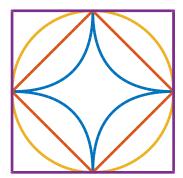
and 
$$c_{i,j} \neq 0 \iff y_i \in S_\ell$$
.

<sup>&</sup>lt;sup>1</sup>If we have enough data points in each subspace and the dimensions of each subspace is comparatively small

## Aside: how to promote sparsity

We would like to minimize the number of non-zero entries (the "zero norm") in the solution to some linear system

$$\label{eq:subject_to} \begin{aligned} & \underset{x \in \mathbb{R}^d}{\text{minimize}} & & \|x\|_0 \\ & \text{subject to} & & \|Ax - b\|_2 \le \varepsilon \end{aligned}$$



## **Proposed Optimization Problem**

This naturally motivates the optimization problem

minimize 
$$\|c_i\|_q$$
 subject to  $y_i = Yc_i, c_{i,i} = 0$   $i = 1, \dots, N.$ 

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(setting q = 1 and re-writing using matrix notation)

$$\label{eq:continuity} \begin{aligned} & \underset{C}{\text{minimize}} & & \|C\|_{1,1} \\ & \text{subject to} & & Y = YC, & & \text{diag}(C) = 0 \end{aligned}$$

## Using C to cluster

To find the clusters:

- $\bullet$  Consider |C| as the adjacency matrix of some graph
- **2** Symmetrize (make undirected):  $W := |C| + |C|^T$
- $\bullet$  Perform spectral clustering on W
  - 1 Build the symmetric nomalized graph Laplacian

$$L = I - D^{-1/2}WD^{1/2}.$$

- **2** calculate the *n* bottom eigenvectors  $U := [u_1, \ldots, u_n]$  of *L*
- **3** Apply k-means to the normalized rows of U.

In the ideal case we will get n connected components

$$W = \begin{bmatrix} W_1 & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & W_n \end{bmatrix} \Gamma.$$

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## Noisy Data Model

We assumed before that we were given N noise-free data points. Let's extend this to a model

$$y_i = y_i^0 + e_i^0 + z_i^0,$$

which can handle:

- sparse outlying entries:  $||e_i^0||_0 \le k$
- noise:  $||z_i^0||_2 \le \zeta$ .

## Practical Optimization Problem (I)

We note that we can sparsely represent  $y_i^0$  as before:

$$y_i^0 = \sum_{j \neq i} c_{i,j} y_j^0.$$

Following our nose, we write

$$y_{i} = \sum_{j \neq i} y_{j} + e_{i} + z_{i}$$

$$e_{i} := e_{i}^{0} - \sum_{j \neq i} c_{i,j} e_{j}^{0}$$

$$z_{i} := z_{i}^{0} - \sum_{j \neq i} c_{i,j} z_{j}^{0}.$$

## Practical Optimization Problem (II)

It is clear now what the optimization program should be:

minimize 
$$||C||_{1,1} + \lambda_e ||E||_{1,1} + \frac{\lambda_z}{2} ||Z||_F^2$$
 subject to 
$$Y = YC + E + Z \quad \text{and} \quad \operatorname{diag}(C) = 0$$

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#### **Definitions and Notation**

#### **Independent Subspaces**

A collection of subspaces  $\{S_{\ell}\}_{\ell=1}^n$  is said to be **independent** if  $\dim (\bigoplus_{\ell=1}^n S_{\ell}) = \sum_{\ell=1}^n \dim S_{\ell}$ 

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#### Disjoint Subspaces

We will call a collection of subspaces  $\{S_{\ell}\}_{\ell=1}^{n}$  disjoint if the intersection of any two of them only contains the origin.

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#### Disjoint Subspaces

We will call a collection of subspaces  $\{S_{\ell}\}_{\ell=1}^{n}$  disjoint if the intersection of any two of them only contains the origin.

Let  $Y_{\ell}$  denote the  $N_{\ell}$  points corresponding to a subspace  $S_{\ell}$  and let  $Y_{-\ell}$  denote the rest of the data points.

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# Independent Subspaces Theory

#### Theorem 1

Assume that we have n independent subspaces of full rank:  $\operatorname{rank}(Y_{\ell}) = d_{\ell}$ . Then for every  $S_{\ell}$  and every non-zero y in  $S_{\ell}$  the  $\ell_q$  minimization problem:

$$\begin{bmatrix} c^* \\ c_-^* \end{bmatrix}. = \arg\min \left\| \begin{bmatrix} c \\ c_- \end{bmatrix} \right\|_q \quad \text{s.t.} \quad y = [Y_\ell \ Y_{-\ell}] \begin{bmatrix} c \\ c_- \end{bmatrix},$$

for  $q < \infty$  recovers a subspace-sparse representation. That is,  $c^* \neq 0$  and  $c^* = 0$ .

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for  $q < \infty$  recovers a subspace-sparse representation. That is,  $c^* \neq 0$  and  $c^* = 0$ .

Note that the condition on independent subspaces is completely unreasonable.

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# Disjoint Subspaces Theory (I)

For a vector x in the intersection of  $S_{\ell}$  with  $\bigoplus_{j\neq \ell} S_j$  define

$$a_{\ell} := \arg\min \|a\|_1$$
 s.t.  $x = Y_{\ell}A$   
 $a_{-\ell} := \arg\min \|a\|_1$  s.t.  $x = Y_{-\ell}A$ 

# Disjoint Subspaces Theory (I)

For a vector x in the intersection of  $S_{\ell}$  with  $\bigoplus_{j\neq\ell} S_j$  define

$$\begin{aligned} a_{\ell} &:= \arg\min \|a\|_1 & \text{s.t.} & x &= Y_{\ell}A \\ a_{-\ell} &:= \arg\min \|a\|_1 & \text{s.t.} & x &= Y_{-\ell}A \end{aligned}$$

We define the exact recover condition:

$$\forall x \in S_{\ell} \cap (\bigoplus_{i \neq \ell} S_i), \quad x \neq 0 \Rightarrow ||a_i||_1 < ||a_{-i}||_1. \tag{ERC}$$

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## Disjoint Subspaces Theory (II)

$$\forall x \in S_{\ell} \cap (\bigoplus_{j \neq \ell} S_j), \quad x \neq 0 \Rightarrow ||a_i||_1 < ||a_{-i}||_1.$$
 (ERC)

## Disjoint Subspaces Theory (II)

$$\forall x \in S_{\ell} \cap (\bigoplus_{j \neq \ell} S_j), \quad x \neq 0 \Rightarrow ||a_i||_1 < ||a_{-i}||_1.$$
 (ERC)

#### Theorem 2

Given n disjoint subspaces of full rank, for every  $S_{\ell}$  and every non-zero y in  $S_{\ell}$  the  $\ell_q$  minimization problem:

$$\begin{bmatrix} c^* \\ c^*_- \end{bmatrix} = \arg \min \begin{bmatrix} c \\ c_- \end{bmatrix} \begin{bmatrix} c \\ c_- \end{bmatrix}, \quad \text{s.t.} \quad y = [Y_i \ Y_{-i}] \begin{bmatrix} c \\ c_- \end{bmatrix},$$

recovers a subspace-sparse representation if and only if (ERC) holds.

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## Face Clustering Example

Given many different pictures of faces with varying illuminations, we seek to identify which faces belong to the same person.



Fig. 2. Face clustering: given face images of multiple subjects (top), the goal is to find images that belong to the same subject (bottom).

(results take from Sparse Subspace Clustering)

## Comparison to other methods

TABLE 1

Clustering error (%) of different algorithms on the Hopkins 155 dataset with the 2F-dimensional data points.

Algorithms	LSA	SCC	LRR	LRR-H	LRSC	SSC
2 Motions						
Mean	4.23	2.89	4.10	2.13	3.69	1.52 (2.07)
Median	0.56	0.00	0.22	0.00	0.29	0.00 (0.00)
3 Motions						
Mean	7.02	8.25	9.89	4.03	7.69	4.40 (5.27)
Median	1.45	0.24	6.22	1.43	3.80	0.56 (0.40)
All						
Mean	4.86	4.10	5.41	2.56	4.59	<b>2.18</b> (2.79)
Median	0.89	0.00	0.53	0.00	0.60	0.00 (0.00)

TABLE 2

Clustering error (%) of different algorithms on the Hopkins 155 dataset with the 4n-dimensional data points obtained by applying PCA.

Algorithms	LSA	SCC	LRR	LRR-H	LRSC	SSC
2 Motions						
Mean	3.61	3.04	4.83	3.41	3.87	1.83 (2.14)
Median	0.51	0.00	0.26	0.00	0.26	0.00 (0.00
3 Motions						
Mean	7.65	7.91	9.89	4.86	7.72	4.40 (5.29
Median	1.27	1.14	6.22	1.47	3.80	0.56 (0.40
All						
Mean	4.52	4.14	5.98	3.74	4.74	2.41 (2.85
Median	0.57	0.00	0.59	0.00	0.58	0.00 (0.00

(results take from Sparse Subspace Clustering)

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**6** Bonus Slides

# Practical Optimization Problem (III), affine subspaces

We can deal with affine spaces by including a single extra linear equality constraint:

minimize 
$$\|C\|_{1,1} + \lambda_e \|E\|_{1,1} + \frac{\lambda_z}{2} \|Z\|_F^2$$
 subject to 
$$Y = YC + E + Z, \quad \operatorname{diag}(C) = 0$$
 and  $1^TC = 1^T,$ 

(simply adds on the equation for an affine plane)

# Prior work on subspace clustering Part 0, iterative approaches

iterative approaches are largely based on generalizations of k-means

- k-subspaces (Tseng, 2000)
- median k-flats (Zhang, Szlam, and Lerman 2009)

# Prior work on subspace clustering Part I, Algebraic

Algebraic approaches to subspace clustering involve either

- Factorization based approaches
  - 1 Costeira and Kanade, 1998
  - 2 Kanatani 2001
- Algebro-geometric approaches (called Generalized PCA)
  - 1 Vidal, Ma, and Sastry (2005)
  - 2 Ma, Yang, Derksen, and Fossum (2008)

# Prior work on subspace clustering Part II: statistical approaches

These approaches generally assume a Gaussian distribution on the data inside each subspace and then use some basic estimation theory to find the spaces

- Mixtures of Probabilistic PCA (Tipping and Bishop, 1999)
- Multi-Stage Learning (EM based), (Gruber and Weiss, 2004)
- Random Sample Consensus (RANSAC, Fischler and Bolles, 1981)

### Prior Work: Part III

Spectral clustering algorithms:

- Sparse Subspace Clustering (Elhamifar and Vidal 2009, Soltanolkotabi and Candes 2012)
- Low-rank recovery (LRR, Liu, Lin, and Yu 2010)
- Spectral Curvature Clustering of Chen and Lerman (2009)
- Local Subspace Affinity (LSA) by Yan and Pollefeys (2006)