

6 Denotational Semantics (MPF)

- (a) For monotone functions $f, f' : P \rightarrow Q$ between posets (P, \sqsubseteq_P) and (Q, \sqsubseteq_Q) , let $f \sqsubseteq f' \stackrel{\text{def}}{\iff} \forall x \in P. f(x) \sqsubseteq_Q f'(x)$.
- (i) Prove that the binary relation \sqsubseteq is a partial order. [3 marks]
- (ii) For monotone functions between posets $p : P' \rightarrow P$, $f, f' : P \rightarrow Q$, and $q : Q \rightarrow Q'$, prove that $f \sqsubseteq f' \implies q \circ f \circ p \sqsubseteq q \circ f' \circ p$. [1 mark]
- (b) An *adjoint pair* $(f : P \rightarrow Q, g : Q \rightarrow P)$ is a pair of monotone functions between posets such that $\text{id}_P \sqsubseteq g \circ f$ and $f \circ g \sqsubseteq \text{id}_Q$.
- (i) Let $f_1, f_2 : P \rightarrow Q$ and $g_1, g_2 : Q \rightarrow P$ be monotone functions between posets such that (f_1, g_1) and (f_2, g_2) are adjoint pairs. Prove that:
- (A) $f_1 \sqsubseteq f_2 \iff g_2 \sqsubseteq g_1$ [4 marks]
- (B) $f_1 = f_2 \iff g_1 = g_2$ [2 marks]
- (ii) Let $(f : P \rightarrow Q, g : Q \rightarrow P)$ be an adjoint pair where the posets P and Q have least elements. Prove that the monotone function f is strict. [2 marks]
- (c) (i) Define the notion of lub (least upper bound) of a countable increasing chain in a poset. [2 marks]
- (ii) Let $(f : D \rightarrow E, g : E \rightarrow D)$ be an adjoint pair where each of the posets D and E is a cpo (chain complete poset). Prove that the monotone function f is continuous. [6 marks]