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ADVANCED COMPUTER ARCHITECTURE

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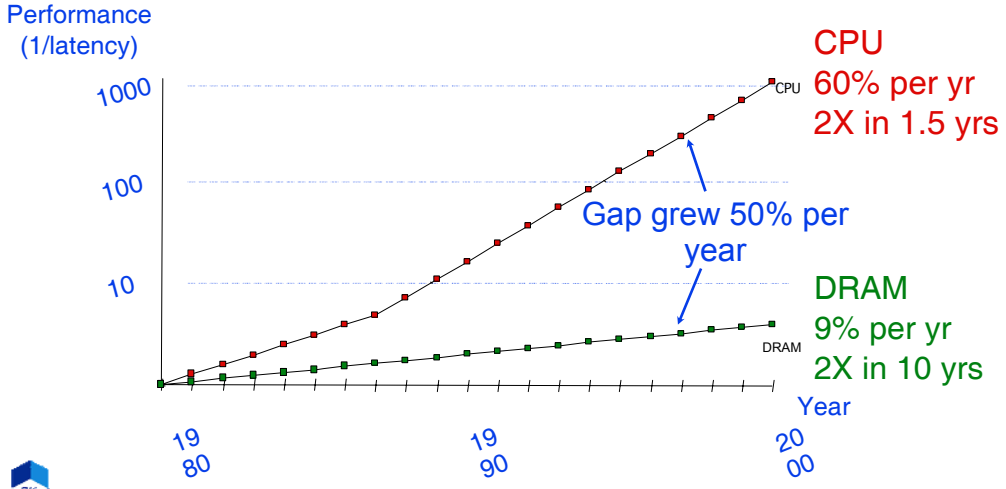
Memory Hierarchy Design

Since 1980, CPU has outpaced DRAM ...

Four-issue 2GHz superscalar accessing 100ns DRAM could execute 800 instructions during time for one memory access!

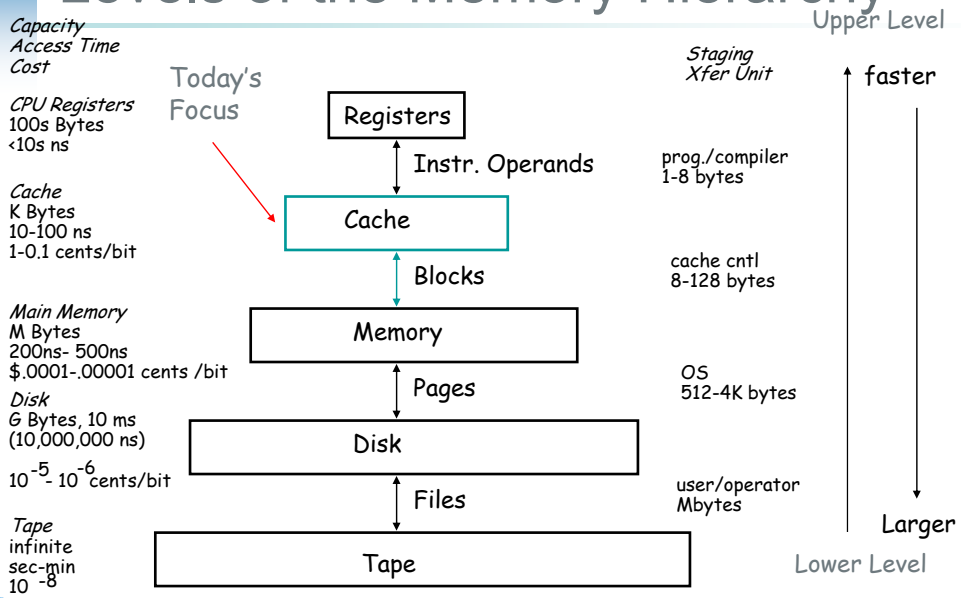


Processor-DRAM Performance Gap Impact

- To illustrate the performance impact, assume a single-issue pipelined CPU with CPI = 1 using non-ideal memory.
- Ignoring other factors, the minimum cost of a full memory access in terms of number of wasted CPU cycles:

Year	CPU speed MHZ	CPU cycle ns	Memory Access ns	Minimum CPU memory stall cycles or instructions wasted
1986:	8	125	190	190/125 - 1 = 0.5
1989:	33	30	165	165/30 -1 = 4.5
1992:	60	16.6	120	120/16.6 -1 = 6.2
1996:	200	5	110	110/5 -1 = 21
1998:	300	3.33	100	100/3.33 -1 = 29
2000:	1000	1	90	90/1 - 1 = 89
2002:	2000	.5	80	80/.5 - 1 = 159
2004:	3000	.333	60	60.333 - 1 = 179

Levels of the Memory Hierarchy



Addressing the Processor-Memory Performance GAP

- Goal: Illusion of large, fast, cheap memory. Let programs address a memory space that scales to the disk size, at a speed that is usually as fast as register access
- Solution: Put smaller, faster "cache" memories between CPU and DRAM. Create a "memory hierarchy".

Common Predictable Patterns

Two predictable properties of memory references:

- Temporal Locality: If a location is referenced, it is likely to be referenced again in the near future (e.g., loops, reuse).
- Spatial Locality: If a location is referenced it is likely that locations near it will be referenced in the near future (e.g., straightline code, array access).

Caches

Caches exploit both types of predictability:

- Exploit temporal locality by remembering the contents of recently accessed locations.
- Exploit spatial locality by fetching blocks of data around recently accessed locations.

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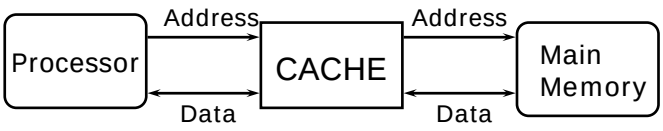
Simple view of cache



- The processor accesses the cache first
- Cache hit: Just use the data
- Cache miss: replace a block in cache by a block from main memory, use the data
- The data transferred between cache and main memory is in blocks, and controlled by independent hardware

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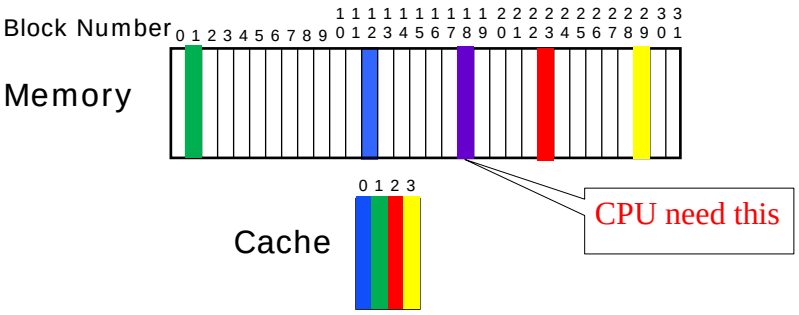
Simple view of cache

- Hit rate: fraction of cache hit ti le
- Miss rate: 1 – Hit rate
- Miss penalty: Time to replace a block + time to deliver the data to the processor

Simple view of cache

- Example: For(i = 0; i < 10; i++) S = S + A[i];
- No cache: At least 12 accesses to main memory (10 A[i] and Read S, write S)
- With Cache: if A[i] and S is in a single block (ex 32-bytes), 1 access to load block to cache, and 1 access to write block to main memory
- Access to S: Temporal Locality
- Access to A[i]: Spatial Locality (A[i])

Replacement



- Cache cannot hold all blocks
- Replace a block by another that is currently needed by CPU

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Basic Cache Design & Operation Issues

- Q1: Where can a block be placed cache?
  (Block placement strategy & Cache organization)
  - Fully Associative, Set Associative, Direct Mapped.
- Q2: How is a block found if it is in cache?
  (Block identification)
  - Tag/Block.
- Q3: Which block should be replaced on a miss?
  (Block replacement)
  - Random, LRU, FIFO.
- Q4: What happens on a write?
  (Cache write policy)
  - Write through, write back.

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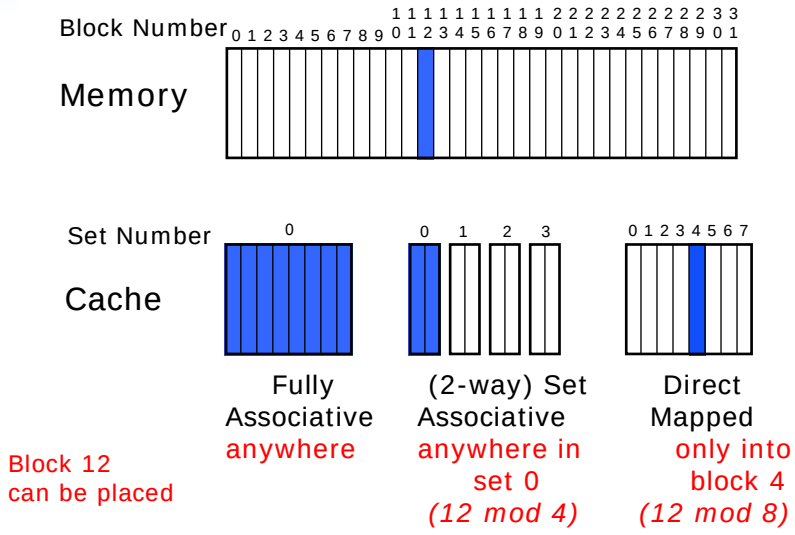
Q1: Where can a block be placed?



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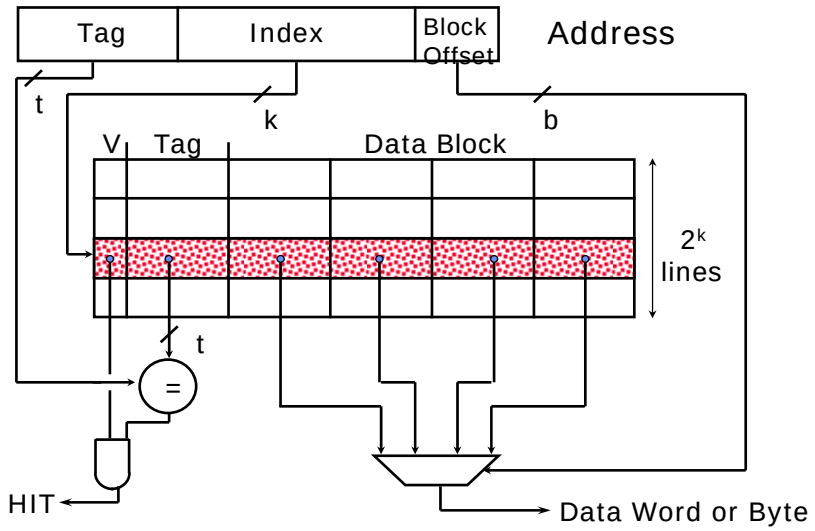
Direct-Mapped Cache



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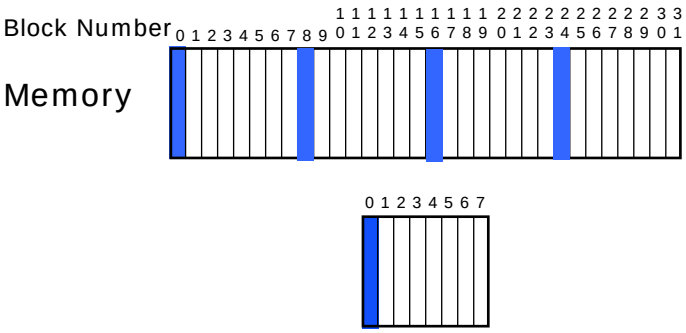
Direct-mapped Cache

- Address: N bits (2^N words)
- Cache has 2^k lines (blocks)
- Each line has 2^b words
- Block M is mapped to the line M % 2^k
- Need t = N-k-bTag bits to identify mem. block
- Advantage: Simple
- Disadvantage: High miss rate
- What if CPU accesses block N0, N1 and N0 % 2^k = N1 % 2^k ?

Direct-mapped Cache



- Access N0, N1 where N0 % 2^k = N1 % 2^k
- Replace a block while there are many rooms available!

4KB Direct Mapped Cache Example

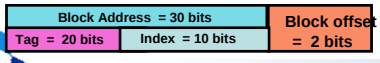
1K = 1024 Blocks
Each block = one word

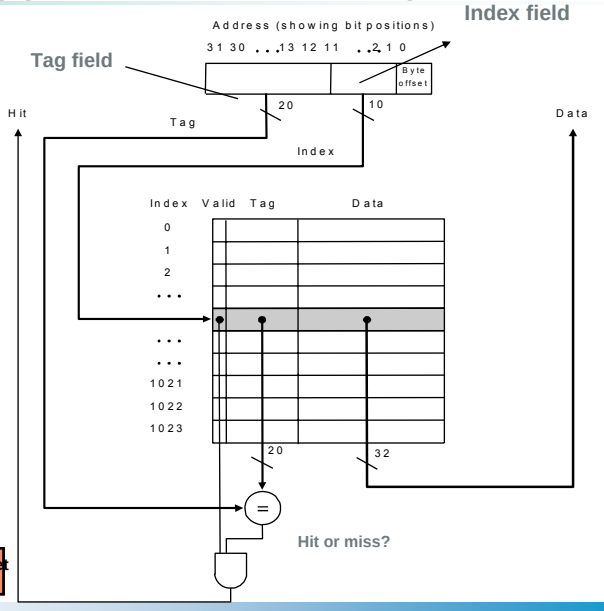
Can cache up to
2^{32} bytes = 4 GB
of memory

Mapping function:

Cache Block frame number = (Block address) MOD (1024)

i.e. index field or
10 low bit of block address





Moi line 2^b word?

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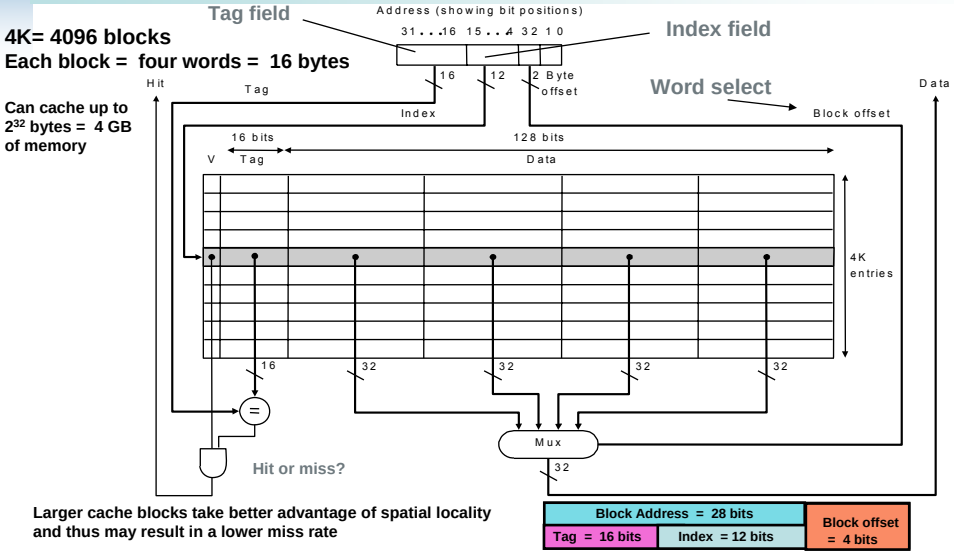
64KB Direct Mapped Cache Example



Mapping Function: Cache Block frame number = (Block address) MOD (4096)

i.e. index field or 12 low bit of block address

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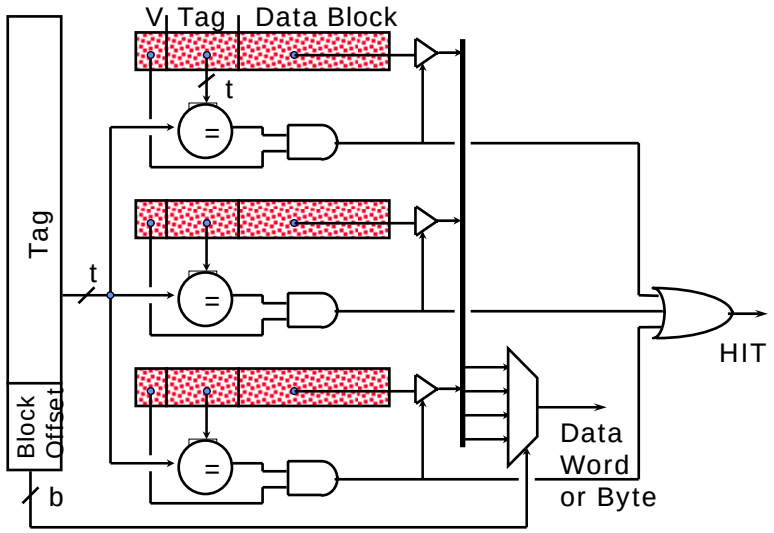
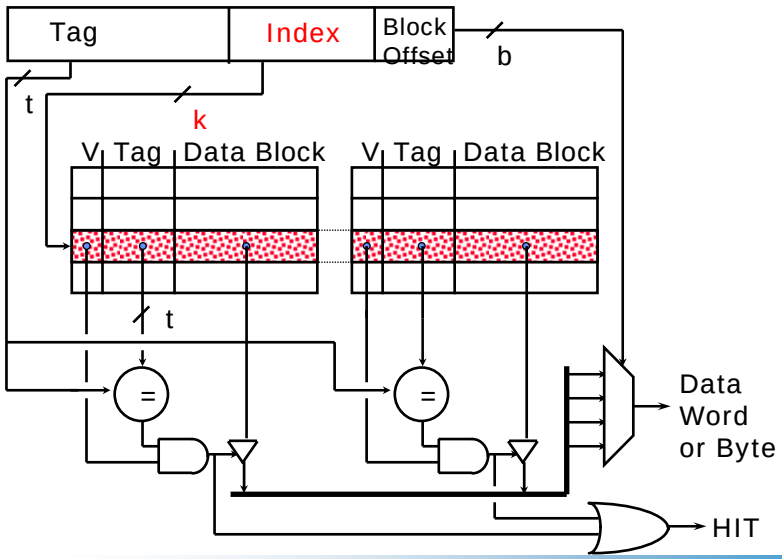
Fully Associative Cache



Fully associative cache

- CAM: Content Addressable Memory
- Each block can be mapped to any lines in cache
- Tag bit: t = N-b. Compared to Tag of all lines
- Advantage: replacement occurs only when no rooms available
- Disadvantage: resource consumption, delay by comparing many elements

Set-Associative Cache



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W-way Set-associative Cache

- Balancing: Direct mapped cache vs Fully associative cache
- Cache has 2^k sets
- Each set has 2^w lines
- Block M is mapped to one of 2^w lines in set M % 2^k
- Tag bit: t = N-k-b
- Currently: widely used (Intel, AMD, …)

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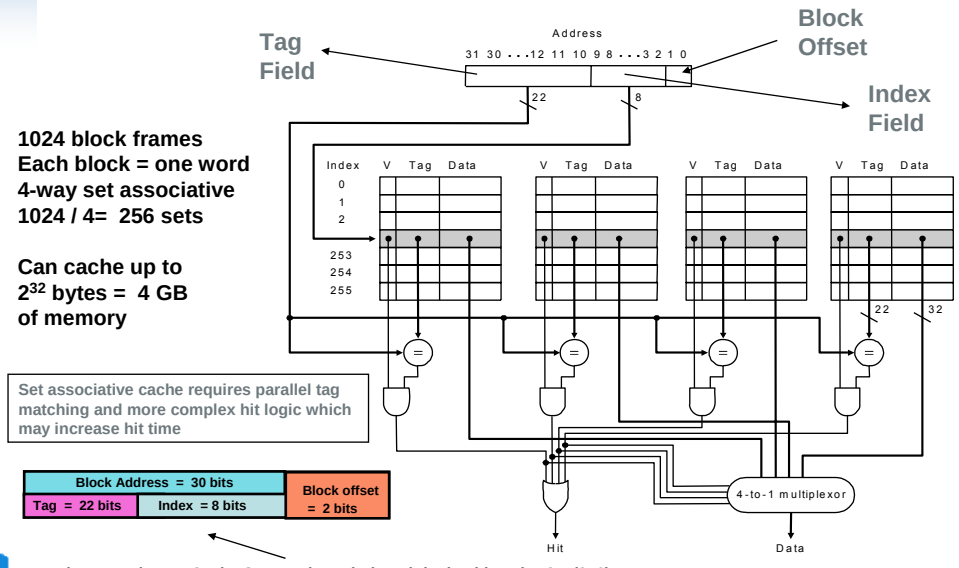
4K Four-Way Set Associative Cache: MIPS Implementation Example



1024 block frames
Each block = one word
4-way set associative
1024 / 4= 256 sets

Can cache up to 2^{32} bytes = 4 GB of memory

Set associative cache requires parallel tag matching and more complex hit logic which may increase hit time

Block Address = 30 bits		Block offset = 2 bits
Tag = 22 bits	Index = 8 bits	

Mapping Function: **Cache Set Number = index= (Block address) MOD (256)**

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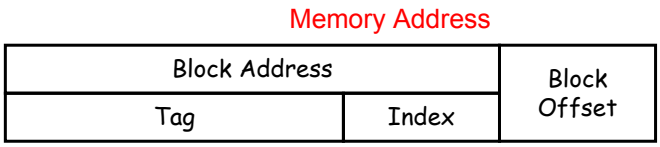
Q2: How is a block found?

- Index selects which set to look in
- Compare Tag to find block
- Increasing associativity shrinks index, expands tag. Fully Associative caches have no index field.
- Direct-mapped: 1-way set associative?
- Fully associative: 1 set?



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What causes a MISS?

- Three Major Categories of Cache Misses:
  - Compulsory Misses: first access to a block
  - Capacity Misses: cache cannot contain all blocks needed to execute the program
  - Conflict Misses: block replaced by another block and then later retrieved - (affects set assoc. or direct mapped caches)
    - Nightmare Scenario: ping pong effect!

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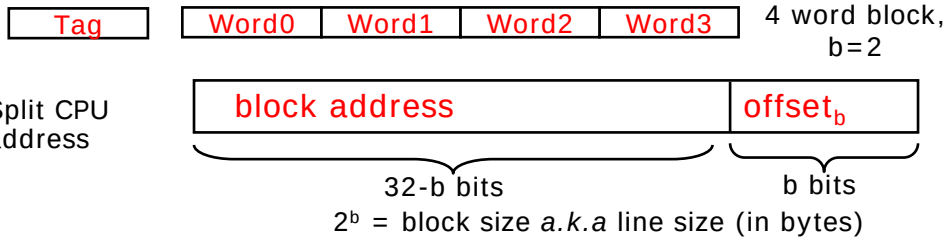
Block Size and Spatial Locality

Block is unit of transfer between the cache and memory



Larger block size has distinct hardware advantages

- less tag overhead
- exploit fast burst transfers from DRAM
- exploit fast burst transfers over wide busses

What are the disadvantages of increasing block size?

Fewer blocks => more conflicts. Can waste bandwidth.

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Q3: Which block should be replaced on a miss?

- Easy for Direct Mapped
- Set Associative or Fully Associative:
  - Random
  - Least Recently Used (LRU)
    - LRU cache state must be updated on every access
    - true implementation only feasible for small sets (2-way, 4-way)
    - pseudo-LRU binary tree often used for 4-8 way
  - First In, First Out (FIFO) a.k.a. Round-Robin
    - used in highly associative caches
- Replacement policy has a second order effect since replacement only happens on misses

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Q4: What happens on a write?

- Cache hit:
  - *write through:* write both cache & memory
    - generally higher traffic but simplifies cache coherence
  - *write back:* write cache only
    (memory is written only when the entry is evicted)
    - a dirty bit per block can further reduce the traffic
- Cache miss:
  - *no write allocate:* only write to main memory
  - *write allocate (aka fetch on write):* fetch into cache
- Common combinations:
  - write through and no write allocate (below example)
  - write back with write allocate (above Example)

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Reading assignment 1

- Cache coherent problem in multicore systems
  - Identify the problem
  - Algorithms for multicore architectures
- Reference
  - *eecs.wsu.edu/~cs460/cs550/**cachecoherence**.pdf*
  - …More on internet

Reading assignment 2

- Cache performance
  - Replacement policy (algorithms)
  - Optimization (Miss rate, penalty, …)
- Reference
  - Hennessy - Patterson - Computer Architecture. A Quantitative
  - *www2.lns.mit.edu/~avinatan/research/**cache**.pdf*
  - *… More on internet*