

# Programming with Effect Handlers in Links

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# WARNING

This talk may contain traces of jargon such as *monads*, *effects*, *algebras* and *handlers*. However the following code examples in this talk may be performed at home.

# What this talk is about

*Handlers for algebraic effects provide a compelling alternative to monads as a basis for effectful programming.*

- **Key idea:** Separate effect signatures from their implementation.
- **“The effect”:** High-degree of modularity.

Definitions will follow later...

## PART 1: Effectively, it's a problem

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# Programs are inherently effectful

Programs may...

- ...halt prematurely
- ...diverge
- ...be stateful (e.g. modify a global state)
- ...communicate via a network
- ...print to standard out

A pure<sup>1</sup> program is not much fun.

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<sup>1</sup>By pure we mean a program that has no effects.

# Fundamental different approaches to effects

**Imperative** Repeatedly performs implicit effects on shared global state.

**Functional** Encapsulates effects in a computational context.

# Fundamental different approaches to effects

**Imperative** Repeatedly performs implicit effects on shared global state.

**Functional** Encapsulates effects in a computational context.

This talk is oriented around *functional* programming with effects.

# Effectful computations (I)

$a \rightarrow b$

Mathematical pure function

$a \rightarrow b$

C/C++ (impure) function

$a \rightarrow b$

ML (impure) function



# Let's be explicit about effects

## Effect annotation

An effect annotation gives a static description of the potential run-time behaviour of a computation.

### Benefits

- Serves as documentation (clarity)
- Compiler can apply specific optimisations
- Possible to reason more precisely about programs

# Enter the Monad

*"Shall I be pure or impure?"*

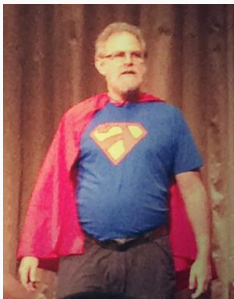


Figure 1: Philip Wadler aka. Lambda Man

- The Essence of Functional Programming [1]
- The Marriage of Effects and Monads [2]

# Effectful computations (II)

$a \rightarrow b$

Mathematical pure function

$a \rightarrow b$

C/C++ (impure) function

$a \rightarrow b$

ML (impure) function

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$a \rightarrow m b$

Haskell impure function

# Effectful computations (II)

$a \rightarrow b$

Mathematical pure function

$a \rightarrow b$

C/C++ (impure) function

$a \rightarrow b$

ML (impure) function

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$a \rightarrow m b$

Haskell impure function

$m$  can be considered an effect annotation

# Monads

## Definition

A monad is a triple  $(m, \text{return}, \text{bind})$  where

- $m$  is a type constructor
- $\text{return} : a \rightarrow m\ a$
- $\text{bind} : m\ a \rightarrow (a \rightarrow m\ b) \rightarrow m\ b$

# Great! Many monads, many effects

A couple of monads and their “effect interpretation”

**IO**  $a$  May perform I/O, returns  $a$

**Reader**  $r$   $a$  May read from  $r$ , returns  $a$

**Writer**  $w$   $a$  May write to  $w$ , returns  $a$

**State**  $s$   $a$  May read/write some state  $s$ , returns  $a$

**Maybe**  $a$  May fail, returns  $a$  on success

# Effectful computations (III)

$a \rightarrow b$

Mathematical pure function

$a \rightarrow b$

C/C++ (impure) function

$a \rightarrow b$

ML (impure) function

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$a \rightarrow m_1 m_2 b$

Haskell impure function

$\not\sim$

$a \rightarrow m_2 m_1 b$

# IO Monad is equivalent to a calzone pizza



Figure 2: “There’s only meat sauce inside”, they said, but you can’t really be sure.



# The importance of effect ordering [3]

Two signatures<sup>2</sup>:

- $A \stackrel{\text{def}}{=} \text{WriterT String Maybe}$ 
  - Returns nothing on failure
- $B \stackrel{\text{def}}{=} \text{MaybeT (Writer String)}$ 
  - Returns a pair on failure

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<sup>2</sup>Elided details about Monad Transformers

## PART 2: Exit the Monad, Enter the Handler

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Figure 3: Gordon Plotkin



Figure 4: John Power

# Algebraic effects and computations

## Definition

**Algebraic effect** An algebraic effect is a collection of abstract operations, e.g.  $\{Op_i : a_i \rightarrow b_i\}$

## Definition

**Abstract computation** An abstract computation is composed from abstract operations. Computations have type

# Nim: A game with sticks

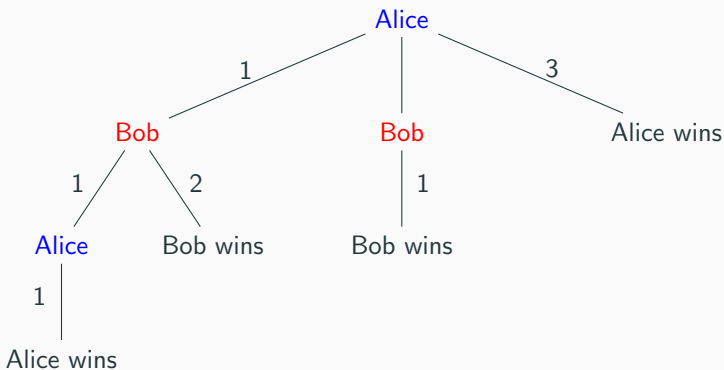


## Set-up

- Two players: Alice and Bob; Alice always starts.
- One heap of  $n$  sticks.
- Turn-based. Each player take between 1-3 sticks.
- The one, who takes the last stick, wins.

We'll demonstrate how to encode strategic behaviour, compute game data, and cheat using handlers.

# Game tree generated by mtGen with $n = 3$



# References



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Links: Linking theory to practice for the web.