

# Programming with Effect Handlers in Links

*PLInG, Autumn 2015, University of Edinburgh*

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November 5, 2015

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# Programs are inherently effectful

Programs may...

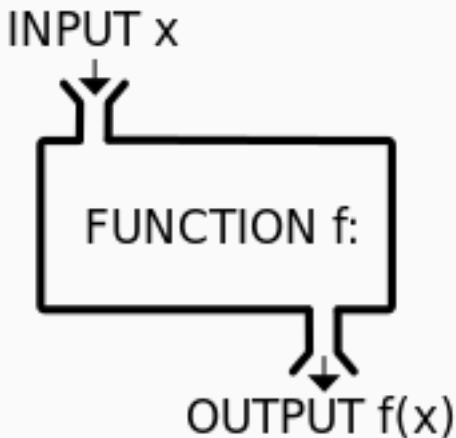
- ...halt prematurely
- ...diverge
- ...be stateful (e.g. allocating memory, throwing exceptions)
- ...communicate via a network
- ...print to standard out

A pure<sup>1</sup> program is not much fun.

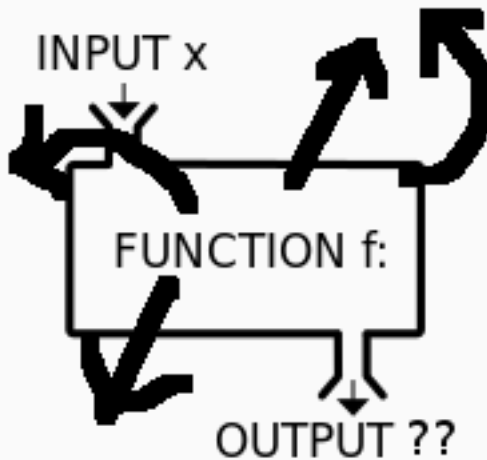
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<sup>1</sup>By pure we mean a program that has no effects.

## Pure vs. impure



## Pure vs. impure



# Function signatures (I)

$f : \mathbb{Z} \rightarrow \mathbb{Z}$

Mathematical pure function

---

`int f(int x)`

C/C++ (impure) function

$f : \text{int} \rightarrow \text{int}$

ML (impure) function

# Let's be explicit about effects

## Effect annotation

An effect annotation gives a static description of the potential run-time behaviour of a computation.

### Benefits

- Serves as documentation (clarity)
- Compiler can apply specific optimisations
- Possible to reason more precisely about programs

# Enter the Monad

“Have you considered using a monad?”



Figure 1: Philip Wadler aka. Lambda Man

- The Essence of Functional Programming (1992)
- The Marriage of Effects and Monads (with Peter Thiemann, 2003)

# Function signatures (II)

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Mathematical pure function

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$f : \text{int} \rightarrow \text{int}$

ML (impure) function

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`f :: Int → IO Int`

Haskell effect annotation



# Function signatures (II)

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Haskell effect annotation

# Function signatures (II)

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ML (impure) function

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`f :: Int → IO Int`

Haskell effect annotation

`int f(int x) throws IOException`

Java effect annotation

# Function signatures (II)

$f : \mathbb{Z} \rightarrow \mathbb{Z}$	Mathematical pure function
<hr/>	
<code>int f(int x)</code>	C/C++ (impure) function
$f : \text{int} \rightarrow \text{int}$	ML (impure) function
<hr/>	
$f :: \text{Int} \rightarrow \text{IO Int}$	Haskell effect annotation
<code>int f(int x) throws IOException</code>	Java effect annotation
$f : (\text{Int}) \xrightarrow{\{Read:String, Write:String \rightarrow ()\}} \text{Int}$	Links effect annotation

# Algebraic effects

Algebraic effects give us a modular way of describing multiple effects.



Figure 2: Gordon Plotkin



Figure 3: John Power

Adequacy for Algebraic Effects (2001)

# Algebraic effects and computations

## Definition (Algebraic effect)

An algebraic effect is a collection of abstract operations. For example  $State \stackrel{\text{def}}{=} \{Get : s, Put : s \rightarrow ()\}$ .

## Definition (Abstract computation)

An abstract computation is composed from abstract operations.

# Effect handlers

Effect handlers assign a particular meaning (semantics) to algebraic effects.



Figure 4: Gordon Plotkin



Figure 5: Matija Pretnar

Handling Algebraic Effects (2003)

# Effect handler

## Definition (Handler)

A handler is an interpreter for abstract computations.

## An example (Links style handler)

Essentially, a handler pattern-matches on operations occurring in a computation  $m$ :

```
handler state(m) {  
  case Get      -> ...  
  case Put(s)  -> ...  
  ...  
}
```

Recall  $State \stackrel{\text{def}}{=} \{Get : s, Put : s \rightarrow ()\}$

# Notion of computation in Links

In Links, we define the type of (handler compatible) computations as

$$\mathit{Comp}(E, a) \stackrel{\text{def}}{=} () \xrightarrow{E} a,$$

where  $E$  is an effect signature, and  $a$  is the return type of the computation. In other words, computations are thunks.



# Handlers in action: Nim – A game with sticks

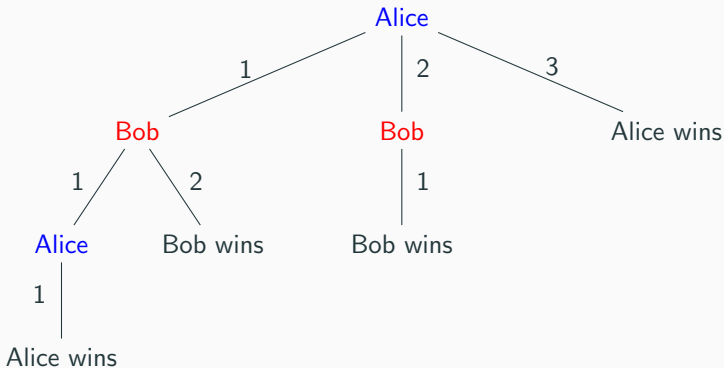


## Set-up

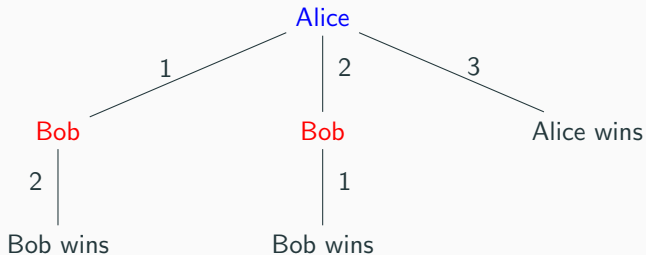
- Two players: Alice and Bob; Alice always starts.
- One heap of  $n$  sticks.
- Turn-based. Each player take between 1-3 sticks.
- The one, who takes the last stick, wins.

We'll demonstrate how to encode strategic behaviour, compute game data, and cheat using handlers.

# Game tree generated by gametree with $n = 3$



# Game tree when Bob plays the perfect strategy, $n = 3$



# Conclusion

Handlers for algebraic effects give us a highly modular basis for effectful programming.

- **Key idea:** Separate effect signatures from their implementation.
- **Consequent:** High-degree of modularity.

# References I



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